

ECE 220

Lecture x0004 - 01/29

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Recap

- Last time we discussed:

- Stacks

- Quarters vs. pancakes

- Implementing PUSH/POP

- Examples of use cases for stacks

Implementation differences in TOS convention

- Current top-most element
- Next available spot

A. Balanced parentheses

B. Palindrome check

C. Stack arithmetic

Lesson objectives

- Understand and explain concepts of infix and postfix notation for arithmetic.
- Understand and explain advantage of postfix notation over infix notation. Be able to evaluate postfix expressions.
- Understand implementation of arithmetic using postfix notation and stack ADT
- Understand steps necessary to implement an RPN calculator in LC3
 - Know how to deal with overflow and underflow on the stack

RPN or postfix arithmetic

- Traditional arithmetic notation is called *infix* notation. Operations are inserted between operands. E.g. $5 + 3$ or 3×4
 - Requires use of parenthesis to indicate order of operations
- An alternative notation is called postfix notation a.k.a Reverse Polish notation (RPN). E.g. $5\ 3\ +$ or $3\ 4\ \times$
 - Implemented properly, does not require parenthesis/brackets

Postfix expressions

- The syntax for a postfix operation is:

`<operand1> <operand2> <operator>`

- $(2 + 5) = 7 \implies 2\ 5\ +$
- Operands may be postfix subexpressions

$2 - (5 + 4) \implies$

2	5	4	+	-
---	---	---	---	---

↓
Operand2

Postfix expressions

- The syntax for a postfix operation is:

```
<operand1> <operand2> <operator>
```

- Advantage: no need for parentheses

• $2 + 5 \times 4 \implies (2 + 5) \times 4$ or $2 + (5 \times 4)$?

$\downarrow$ $\downarrow$

$2\ 5\ +\ 4\ \times$ $2\ 5\ 4\ \times\ +$

Practice RPN - MP2 material

- Note: $53 - \mapsto 5 - 3$
- Consider: $3\ 4\ *\ 7\ 2\ -\ 3\ *\ +$
 - What does it evaluate to?
 - What is the *infix* version of the above?
 - Can we evaluate it using a stack?

Example

+
x
3
-
2
7
x
4
3

□ = 27

□ = 15

□ = 5

□ = 12

POP FROM STACK

Postfix expressions

- Rewrite the following infix expressions in RPN:
 - $(8 + 4)^2$
 - $7 + (9 - 6)/3$
 - $(5 + (1 + 2) \times 4) - 3$

Postfix expressions

- Now evaluate them
 - $8\ 4\ +\ 2\ \wedge$
 - $7\ 9\ 6\ -\ 3\ \div\ +$
 - $5\ 1\ 2\ +\ 4\ \times\ +\ 3\ -$

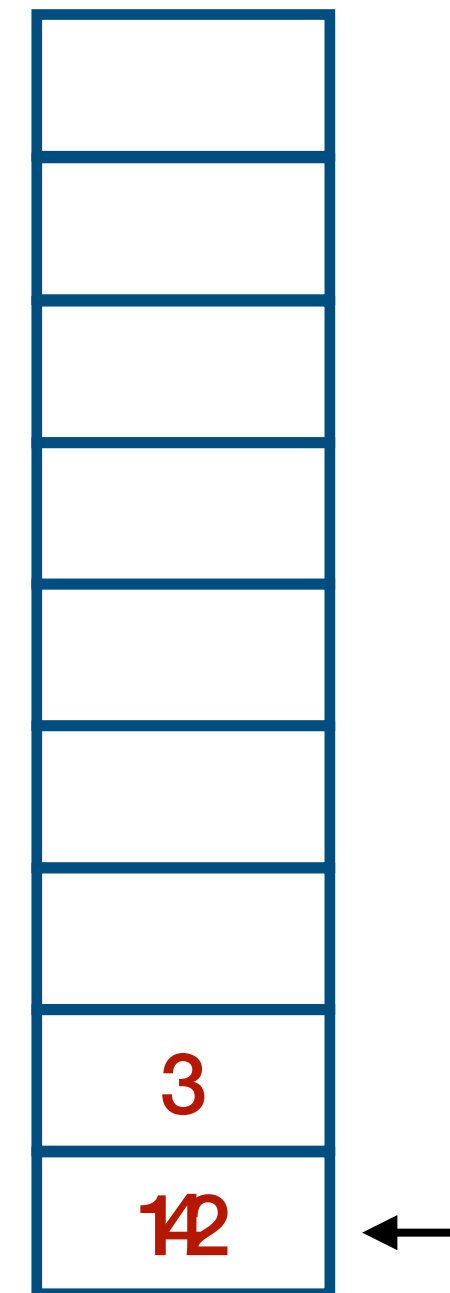
Redo example - single pass?



• Expression: 4 3 $\times$ 7 2 $-$ 3 $\times$ +

• Strategy:

- Numbers $\rightarrow$ Push
- Operator:
 - Pop two elements
 - Perform operation
 - Push on stack

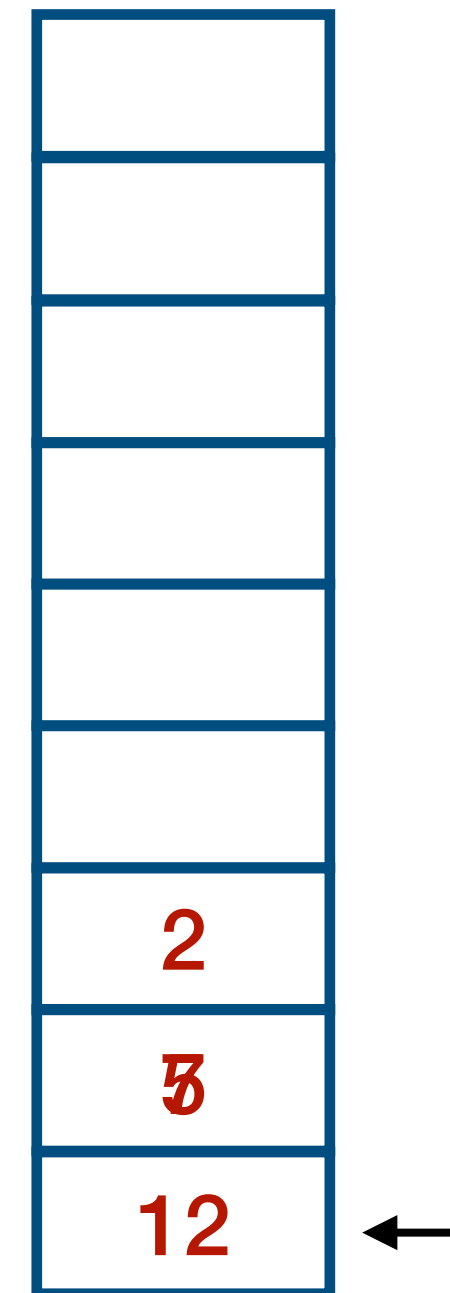


Redo example - single pass?

↓
• Expression: 4 3 × 7 2 − 3 × +

• Strategy:

- Numbers → Push
- Operator:
 - Pop two elements
 - Perform operation
 - Push on stack

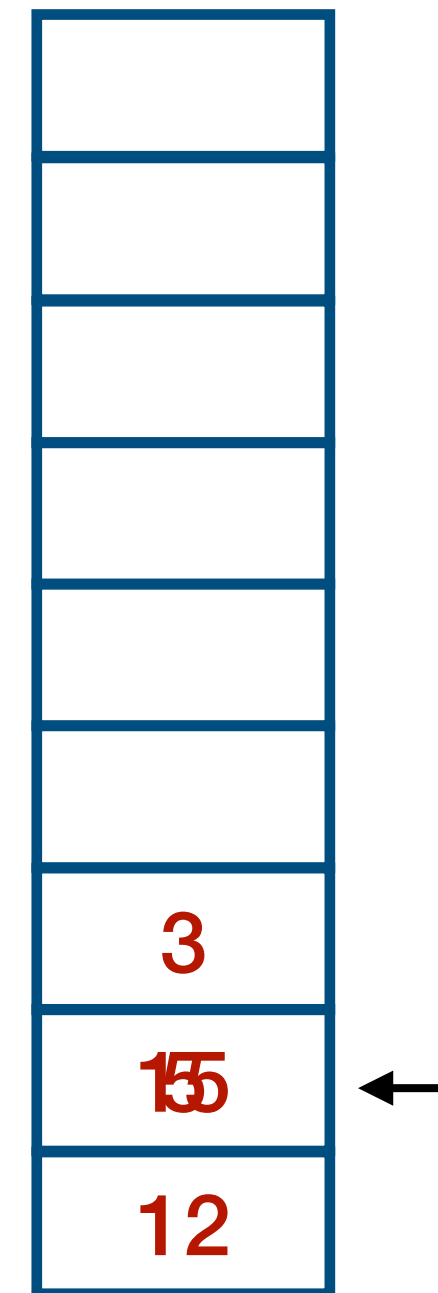


Redo example - single pass?

• Expression: 4 3 $\times$ 7 2 $-$ 3 $\times$ +
↓

• Strategy:

- Numbers $\rightarrow$ Push
- Operator:
 - Pop two elements
 - Perform operation
 - Push on stack



Redo example - single pass?

• Expression: 4 3 $\times$ 7 2 $-$ 3 $\times$ +
↓

• Strategy:

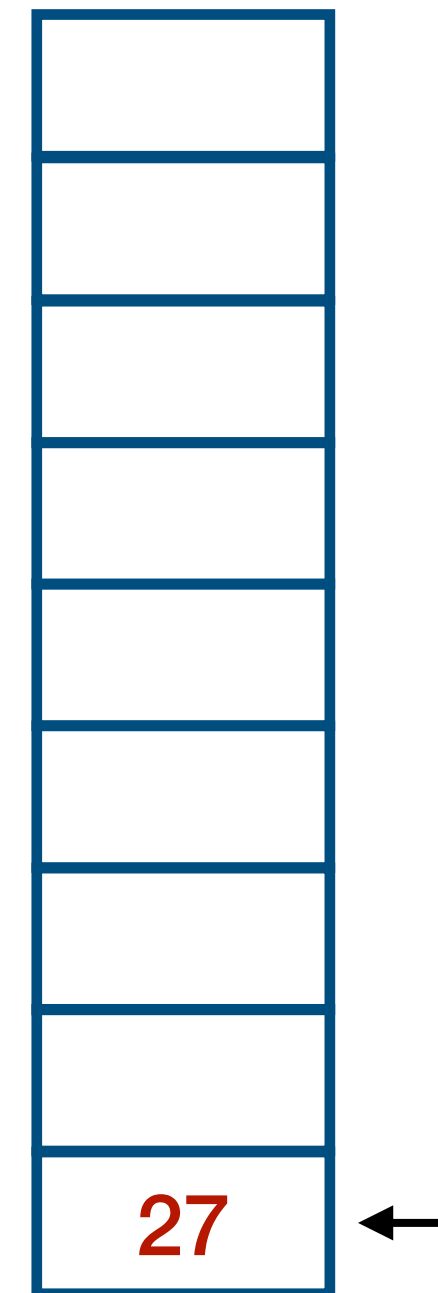
- Numbers $\rightarrow$ Push

- Operator:

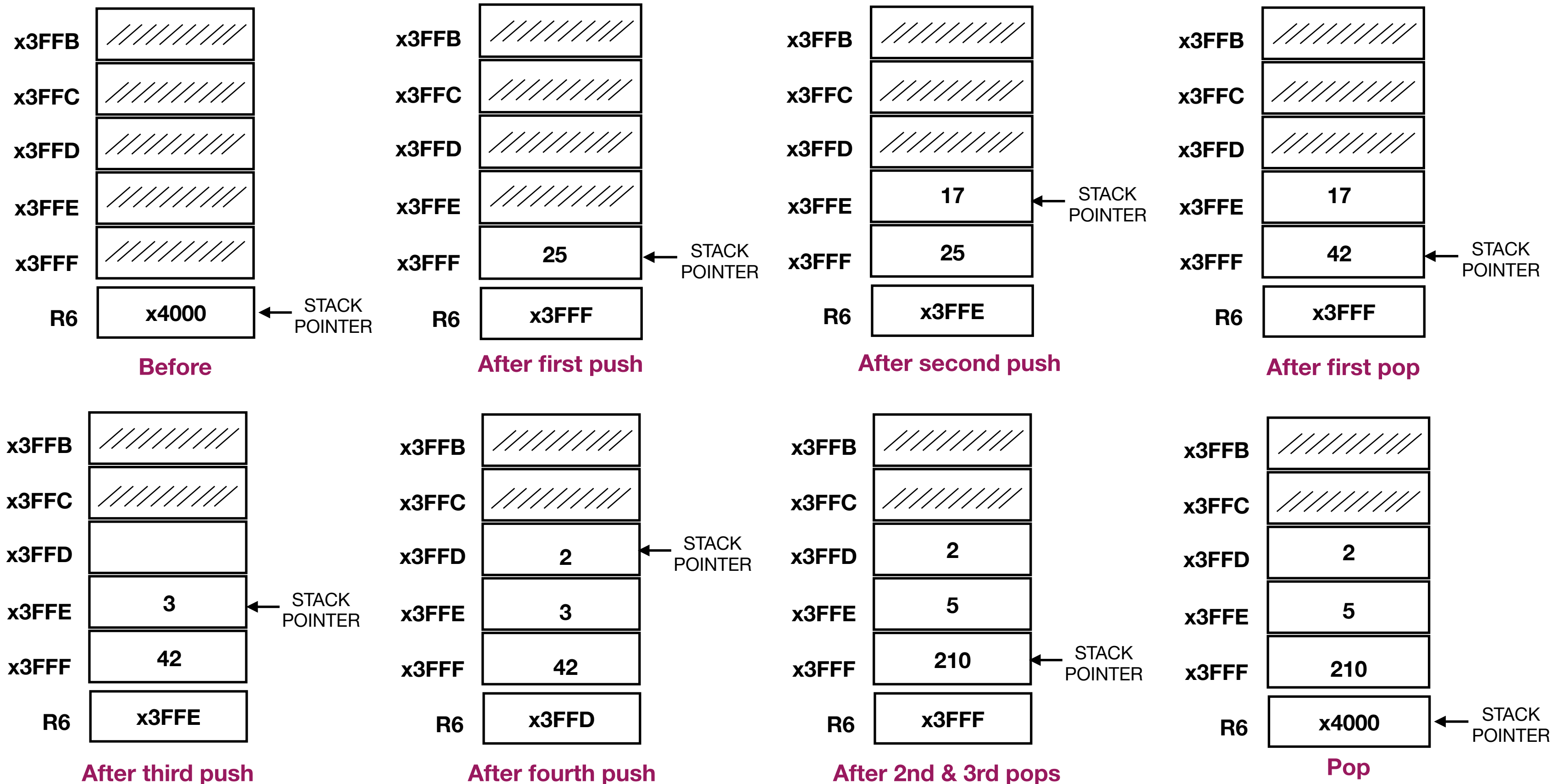
- Pop two elements

- Perform operation

- Push on stack



Stack usage - memory



Arithmetic using stack - LC3

- Compute $(A + B) \times (C + D)$ and store result in R0
- Compute $A \times B + C \times D$ and store result in R0

;Implementation using registers

```
LD R0, A
LD R1, B
ADD R1, R0, R1
LD R2, C
LD R3, D
ADD R3, R2, R3
JSR MULT
HALT
```

MULT: subroutine such that
Input: R1, R3 and Output: R0

MULTIPLY: **POP** two
numbers, compute and
then **PUSH** result back

;Implementation using stack

```
LD R0, A
PUSH
LD R0, B
PUSH
JSR ADD ;Assuming ADD exists
LD R0, C
PUSH
LD R0, D
PUSH
JSR ADD
JSR MULTIPLY ;Assuming MULTIPLY exists
POP ;RESULT in R0
```

ADD: **POP** two numbers,
compute and then **PUSH**
result back

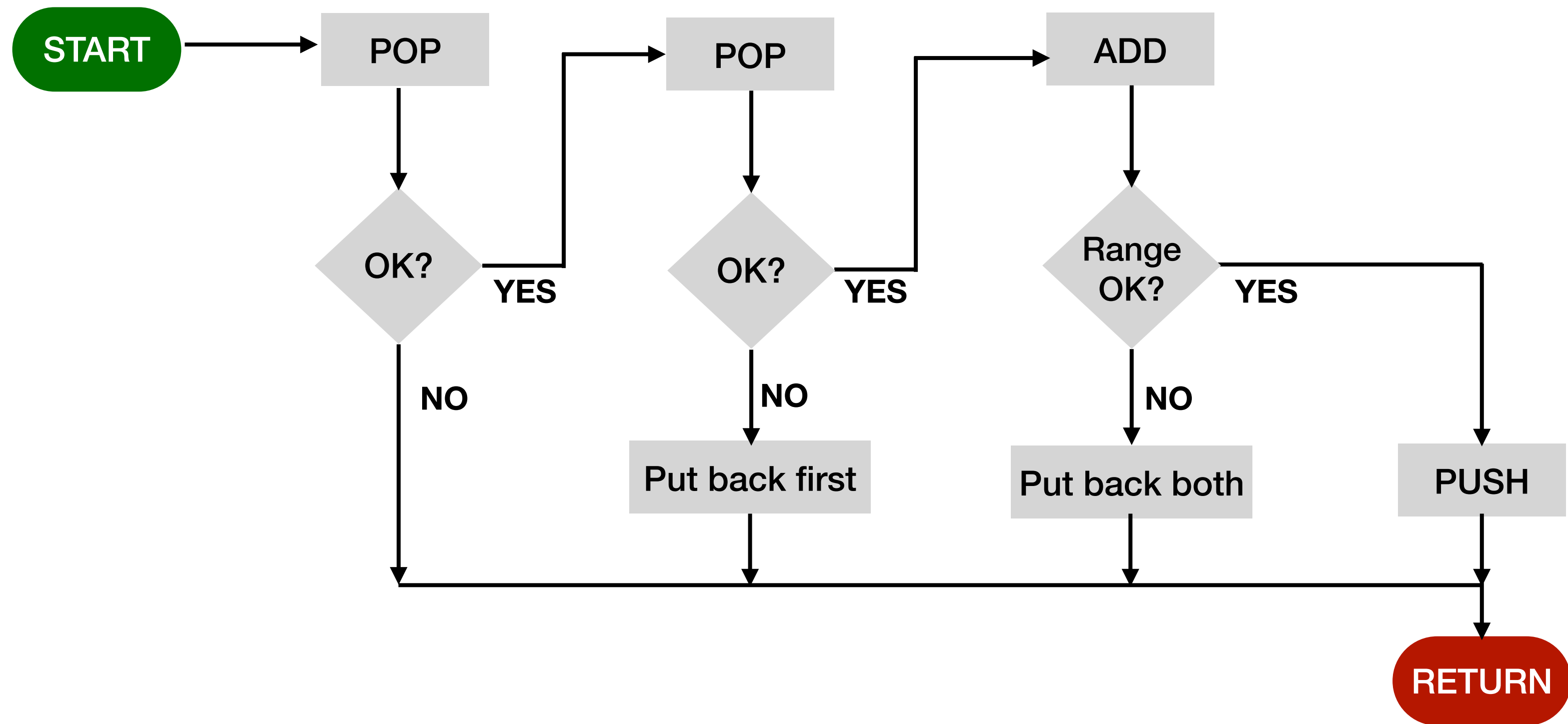
Postfix evaluation

Problem: Given a postfix expression with numerals and '+', '-', '*' in the form of a string, evaluate it and store the answer in R5. Each numeral is a single character.

Algorithm:

- Read the string (postfix expression) left to right
- Push the numbers in the expression on the stack
- For an operator, pop the top two elements, compute the answer and push it on the stack

Flowchart - ADD subroutine



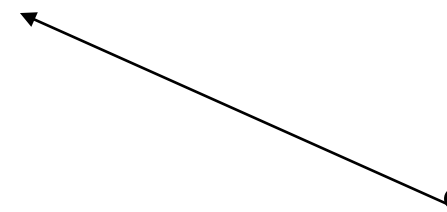
Implement ADD subroutine

```
;PUSH
;Input: R0 (value to store on stack)
;Output: R5 (0-success, 1-fail)
```

```
;POP
;Output: R0 (value to load from stack)
;Output: R5 (0-success, 1-fail)
```

```
;CHECK_RANGE
;Input: R0 (value to be checked)
;Output: R5 (0-success, 1-fail)
```

- Save R7 before calling other subroutines.
- Save registers that will be altered in this subroutine
- R6 is stack pointer (points to the next available spot on the stack)
- Assume PUSH, POP and CHECK_RANGE subroutines are provided to you



```
; ADD subroutine – pop two numbers from stack,  
; perform '+' operation and then push result back to  
the stack
```

```
ADD_OP  
; save registers
```

```
; initialize R5
```

```
; first pop
```

```
; check return value of first pop, go to EXIT if it  
failed (R5 = 1)
```

```
;save value in R1 before second pop
```

```
; second pop
```

```
; check result of second pop, go to RESTORE_1 if it  
failed
```

```
; add two numbers: R0 <- R0 + R1
```

```
; check range of sum, go to RESTORE_2 if it failed
```

```
; everything is good, push sum (already in R0) to  
stack  
;
```

```
RESTORE_1      ; put back first number
```

```
    ; Load STACK_TOP  
    ; Put back item  
    ; Update STACK_TOP  
    ; Go to exit
```

```
RESTORE_2      ; put back both numbers
```

```
    ; Load STACK_TOP  
    ; Put back item(s)  
    ; Update STACK_TOP  
;  
EXIT  
; update stack top pointer  
; restore registers
```

```
RET
```

Next time

- Introduction to C
 - Compiling a C program on EWS
 - Running the GNU debugger etc.
 - Bring your laptop!