

### Homework 3 (due Nov 10 Wednesday 5pm)

**Instructions:** see previous homework.

1. [50 pts] Given a set  $P$  of  $n$  points in  $\mathbb{R}^d$  for a constant dimension  $d$ , a  $k$ -hop  $c$ -spanner is a subgraph  $H$  of the complete graph with vertex set  $P$  such that for every  $p, q \in P$ , there exists a path from  $p$  to  $q$  in  $H$  with at most  $k + 1$  edges and total length at most  $c\|p - q\|$  (in the Euclidean metric). Describe an algorithm to construct a 1-hop  $c$ -spanner with  $O(n \log n)$  edges for some constant  $c$ .

[Hint: one approach is to use shifted, balanced quadtrees; another approach is to use shifted Z-ordering.]

[Bonus: make  $c$  arbitrarily close to 1. Also, how many edges for a 2-hop spanner?]

2. [50 pts] In the *discrete 1-median* problem, we are given a set  $P$  of  $n$  points in 2D, and want to find a point  $q^* \in P$  that minimizes the sum of Euclidean distances  $\sum_{p \in P} \|p - q^*\|$ . (Note that  $q^*$  must be in  $P$ .) There is a naive  $O(n^2)$ -time exact algorithm. Describe a more efficient approximation algorithm to solve the problem.

[Note:  $O(n \log n)$  run time and constant approximation factor would be fine, though  $O(n)$  time and  $1 + \varepsilon$  approximation would be even better (and worth bonus points).]

[Hint: one approach is to change the metric...]