ECE 220

Lecture x000D

Recap

- Formal introduction to recursion
  - Factorial
  - Binary search
  - Towers of Hanoi
  - LC3 implementation
- Today: More recursion & problem solving
  - N - Queens problem
  - Maze solving
  - Exercise(s)

Lesson objectives

- Understanding recursion vs. iteration tradeoff
  - Introduce ***memoization*** as a speedup technique
- Use recursion with backtracking to produce elegant solutions to problems
- Be able to implement recursion with bactracking to solve puzzles, problems, etc.

Good recursion vs. bad recursion

- Consider the recursive Fibonacci function from last time.

```
long long fib(long long n){
     long long sum;

     if (n == 0 || n == 1)
          return 1;
     else {
          sum = (fib(n-1) + fib(n-2));
          return sum;
}
}
```

- Let's do an activity
- Convert this function to an iterative version.
- Compare run times.
- Can recursion be made faster?

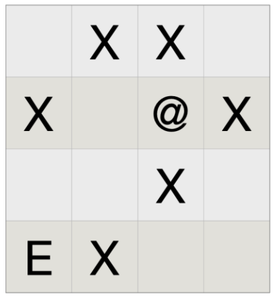
Solving a maze

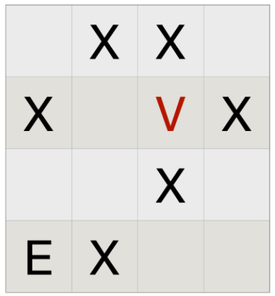
- We represent a maze by a 2D grid of size $N \times M$
- Walls are marked with X and the exit with E.
- Given starting point (i, j) marked with @, find a path to E (if it exists).
  - Do not go outside grid
  - Avoid going around in circles.
  - Mark valid path with P.
- Write a recursive function `int ExitMaze` solving problem.

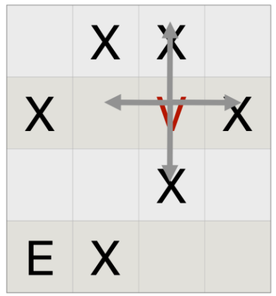
	X	X	
X		@	X
		X	
E	X		

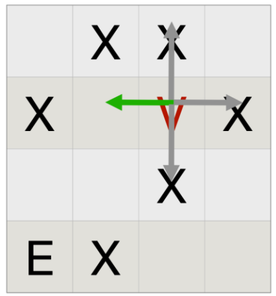
Solving a maze - strategy

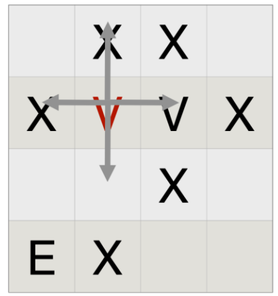
Strategy: Mark current cell as visited and explore solution space. Exploration defined by four possible moves (U, D, L, R).

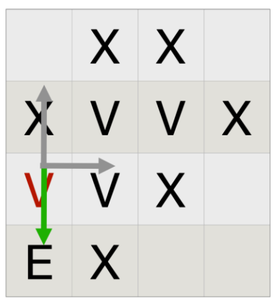


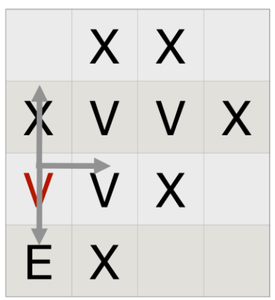


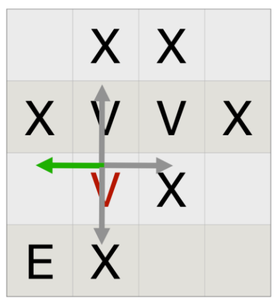


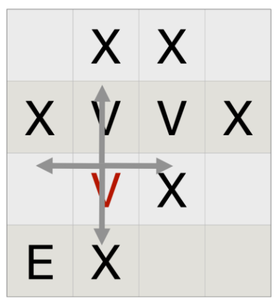


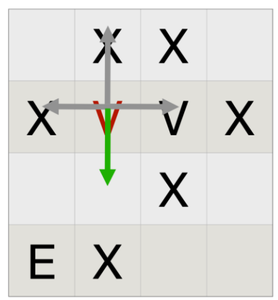












Solving a maze - setup

- Recursion needs base case. What should be the base case?
  - Found exit (return “good”) OR hit X or hit V or out-of-bounds (return “bad”)
  - Let `xpos` and `ypos` be the *row* and *column* index.

```
if (xpos < 0 || xpos >= MAZE_HEIGHT || ypos < 0 || ypos >= MAZE_WIDTH)
  return 0;

if (maze[xpos][ypos] == 'E')    /* Found the Exit! *
  return 1;

if (maze[xpos][ypos] != ' ')    /* Space is not empty (possibly X or V)*/
  return 0;
```

Solving a maze - code

- What should be the recursive call?
  - Go down, up, left or right.
  - `int ExitMaze(maze, xpos, ypos)`
    - Function exploring the solution space.

```
/* Go Down */
if (ExitMaze(maze, xpos + 1, ypos)) {
  maze[xpos][ypos]='P';
  return 1;
}

/* Go Right */
if (ExitMaze(maze, xpos, ypos + 1)) {
  maze[xpos][ypos]='P';
  return 1;
}
```

```
/* Go Up */
if (ExitMaze(maze, xpos - 1, ypos)) {
  maze[xpos][ypos]='P';
  return 1;
}

/* Go Left */
if (ExitMaze(maze, xpos, ypos - 1)) {
  maze[xpos][ypos]='P';
  return 1;
}
```

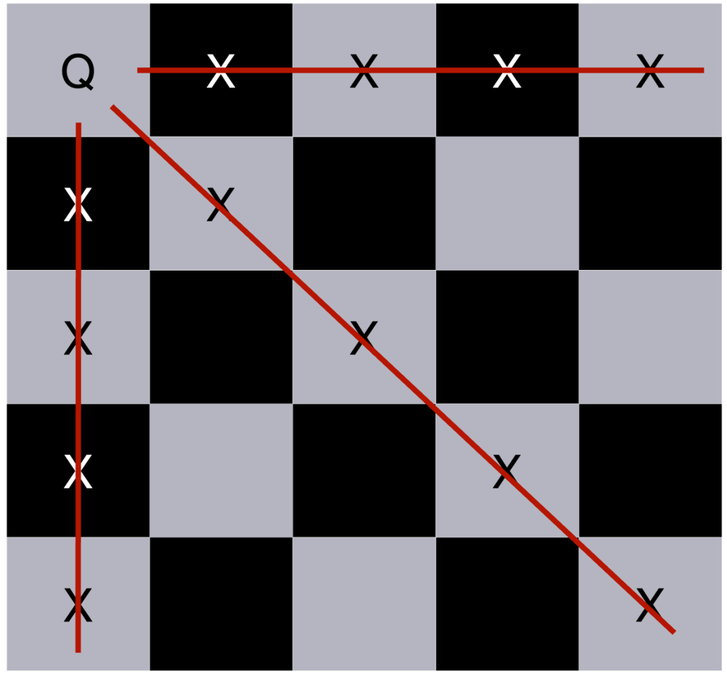
Exercise

- There is an `ExitMaze` function on Github which I tested to work.
- Modify it by adding a `main` function, board definition and try it on this maze.

	X					X
	X			X		
			@	X	X	X
	X					
	X	X	X	X		
X			E	X		
X	X		X	X		X
						X

N - Queens Problem

- In chess, a Queen can attack another piece within its line of sight as long as that piece is in the same: **row**, **column** or **diagonal**.

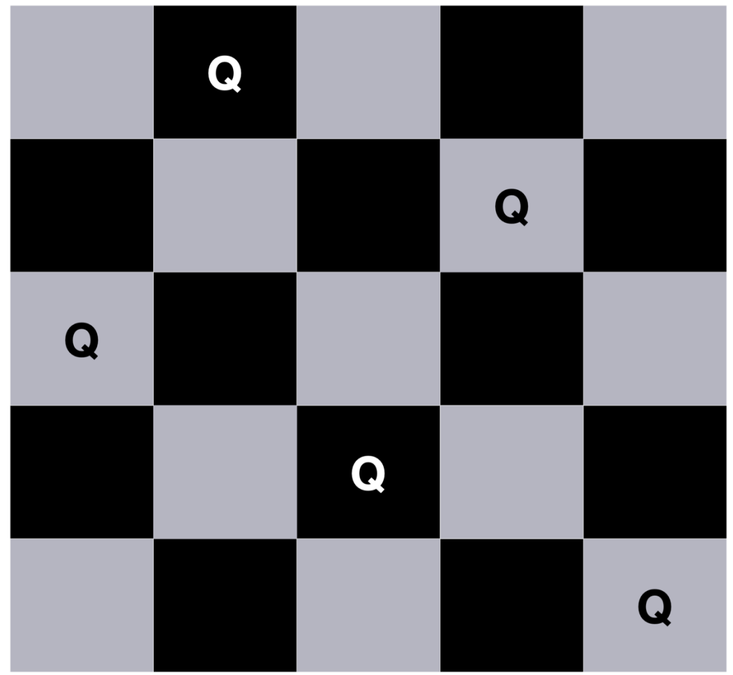


- **Question**: Given an N x N grid, is it possible to place N Queens in the grid so that no two Queens can attack each other ?

- **Answer:** Yes for N>=4.

N - Queens

- Here is a possible solution for the 5 x 5 grid.



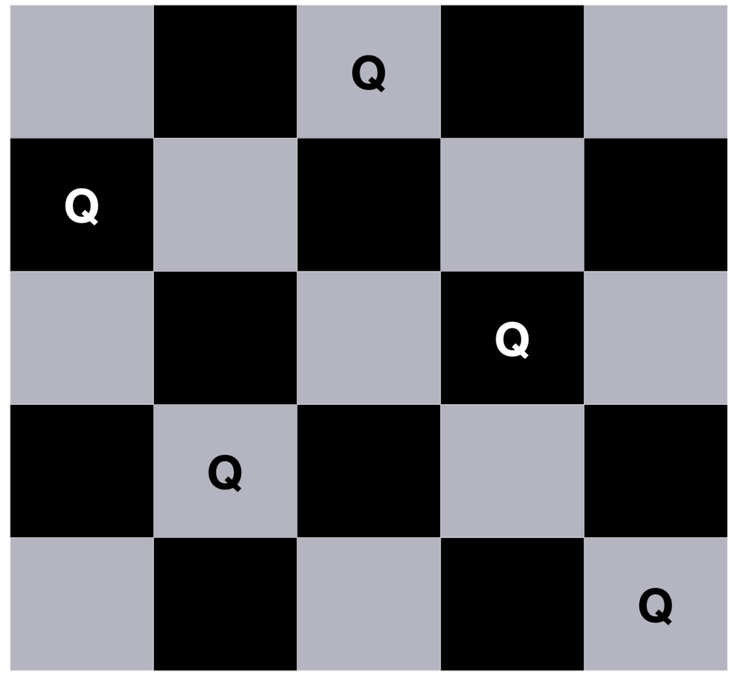
  - Not unique

- Can we make the computer solve it for any given N?

  - Solution: Recursion with *backtracking.*

N - Queens Solution

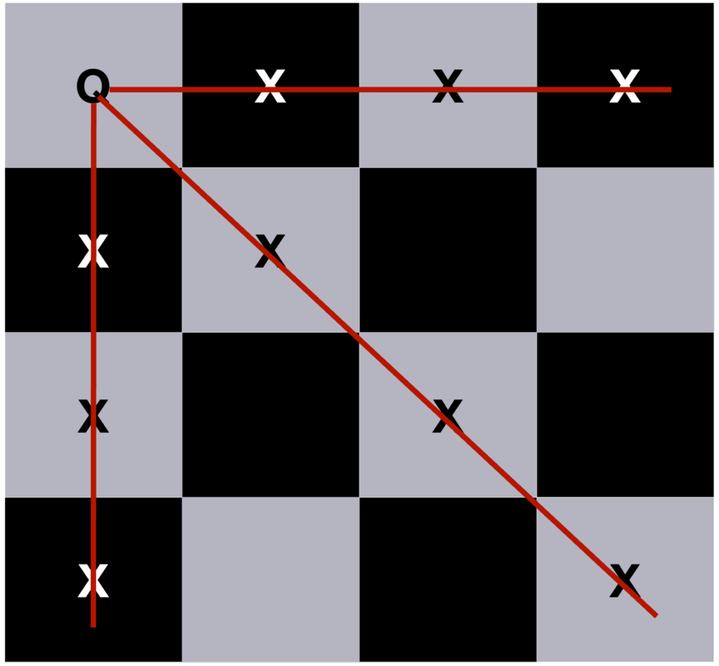
- Here is a possible solution for the 5 x 5 grid.
  - Not unique
- Can we make the computer solve it for any given N?
  - Solution: Recursion with *backtracking.*

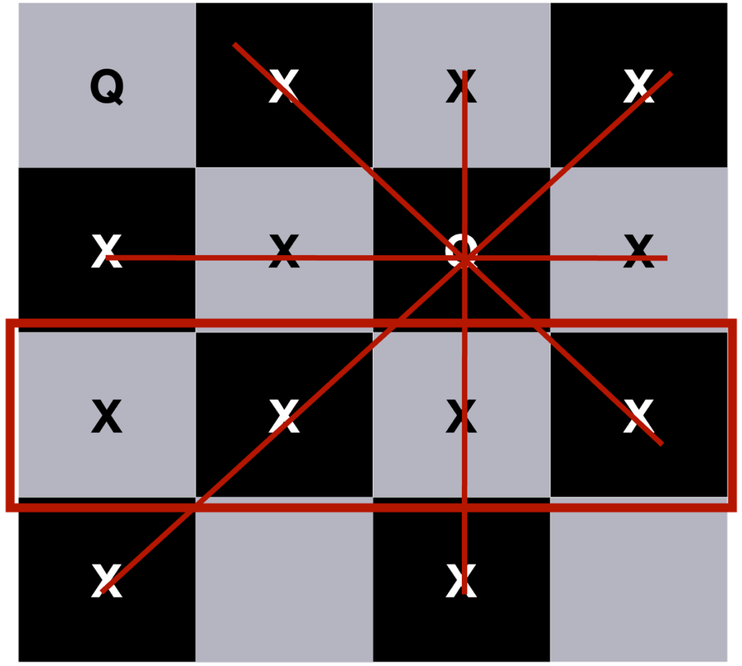


N - Queens Search

- **Back-tracking**: Make a choice and search the solution space. If solution space is empty, return and make a different choice.

Choice #1





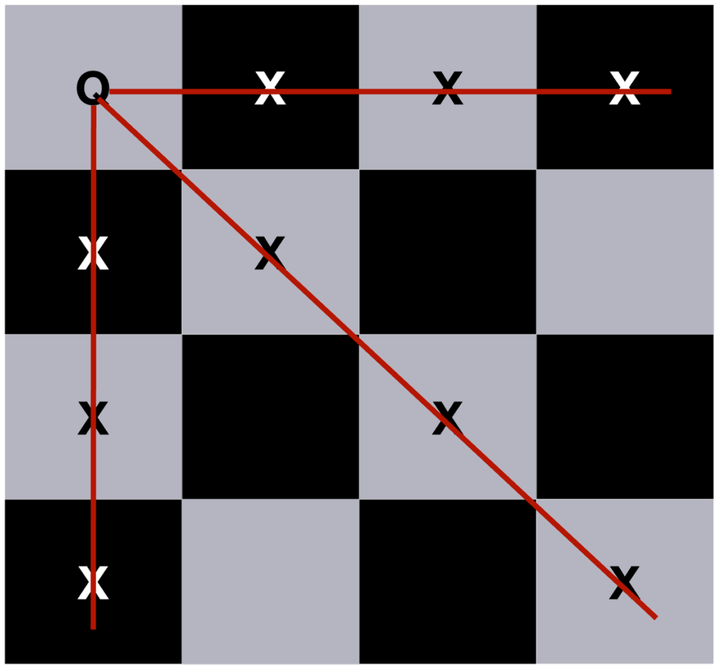
Choice #1.1

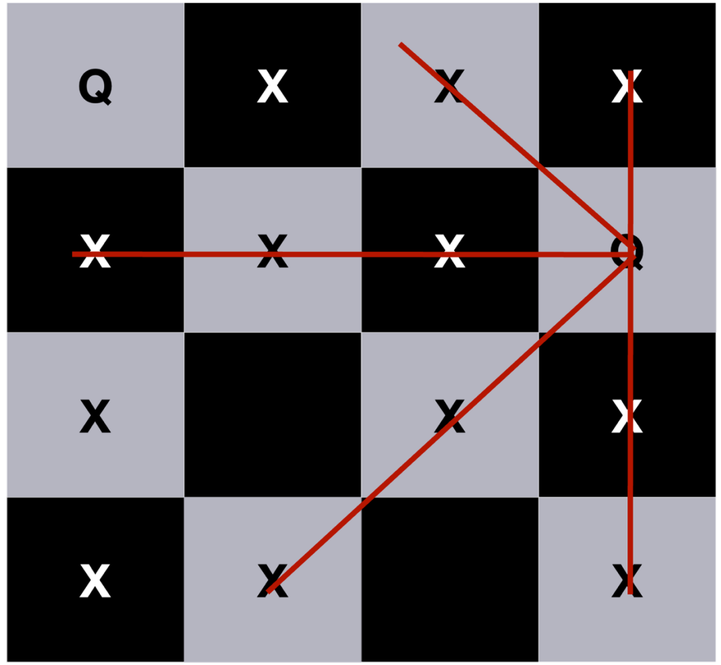
Not a solution!

N - Queens Search Tree

- **Back-tracking**: Make a choice and search the solution space. If solution space is empty, return and make a different choice.

Choice #1



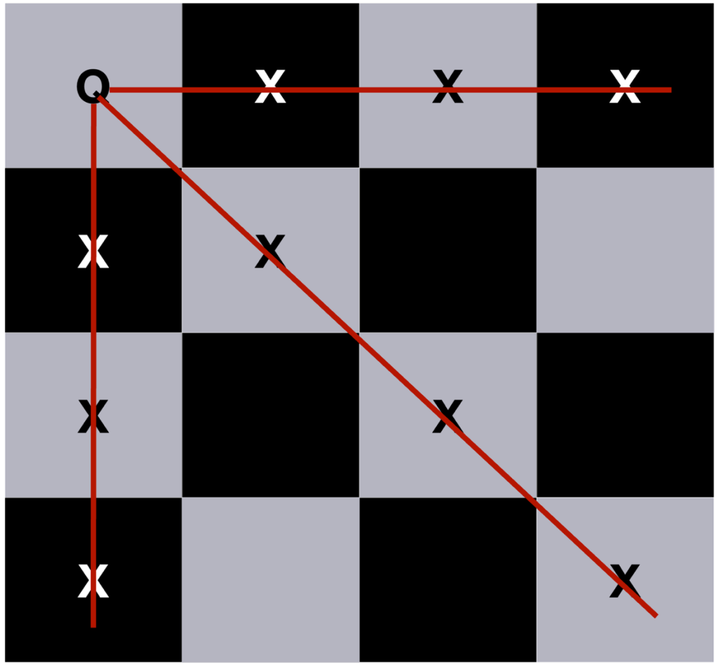


Choice #1.2

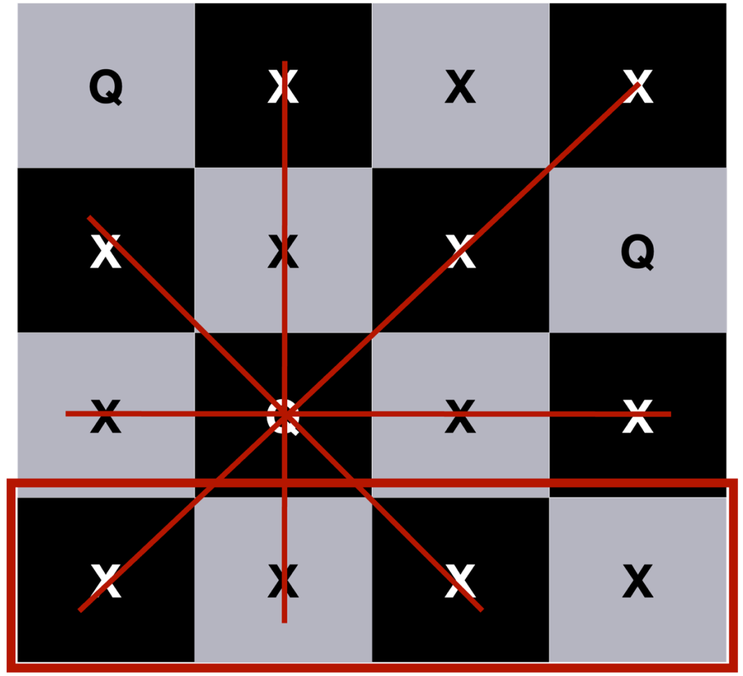
N - Queens Failed Search

- **Back-tracking**: Make a choice and search the solution space. If solution space is empty, return and make a different choice.

Choice #1



Choice #1.2.1



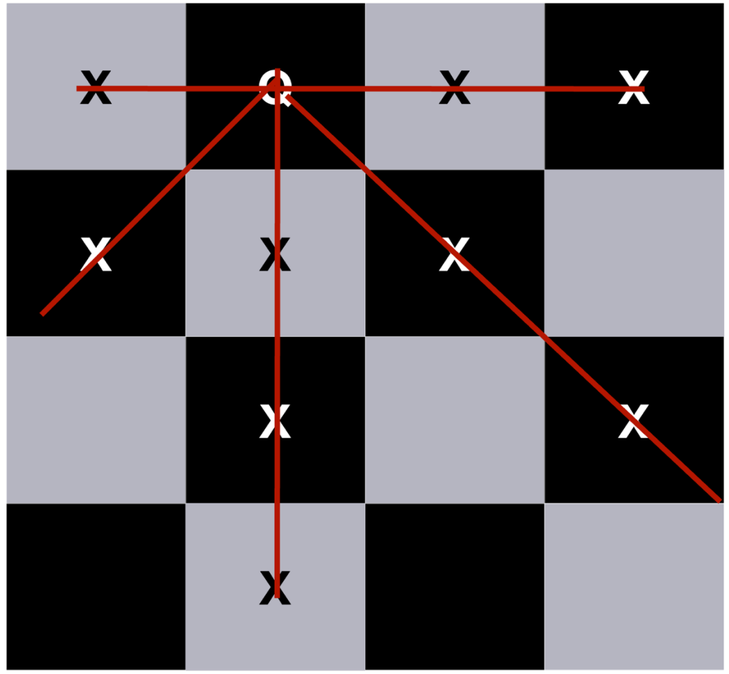
Choice #1.2

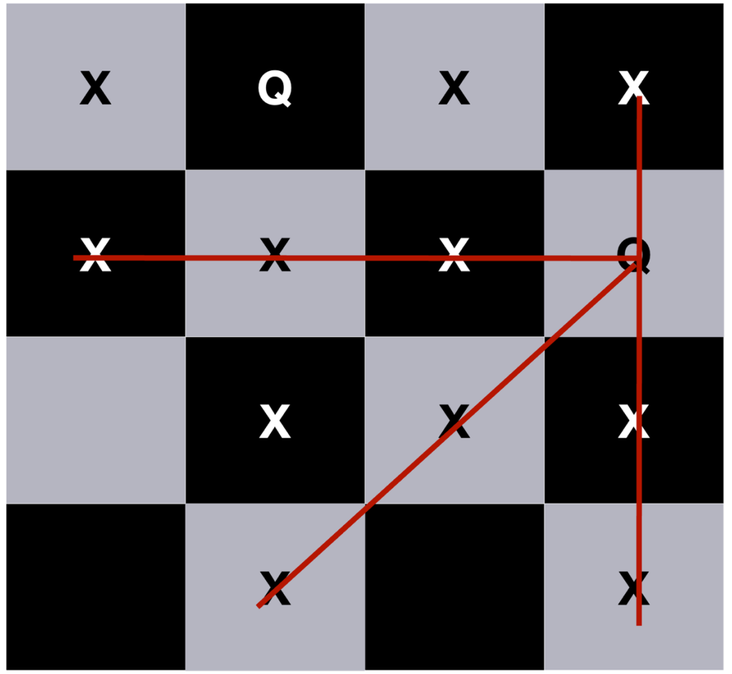
Not a solution!

N - Queens Backtrack

- **Back-tracking**: Make a choice and search the solution space. If solution space is empty, return and make a different choice.

Choice #2

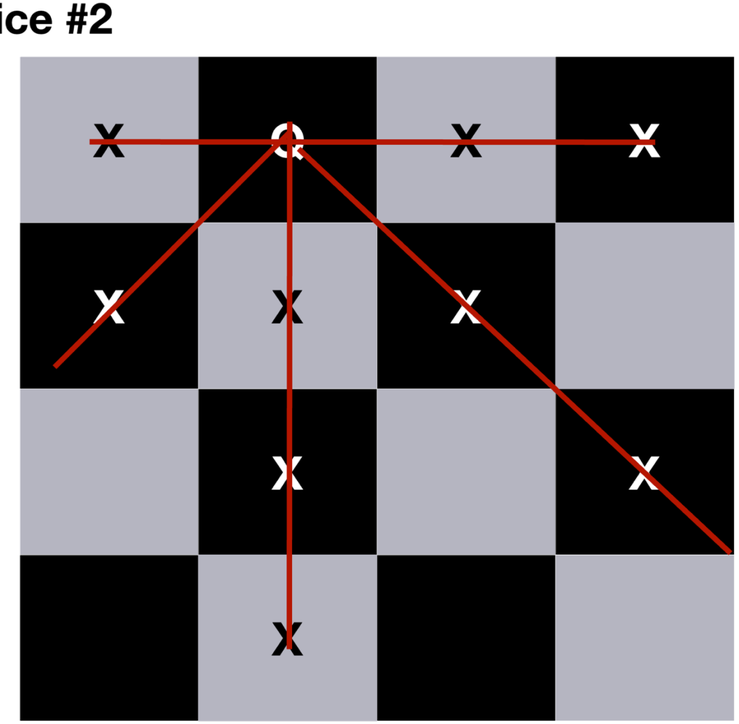


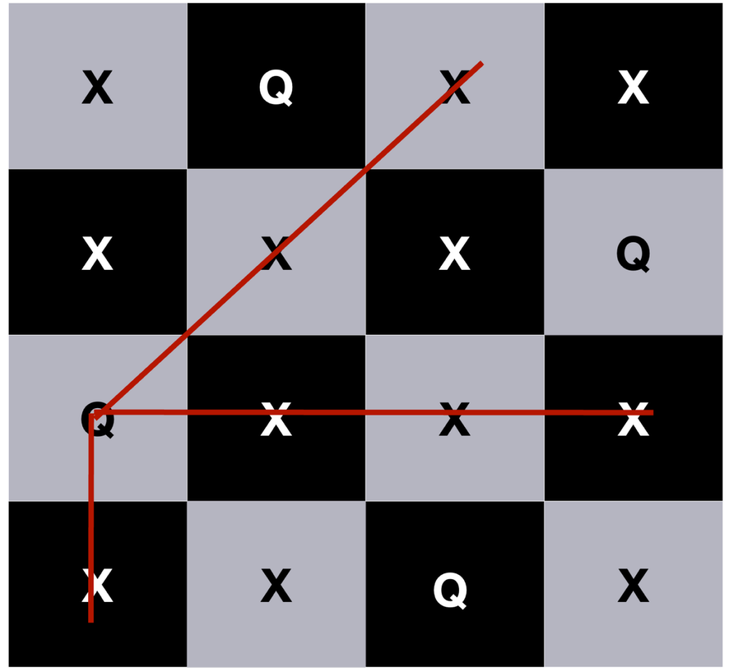


Choice #2.1

N - Queens Successful Search

- **Back-tracking**: Make a choice and search the solution space. If solution space is empty, return and make a different choice.





N - Queens implementation?

- **Question:** Can we set this up as a recursive problem?
  - What is the action/sub-problem that we want to repeat?

    → Placing a Queen in a row
    - If not successful how do we *backtrack?*

      → Undo placing a queen
- How do we know we have reached an end case?

  → No more rows to fill.

N - Queens set-up?

- We represent the configuration space with a grid.
  - We will denote with digit **zero** an empty spot (maybe safe or unsafe, but its unoccupied).
  - We will denote with the digit **one** a space occupied by a queen.
  - We will fill in rows starting with the first row and proceeding downward.

N - Queens implementation

```
int is_safe(int board[N][N], int rnum, int cnum);

/*Function places a queen in row rnum */
int place_queen(int board[N][N], int rnum){
  if (           )  /* Finished all rows */
    return 1;       /* Found a solution  */
  else{
    /* Iterate over possible columns */
    for(int cnum=0;           ; cnum++)
      if (is_safe(                )==1){
        board[rnum][cnum] = 1; /* Place a queen there */
        /* Update row number and recurse */
        if (                                ==1)
          return 1;
        else /* Hit a road block down the line */
                        ; /* Remove queen */
    } /* Try next column along row */
  }   /* For loop finished without hitting a return */
  return 0;   /* Solution doesn't exist. */
}
```

`is_safe` checks whether it is possible to place a queen on position `(rnum, num)` given the configuration of the board at some given time. It returns 1 if safe or 0 if unsafe

`place_queen` fills the board with a valid solution and returns 1 or returns 0 if no solution found.

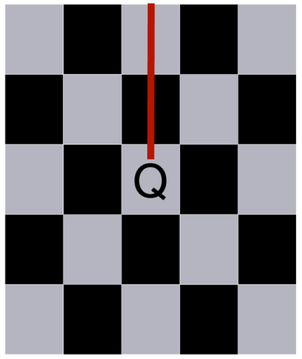
Is it safe/unsafe?

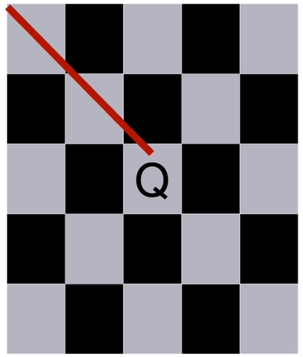
- On the N-th row when we place a queen on a square `(i,j)` what do we need to check?

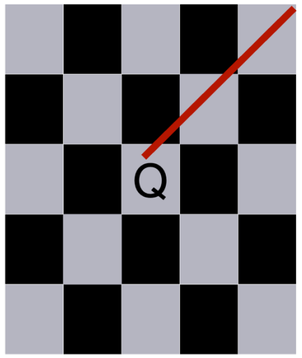
  → Are we in the line of sight (LOS) for any previous Queen.

  - We are in LOS if
    - The column `i` contains any Queen OR
    - The diagonals to the top-left of `(i,j)` contains a Queen OR
    - The diagonals to the top-right of `(i,j)` contains a Queen
    - What about diagonals to the bottom left or bottom right?

Is it safe/unsafe - code








```
int is_safe(int board[N][N], int row, int col){
  int i, j;
  for (                    ){/*Check along column*/
    if (board[i][col]==1)
      return 0;
  }
  /* Check diagonal to upper left */
  for (                    ; i>=0 && j>=0; i--, j--){
    if (board[i][j] == 1)
      return 0;
  }
  /* Check diagonal to upper right */
  for (i=row-1, j=col+1;           ;              ){
    if (board[i][j]==1)
      return 0;
  }
  return 1;
}
```

Exercise for fun outside lecture

- Exercise for the *curious/mighty/brave*:
  - Modify the source of `queens.c` so that it keeps a static variable to keep track of the recursive calls.
  - Varying `N`, generate a plot of `N` vs number of recursive calls. Try `N=4, 5, …, 15`. What kind of growth is it?

Exercise - practice, practice, p….

- You have a pile of wood sticks with 3 different lengths: 3, 7, and 10 feet. You want to connect them and make an X-feet long stick using **at most** 10 sticks.

- To make a stick 33 feet long you can do:
  - 4 x 3F + 3 x 7F ✓
  - 11 x 3F ✗

- Use recursion with backtracking to find a solution

Exercise - sticks

```c
#define N 10 /* Number allowed   */
#define M 3  /* Types of lengths */

/* Implement this function
    solution[N]: stores the solution
    idx: index for the solution matrix
    total: remaining length */
int solve(int solution[N], int idx, int total);

const int set[M] = {3,7,10};

int main(){
  int solution[N] = {0, 0, 0, 0, 0, 0, 0, 0, 0, 0};
  int total;
  printf("Enter total length: ");
  scanf("%d", &total);
  /* Write your code here */
}
```

Time permitting

- Using the `gdb` debugger with gdb dashboard.