

Discrete Random Variables

ECE 313

Probability with Engineering Applications

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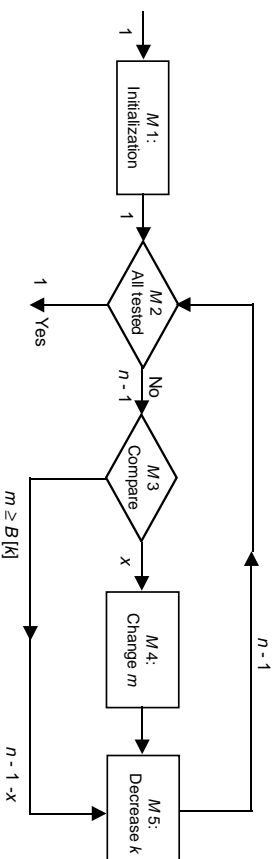
Analysis of Program Max

- Consider the **Pascal program MAX** that finds the largest element in the given array B .
- Given an array of n elements, $B[1], B[2], \dots, B[n]$, we will find m and j such that $m = B[j] = \max\{B[k] \mid 1 \leq k \leq n\}$, and for which j is as large as possible.

```
program MAX (input, output);  
  label 1, 2, 3, 4, 5, 6;  
  const n = 100;  
  var j, k, m: integer;  
      B: array [1..n] of integer;  
  
begin  
  1: j := n; k := n - 1; m := B[n]  
  2: while (k > 0) do  
    begin  
  3:       if B[k] > m  
        then  
  4:           begin  
              j := k;  
              m := B[k]  
            end;  
  5:       k := k - 1  
    end;  
  6: writeln (j, m)  
end.
```

Analysis of Program Max (cont.)

- Analyze the time required for its execution.
- The time of execution depends on:
 - machine
 - compiler
 - input data
- Examine the effect of input data on execution time.
- We study frequency counts for each step.



Analysis of Program Max (cont.)

- X is the number of times we must change the value of the current maximum.
- The value of X depends on the pattern of numbers constituting the elements of the array B .
- Each such pattern may be considered a sample point with a fixed assigned probability.
- X can be thought of as a random variable over the sample space.

Step number	Frequency count	Number of statements
M1	1	3
M2	n	1
M3	$n-1$	1
M4	X	2
M5	$n-1$	1
M6	1	1

Analysis of Program Max (cont.)

- Determine the distribution of the random variable X for a given assignment of probabilities over the sample space.
- The image of the random variable X is $\{0, 1, \dots, n - 1\}$:
 - minimum value of X occurs when $B[n] = \max \{B[k] \mid 1 \leq k \leq n\}$
 - maximum value of X occurs when $B[1] > B[2] > \dots > B[n]$
- For simplicity, assume that $B[k]$ are distinct values. Without a loss of generality, also assume the vector of elements $(B[1], B[2], \dots, B[n])$ is any one of the permutations of the integers $\{1, 2, \dots, n\}$, $n!$
- Sample space S_n : set of all permutations of n integers $\{1, 2, \dots, n\}$.
- Assume that all $n!$ permutations are equally likely.
- For all s in S_n , $P(s) = 1/n!$.

We define a random variable X_n as a function with domain S_n and the image $\{0, 1, \dots, n - 1\}$.

Analysis of Program Max (cont.)

- As n changes, we have a sequence of random variables X_1, X_2, \dots , where X_i is defined on the sample space S_i .

The probability mass function of X_n , $P_{X_n}(k)$ is denoted by P_{nk} . Then:

$$\begin{aligned} P_{nk} &= P(X_n = k) \\ &= \frac{\text{numbers of permutations of } n \text{ objects for which } X_n = k}{n!} \end{aligned}$$

Analysis of Program Max (cont.)

- Now establish a recurrence relation for p_{nk} :

Consider a sample point $s = (b_1, b_2, \dots, b_n)$, a permutation on $\{1, 2, \dots, n\}$.

Consider two mutually exclusive and collectively exhaustive events:

$$A = "b_1 = n"$$

$$\bar{A} = "b_1 \neq n"$$

- If event A occurs, then a comparison with b_1 (in program MAX) will force a change in the value of m .

Analysis of Program Max (cont.)

- The value obtained for X_n will be one higher than a similar value obtained while examining the previous $n-1$ elements (b_2, \dots, b_n) .
- Note that (b_2, \dots, b_n) is a permutation on $\{1, 2, \dots, n-1\}$.
- The number of times the value of m gets changed while examining (b_2, \dots, b_n) is X_{n-1} .

$$\begin{aligned} \text{Thus: } P(X_n = k \mid A) &= P(X_{n-1} = k - 1) \\ &= p_{n-1, k-1} \end{aligned}$$

- The occurrence of event \bar{A} implies that the count of exchanges does not change when we examine b_1 :

$$\begin{aligned} P(X_n = k \mid \bar{A}) &= P(X_{n-1} = k) \\ &= p_{n-1, k} \end{aligned}$$

Analysis of Program Max (cont.)

- By assumption of equiprobable sample space, we have:

$$P(A)=1/n \text{ and } p(A)=(n-1)/n$$

By the theorem of total probability:

$$\begin{aligned} P_{nk} &= P(X_n = k) \\ &= P(X_n = k \mid A)P(A) + P(X_n = k \mid \bar{A})P(\bar{A}) \\ &= \frac{1}{n}P_{n-1,k-1} + \frac{n-1}{n}P_{n-1,k} \end{aligned}$$

- Thus, we can recursively compute p_{nk} if we provide the initial conditions.

Since the image of X_n is $\{0, 1, \dots, n-1\}$, we know that: $P_{nk} = 0$ if $k < 0$. With $n = 1$, the **while** loop in program MAX will never be executed.

Analysis of Program Max (cont.)

- Therefore $X_1=0$, i.e:

$$P(X_1=0) = p_{1,0} = 1 \text{ and } P(X_1=1) = p_{1,1} = 0$$

- The complete specification to evaluate p_{nk} is:

$$p_{1,0} = 1$$

$$p_{1,1} = 0$$

$$p_{nk} = \begin{cases} \frac{1}{n}p_{n-1,k-1} + \frac{n-1}{n}p_{n-1,k}, & 0 \leq k \leq n-1, n \geq 2, \\ 0, & \text{otherwise.} \end{cases}$$