

# ★ External Regret Model.

Decision / player / agent  
Market

VS

Adversary

$A$ : set of actions.

★  $t=1, \dots, T$

1. Agent chooses  $p^t \in \Delta(A)$ .

2. Adversary (after looking at  $p^t$ ) chooses  $c^t: A \rightarrow [0, 1]$

3. Agent draws  $a^t \sim p^t$  (at w.p.  $p^t(a^t)$ )

incurs cost  $c^t(a^t)$

learns the entire  $c^t$ .

Learns  $c^t(a^t) \equiv$  Bandit model.  
Guarantees are similar w/ some loss

Q: How bad  $\sum_{t=1}^T c^t(a^t)$  (cost to the agent)

compared to "best possible"?

$$\sum_{t=1}^T \min_{a \in A} c^t(a)$$

eg.  $A = \{1, 2\}$

$t = 1, \dots, T$

After looking at  $p^t$  adversary chooses  $c^t$  as follows  
 $c^t = \begin{cases} (1, 0) & \text{if } p^t(1) \geq p^t(2) \\ (0, 1) & \text{o.w.} \end{cases}$

Expected cost of the agent =  $\sum_{t=1}^T \max\{p^t(1), p^t(2)\} \geq \frac{T}{2}$

"best possible" = 0

★ Compare to the "best-action-in-hindsight".

argmin  $a \in A \sum_{t=1}^T c^t(a)$

★ Def Algo "J" is no-regret if it's

Time-avg-regret =  $\frac{1}{T} \left[ E[\text{cost of the algo}] - \text{cost of Best-action-in-hindsight} \right]$

cost of Best-action-in-hindsight

lim  $T \rightarrow \infty$

0

(Bad eg.)  $A = \{1, 2\}$

Oblivious adversary:

$c^t(1)$	$c^t(2)$	
1	0	w.p. $\frac{1}{2}$
0	1	w.p. $\frac{1}{2}$

$$\text{Expected cost - regret in round } t = \frac{p^t(1)}{2} + \frac{p^t(2)}{2}$$

$$= \frac{1}{2}$$

$$\mathbb{E}[\text{cost - regret}] = \frac{1}{2} \left[ \sum_{t=1}^T c^t(1) + \sum_{t=1}^T c^t(2) \right]$$

$$\mathbb{E}[\text{cost of best action}] = \min \left\{ \begin{array}{l} \# \text{ times } (1,0) \\ \# \text{ times } (0,1) \end{array} \right\}$$

$$= \frac{T}{2} - \Theta(\sqrt{T})$$

↑ variance.

$$T\text{-A-Regret} = \frac{1}{T} \left[ \frac{T}{2} - \frac{T}{2} + \Theta(\sqrt{T}) \right]$$

$$= \Theta\left(\frac{1}{\sqrt{T}}\right) \rightarrow 0 \text{ as } T \rightarrow \infty$$

↓ extend

n-action  $\text{Regret} = \Theta\left(\sqrt{\frac{\ln n}{T}}\right)$

n-action

regret

$$\sqrt{\frac{\ln n}{T}}$$

Thm:  $\exists$  no-regret Algorithm with regret  $\leq O\left(\sqrt{\frac{\ln n}{T}}\right)$

(If worst regret =  $\epsilon$  then  $T = \frac{\ln n}{\epsilon^2}$ )

⊗ Multiplicative - Weight - Update (MWU)

OR Hedge

OR expert advice

Intuition: Increase prob. of "good actions"

$\equiv$  aggressively punish bad actions.

★ Player's MWU Algo:

★  $w^t(a) = 1, \forall a \in A$  (init.)

★ For  $t = 1, \dots, T$

①  $p^t$  prop. to  $w^t$

$$p^t(a) = \frac{w^t(a)}{\Gamma^t}, \forall a \in A$$

where  $\Gamma^t = \sum_{a \in A} w^t(a)$

$$1 = \sum_{a \in A} w^t(a)$$

② After observing  $c^t: A \rightarrow [0, 1]$ , update

$$\forall a \in A, \quad w^{t+1}(a) = w^t(a) (1 - \epsilon)^{c^t(a)}$$

$\epsilon \rightarrow 0$  :  $w^{t+1}(a) \sim w^t(a)$  explore

$\epsilon \rightarrow \frac{1}{2}$  :  $w^{t+1}(a) = \text{B.R.}$  exploit.

$\epsilon \in [0, \frac{1}{2}]$  pick later.

pf : 
$$r^t = \sum_{a \in A} w^t(a)$$

Goal: Relate  $\mathbb{E}[\text{cost of agent}] = \sum_{t=1}^T \underbrace{\sum_{a \in A} p^t(a) c^t(a)}_{\text{expected cost in round } t}$

$\uparrow$

$$\mathbb{E}[\text{cost of best action}] = \underset{a \in A}{\text{argmin}} \underbrace{\sum_{t=1}^T c^t(a)}_{\text{OPT}}$$

Via  $r^t$ .

$$r^T = \sum_{a \in A} w^T(a) \geq w^T(a^*) = \sum_{t=1}^T c^t(a^*)$$

$$\begin{aligned}
 \Gamma^t &= \sum_{a \in A} w^t(a) = w^t(a^*) \prod_{t=1}^t (1-\epsilon) \\
 &= w^t(a^*) \frac{1}{\sum_{t=1}^t c^t(a^*)} \\
 &= 1 \cdot (1-\epsilon)^t \\
 &= (1-\epsilon)^t \text{ opt.} \rightarrow \textcircled{1}
 \end{aligned}$$

- Expected cost of agent in time  $t$ .

$$V^t = \sum_{a \in A} p^t(a) c^t(a) = \sum_{a \in A} \frac{w^t(a)}{\Gamma^t} c^t(a) \rightarrow \textcircled{2}$$

$$\begin{aligned}
 \Gamma^{t+1} &= \sum_{a \in A} w^{t+1}(a) \\
 &= \sum_{a \in A} w^t(a) (1-\epsilon) c^t(a)
 \end{aligned}$$

$\left( \begin{array}{l} \epsilon \in [0, 1/2] \\ c^t \in [0, 1] \end{array} \right)$

$$\begin{aligned}
 &\leq \sum_{a \in A} w^t(a) (1 - \epsilon c^t(a)) \\
 &= \Gamma^t \left( \frac{\sum_{a \in A} w^t(a)}{\Gamma^t} - \epsilon \sum_{a \in A} \frac{w^t(a)}{\Gamma^t} c^t(a) \right)
 \end{aligned}$$

$$= \underbrace{\Gamma^t}_{t-1} (1 - \epsilon \nu^t) (1 - \epsilon \nu^t)$$

$$\leq \underbrace{\prod_{t=1}^T (1 - \epsilon v^{t-1})}_{\text{expected cost of agent in time } k} (1 - \epsilon v^T)$$

$$\prod_{t=1}^{t+1} \leq \prod_{k=1}^t (1 - \epsilon v^k)$$

expected cost of agent in time k.

$$(1 - \epsilon)^{\text{OPT}} \stackrel{\text{①}}{\leq} \underbrace{\prod_{t=1}^T}_{\text{②}} (1 - \epsilon v^t) \leftarrow$$

Apply  $\ln$

$$\text{OPT} \ln(1 - \epsilon) \leq \ln n + \sum_{t=1}^T \ln(1 - \epsilon v^t)$$

$$\ln(1 - x) = -x - \frac{x^2}{2} - \frac{x^3}{3} - \dots$$

$$-x - x^2 \leq \ln(1 - x) \leq -x$$

$$\frac{\text{OPT} (1 - \epsilon + \epsilon^2)}{\epsilon} \geq \frac{-\ln n}{\epsilon} + \sum_{t=1}^T \frac{(1 - \epsilon v^t)}{\epsilon}$$

$$\Rightarrow \mathbb{E} [\text{cost of best among MWU}] = \frac{\sum_{t=1}^T v^t}{1 - \text{OPT}} \leq \frac{\ln n}{\epsilon} \xrightarrow{+ \epsilon \text{OPT}} \frac{1 + \epsilon \text{OPT}}{\epsilon}$$

$$\Rightarrow E[\text{cost}_{\text{MWC}}] = \frac{\text{cost} - \text{OPT}}{\epsilon}$$

$$\epsilon = \sqrt{\frac{\ln n}{T}}$$

$$\begin{aligned} \text{avg Regret} &= \frac{1}{T} \left[ \sum_{t=1}^T v^t - \text{OPT} \right] \leq \frac{1}{T} \left[ \frac{\sqrt{\ln n}}{\sqrt{T}} + \sqrt{\frac{\ln n \cdot \text{OPT}}{T}} \right] \\ &\leq \frac{1}{\sqrt{T}} \left[ \sqrt{\ln n} \sqrt{T} + \sqrt{\ln n} \sqrt{T} \right] \\ &= 2 \sqrt{\frac{\ln n}{T}} \end{aligned}$$

N players:

$i \in N$   $S_i$ : action set.

\* If each player plays as per MWC.

after  $T = \frac{\ln n}{\epsilon^2}$  rounds  $\rightarrow$  regret  $\leq \epsilon$

the "avg play" is  $\epsilon$ -CCE.



That is,  
 Suppose player  $i$  plays  $p_i^t \in \Delta(S_i)$

$$\sigma_i^t = \prod_{i \in N} p_i^t$$

$$\sigma_i^t(s) = \prod_{i \in N} p_i^t(s_i) \quad \forall s \in \prod_{i \in N} S_i$$

(s<sub>1</sub>, ..., s<sub>n</sub>)

$$\sigma = \frac{1}{T} \sum_{t=1}^T \sigma_i^t \quad \text{is } \underline{\epsilon\text{-CCE}}$$

$\forall i, \forall s_i' \in S_i$

$$\mathbb{E}_{s \sim \sigma} [c_i(s)] \leq \mathbb{E}_{s \sim \sigma} [c_i(s_i', s_{-i})] + \epsilon$$

$\mathbb{E} [$  <sup>avg.</sup> cost of the agent  $]$

$\mathbb{E} [$  <sup>avg</sup> cost of the best-action-in-hindsight  $]$