

ECE 220: Computer Systems & Programming

Lecture 14: Recursion

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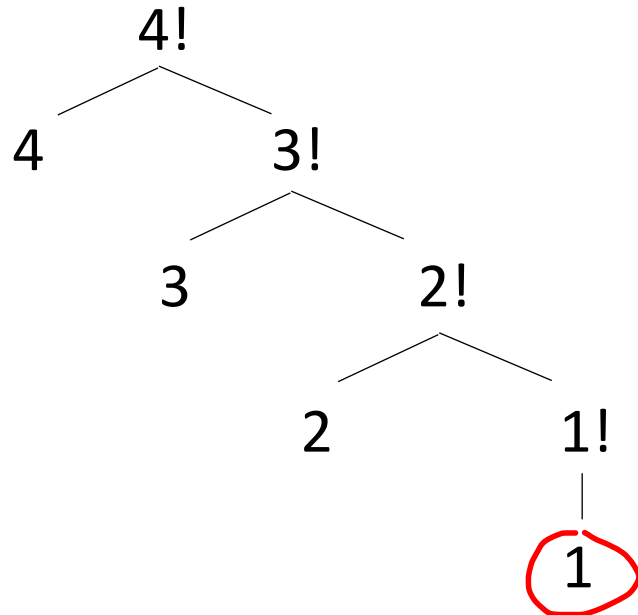
Example: Factorial

- Mathematical definition

$$n! = n \cdot (n - 1) \cdot (n - 2) \dots \cdot 2 \cdot 1 \quad \text{for } n > 0$$

- Recursive form

$$\underline{n!} = \begin{cases} 1, & \text{if } n = 1 \leftarrow \\ n \cdot \underline{(n - 1)!}, & \text{else} \end{cases}$$



Recursion



- A recursive function solves its task by calling **itself** on smaller pieces of data.
 - Each step 1) calls the same function with 2) a variant of the input until 3) a certain condition is satisfied.
 - Must have at least one **base case** (terminating case) that ends the recursive process.
 - Sometimes recursion results in a simpler solution than iteration version (can be used interchangeably).

Recursive Function

1. Base case (= Terminating case)

- This case is mandatory as it provides the way out from the recurring loop.
- The base case must provide a condition that will eventually become true and returns from the function. Otherwise, the run-time stack will overflow.

2. Recursive case (= Induction case)

- This case returns a recursive call to itself. It breaks down the problem into smaller chunks that can be solved over and over by the same function.
- The input to the next call gets induced gradually.

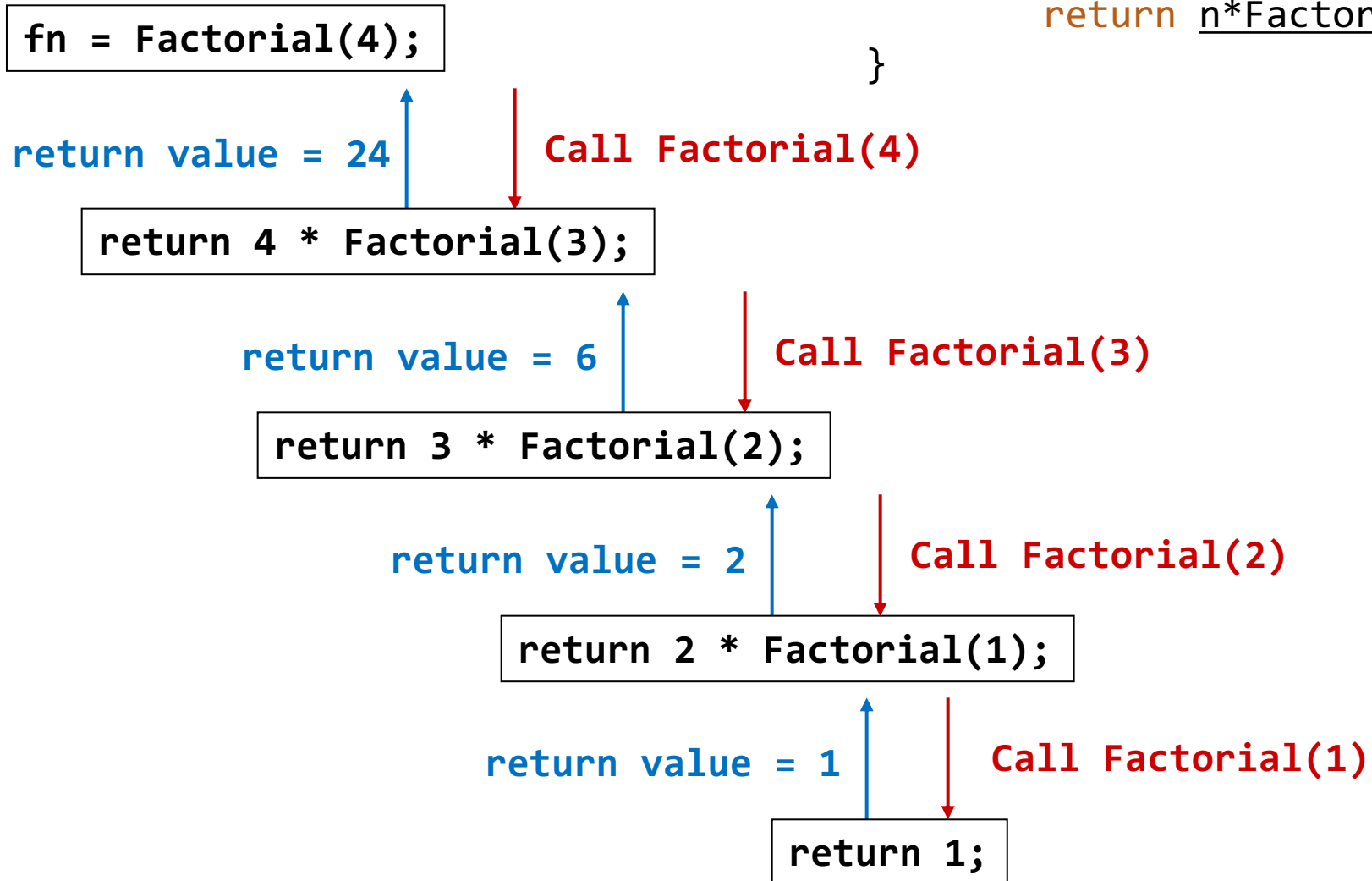
Example: Factorial

```
int Factorial_iter(int n){
    int i, result = 1;
    for(i=0;i<n;i++)
        result = result*(i+1);

    return result;
}
int Factorial(int n){
    if( n <= 1)
        return 1;
    else
        return n*Factorial(n-1);
}
```

Tracing the Factorial Calls

```
int Factorial(int n){  
    if( n <= 1)  
        return 1;  
    else  
        return n*Factorial(n-1);  
}
```



Exercise- Binary Search

```
int binarySearch(int array[], int l, int h, int key)
{
    int mid;
    while( l <= h){
        mid = (l+h)/2;
        if(key < array[mid])
            h = mid -1;
        else if(key > array[mid])
            l = mid +1;
        else
            return mid;
    }
    return -1;
}
```

→ Can we implement binary search in **recursive** way?

Binary Search

	0	1	2	3	4	5	6	7	8	9
Search 23	2	5	8	12	16	23	38	56	72	91
	L=0	1	2	3	M=4	5	6	7	8	H=9
23 > 16 take 2 nd half	2	5	8	12	16	23	38	56	72	91
	0	1	2	3	4	L=5	6	M=7	8	H=9
23 > 56 take 1 st half	2	5	8	12	16	23	38	56	72	91
	0	1	2	3	4	L=5, M=5	H=6	7	8	9
Found 23, Return 5	2	5	8	12	16	23	38	56	72	91

Exercise- Binary Search: in Recursive way

```
int bsr(int array[], int l, int h, int key)
{
    int mid = (l+h)/2;

    if (l > h) // base case1
        return -1;
    else if(key == array[mid]) // base case2
        return mid;
    else if(key > array[mid])
        return bsr(array, mid+1 ,h , key);
    else //key < array[mid]
        return bsr(array, l , mid-1 , key);
}
```


0	1	2	3	4	5
11	12	22	25	34	64

↑ ↑

`found = bsr(array, 0, 5, 64);`

return value = 5 Call bsr(array, 0, 5, 64)
↑ ↓
mid=2

`return bsr(array, 3, 5, 64);`

return value = 5 Call bsr(array, 3, 5, 64)
↑ ↓
mid=4

`return bsr(array, 5, 5, 64);`

return value = 5 Call bsr(array, 5, 5, 64)
↑ ↓
mid=5

`return 5;`

```

int bsr(int array[], int l, int h, int key)
{
    int mid = (l+h)/2;

    if (l > h) // base case1
        return -1;
    else if(key == array[mid]) // base case2
        return mid;
    else if(key > array[mid])
        return bsr(array, mid+1, h, key);
    else //key < array[mid]
        return bsr(array, l, mid-1, key);
}

```

Exercise: ???

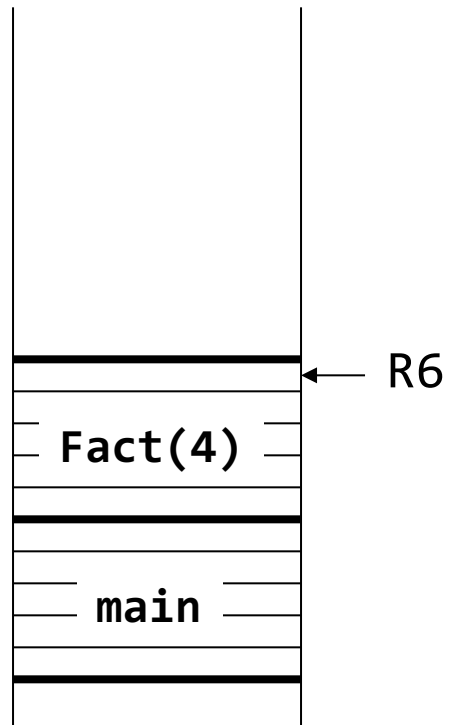
```
int foo(char *c){  
    if(*c==0)  
        return 0;  
    else  
        return foo(c+1) + 1;  
}
```

```
int main(){  
    char str[100] = "ABC";  
    printf("%d", foo(str));  
}
```

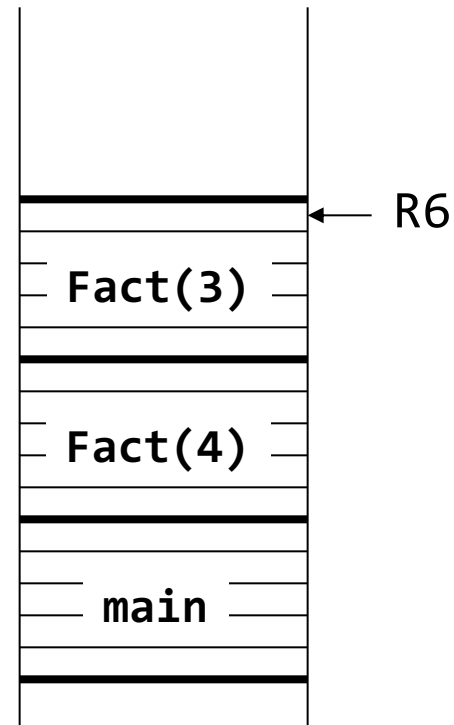
- What does the “foo” function return?
- How many function call?

Run-time Stack during Execution of Factorial

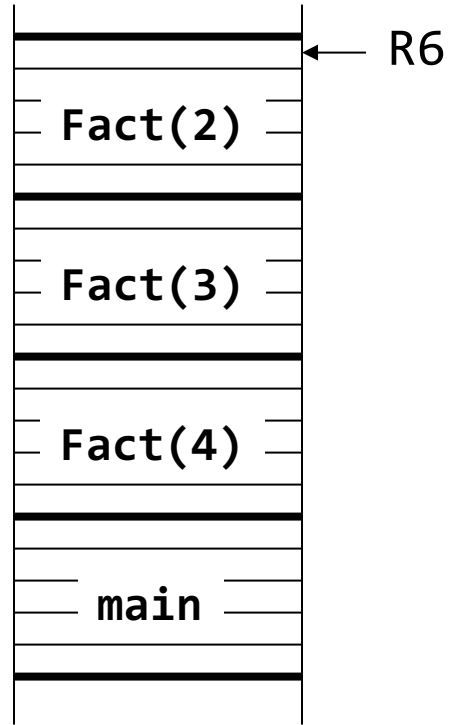
main calls
Factorial(4)



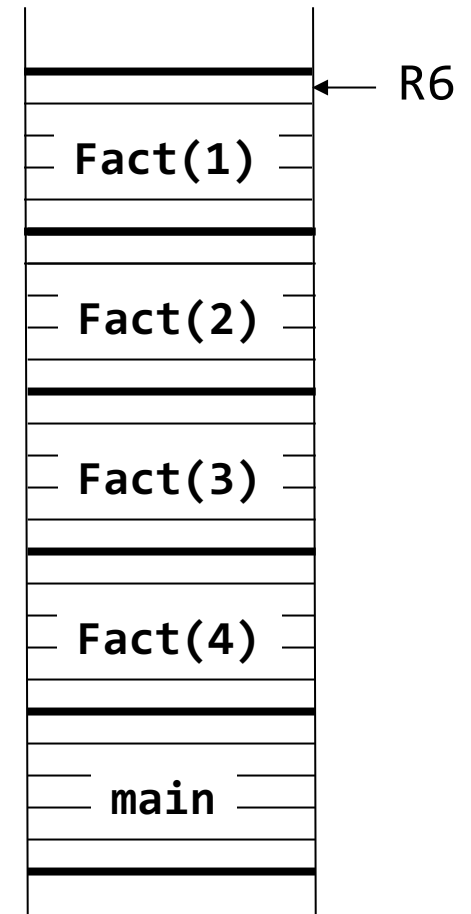
Factorial(4) calls
Factorial(3)



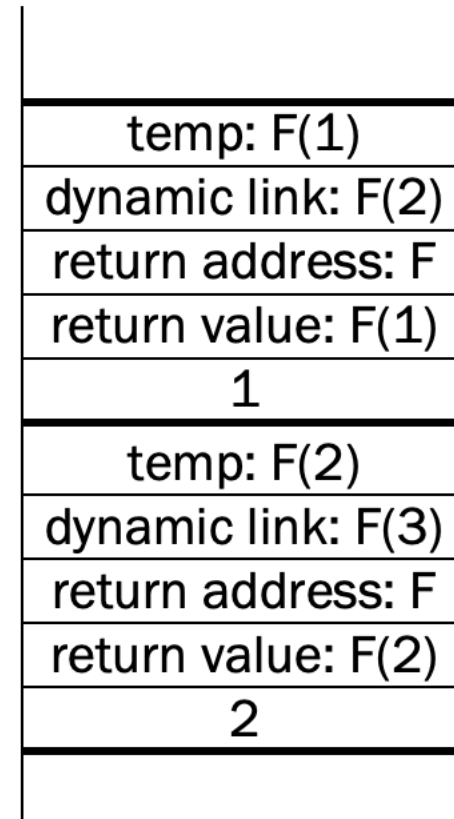
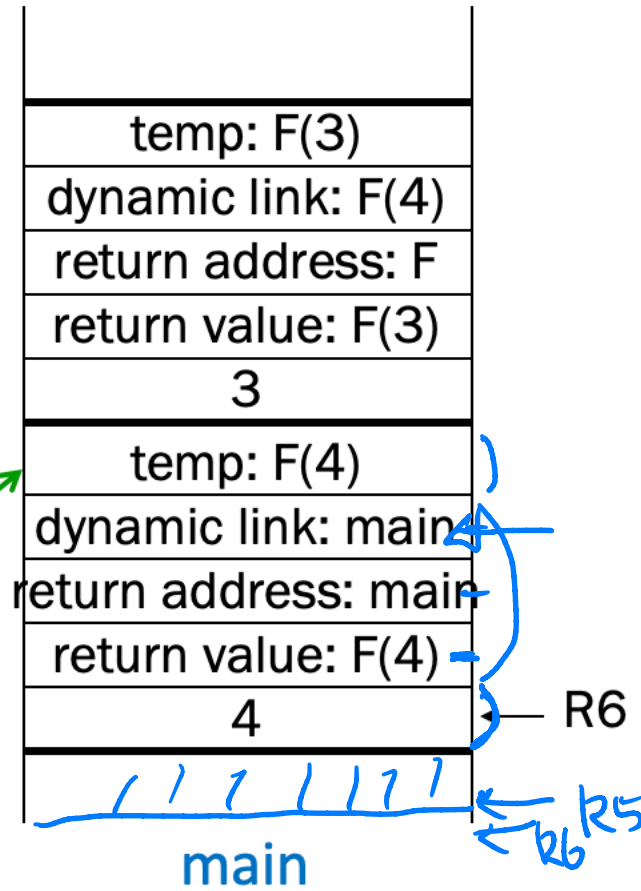
Factorial(3) calls
Factorial(2)



Factorial(2) calls
Factorial(1)

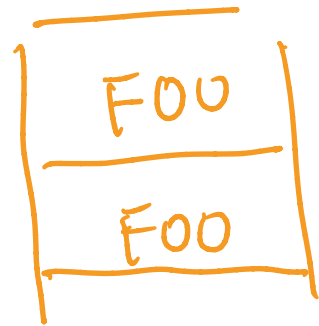


Activation Record



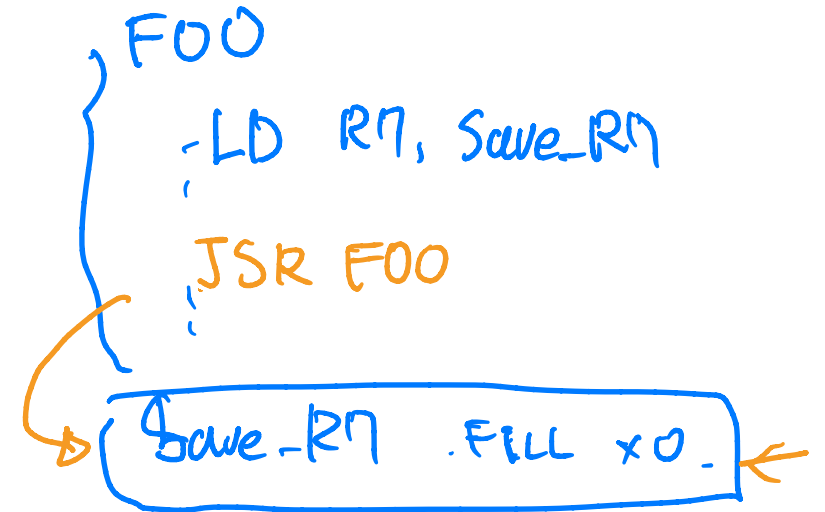
Compiler generates temporary variable to hold result.

C to LC3: Recursion with Run-time Stack



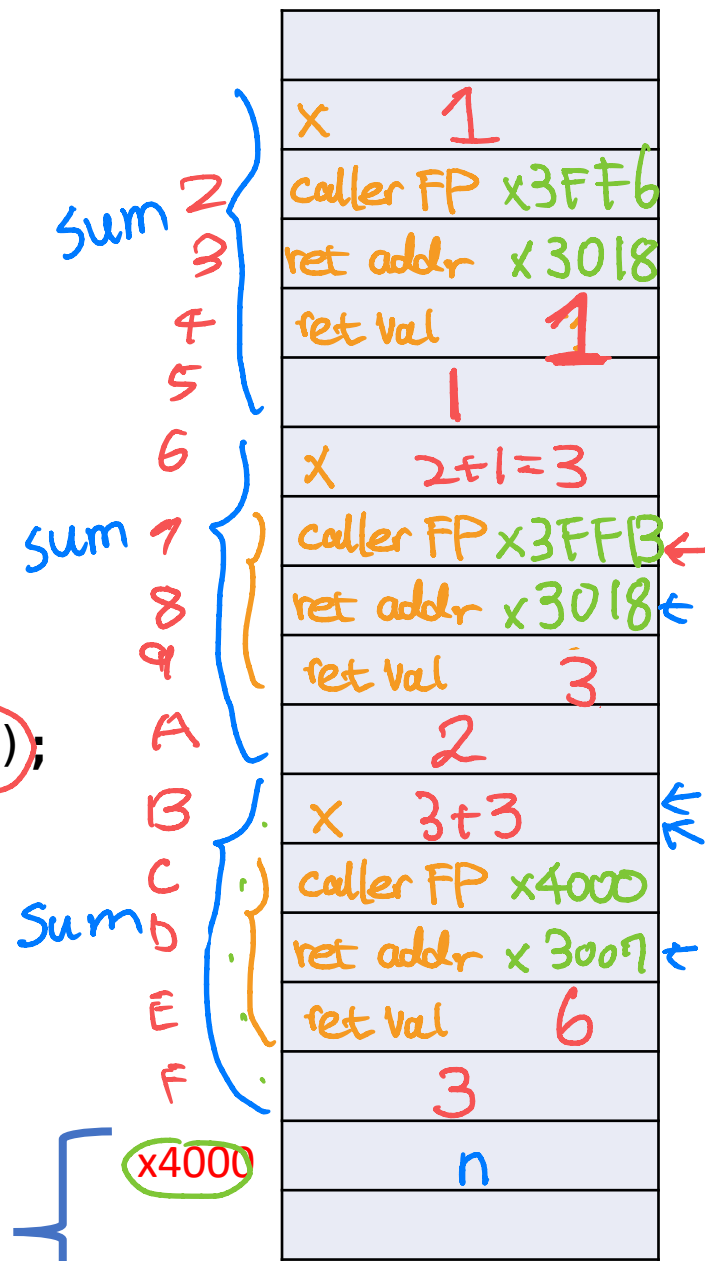
```
int main(){
    int n;
    n = sum(3);
}

int sum(int n){
    int x;
    if(n<=1){
        x = 1;
        return x;
    }
    else{
        x = n + sum(n-1);
        return x;
    }
}
```



```
int main(){
  int n;
  n = sum(3);
}
```

```
int sum(int n){
  int x;
  if(n<=1){
    x = 1;
    return x;
  }
  else{
    x = n + sum(n-1);
    return x;
  }
}
```



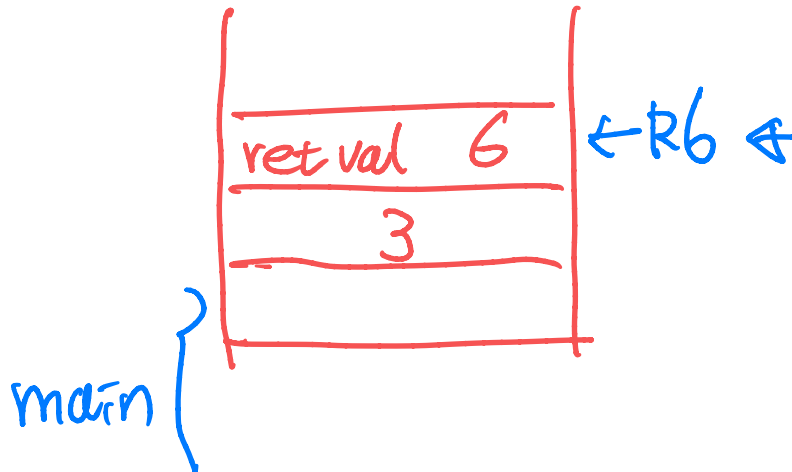
activation record of main

```
.ORIG x3000
; push argument ) push '3'
; ...
x3006 JSR SUM
; pop return value and tear-down
; ...
HALT
```

```
SUM
; push bookkeeping & local }
; ...
; if-else )
; ...
; push n-1 )
; ...
; recursive call
x3017 JSR SUM
; set return value n+sum(n-1)
; and tear-down
; ...
RET
```

C to LC3: Recursion

```
int main(){  
    int n;  
    n = sum(3);  
}
```



```
.ORIG x3000  
; init R6 and R5  
    LD R6, STACK_TOP  
    ADD R5, R6, #0  
; push argument  
    ADD R6, R6, #-1  
    AND R0, R0, #0  
    ADD R0, R0, #3  
    STR R0, R6, #0  
; call subroutine sum(n);  
    JSR SUM  
; pop return value from run-time stack  
    LDR R0, R6, #0  
    ADD R6, R6, #-2  
    STR R0, R5, #0  
    HALT
```

C to LC3: Recursion

```

int sum(int n){
    int x;
    if(n<=1){
        x = 1;
        return x;
    }
    else{
        x = n + sum(n-1);
        return x;
    }
}

```

SUM

```

; push callee's bookkeeping info onto the run-time stack
; . . .

; push n-1 onto run-time stack
    ADD R6, R6, #-1
    STR R2, R6, #0 ; R2 = n-1 } n-1
; call sum subroutine
    JSR SUM
; pop return value from run-time stack (R0 = sum(n-1))
    LDR R0, R6, #0 }
    ADD R6, R6, #1
; pop function argument from the run-time stack
    ADD R6, R6, #1
; add n by the return value (n+sum(n-1))
    LDR R1, R5, #4
    ADD R0, R0, R1
; store result in local var
    STR R0, R5, #0

```

