

Use cases of Programmable Dataplane (P4)

ECE/CS598HPN

Radhika Mittal

Which paper(s) did you read?

- (A) BeauCoup: Network Monitoring
- (B) Elmo: Multicast
- (C) Both
- (D) Neither

Network Monitoring

- Most popular usecase of programmable dataplanes.
- Lots of recent papers!
- Key challenges:
 - Dealing with small amount of memory.
 - Ensuring high line rate (small processing capability, limited memory access)
 - Supporting a wide variety of queries.

¹[bo'ku] Adv. many, a lot.

BeauCoup:¹

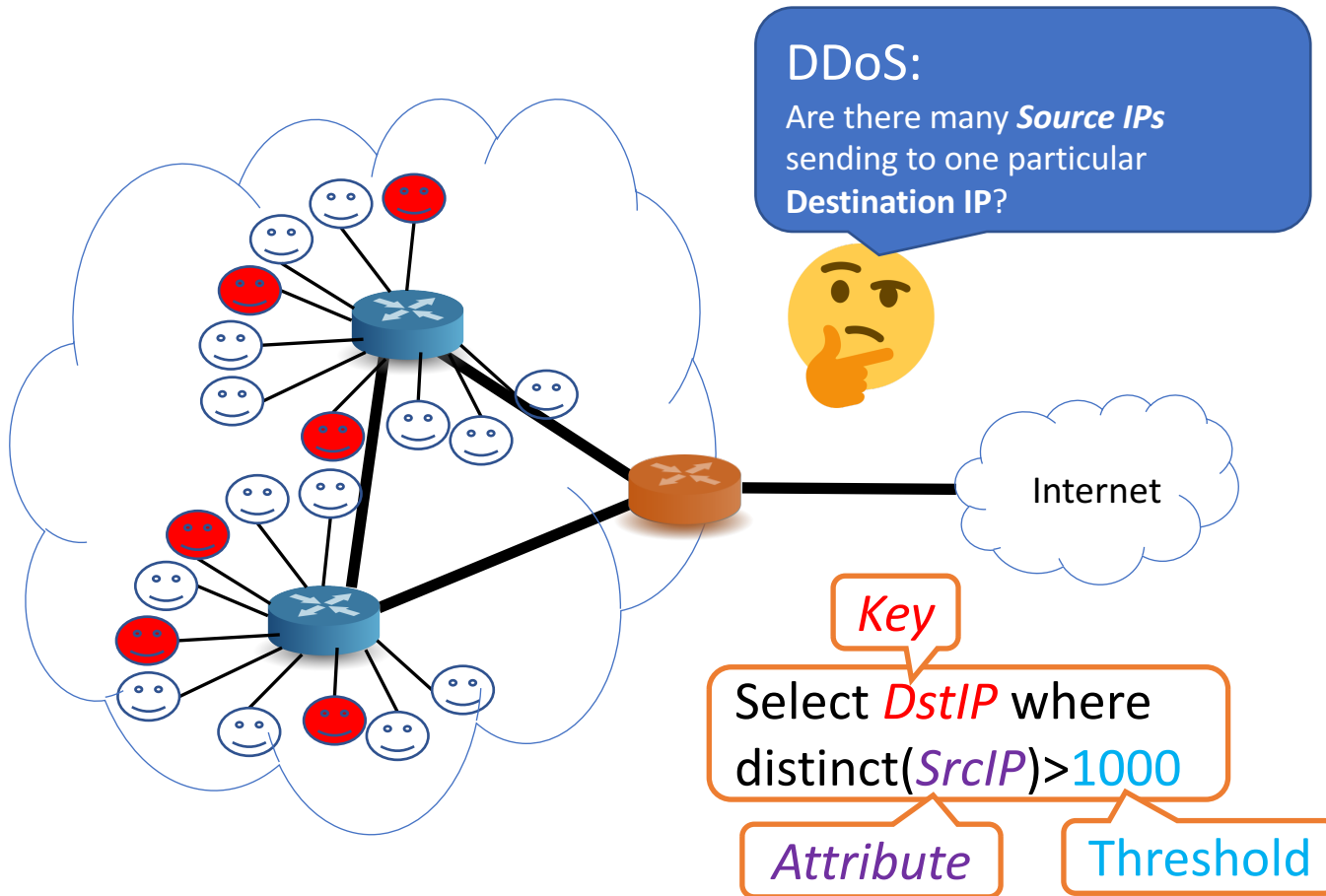
Answering *many* network traffic queries,
one memory update at a time!

Xiaoqi Chen, Shir Landau-Feibish, Mark Braverman, Jennifer Rexford

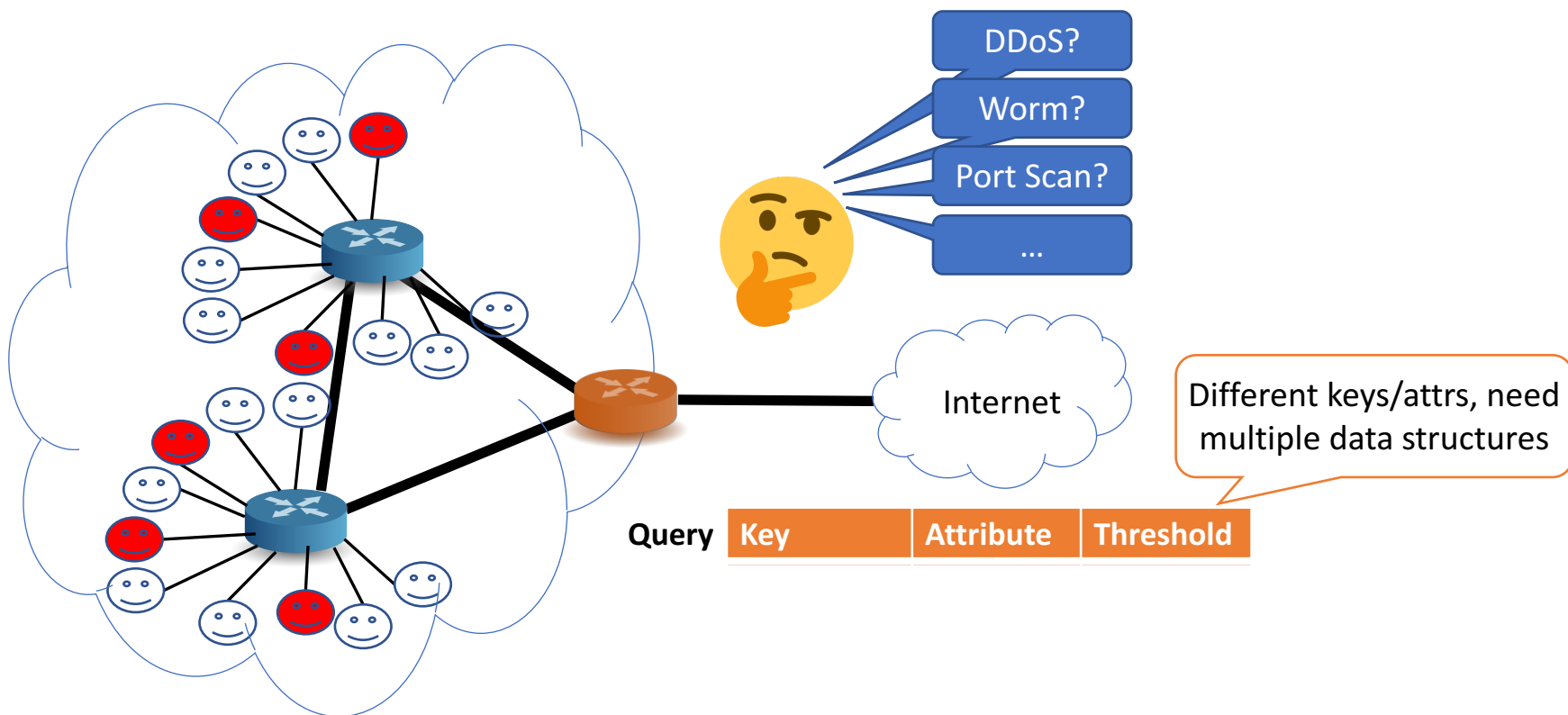


Slides from Xiaoqi Chen

Network traffic query



Many network traffic queries



Many network traffic queries

Run 42 data structures?



I have 42 queries



I can't...

Spec for today's commodity programmable switch:

- **XX Tbps** aggregated throughput
- **YY MB** data-plane memory
- Can only access **ZZ bytes** of memory **per packet**

One memory update at a time?

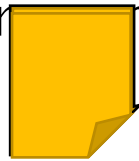
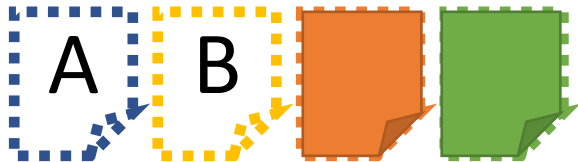
- Constant memory update per packet, regardless of the number of queries?
- Game plan:
 1. Each query uses only $\mathbf{o(l)}$ memory update per packet **on average**
 2. Combine many different queries, on average uses $\mathbf{O(l)}$
 3. Coordinate, **at most $O(l)$** per packet

BeauCoup's Approach

- Challenge:
many queries, few memory updates
- Achieving $\mathcal{O}(1)$ memory access:
coupon collectors

The coupon collector problem

- 4 different coupons, collect all of them
- Random draws
- How many total draws are required?




Naïve Approach

Query: Select *DstIP* where $\text{distinct}(\text{SrcIP}) > 130$

- Map each SrcIP to a coupon
 - How many total coupons?
 - How many do you need to collect?
- Issues with this approach:
 - Too much memory
 - Each packet results in a coupon collection.
 - Exceed $O(1)$ access when multiple such queries are combined.

BeauCoup coupon collector

$f(\text{SrcIP}) \rightarrow \text{Coupon}$ 

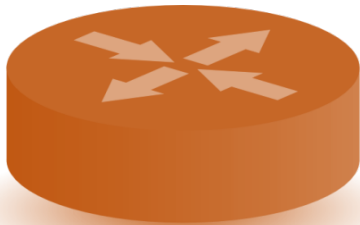
Mapping

Select *DstIP* where
 $\text{distinct}(\text{SrcIP}) > 100$

Collect different
coupons


Key: 162.249.4.107

Coupons:    



- $f(10.0.1.15) \rightarrow \text{Coupon } \text{C}$
- $f(10.0.1.33) \rightarrow \text{Coupon } \text{B}$
- $f(10.0.1.15) \rightarrow \text{Coupon } \text{C}$
- $f(10.0.1.42) \rightarrow \text{No Coupon}$

BeauCoup coupon collector

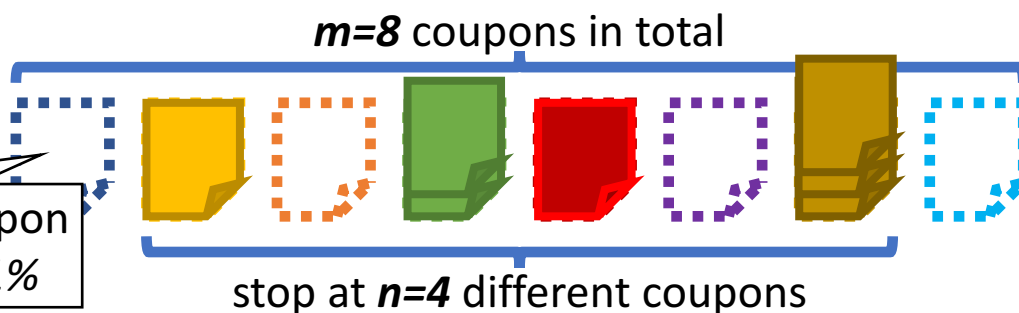
$f(\text{SrcIP}) \rightarrow \text{Coupon}$ 

- Generalization: (m, p, n) -coupon collector
- $m * p < 1$, most packets collect no coupon

Example:

$(m=8, p=1\%, n=4)$

Given a new SrcIP, each coupon is drawn with probability 1%

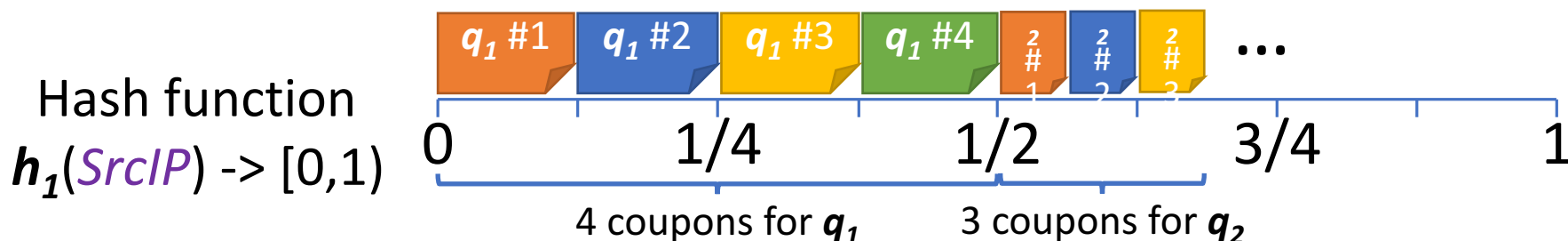


Stacking queries: same attribute


$q_1: f(\text{SrcIP}) \rightarrow \text{Coupon } \boxed{?}$
 $m_1=4, p_1=1/8$


$q_2: f(\text{SrcIP}) \rightarrow \text{Coupon } \boxed{?}$
 $m_2=3, p_2=1/16$

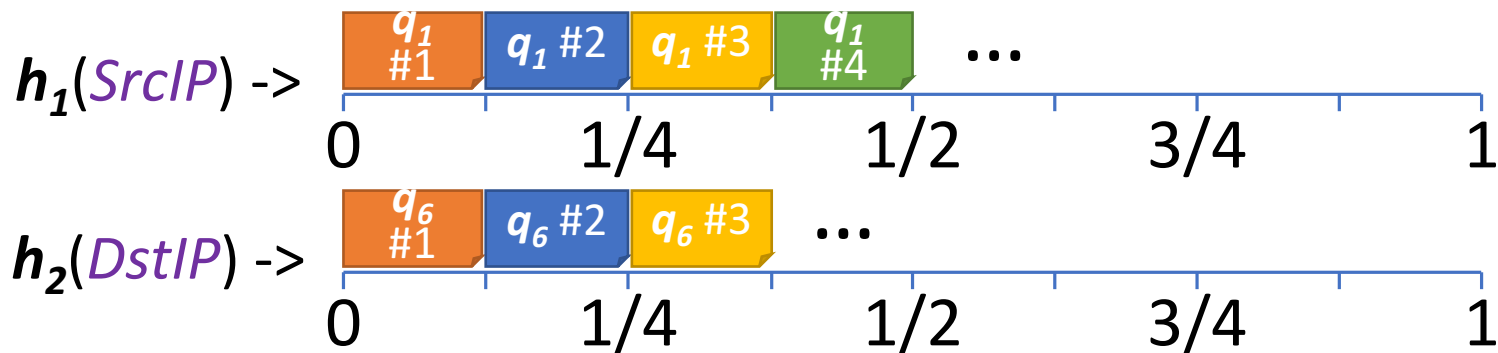
⋮



One hash function for each attribute

$q_1: f(\text{SrcIP}) \rightarrow \text{Coupon}$ 
 $m_1=4, p_1=1/8$

$q_6: g(\text{DstIP}) \rightarrow \text{Coupon}$ 
 $m_6=3, p_6=1/8$



Randomly break ties if a coupon needs to be collected for two different attributes

System design

- Query compiler: finds coupon collector configurations
 - Stops near query thresholds, minimize error
 - Hardware limits (e.g., memory access limit)
 - Fairness across queries
- Data plane program: collect coupons into in-memory table
 - Simultaneously run **many** queries
 - At most **one** coupon per packet
 - Update queries on-the-fly

Query compiler

Query set
 $Q = \{q_1, q_2, \dots\}$

Total memory update
limit: Γ per packet

Per-query limit:
 γ_q per packet

Compiler

$\gamma_q = \Gamma / |Q|$
(fair allocation)

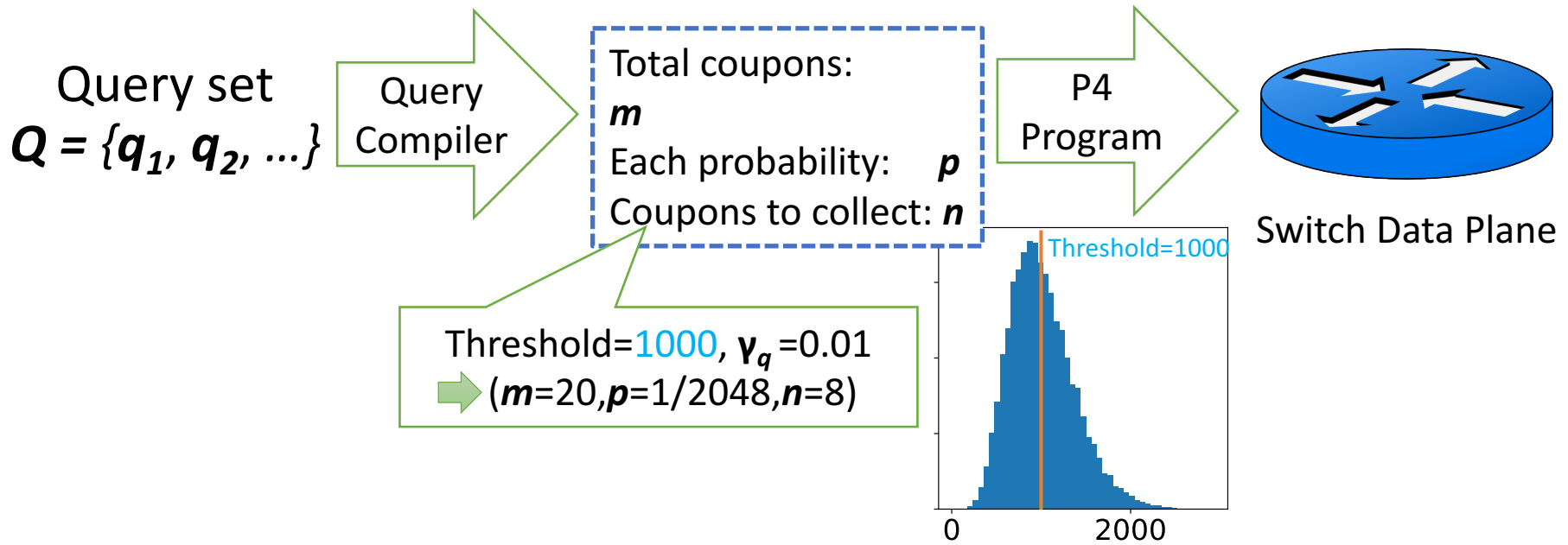
Query q_i
Key, Attribute,
Threshold

q_i 's Collector Configuration
Total coupons: m
Each probability: p
Coupons to collect: n

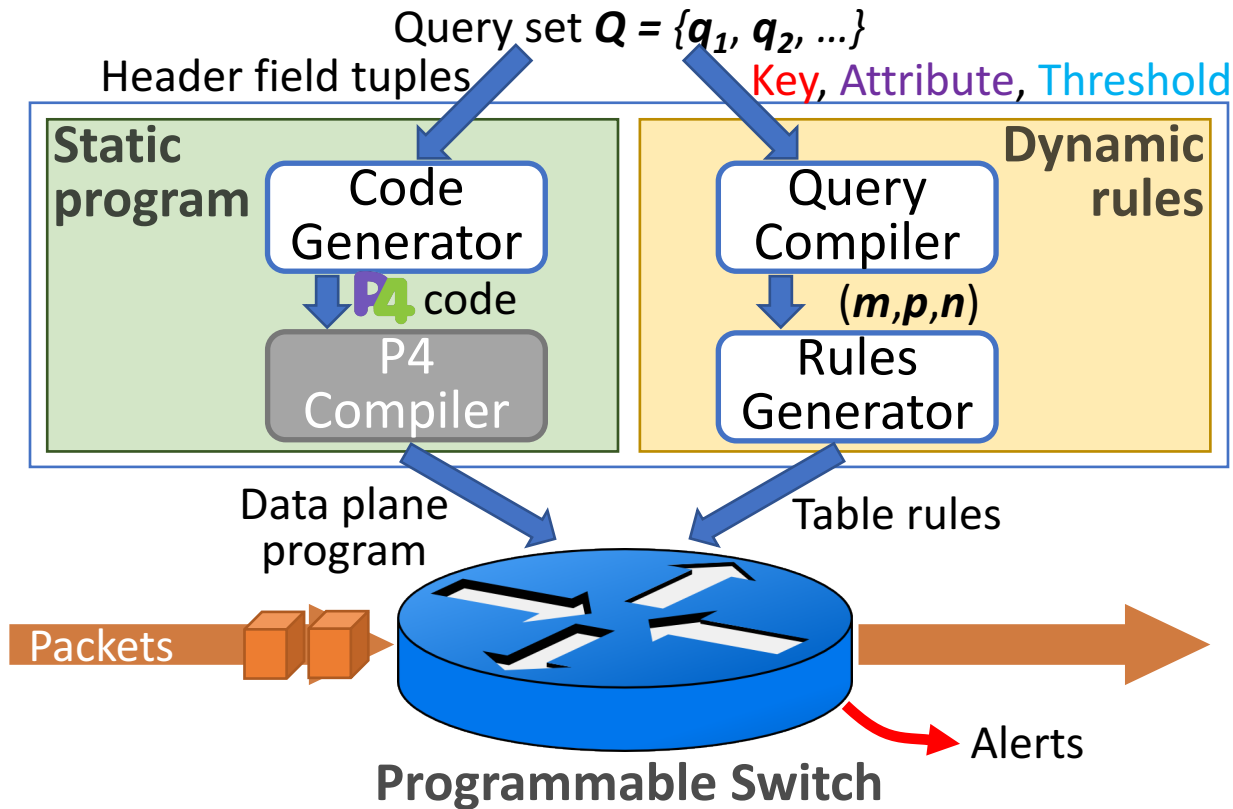
Goal:

- I. Stop near **Threshold**
- II. Update limit $m * p \leq \gamma_q$
- III. HW limit, e.g., $m \leq 32$

Query compiler



Installing queries into switches

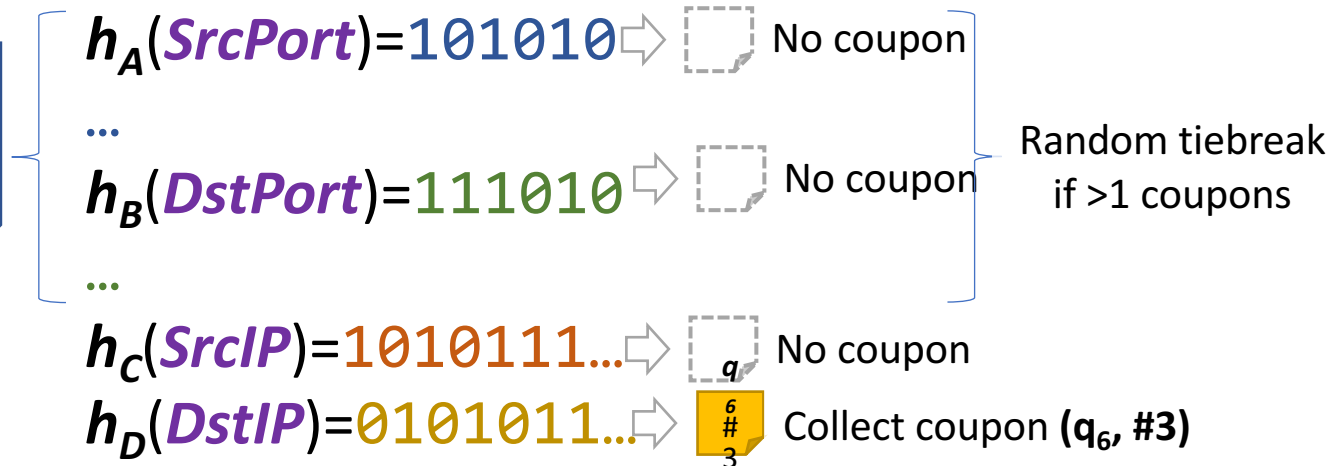


- The installed rules represent query set Q
- Update queries **on the fly**, without recompiling P4

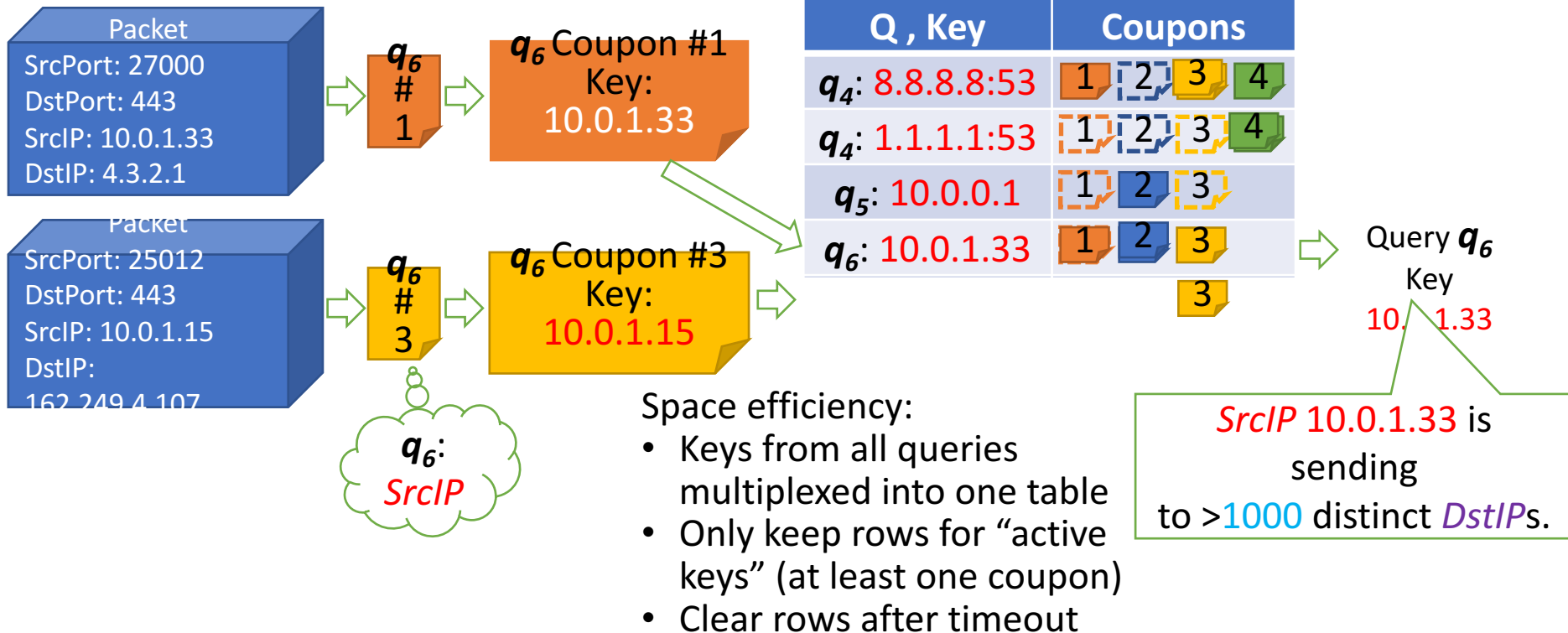
TCAM for selecting a coupon

Match $h_A(\text{SrcPort})$		Query#,Coupon#
Match $h_B(\text{DstPort})$		Query#,Coupon#
Match $h_C(\text{SrcIP})$		Query#,Coupon#
Match $h_D(\text{DstIP})$		Query#,Coupon#
000*****		(6,1)
001*****		(6,2)
010*****		(6,3)
01101***		(8,1)
...		...

Packet
 SrcPort: 25012
 DstPort: 443
 SrcIP: 10.0.1.15
 DstIP:
 162.249.4.107



Coupon collector table in SRAM



Evaluation highlights

- How efficient is BeauCoup?

Uses **4x~10x fewer memory access** than the state-of-the-art to achieve the same accuracy.

- How much hardware resource?

On the Barefoot Tofino programmable switch, BeauCoup occupies **<50% of each resource**

BeauCoup:

Answering *many* network traffic queries,
one memory update at a time!

- **Scalable:** built upon *coupon collectors*, runs many queries simultaneously
- **Versatile:** change queries on the fly, without recompiling P4 program
- **Efficient:** achieve the same accuracy using 4x-10x fewer memory accesses

Is this a good usecase of
programmable dataplanes?

What are the limitations?

Elmo: Source Routed Multicast for Public Clouds

Muhammad Shahbaz

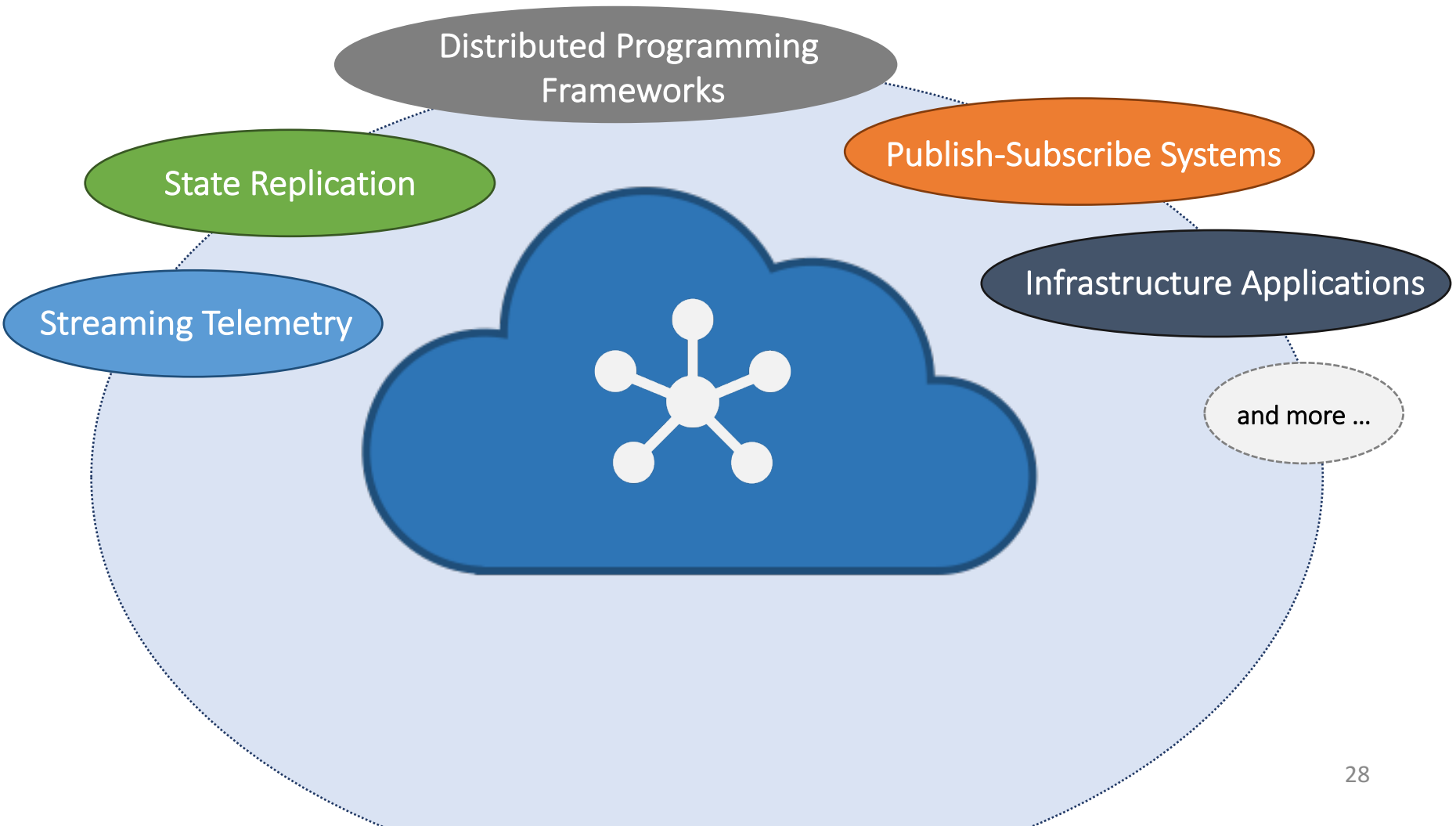
Lalith Suresh, Jennifer Rexford, Nick Feamster,
Ori Rottenstreich, and Mukesh Hira



I-to-Many Communication in Cloud



I-to-Many Communication in Cloud

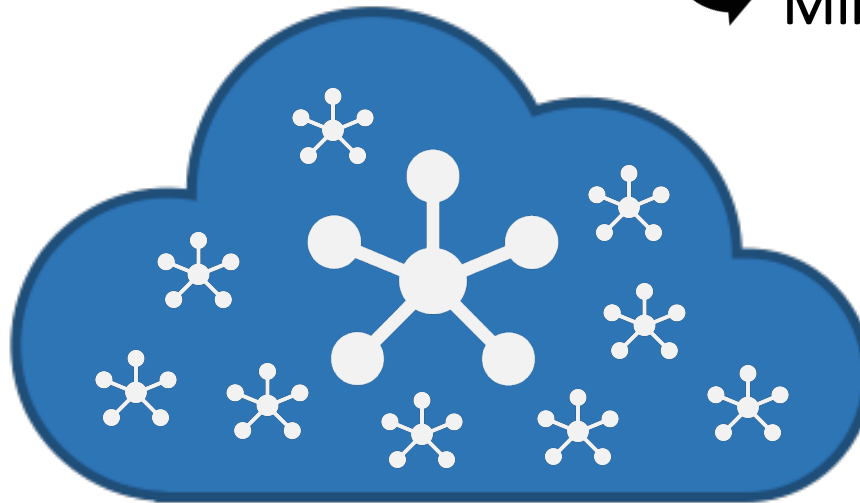


I-to-Many Communication in Cloud

10,000s of tenants

➔ 100s of workloads

➔ Millions of groups



amazon

Google

Microsoft

I-to-Many Communication in Cloud

10,000s of tenants

➔ 100s of workloads

➔ Millions of groups



Multicast

amazon

Google

Microsoft

I-to-Many Communication in Cloud

10,000s of tenants

➔ 100s of workloads

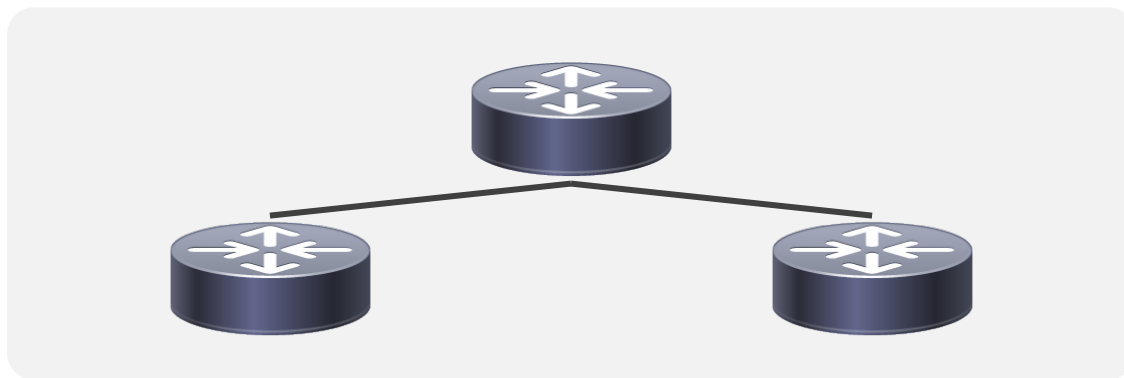
➔ Millions of groups



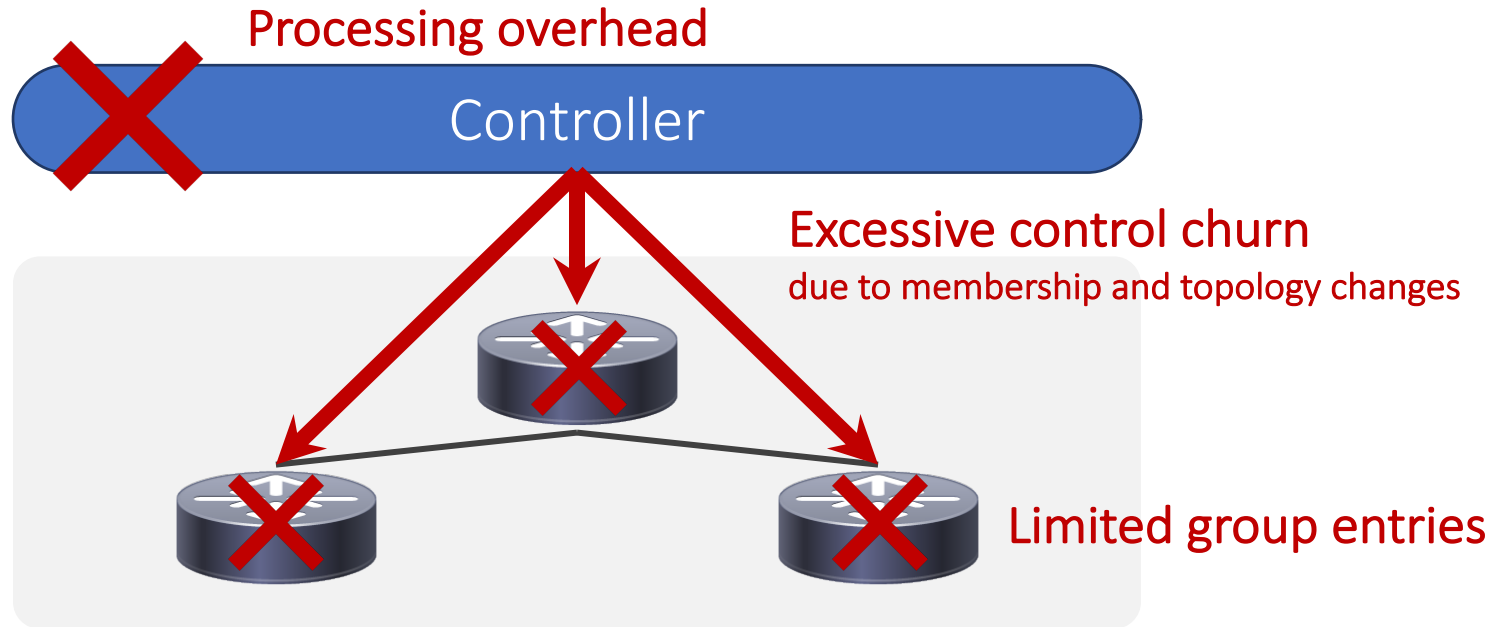
amazon Google Microsoft

Limitations of Native Multicast

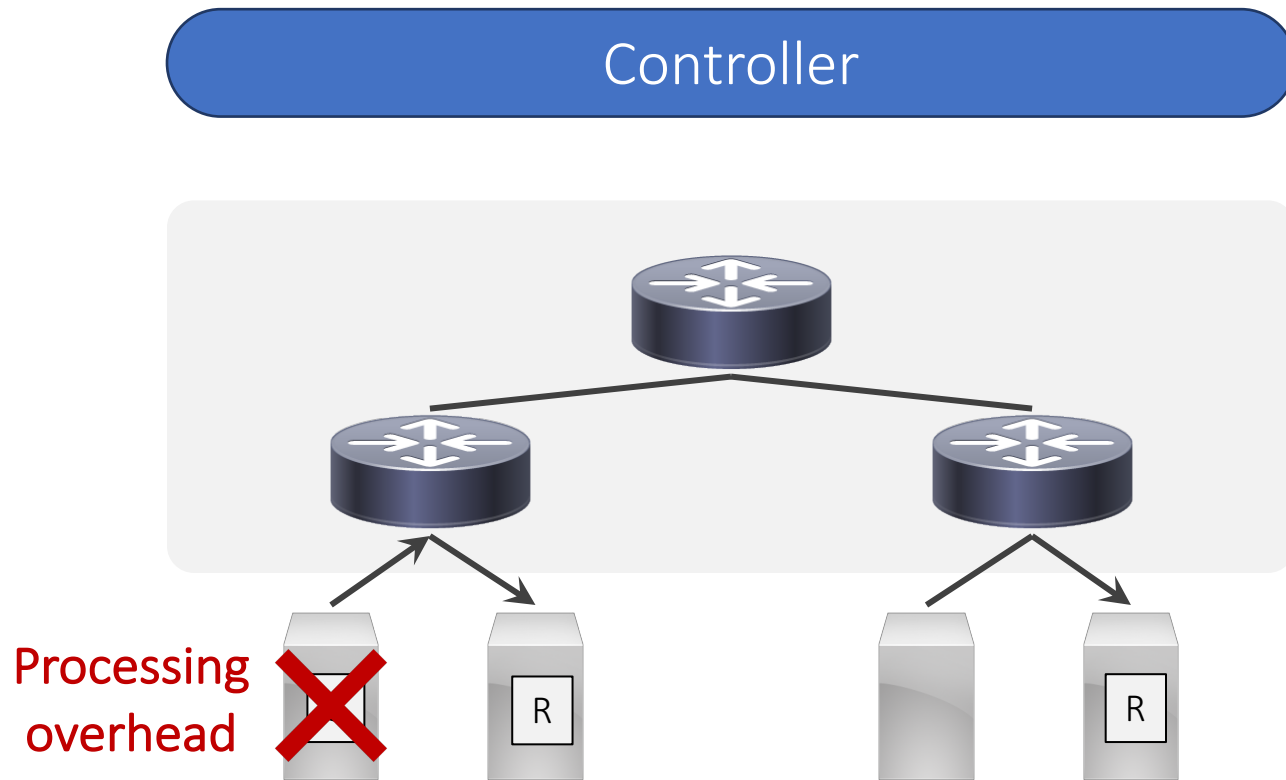
Controller



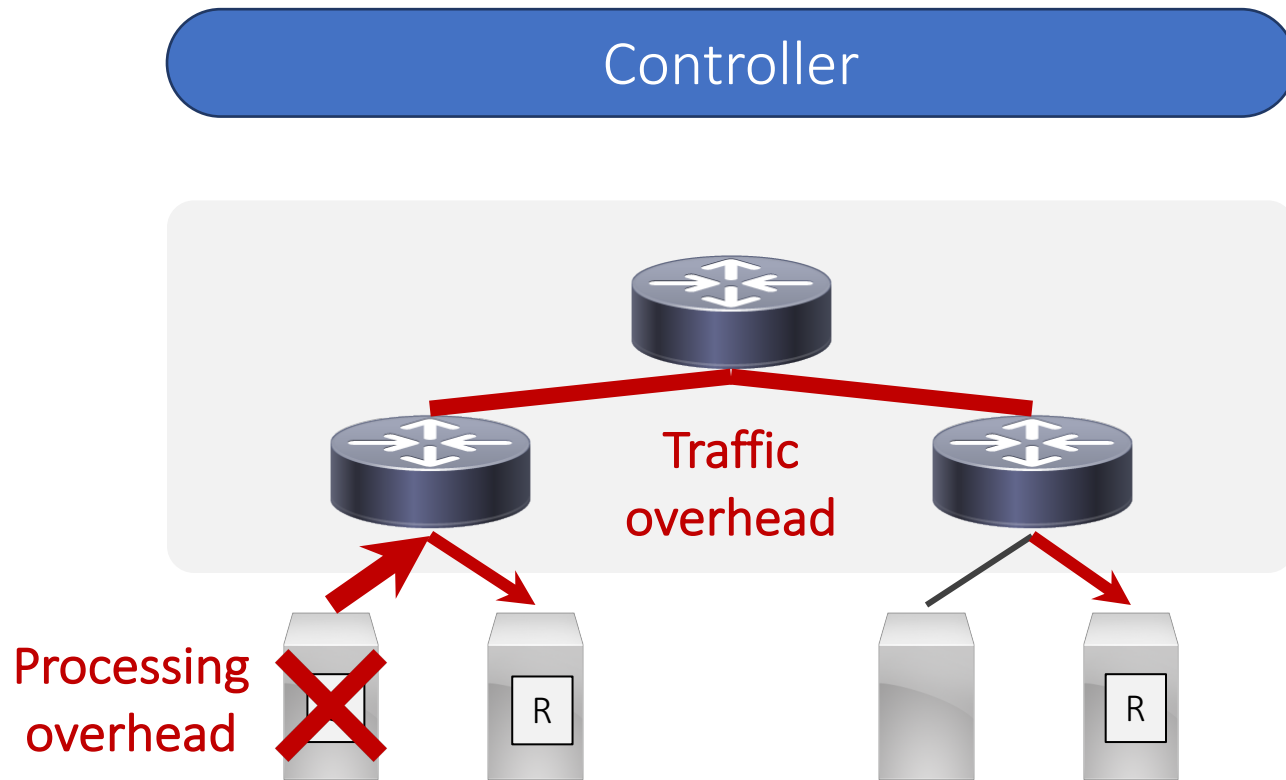
Limitations of Native Multicast



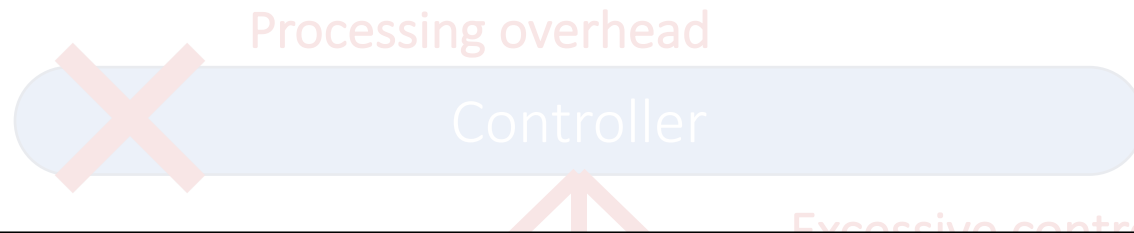
Restricted to Unicast-based Alternatives



Restricted to Unicast-based Alternatives



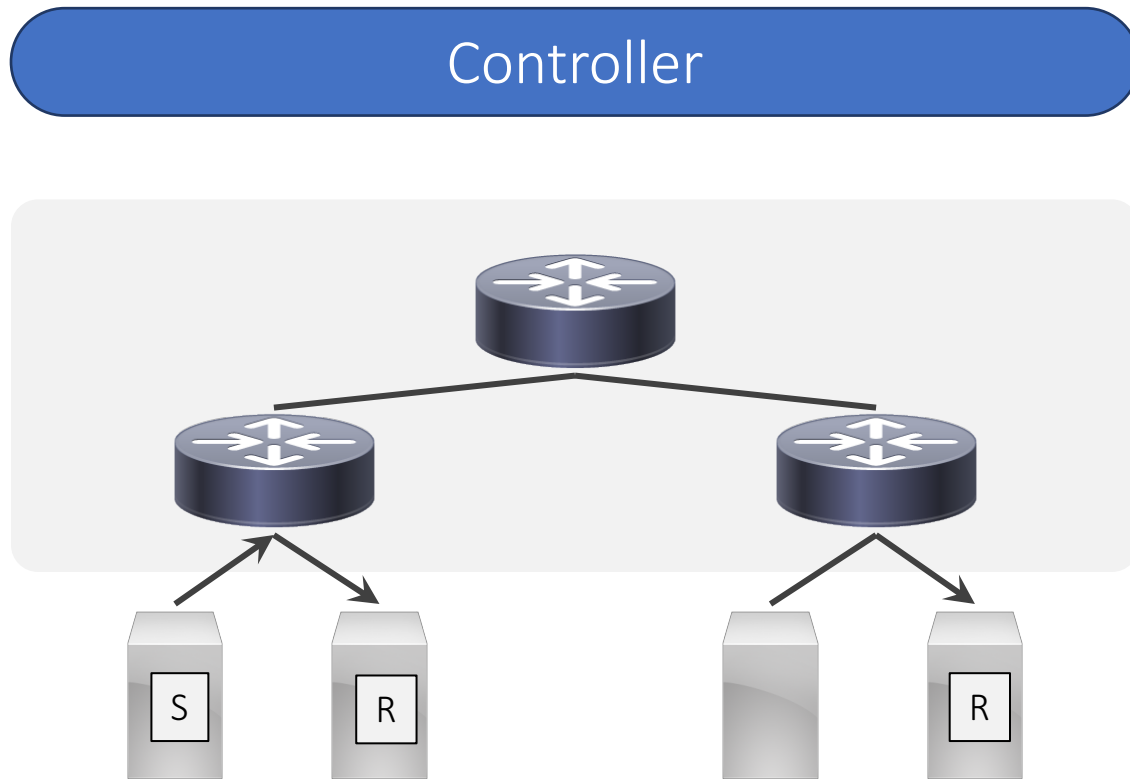
I-to-Many Communication in the Cloud



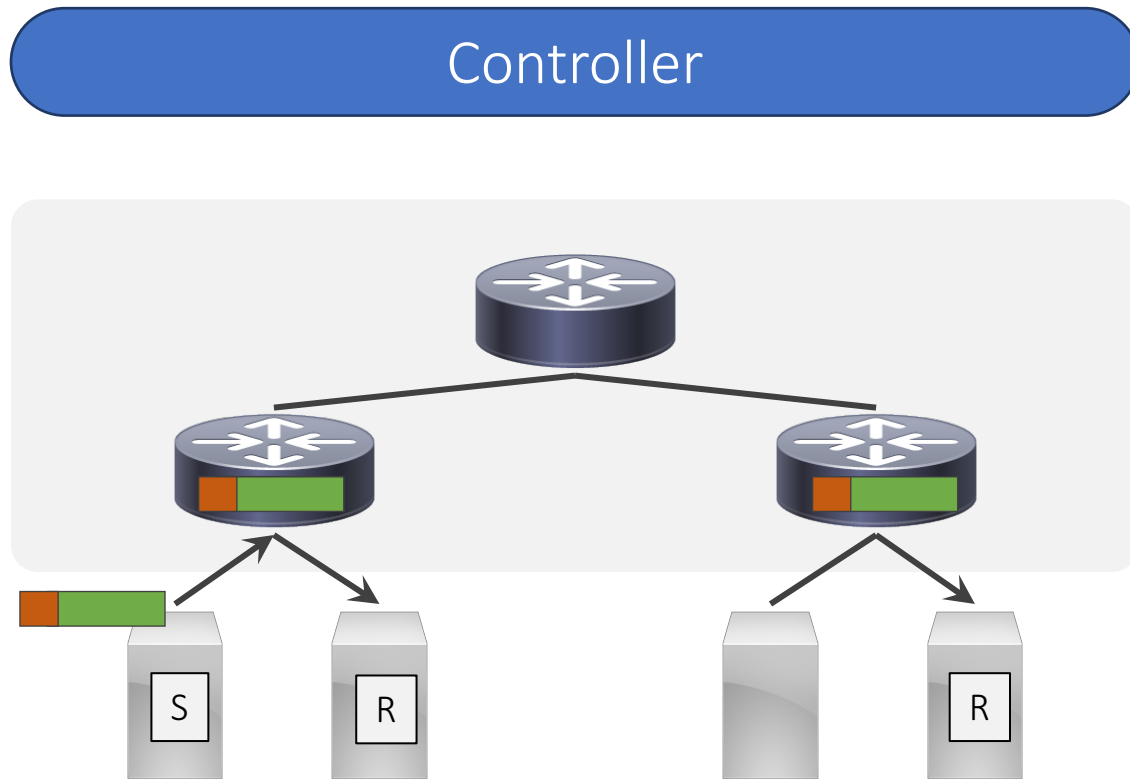
Need a scheme that scales to millions of groups without excessive control, end-host CPU, and traffic overheads!



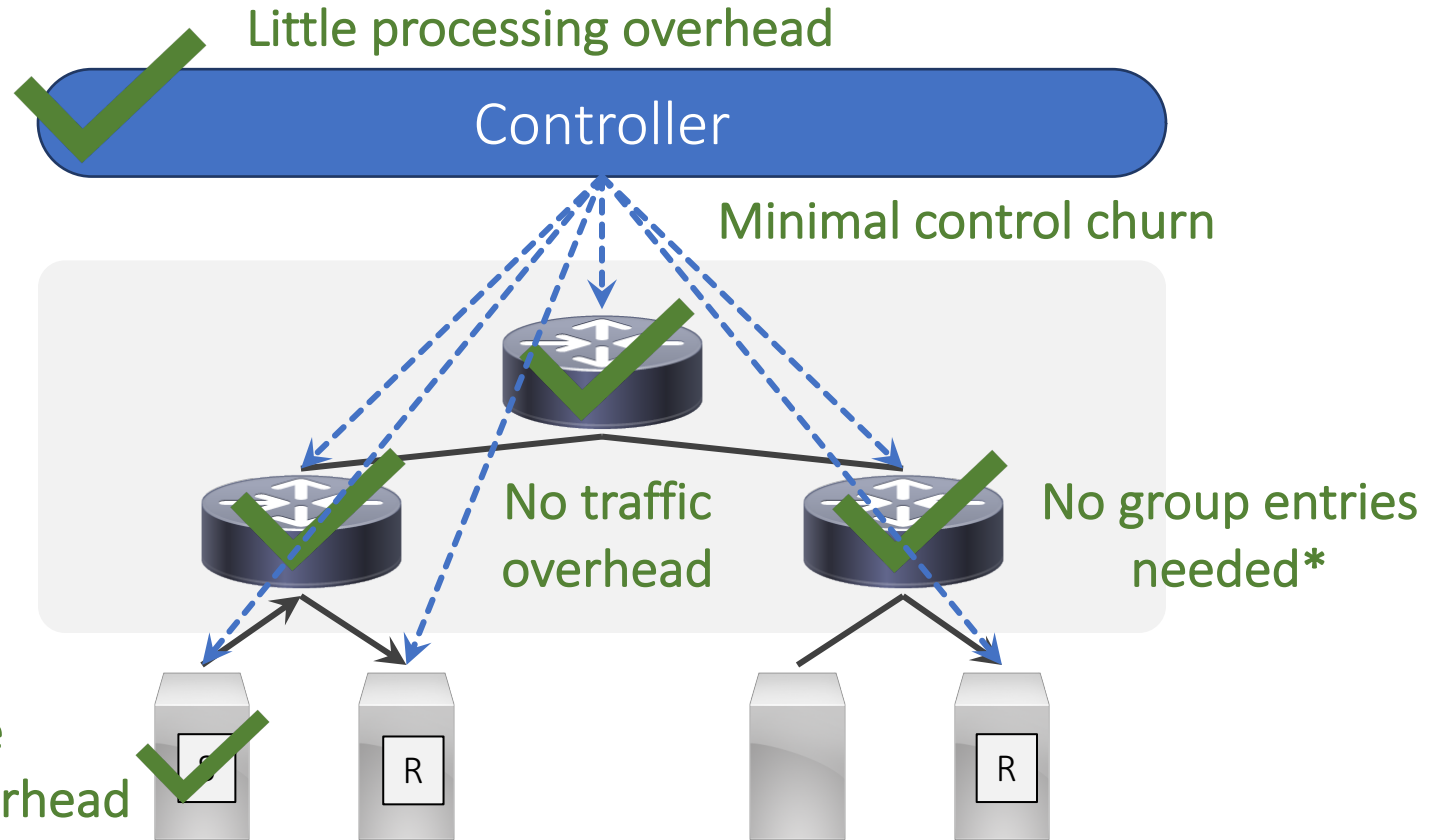
Proposal: Source Routed Multicast



Proposal: Source Routed Multicast



Proposal: Source Routed Multicast



A Naive Source Routed Multicast

A multicast group encoded as a list of (Switch, Ports) pairs

Switch 1: [Ports]
Switch 2: [.. .. .]
Switch 3: [.. .. .]
Switch 4: [.. .. .x ..]
Switch 5: [.x]

For a data center with:

- 1000 switches
- 48 ports per switch



O(30) bytes per switch



O(30,000) bytes header

for a group spanning 1000 switches

20x the packet size!

Enabling Source Routed Multicast in Public Clouds

Key attributes:

- Efficiently encode multicast forwarding policy inside packets
- Process this encoding at hardware speed in the switches
- Execute tenants' applications without modification

Encoding a Multicast Policy in Elmo

A multicast group encoded as a list of (Switch, Ports) pairs

```
Switch 1: [Ports ]  
Switch 2: [.. .. ..]  
Switch 3: [.. .. ..]  
Switch 4: [.. .. .x ..]  
Switch 5: [.x .. .. ..]
```

Encoding a Multicast Policy in Elmo

A multicast group encoded as a list of (Switch, Ports) pairs

```
Switch 1: [Bitmap] ←  
Switch 2: [.. .. ..]  
Switch 3: [.. .. ..]  
Switch 4: [.. .. .x ..]  
Switch 5: [.x .. .. ..]
```

1 Encode switch ports as a **bitmap**

Bitmap is the internal **data structure** that switches use for **replicating packets**

Encoding a Multicast Policy in Elmo

A multicast group encoded as a list of (Switch, Ports) pairs

```
Switch 1: [Bitmap]
Switch 2: [.. .. ..]
Switch 3: [.. .. ..]
Switch 4: [.. .. .x ..]
Switch 5: [.x .. .. ..]
```

2 Group switches into layers

Encoding a Multicast Policy in Elmo

A multicast group encoded as a list of (Switch, Ports) pairs

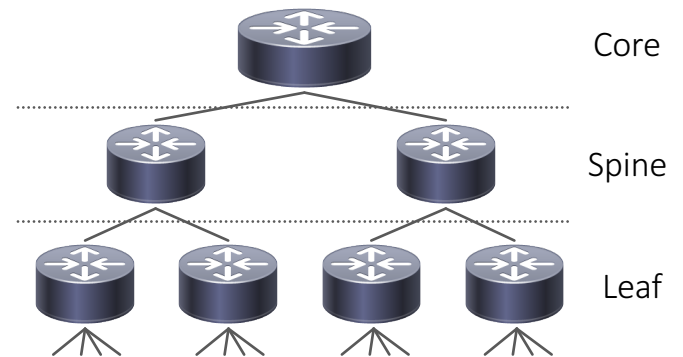
Switch 1: [Bitmap]
Switch 2: [.. .. .]
Switch 3: [.. .. .]
Switch 4: [.. .. .x ..]
Switch 5: [.x]

Core

Spine

Leaf

2 Group switches into layers



More precisely: *upstream leaf, upstream spine, core, downstream spine, downstream leaf*

Encoding a Multicast Policy in Elmo

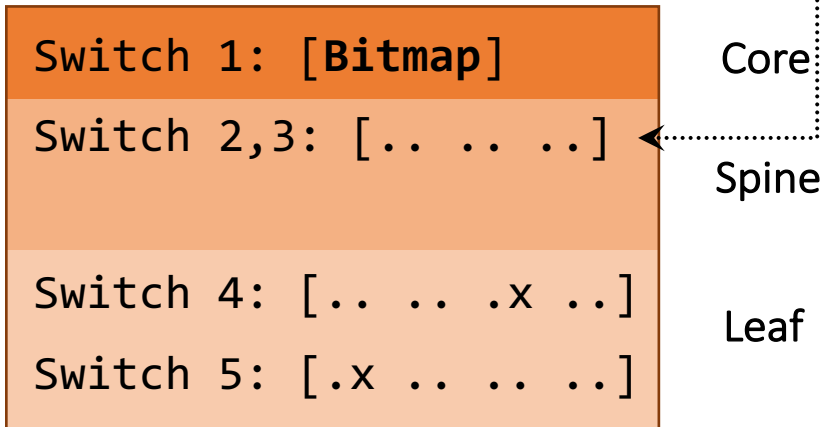
A multicast group encoded as a list of (Switch, Ports) pairs

```
Switch 1: [Bitmap]
Switch 2: [.. .. ..]
Switch 3: [.. .. ..]
Switch 4: [.. .. .x ..]
Switch 5: [.x .. .. ..]
```

3 Switches within a layer with same ports share a bitmap

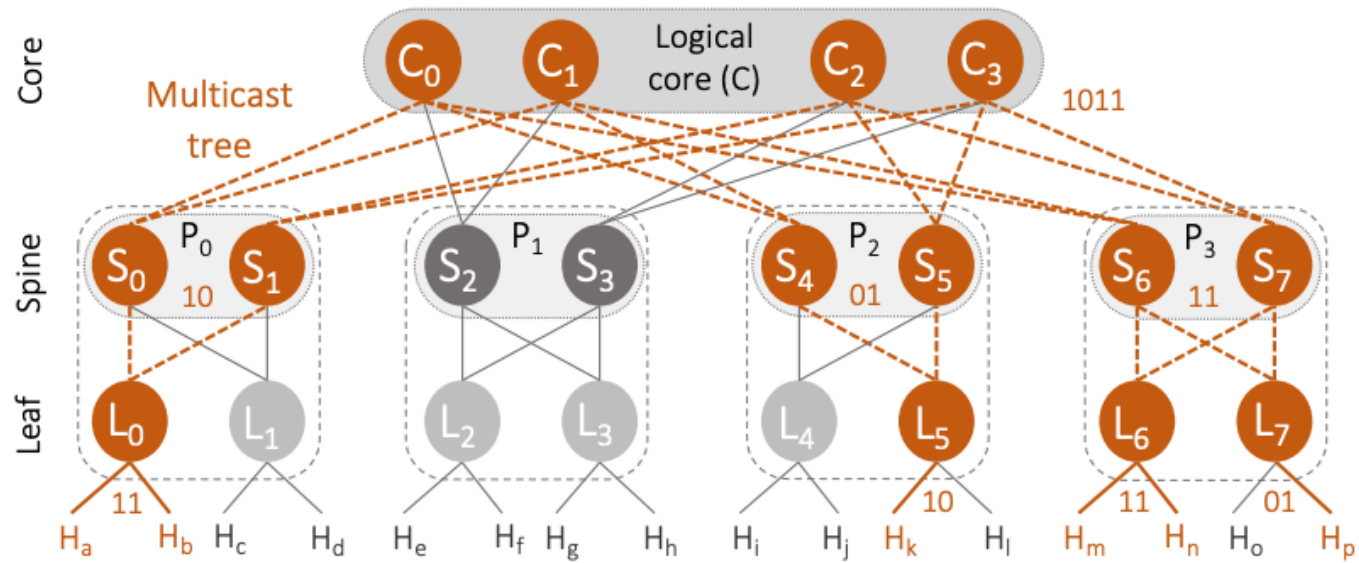
Encoding a Multicast Policy in Elmo

A multicast group encoded as a list of (Switch, Ports) pairs



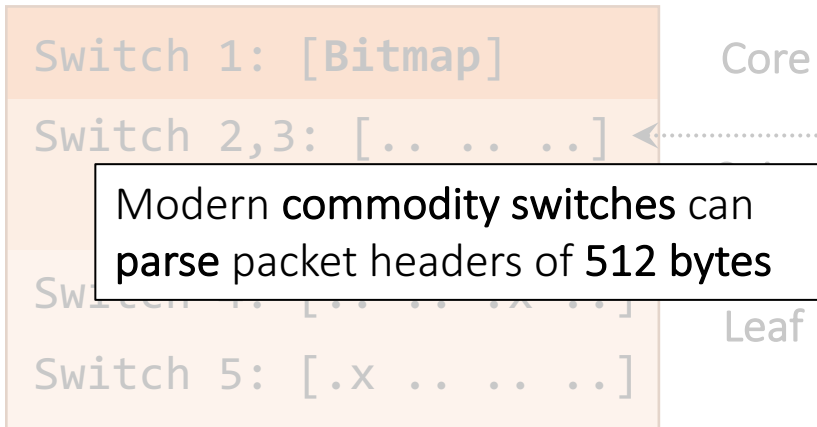
3 Switches within a layer with same ports share a bitmap

Encoding a Multicast Policy in Elmo

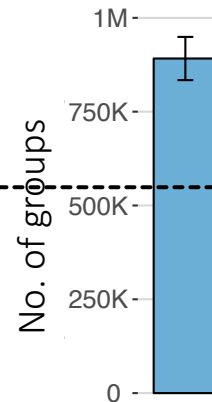


Encoding a Multicast Policy in Elmo

A multicast group encoded as a list of (Switch, Ports) pairs



3 Switches within a layer with same ports share a bitmap



For a data center with:
- 628 switches
- 325 bytes header space

Supports **890,000** groups!

Encoding a Multicast Policy in Elmo

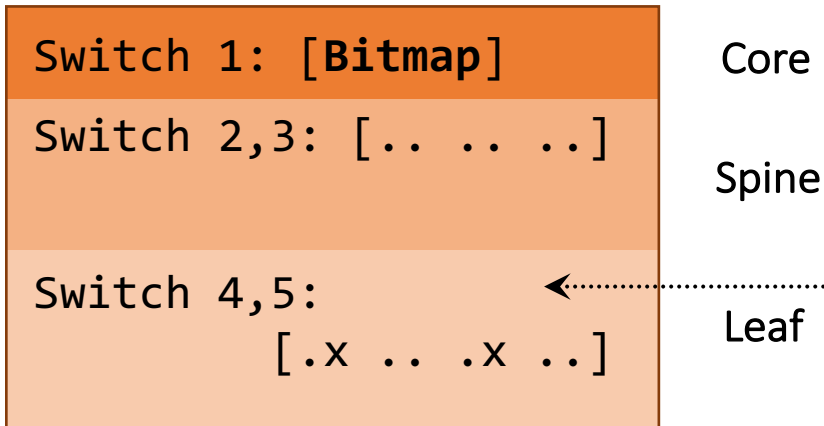
A multicast group encoded as a list of (Switch, Ports) pairs

Switch 1: [Bitmap]	Core
Switch 2,3: [..]	Spine
Switch 4: [.. .. .x ..]	Leaf
Switch 5: [.x]	

4 Switches within a layer with **N** different ports share a bitmap

Encoding a Multicast Policy in Elmo

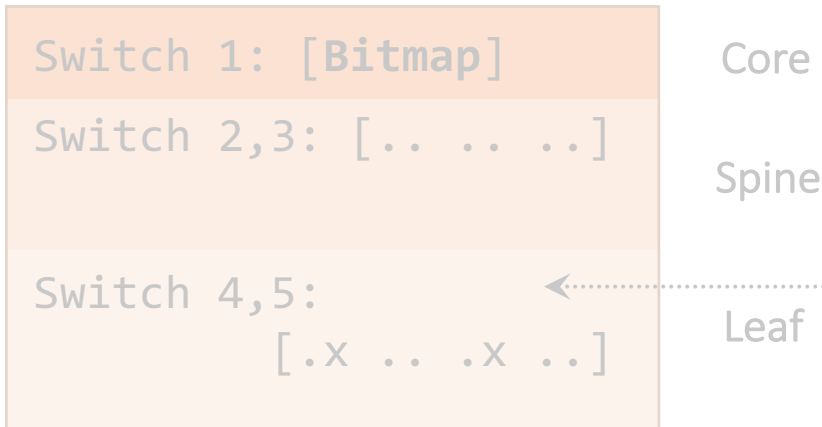
A multicast group encoded as a list of (Switch, Ports) pairs



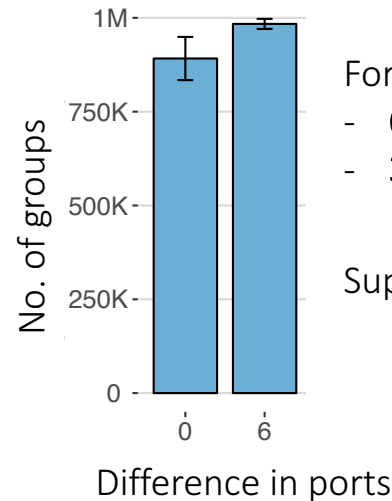
4 Switches within a layer with N different ports share a bitmap

Encoding a Multicast Policy in Elmo

A multicast group encoded as a list of (Switch, Ports) pairs



4 Switches within a layer with N different ports share a bitmap



For a data center with:
- 628 switches
- 325 bytes header space
Supports **980,000** groups!

Encoding a Multicast Policy in Elmo

A multicast group encoded as a list of (Switch, Ports) pairs

Fixed Header Size

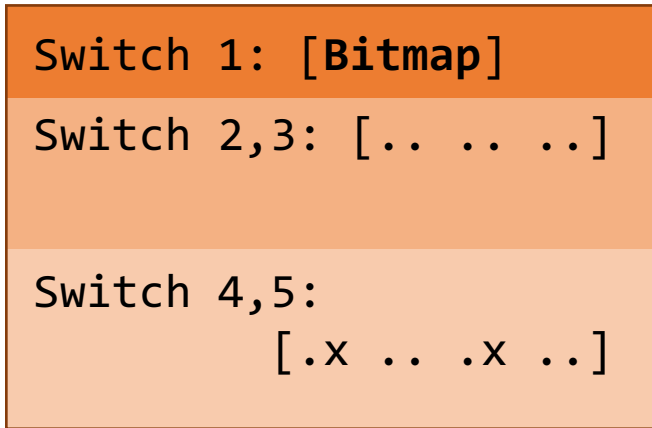
Switch 1: [Bitmap]	Core
Switch 2,3: [..]	Spine
Switch 4,5: [.x .. .x ..]	Leaf

- 4 Switches within a layer with N different ports share a bitmap

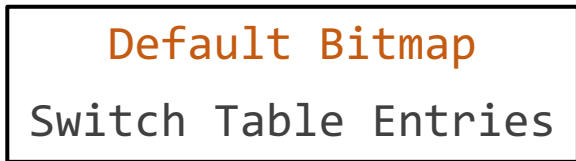
Encoding a Multicast Policy in Elmo

A multicast group encoded as a list of (Switch, Ports) pairs

Fixed Header Size



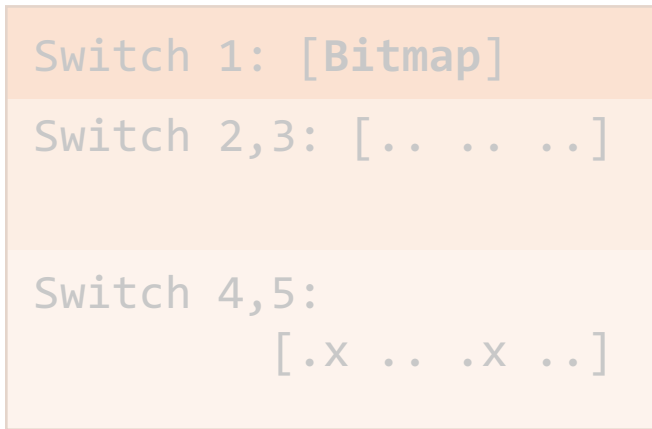
5 Use switch entries and a default bitmap for larger groups



Encoding a Multicast Policy in Elmo

A multicast group encoded as a list of (Switch, Ports) pairs

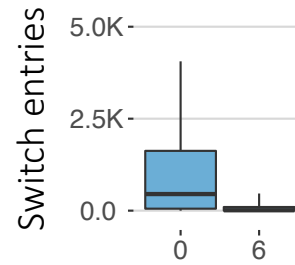
Fixed Header Size



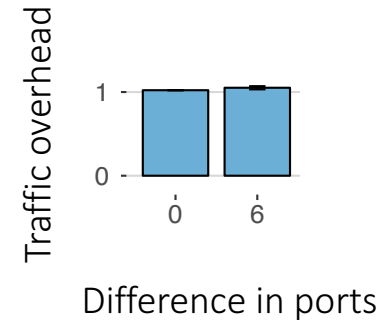
Core
Spine
Leaf

Default Bitmap
Switch Table Entries

5 Use switch entries and a default bitmap for larger groups



For a data center with:
- 628 switches
- 325 bytes header space



Encoding a Multicast Policy in Elmo

A multicast group encoded as a list of (Switch, Ports) pairs

Fixed Header Size

Switch 1: [Bitmap]
Switch 2,3: [... ..]
Switch 4,5: [.x .. .x ..]

Core

Spine

Leaf

Default Bitmap
Switch Table Entries

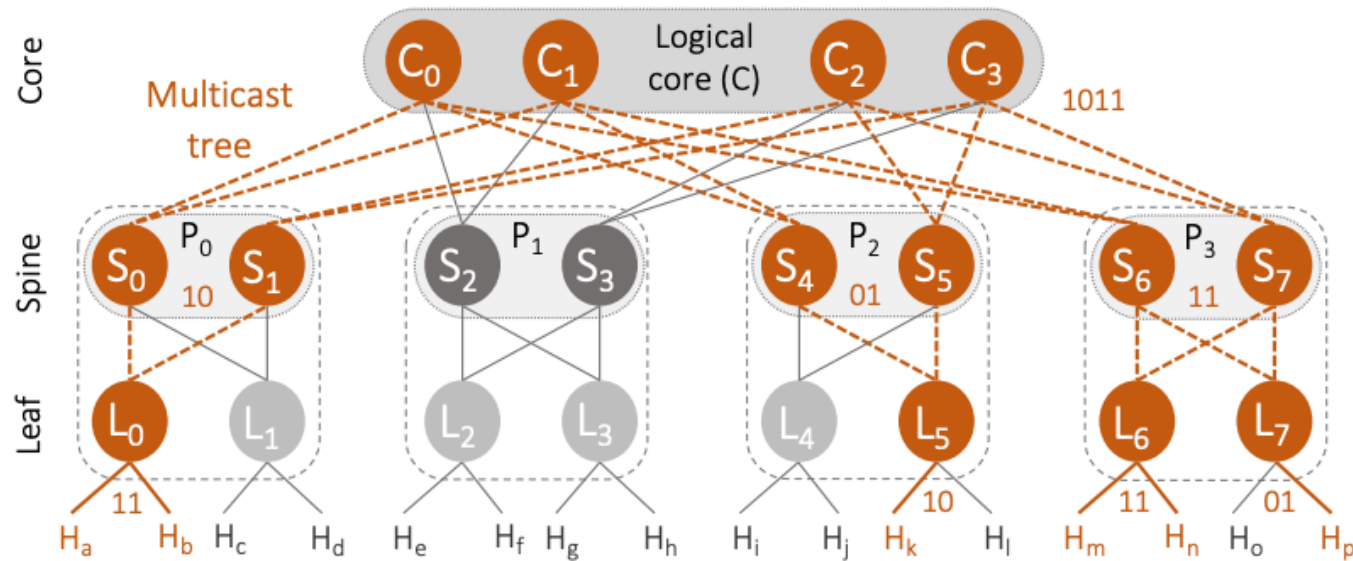
- 1 Encode switch ports as a **bitmap**
- 2 Group switches into **layers**
Switches within a layer with:
 - 3 - same ports share a **bitmap**
 - 4 - N different ports share a **bitmap**
- 5 Use **switch entries** and a **default bitmap** for larger groups

For a data center with:

- 628 switches
- 325 bytes header space

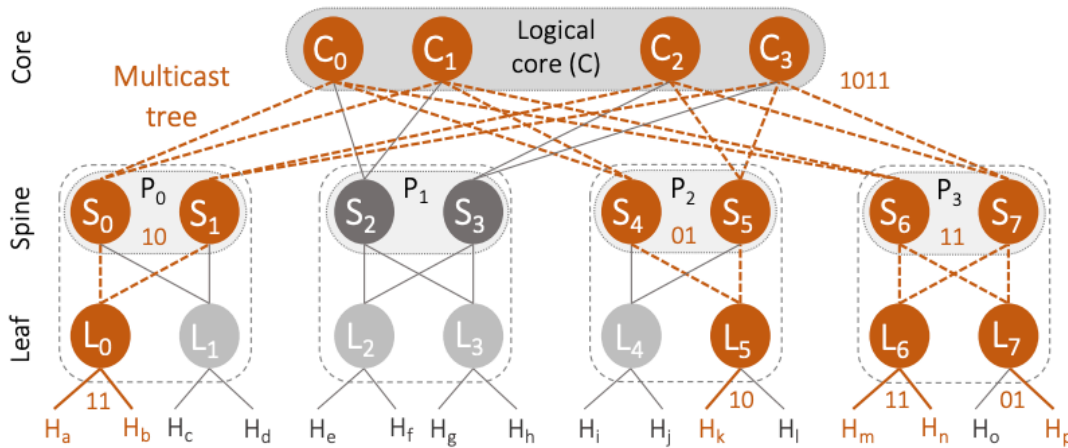
Supports **a Million** groups!

Encoding a Multicast Policy in Elmo



Sender H_a		type	Sender-specific leaf, spine, and core p -rules			Common downstream spine and leaf p -rules				
Outer header(s)	VXLAN	u	u -leaf	u -spine	d -core	d -spine		d -leaf		Packet body
			01 M	00 M	0011	10:[P ₀]	11:[P ₃]	11:[L ₀ ,L ₆]	01:[L ₇]	Packet body
			At L ₀ : forward to H _b and multipath to P ₀			01:[P ₂]	Default	10:[L ₅]	Default	
Sender H_k		type	Sender-specific leaf, spine, and core p -rules			Common downstream spine and leaf p -rules				
Outer header(s)	VXLAN	u	00 M	00 M	1001	d -spine		d -leaf		Packet body
			At L ₅ : multipath to P ₂			P ₀ : forward to L ₀ P ₂ : forward to L ₅ P ₃ : forward to L ₆ , L ₇		L ₀ : forward to H _a , H _b L ₅ : forward to H _k L ₆ : forward to H _m , H _n L ₇ : forward to H _p		

Encoding a Multicast Policy in Elmo



Downstream spine and leaf assignments with varying degree of redundancy (R) and *s*-rules.

	R = 0	R = 2	
	# <i>s</i> -rules = 0	# <i>s</i> -rules = 1	
	# <i>s</i> -rules = 0, 1	# <i>s</i> -rules = 0, 1	
	(<i>p</i>) 10:[P ₀] (<i>p</i>) 01:[P ₂] (<i>d</i>) 11:[P ₃]	(<i>p</i>) 10:[P ₀] (<i>p</i>) 01:[P ₂] (<i>s</i>) 11:[P ₃]	(<i>p</i>) 10:[P ₀] (<i>p</i>) 11:[P ₂ , P ₃]
	(<i>p</i>) 11:[L ₀ , L ₆] (<i>p</i>) 10:[L ₅] (<i>d</i>) 01:[L ₇]	(<i>p</i>) 11:[L ₀ , L ₆] (<i>p</i>) 10:[L ₅] (<i>s</i>) 01:[L ₇]	(<i>p</i>) 11:[L ₀ , L ₆] (<i>p</i>) 11:[L ₅ , L ₇]

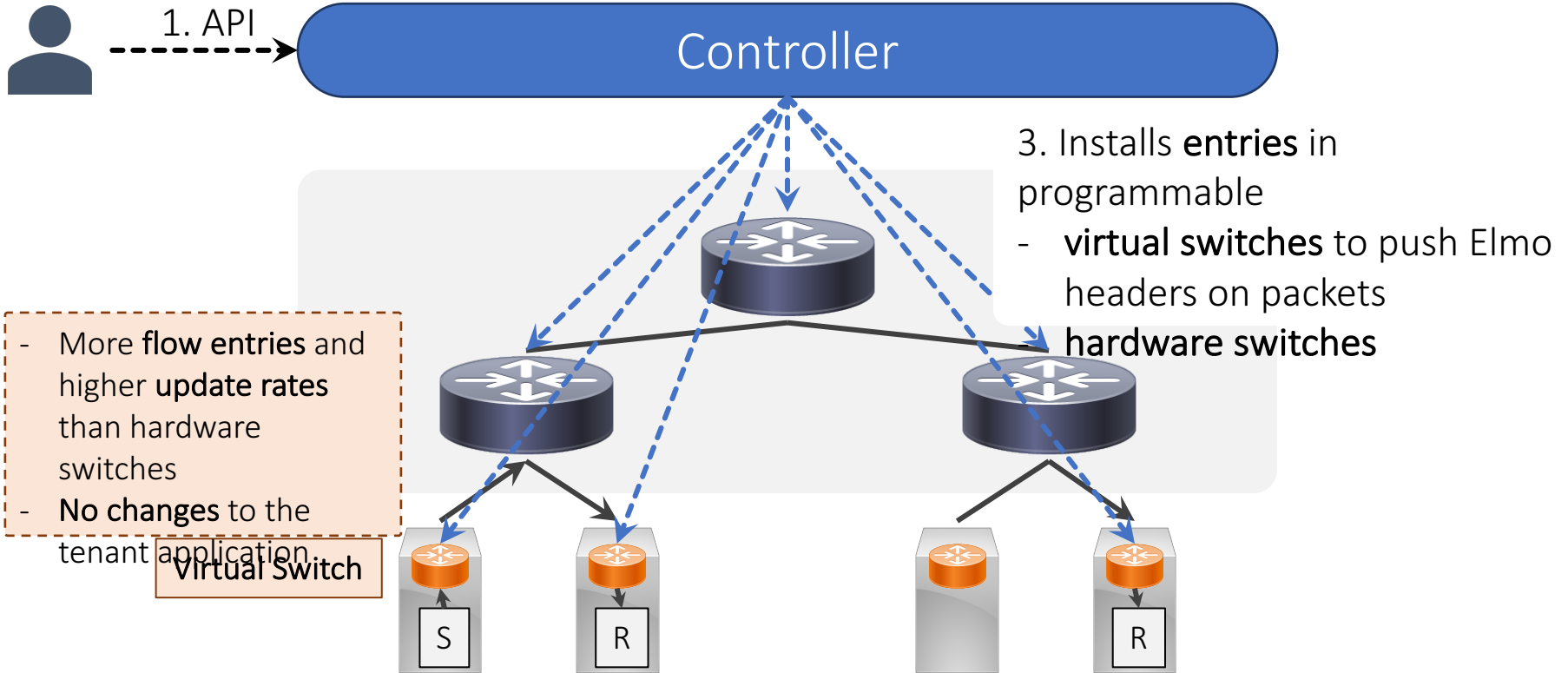
#*p*-rules = 2 (and max two switches per *p*-rule)

legends: *p*-rule (*p*), *s*-rule (*s*), and default *p*-rule (*d*)

Sender-specific leaf, spine, and core <i>p</i> -rules					Common downstream spine and leaf <i>p</i> -rules					
Sender H _a	type	<i>u</i> -leaf	<i>u</i> -spine	<i>d</i> -core	<i>d</i> -spine		<i>d</i> -leaf		Packet body	
Outer header(s)	VXLAN	<i>u</i>	01 M	00 M	10:[P ₀]	11:[P ₃]	11:[L ₀ , L ₆]	01:[L ₇]	Packet body	
			At L ₀ : forward to H _b and multipath to P ₀	P ₀ : multipath to C	C: forward to P ₂ , P ₃	01:[P ₂]	Default	10:[L ₅]		Default
Outer header(s)	VXLAN	<i>u</i>	00 M	00 M	1001					Packet body
			At L ₅ : multipath to P ₂	P ₂ : multipath to C	C: forward to P ₀ , P ₃					

Processing a Multicast Policy in Elmo

2. Computes the multicast policy

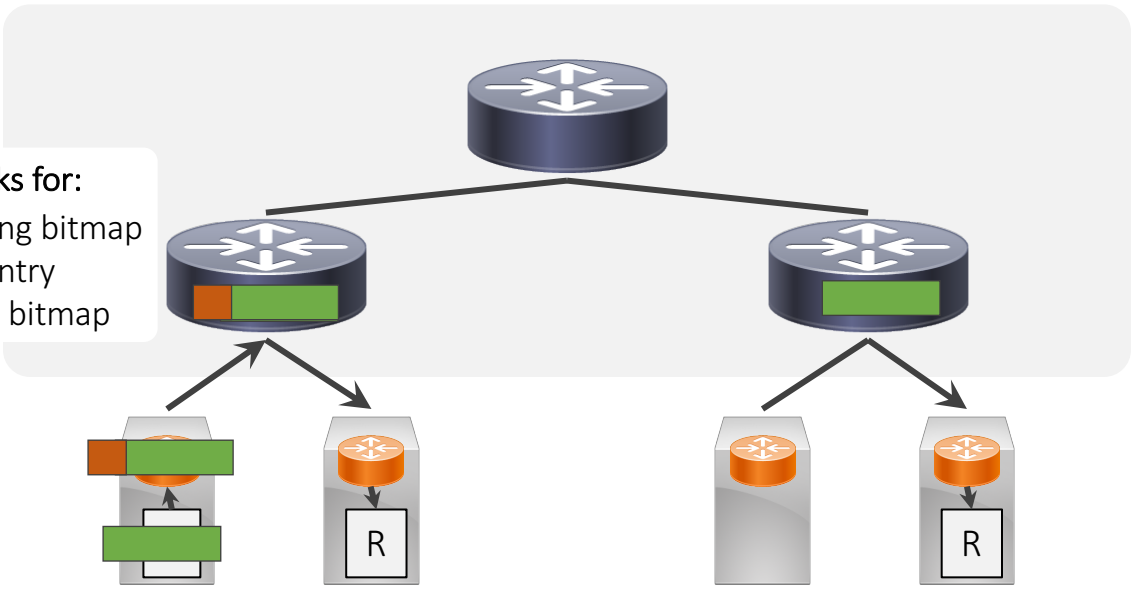


Processing a Multicast Policy in Elmo



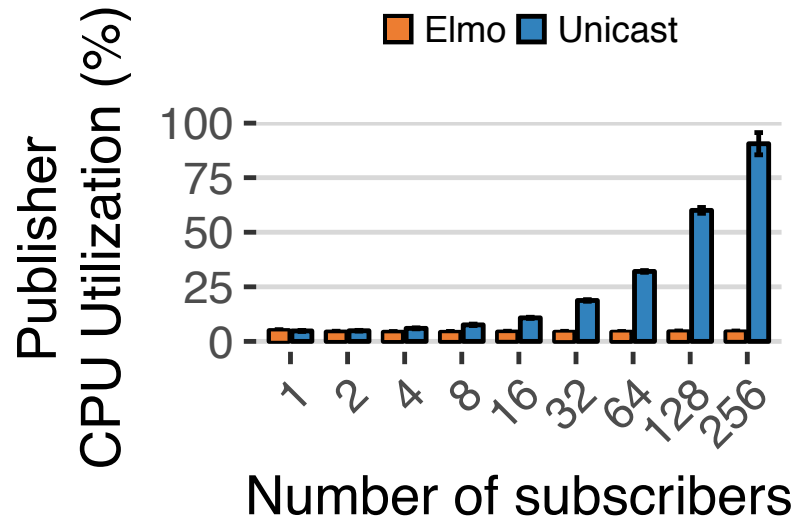
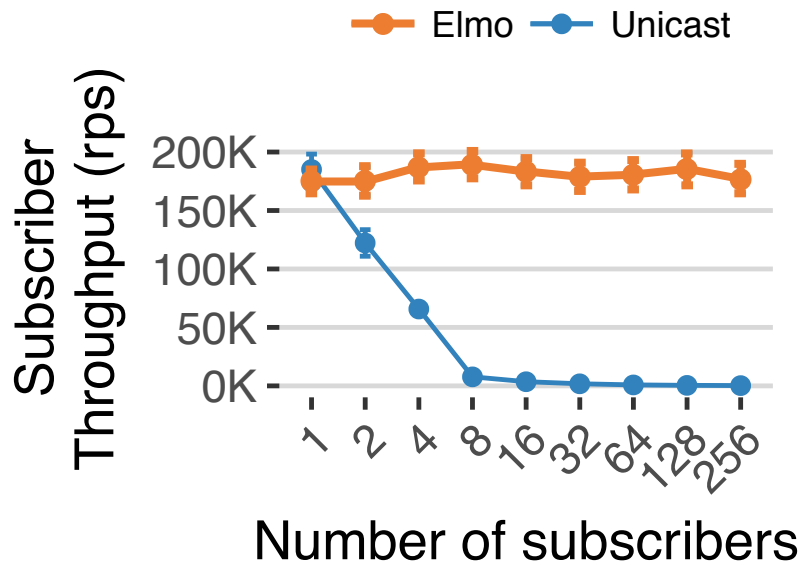
Controller

Switch looks for:
Matching bitmap
or Table entry
or Default bitmap



Implemented using P4 on a Barefoot Tofino Switch

Applications Run Without Performance Overhead



Conclusion

Elmo

Source Routed
Multicast
for Public Clouds

- Designed for multi-tenant data centers
- Compactly encodes multicast policy inside packets
- Operates at hardware speed using programmable data planes

Is this a good usecase of
programmable dataplanes?

What are the limitations?

Other networking usecases

- Load balancing:
 - HULA: Scalable Load Balancing Using Programmable Data Planes, SOSR'16
- Congestion control:
 - Evaluating the Power of Flexible Packet Processing for Network Resource Allocation, NSDI'17
 - *Support RCP and XCP on programmable switches*
 - HPCC: High Precision Congestion Control, SIGCOMM'19
 - *Obtain precise link information for congestion control*
- A new protocols for more efficient L2 switching
 - The Deforestation of L2, SIGCOMM'16
- And others...