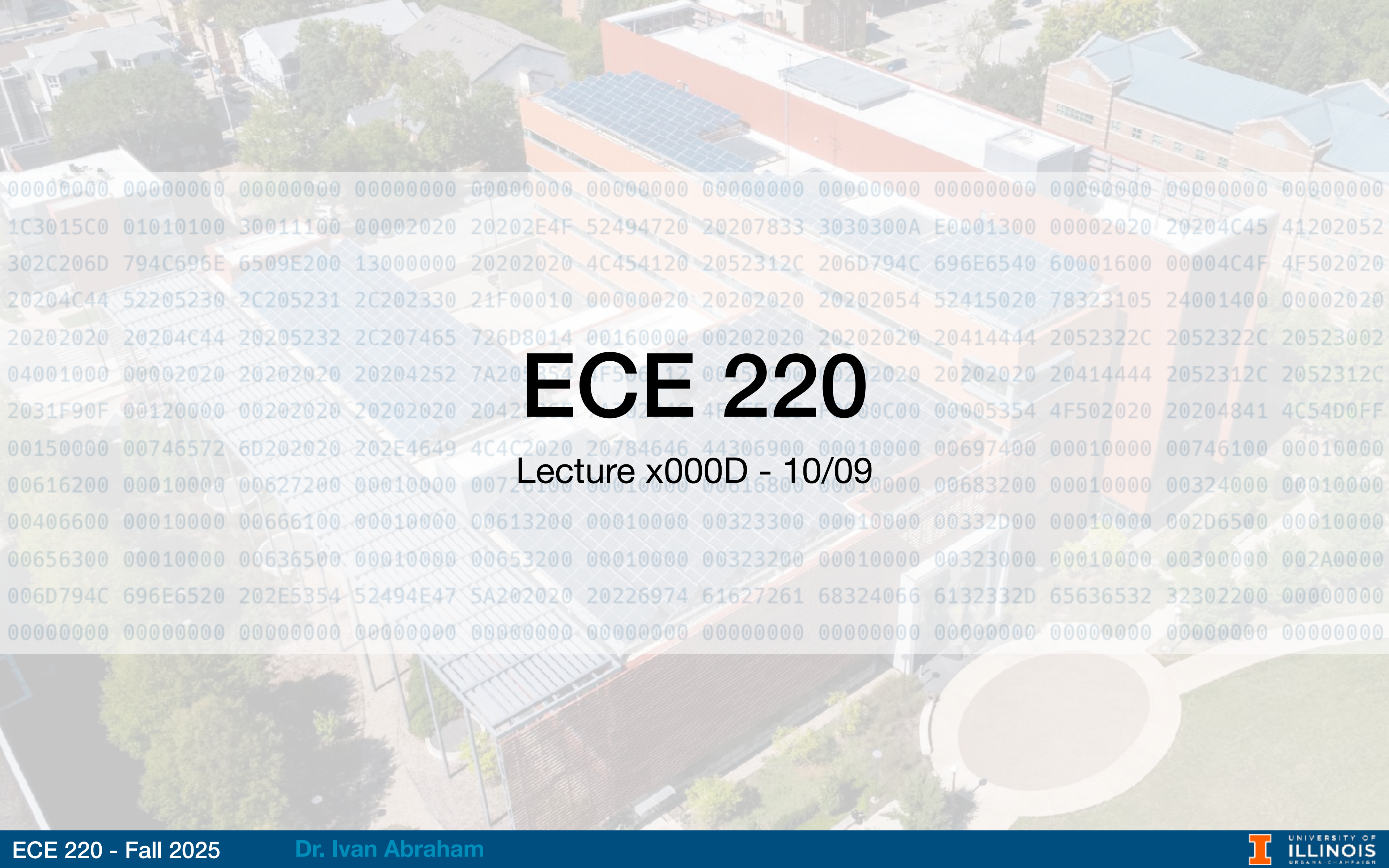




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# ECE 220

Lecture x000D - 10/09

# Recap

- Formal introduction to recursion
  - Factorial
  - Binary search
  - Towers of Hanoi
  - LC3 implementation
- Today: More recursion & problem solving
  - N - Queens problem
  - Maze solving
  - Exercise(s)

# Lesson objectives

- Understanding recursion vs. iteration tradeoff
  - Introduce *memoization* as a speedup technique
- Use recursion with backtracking to produce elegant solutions to problems
- Be able to implement recursion with backtracking to solve puzzles, problems, etc.

# Good recursion vs. bad recursion

- Consider the recursive Fibonacci function from last time.

```
long long fib(long long n){  
    long long sum;  
  
    if (n == 0 || n == 1)  
        return 1;  
    else {  
        sum = (fib(n-1) + fib(n-2));  
        return sum;  
    }  
}
```

- Let's do an activity
- Convert this function to an iterative version.
- Compare run times.
- Can recursion be made faster?

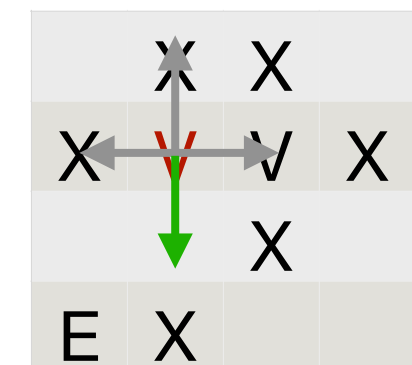
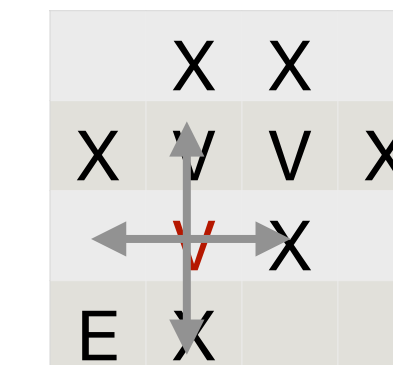
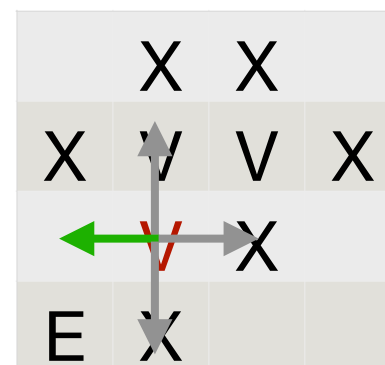
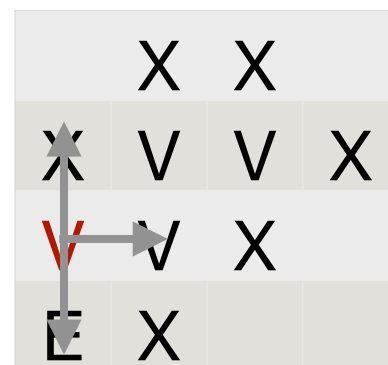
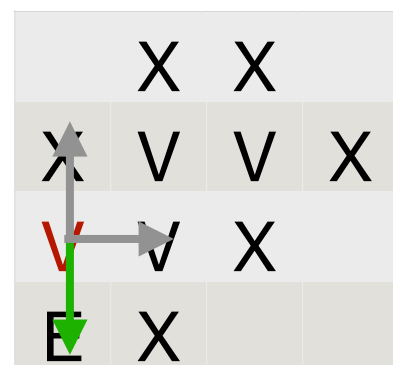
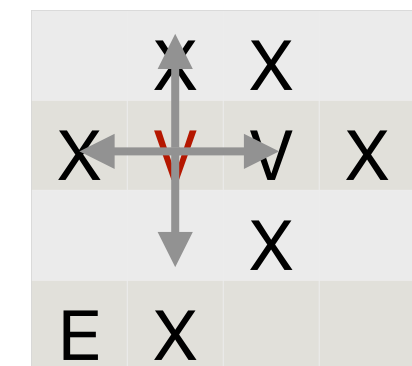
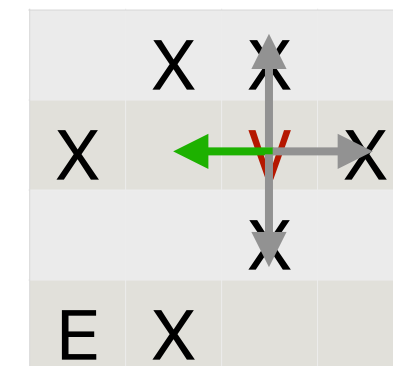
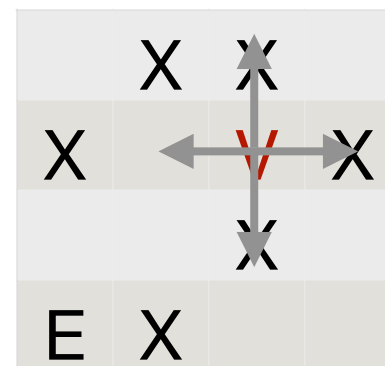
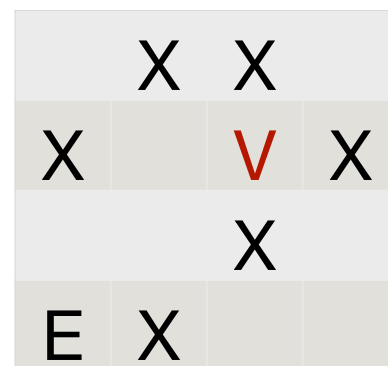
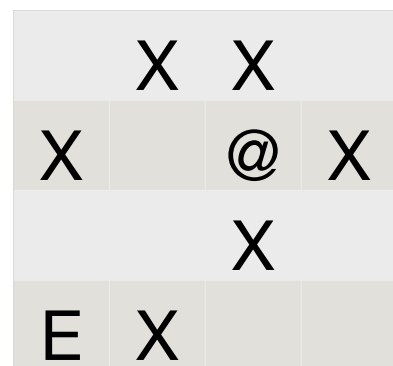
# Solving a maze

- We represent a maze by a 2D grid of size  $N \times M$
- Walls are marked with X and the exit with E.
- Given starting point  $(i, j)$  marked with @, find a path to E (if it exists).
  - Do not go outside grid
  - Avoid going around in circles.
  - Mark valid path with P.
- Write a recursive function `int ExitMaze` solving problem.

|   |   |   |   |
|---|---|---|---|
|   | X | X |   |
| X |   | @ | X |
|   |   | X |   |
| E | X |   |   |

# Solving a maze

**Strategy:** Mark current cell as visited and explore solution space. Exploration defined by four possible moves (U, D, L, R).



# Solving a maze

- Recursion needs base case. What should be the base case?
  - Found exit (return “good”) **OR** hit X **or** hit V **or** out-of-bounds (return “bad”)
  - Let **xpos** and **ypos** be the *row* and *column* index.

```
if (xpos < 0 || xpos >= MAZE_HEIGHT || ypos < 0 || ypos >= MAZE_WIDTH)
    return 0;

if (maze[xpos][ypos] == 'E')                // Found the Exit!
    return 1;

if (maze[xpos][ypos] != ' ')                // Space is not empty (possibly X or V)
    return 0;
```

# Solving a maze

- What should be the recursive call?
  - Go down, up, left or right.
  - `int ExitMaze(maze, xpos, ypos)`
    - Function exploring the solution space.

```
// Go Down
if (ExitMaze(maze, xpos + 1, ypos)) {
    maze[xpos][ypos]='P';
    return 1;
}
```

```
// Go Right
if (ExitMaze(maze, xpos, ypos + 1)) {
    maze[xpos][ypos]='P';
    return 1;
}
```

```
// Go Up
if (ExitMaze(maze, xpos - 1, ypos)) {
    maze[xpos][ypos]='P';
    return 1;
}
```

```
// Go Left
if (ExitMaze(maze, xpos, ypos - 1)) {
    maze[xpos][ypos]='P';
    return 1;
}
```

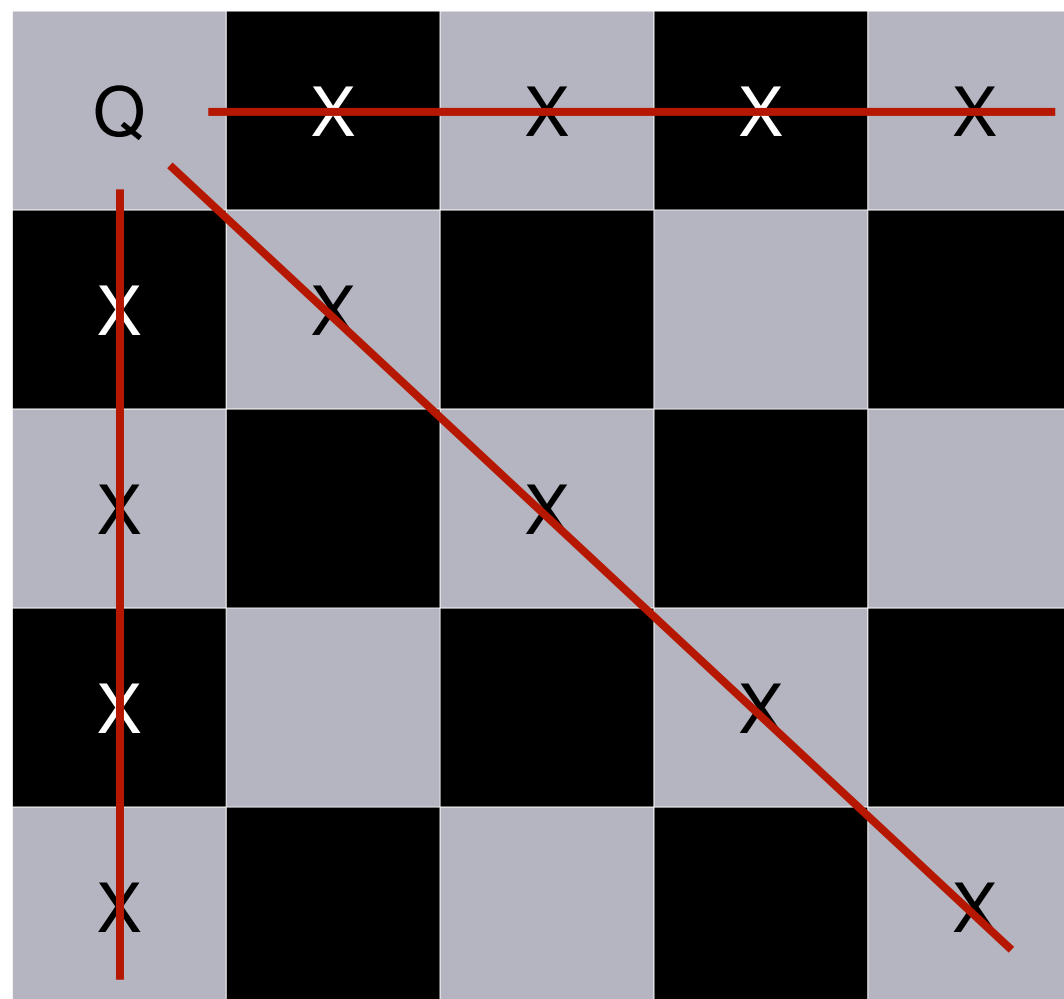
# Exercise

- There is an `ExitMaze` function on Github which I tested to work.
- Modify it by adding a `main` function, board definition and try it on this maze.

|   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|
|   | X |   |   |   |   | X |
|   | X |   |   | X |   |   |
|   |   |   | @ | X | X | X |
|   | X |   |   |   |   |   |
|   | X | X | X | X |   |   |
| X |   |   | E | X |   |   |
| X | X |   | X | X |   | X |
|   |   |   |   |   |   | X |

# N - Queens Problem

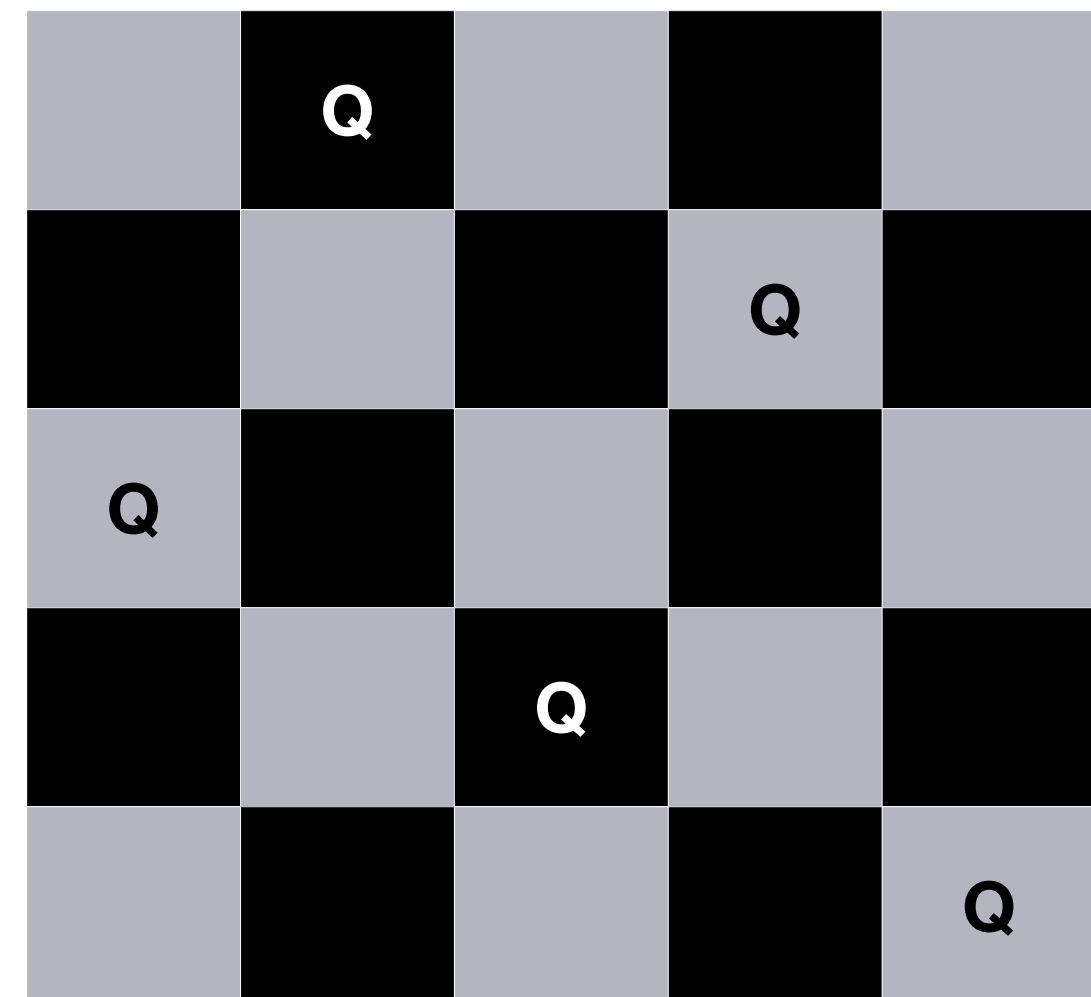
- In chess, a Queen can attack another piece within its line of sight as long as that piece is in the same: **row**, **column** or **diagonal**.



- Question:** Given an  $N \times N$  grid, is it possible to place  $N$  Queens in the grid so that no two Queens can attack each other ?
- Answer:** Yes.

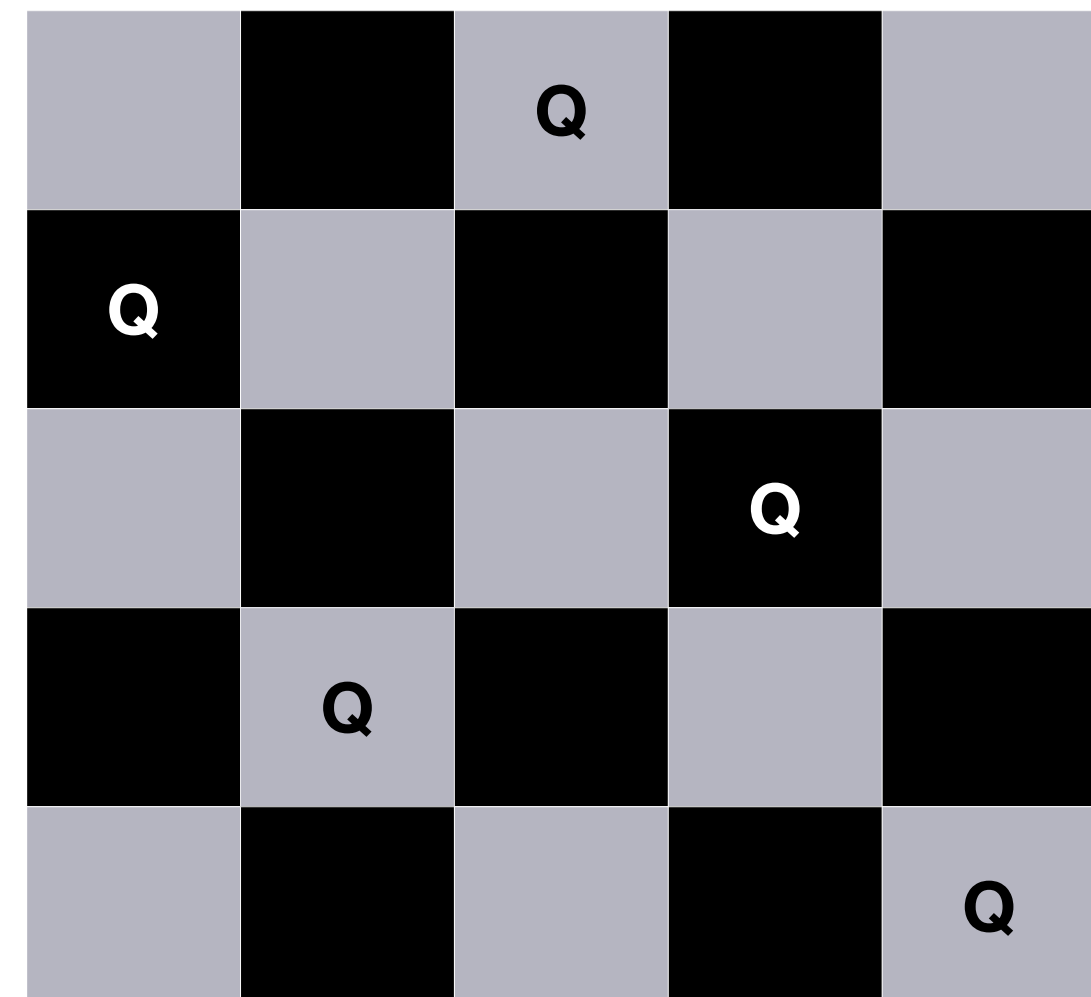
# N - Queens Problem

- Here is a possible solution for the 5 x 5 grid.
  - Not unique
- Can we make the computer solve it for any given N?
  - Solution: Recursion with *backtracking*.



# N - Queens Problem

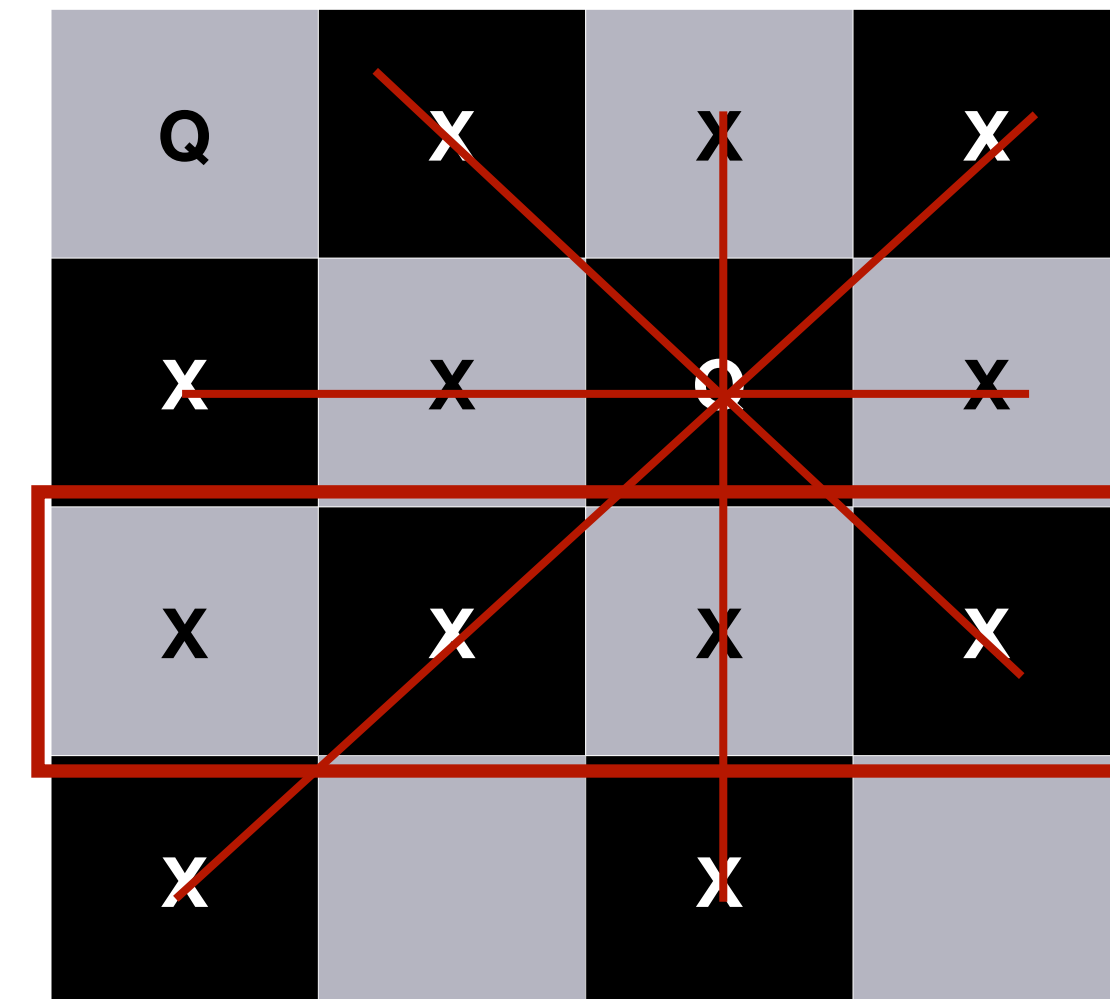
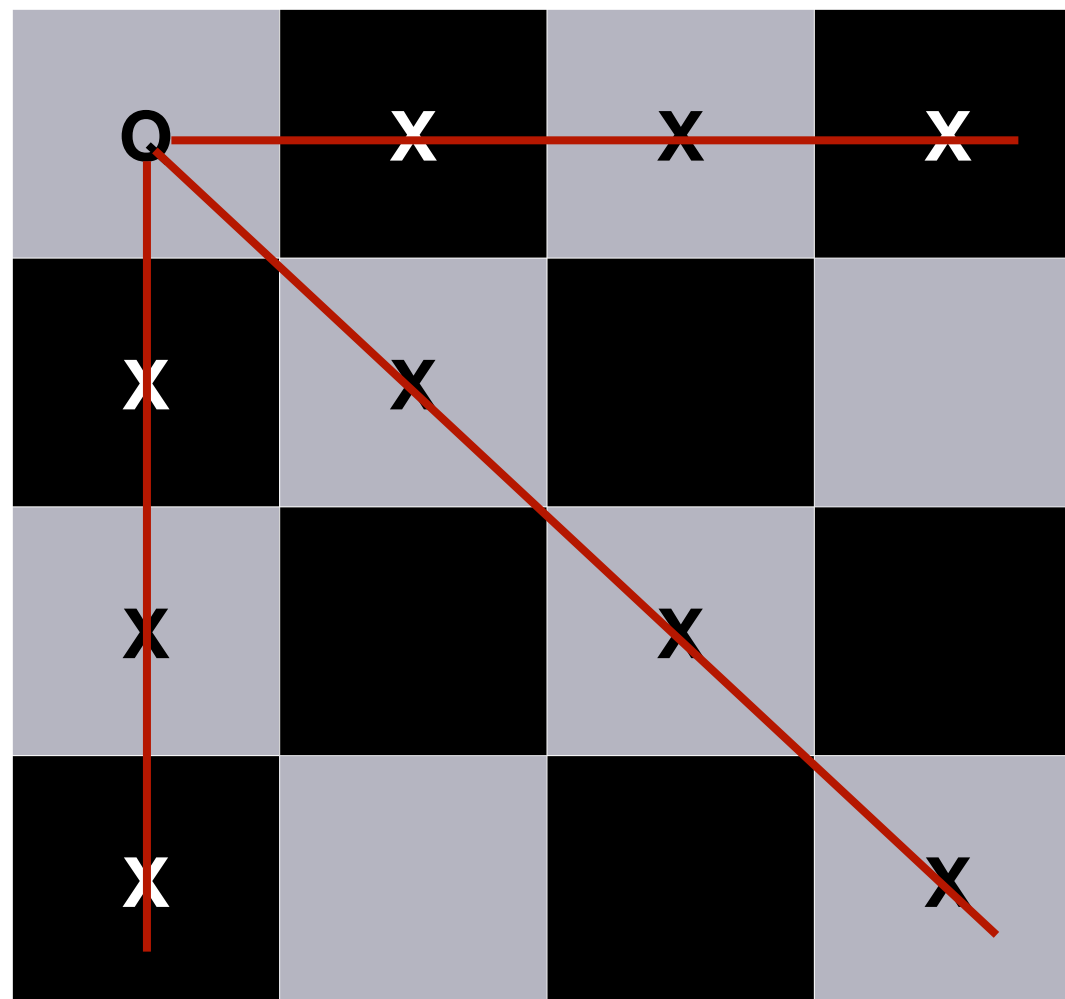
- Here is a possible solution for the 5 x 5 grid.
  - Not unique
- Can we make the computer solve it for any given N?
  - Solution: Recursion with *backtracking*.



# N - Queens Problem

- **Back-tracking:** Make a choice and search the solution space. If solution space is empty, return and make a different choice.

Choice #1



Choice #1.1

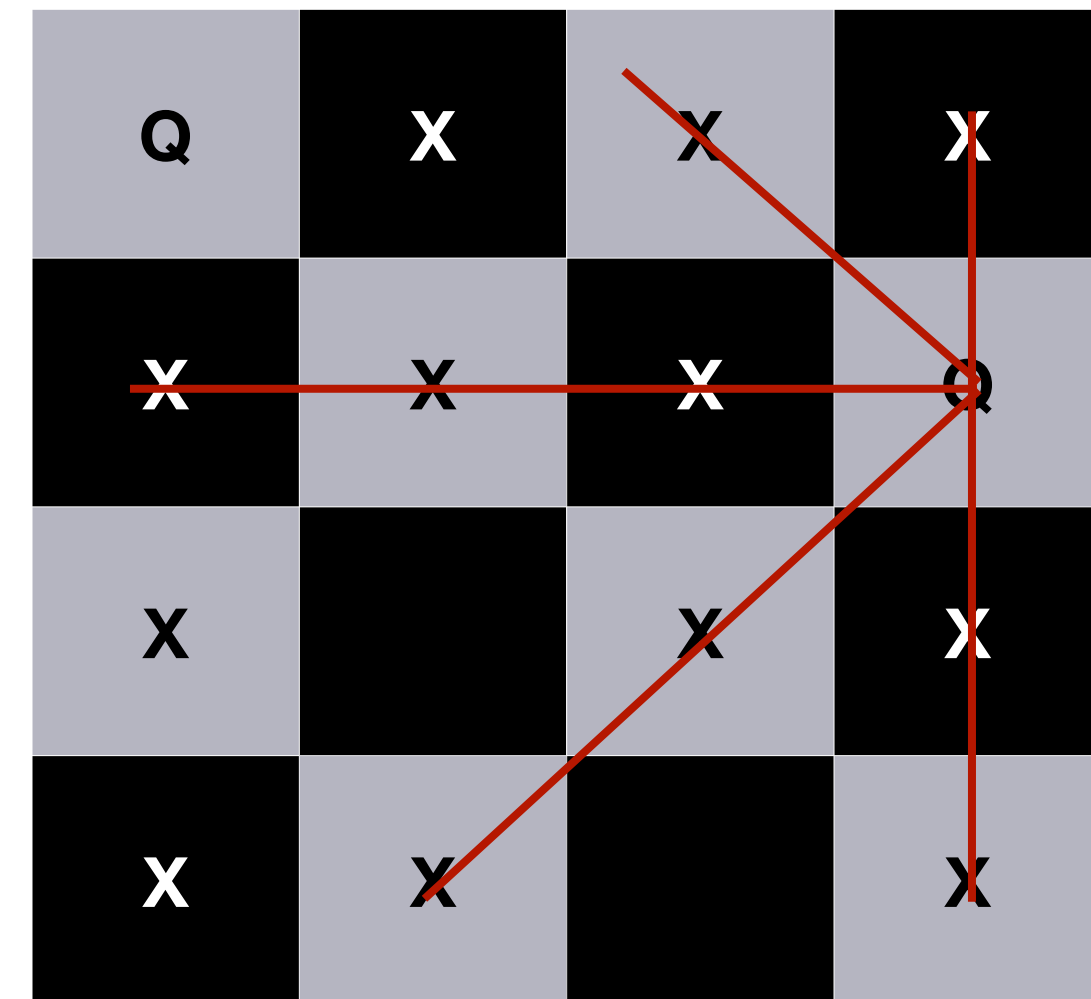
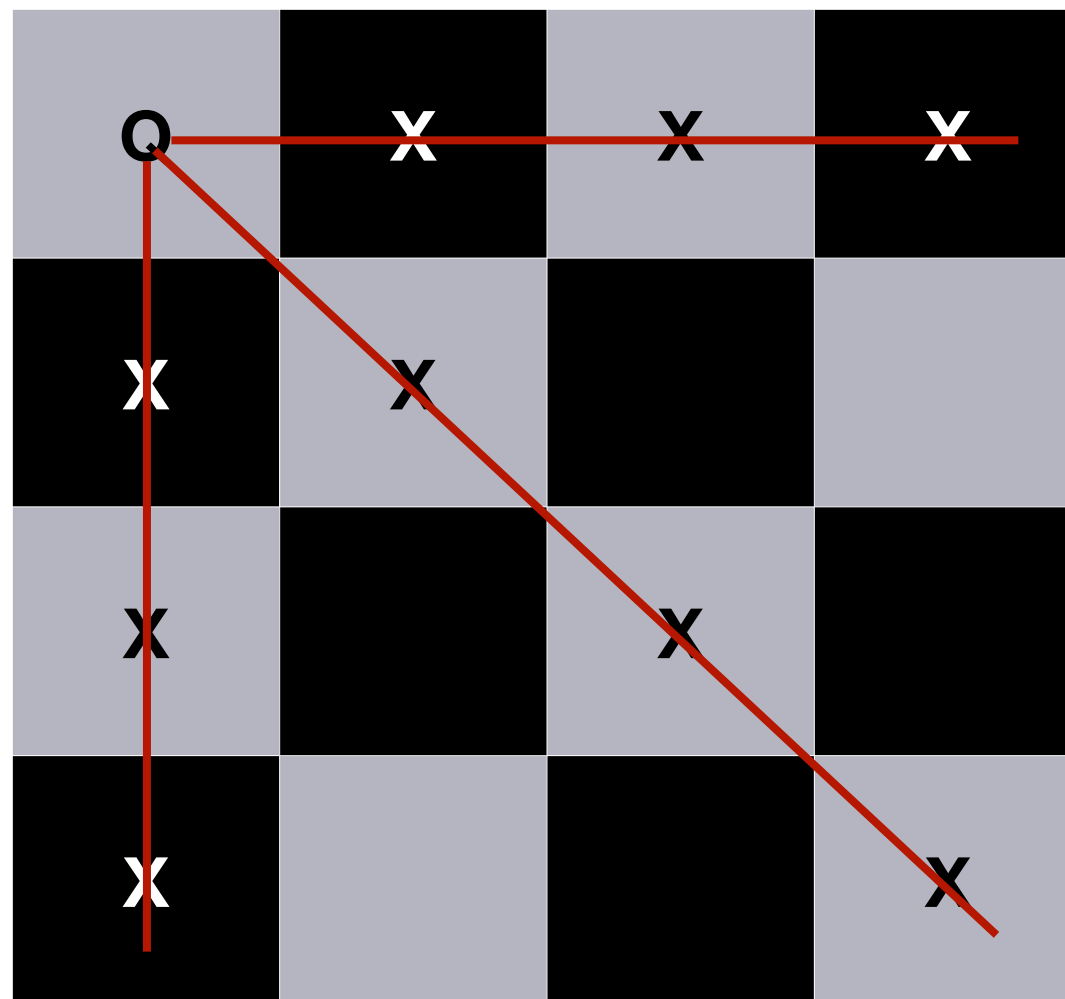
Not a solution!



# N - Queens Problem

- **Back-tracking:** Make a choice and search the solution space. If solution space is empty, return and make a different choice.

Choice #1

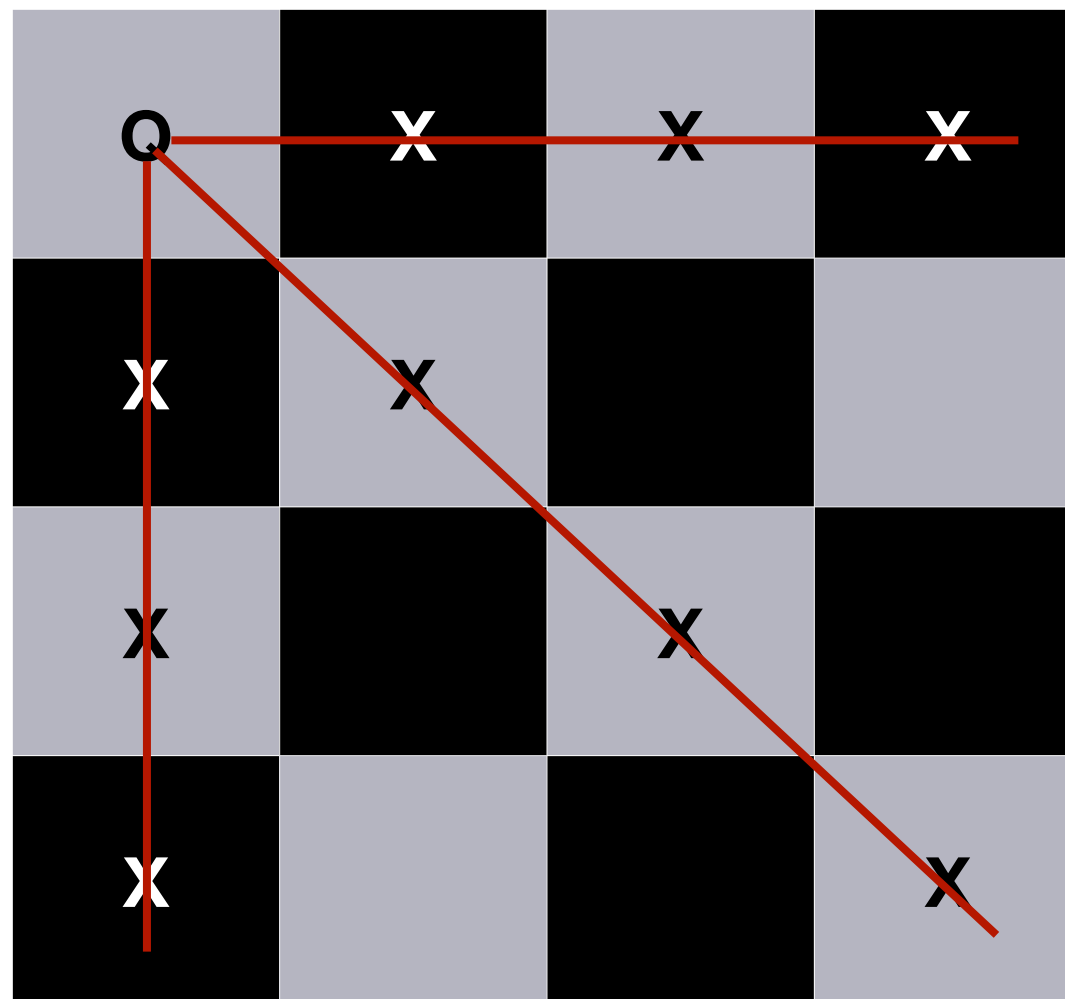


Choice #1.2

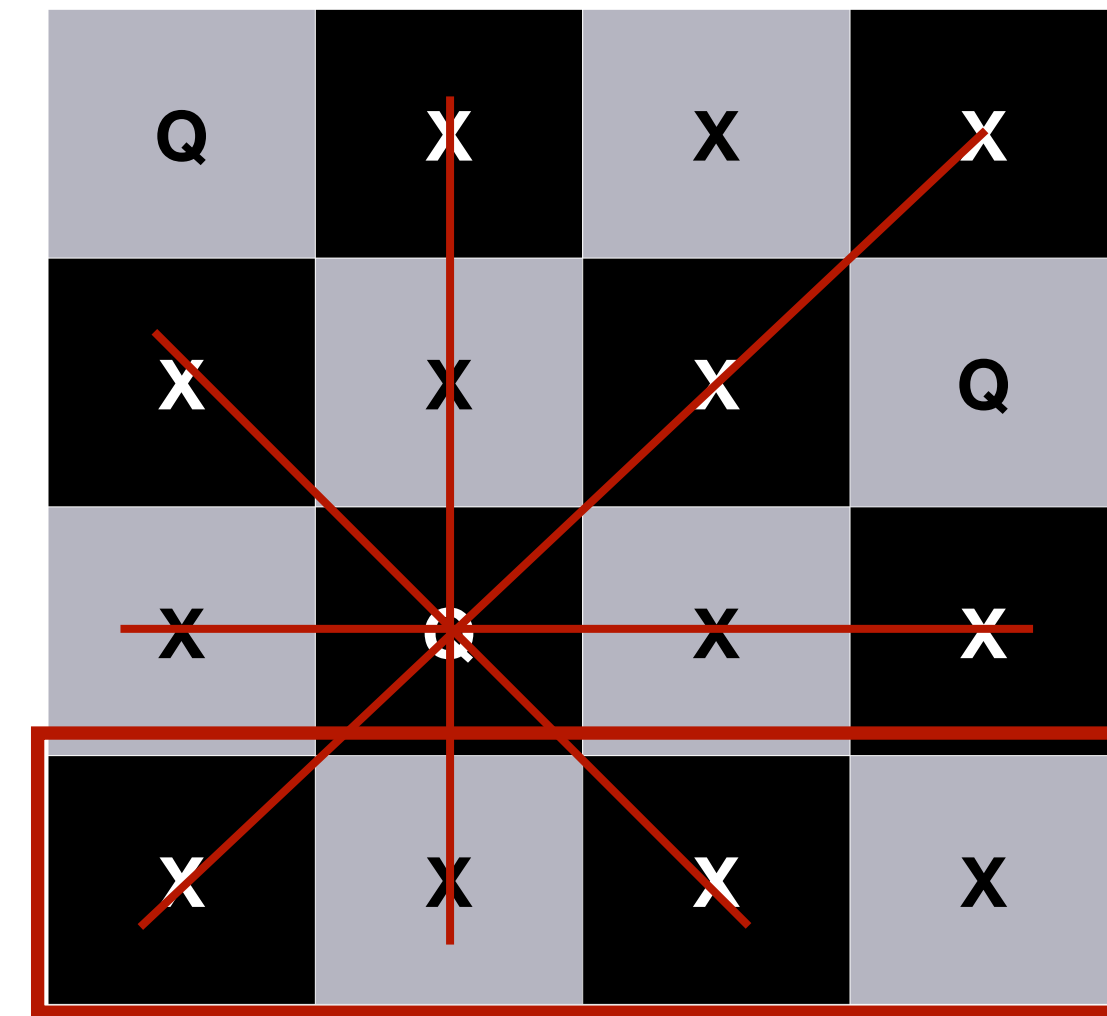
# N - Queens Problem

- **Back-tracking:** Make a choice and search the solution space. If solution space is empty, return and make a different choice.

Choice #1



Choice #1.2.1



Choice #1.2

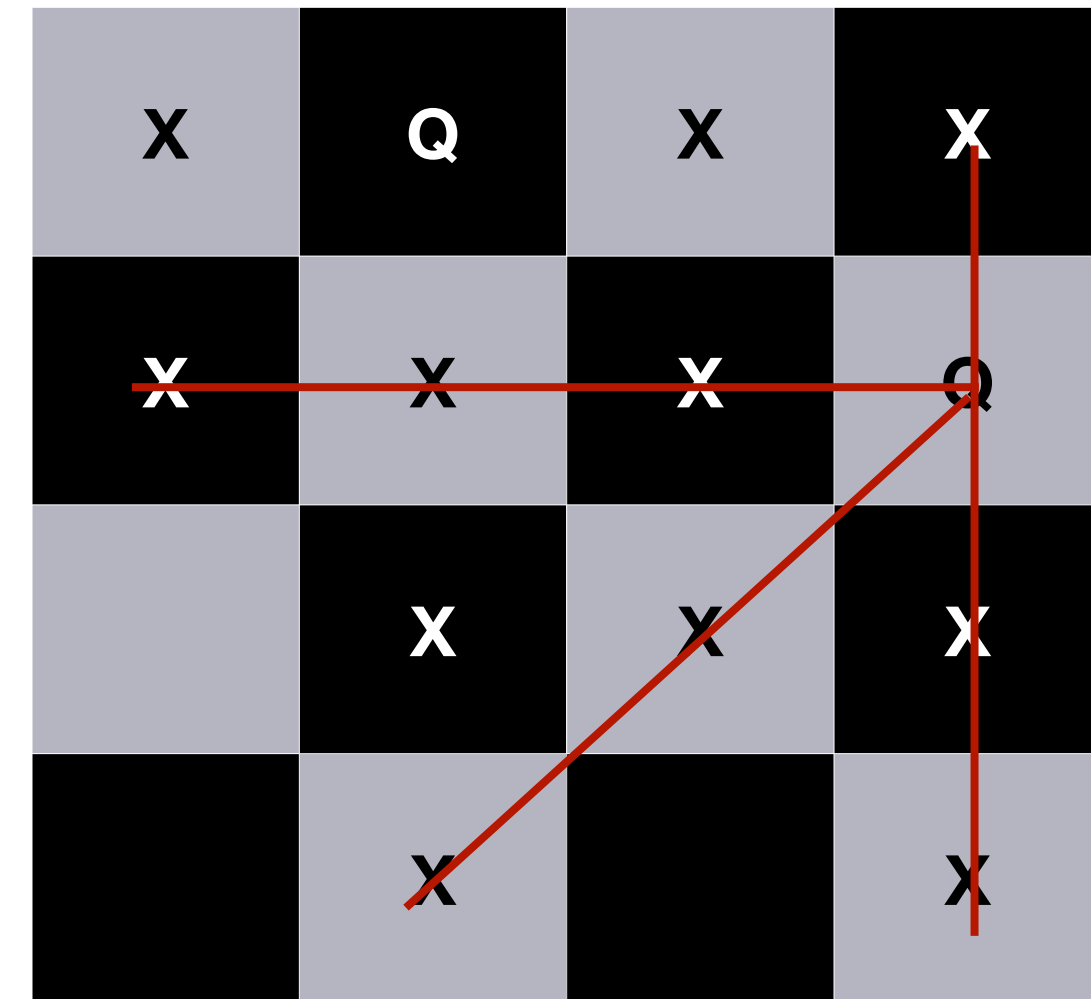
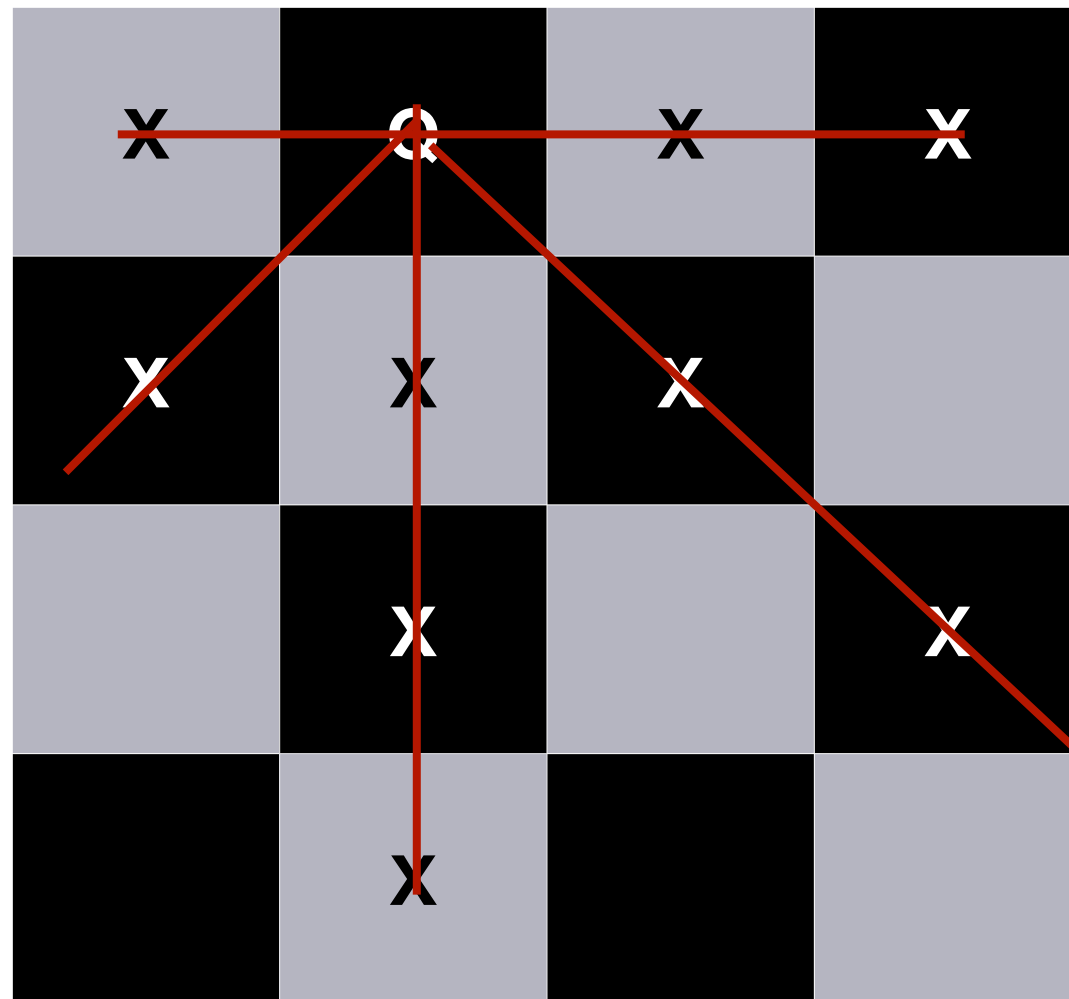
Not a solution!



# N - Queens Problem

- **Back-tracking:** Make a choice and search the solution space. If solution space is empty, return and make a different choice.

Choice #2

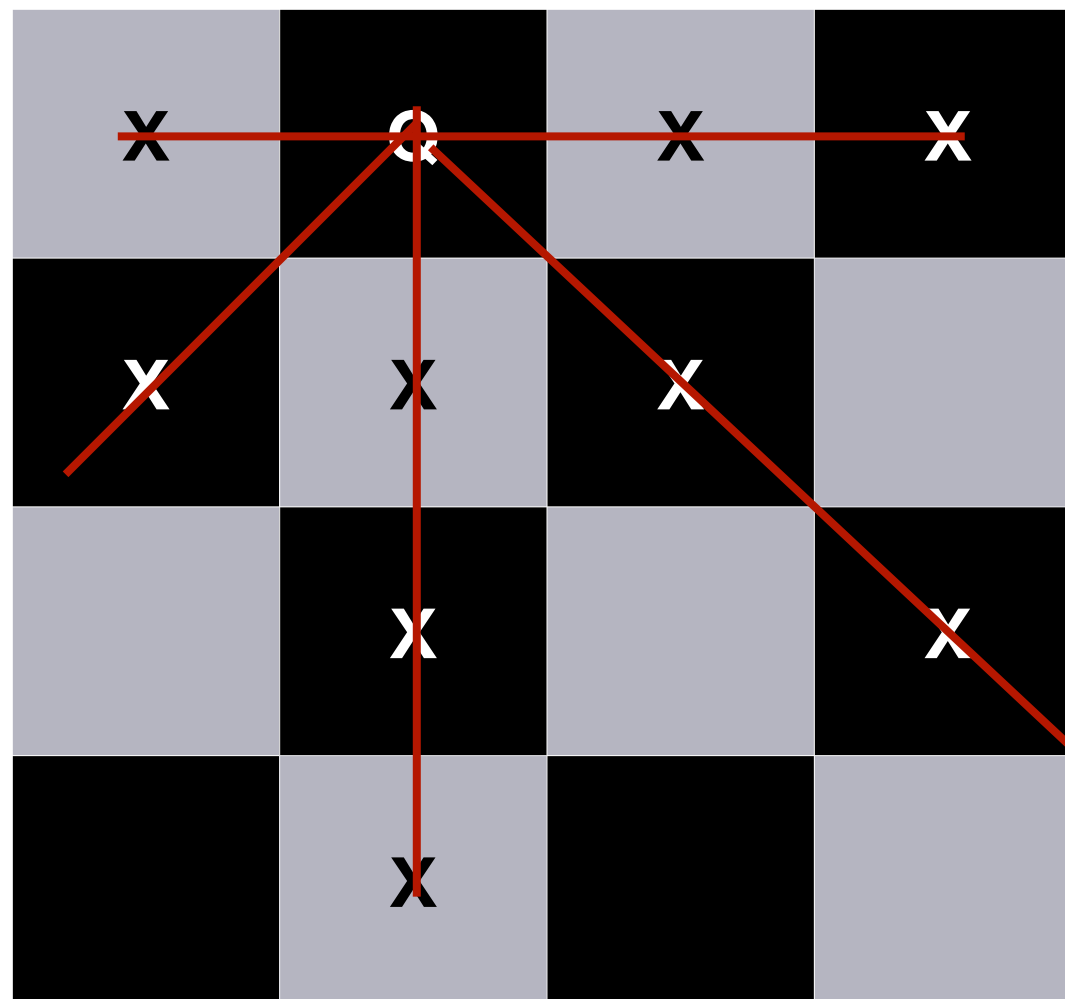


Choice #2.1

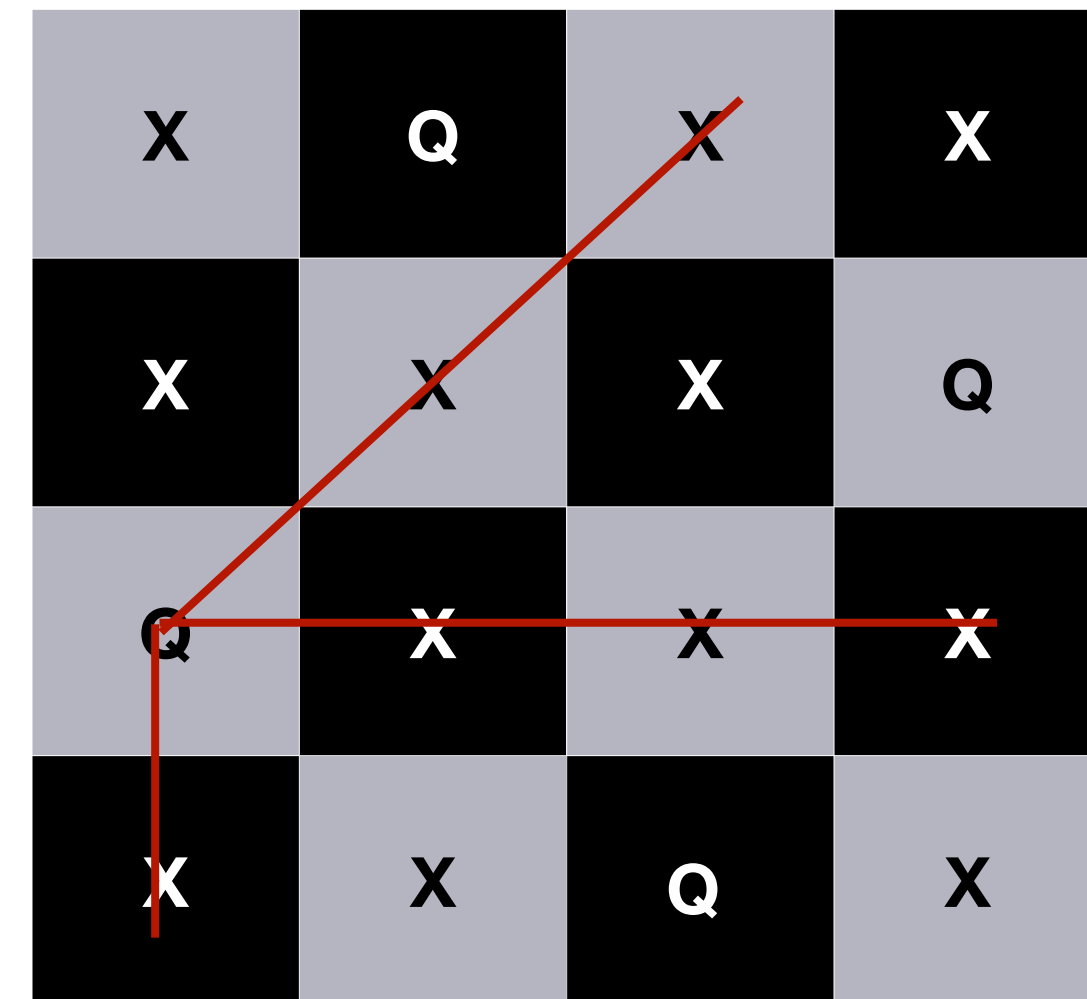
# N - Queens Problem

- **Back-tracking:** Make a choice and search the solution space. If solution space is empty, return and make a different choice.

Choice #2



Choice #2.1.1



Choice #2.1

Valid solution



Choice #2.1.1.1

# N - Queens implementation?

- **Question:** Can we set this up as a recursive problem?
  - What is the action/sub-problem that we want to repeat?
    - Placing a Queen in a row
      - If not successful how do we *backtrack*?
        - Undo placing a queen
  - How do we know we have reached an end case?
    - No more rows to fill.

# N - Queens set-up?

- We represent the configuration space with a grid.
  - We will denote with digit **zero** an empty spot (maybe safe or unsafe, but its unoccupied).
  - We will denote with the digit **one** a space occupied by a queen.
  - We will fill in rows starting with the first row and proceeding downward.

# N - Queens implementation

```
int is_safe(int board[N][N], int rnum, int cnum);

/*Function places a queen in row rnum */
int place_queen(int board[N][N], int rnum){
    if ( ) // Finished all rows
        return 1; // Found a solution
    else{
        // Iterate over possible columns
        for(int cnum=0; ; cnum++){
            if (is_safe( )==1){
                board[rnum][cnum] = 1; // Place a queen there
                // Update row number and recurse
                if ( ==1)
                    return 1;
                else // Hit a road block down the line
                    // Remove queen
            } // Try next column along row
        } // For loop finished without hitting a return
        // Solution doesn't exist.
    }
}
```

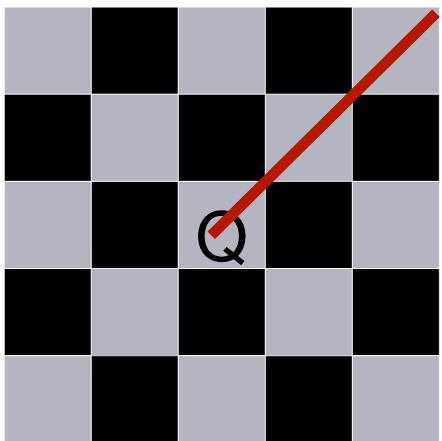
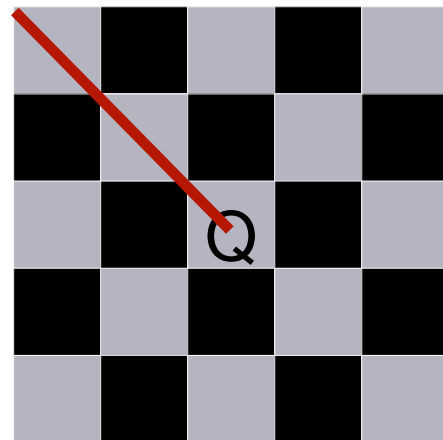
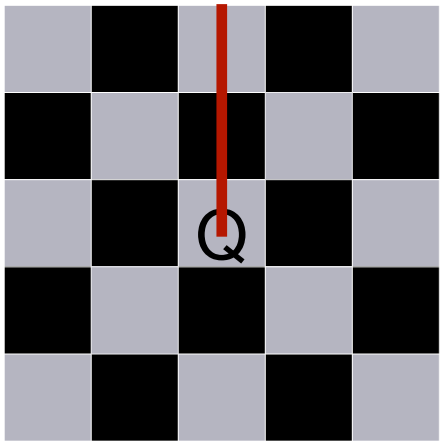
**is\_safe** checks whether it is possible to place a queen on position (**rnum**, **num**) given the configuration of the board at some given time. It returns 1 if safe or 0 if unsafe

**place\_queen** fills the board with a valid solution and returns 1 or returns 0 if no solution found.

# Is it safe/unsafe?

- On the N-th row when we place a queen on a square  $(i, j)$  what do we need to check?
  - Are we in the line of sight (LOS) for any previous Queen.
  - We are in LOS if
    - The column  $i$  contains any Queen OR
    - The diagonals to the top-left of  $(i, j)$  contains a Queen OR
    - The diagonals to the top-right of  $(i, j)$  contains a Queen
    - What about diagonals to the bottom left or bottom right?

# Is it safe/unsafe?



```
int is_safe(int board[N][N], int row, int col){
    int i, j;
    for (                ){ //Check along column
        if (board[i][col]==1)
            return 0;
    }
    // Check diagonal to upper left
    for (                ; i>=0 && j>=0; i--, j--){
        if (board[i][j] == 1)
            return 0;
    }
    // Check diagonal to upper right
    for (i=row-1, j=col+1;                ;                ){
        if (board[i][j]==1)
            return 0;
    }
    return 1;
}
```

# Exercise for fun outside lecture

- Exercise for the *curious/mighty/brave*:
  - Modify the source of `queens.c` so that it keeps a static variable to keep track of the recursive calls.
  - Varying `N`, generate a plot of `N` vs number of recursive calls. Try `N=4, 5, ..., 15`. What kind of growth is it?

# Exercise - practice, practice, p....

- You have a pile of wood sticks with 3 different lengths: 3, 7, and 10 feet. You want to connect them and make an X-feet long stick using **at most** 10 sticks.
- To make a stick 33 feet long you can do:
  - $4 \times 3F + 3 \times 7F$  ✓
  - $11 \times 3F$  ✗
- Use recursion with backtracking to find a solution

# Exercise

```
#define N 10 // Number allowed
#define M 3  // Types of lengths

// Implement this function
// solution[N]: stores the solution
// idx: index for the solution matrix
// total: remaining length
int solve(int solution[N], int idx, int total);

const int set[M] = {3,7,10};

int main(){
    int solution[N] = {0, 0, 0, 0, 0, 0, 0, 0, 0, 0};
    int total;
    printf("Enter total length: ");
    scanf("%d", &total);
    // Write your code here
}
```

# Time permitting

- Using the `gdb` debugger with gdb dashboard.