



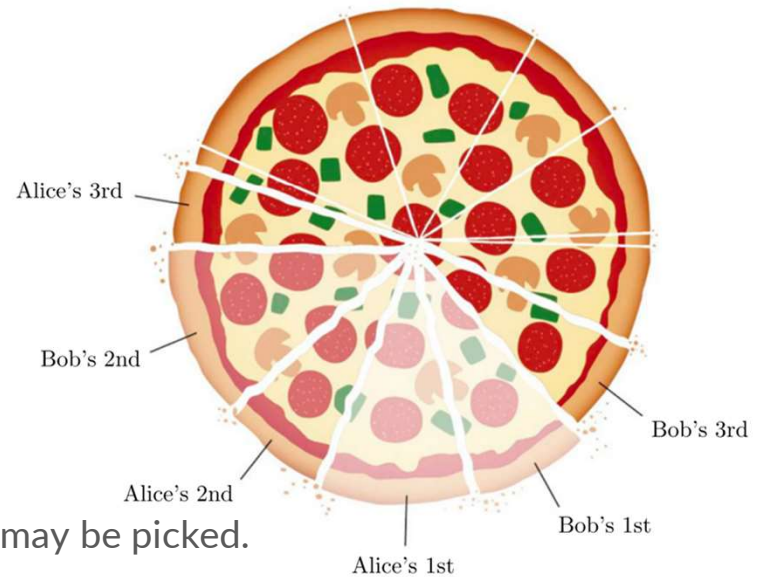
How to eat $4/9$ of a pizza

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The problem

Alice and Bob share a pizza.

- (i) They pick pieces in an alternating fashion;
- (ii) Alice starts by eating any piece of the pizza;
- (iii) Afterward only pieces adjacent to already eaten pieces may be picked.





Observation 1: Even count of pieces

Claim: Alice can get at least $\frac{1}{2}$ of the pizza (For even pieces).

Proof idea: color the pizza alternatively by red and green, then Alice can decide to pick all the green (or red) pieces.



Some definitions/ notations:

Def: 'interval' of a pizza = set of consecutive pieces. Odd intervals= intervals with odd pieces(Same for even) Two cuts, say C_1 and C_2 , encloses an interval.

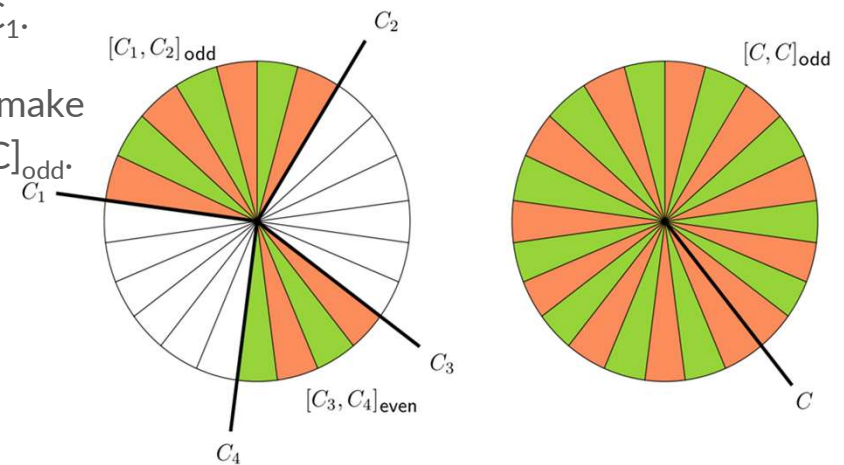
Def: ' $[C_1, C_2]_{\text{odd}}$ ' the odd interval between C_1 and C_2

' $[C_1, C_2]_{\text{even}}$ ' the even interval between C_1 and C_2

Coloring the intervals

Just like what we did in Observation 1, we want to color each interval by red and green. For an interval $[C_1, C_2]$, we make sure that red is always adjacent to C_1 .

Since now we are only focusing on odd pizzas, we can make a random cut, calling it C , then color the pizza w.r.t $[C, C]_{\text{odd}}$.





Follow Bob strategy (fB)

After the first pick, if Alice decides to pick the piece that was just revealed by Bob, we call this strategy 'Follow Bob'(fB).



Prop 2: Easy threshold for odd cases

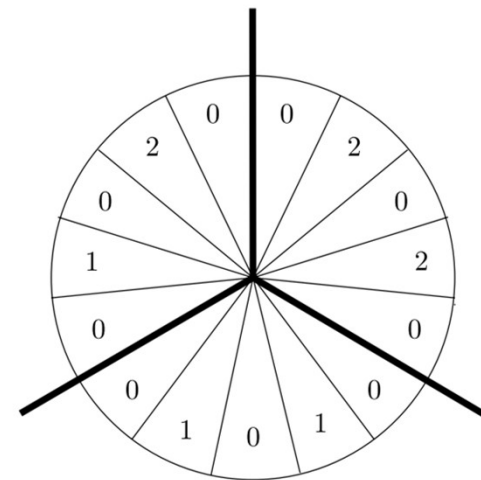
Claim: Alice can get at least $\frac{1}{3}$ of the pizza (For odd pieces).

Proof idea: Randomly pick a cut C of the pizza, then color $[C, C]_{\text{odd}}$ with reds and greens. We observe that Alice can always pick all the red pieces, then we assume that $R([C, C]_{\text{odd}}) < \frac{1}{3}$. Then the pizza has more greens. Starting with the 'mid-point green piece' opposing C , one can make sure that Alice gets at least half of the green areas, which is greater than $\frac{1}{3}$.

Observation 3:

Claim: There are some pizzas where Alice can only get $4/9$

Proof: Picture



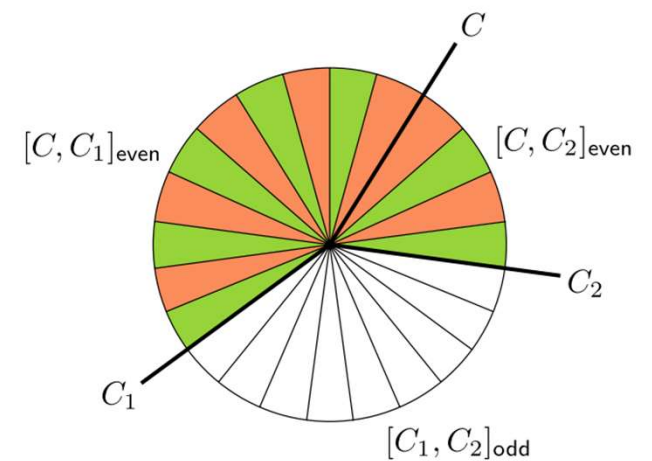


Now it's time to start the proof!

Assume that $[C_1, C_2]_{\text{odd}}$ has already been eaten. Now it's Bob's turn, and Alice is using fB-strategy.

We analyze Bob's behavior: Then Bob is trying to find a cut C such that $[C, C_1]_{\text{even}}$ and $[C, C_2]_{\text{even}}$ would have maximized green area.

We call C 'the (Bob's) *best answer to* $[C_1, C_2]_{\text{odd}}$ '





Definition: Heavy Greens Property

For an even interval $[C_1, C_2]_{\text{even}}$, if for every $[C_1, C]_{\text{even}}$ in $[C_1, C_2]_{\text{even}}$, the green pieces have larger (or equal) size.

One can prove that a cut C with $[C, C_1]_{\text{even}} \subseteq [C_1, C_2]_{\text{even}}$ is a best answer to $[C_1, C_2]_{\text{odd}}$ if and only if $[C, C_1]_{\text{even}}$ and $[C, C_2]_{\text{even}}$ have the heavy greens property.

Proof idea: We find C' that's also best answer, then $[C, C']$ would have equal reds and greens. Otherwise contradiction.

C best answer \Leftrightarrow both partitions have heavy greens property.



Definition: Easy pizzas and hard pizzas

A pizza is *easy* if fB-strategy for Alice would yield at least $\frac{1}{2}$ of the pizza.

Pizza is *hard* otherwise.

Theorem 4:

A hard pizza can be partitioned into three odd intervals each with heavy greens property.

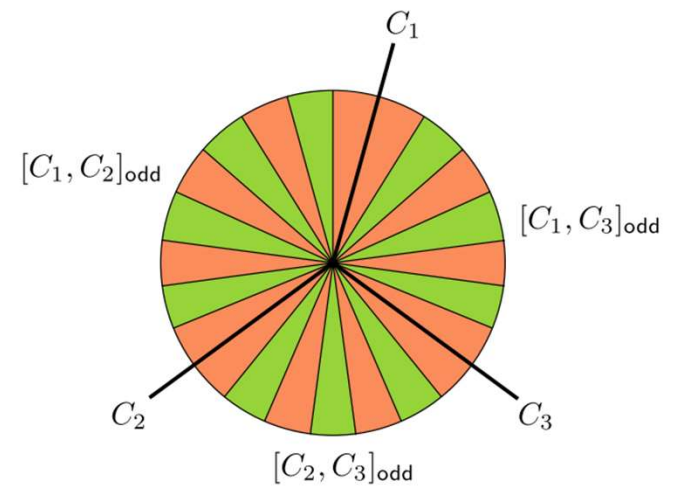
Proof: Very long, will discuss if time allows.

Remarks: There's a stronger version of the theorem, but for now we can just assume that there are some properties for the chosen cuts:

$$\text{i) } g_1 + r_2 + r_3 \leq r_1 + r_2 + g_3 \leq r_1 + g_2 + r_3$$

where r_i means the sum of red areas in the interval opposing C_i . Similar for g_i

ii) Each cut is the best answer to some piece p . (best answer w.r.t one piece).





'Partial pizza'

We partition the pizza into 3 odd intervals, where the three cuts are named C_1 , C_2 and C_3 .

Consider $[C_2, C_3]_{\text{odd}}$, we can think of it as a self-contained pizza, which glues together C_2 and C_3 . We call this $C_{2,3}$, and the resulting pizza 'partial pizza'.

Note. The reason why 2 and 3 were picked are because of the stronger version of Theorem 4.



strategies on the partial pizza

A strategy is said to be *plugged into the whole pizza* (call it P_{strat}) if:

- 1) If none of C_2 and C_3 is revealed, Alice follow the good strategy.
- 2) If Bob reveals C_2 or C_3 , Alice does not follow Bob. Instead, pick the other option.
- 3) Alice follows Bob if Bob picks something outside the partial pizza.
- 4) Inside the partial pizza, Alice has some specific strategy.

A P_{strat} ensures the outcome of Alice to be at least the strategy inside the partial pizza, plus $r_2 + g_3$.

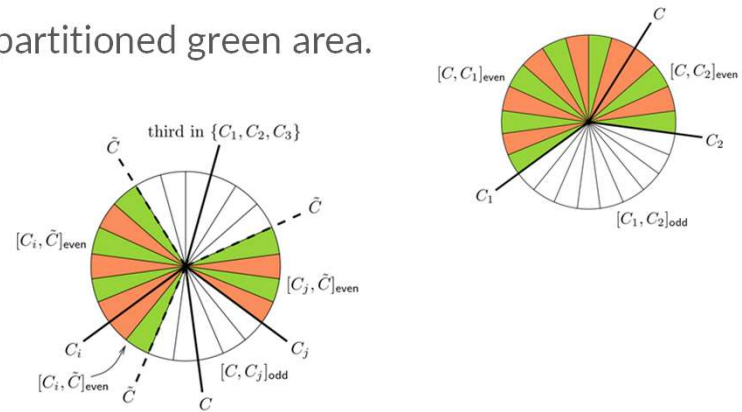
Proof on next page

Lemma 5: Best answer on the edge.

Claim: Let C_i, C_j be two cuts in our 3 cuts. Then if $[C, C_j]_{\text{odd}}$ is in $[C_i, C_j]_{\text{odd}}$, then either C_i or C_j is a best answer to it.

Reminder: best answer means that it maximizes the the partitioned green area.

Proof: Consider \tilde{C} , a best answer to $[C, C_j]_{\text{odd}}$, then clearly it's not in $[C_i, C_j]_{\text{odd}}$, then at least one of $[\tilde{C}, C_j]_{\text{ev}}$ and $[\tilde{C}, C_i]_{\text{odd}}$ is a subset of an interval with the heavy green property. WLOG suppose that's i .



-continue: We know $\sim C$ is best answer, so $[\sim C, C_j]_{\text{even}}$ is heavy green. We just mentioned that $[C_i, \sim C]_{\text{even}}$ is also heavy green. So it has equal reds and greens. Therefore C_i is also a best answer. □

After proving this we prove

Theorem 6: A Pstrat ensures the outcome of Alice to be at least the strategy inside the partial pizza, plus $r_2 + g_3$.

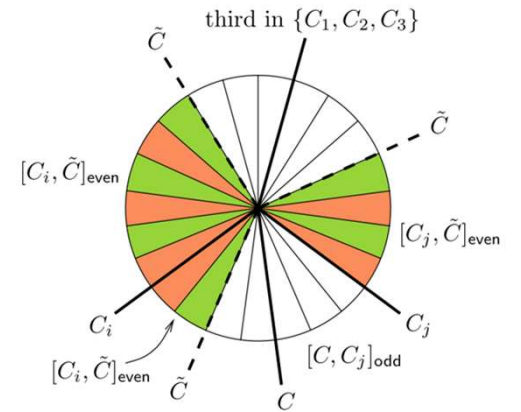
Note that Alice follow Bob outside the partial pizza when some

$[C, C_2]_{\text{odd}} \subseteq [C_2, C_3]_{\text{odd}}$ or $[C, C_3]_{\text{odd}} \subseteq [C_2, C_3]_{\text{odd}}$ is eaten.

By lemma 5 we know that the best answer of it is given by C_2 or C_3 .

In the worst case Alice gets $r_2 + g_3$ or $r_3 + g_2$, by our Observation 1.

$r_2 + g_3 \leq r_3 + g_2$ by our assumption in Theorem 4 proof.





Last before final: calculating some strategies outcome.

Our final result is the combination of a list of strategies, taking the optimal one.

-fB

-mfB

-Pstrat combined with fB or mfB.

fB-strategy associated with cut C

Just like what we assumed in the proof of Theorem 4,

$$g_1 + r_2 + r_3 \leq r_1 + r_2 + g_3 \leq r_1 + g_2 + r_3$$

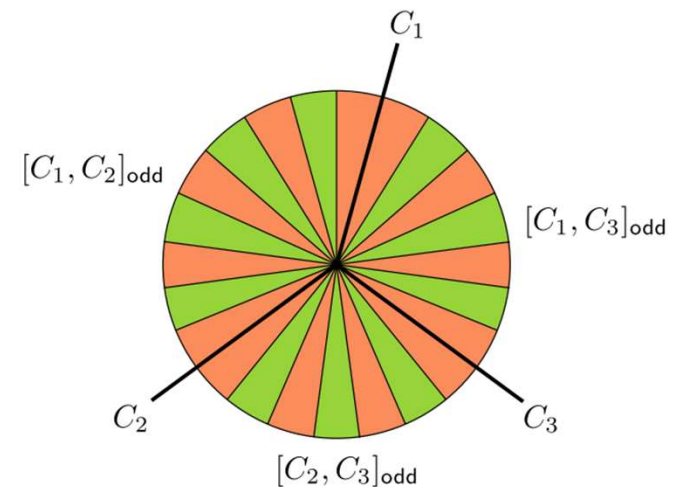
These values are actually the least outcomes of the fB-strategy:

Let C be the cut we pick among the three, and p the piece which

has the best answer of C, then if Alice starts with p and follows Bob

afterwards, she gets all the red pieces of [C, C], which corresponds to

to the values above.



Cut	Alice's outcome	Bob's outcome
C_1	$g_1 + r_2 + r_3$	$r_1 + g_2 + g_3$
C_2	$r_1 + g_2 + r_3$	$g_1 + r_2 + g_3$
C_3	$r_1 + r_2 + g_3$	$g_1 + g_2 + r_3$



modified follow-Bob strategy(mfB) associated with cut C

Idea: eat more green pieces.

(i) Alice starts with eating $p_i \in G_i$.

(ii) As long as Bob's moves reveal pieces in G_i Alice picks them, i.e., follows Bob.

(iii) At the moment Bob's move reveals the first red piece from another of the three odd intervals, Alice makes a single move that does not follow Bob. This means she picks a piece from R_i .

(iv) Alice follows Bob from then on.

If the starting piece is the middle-piece of the green pieces in G_i , then Alice gets at least half of the size of the green pieces in G_i . One can show that the following table is t

mfB-strategy	Alice's outcome
mfB ₁	$\frac{g_1}{2} + g_2 + r_3$
mfB ₂	$g_1 + \frac{g_2}{2} + r_3$
mfB ₃	$g_1 + r_2 + \frac{g_3}{2}$



Final proof: 4/9.

Consider two cases, the partial pizza is easy or hard.



First case: partial is easy

Use the following strategies:

(i) The fB-strategy associated to C2, which yields at least $r_1 + g_2 + r_3$;

(ii) The mfB-strategy mfB2, which yields at least $g_1 + g_2/2 + r_3$;

(iii) The fB-strategy plugged into the whole pizza, which yields at least $(g_1+r_1)/2 + r_2 + g_3$

Summing them up with weights we have:

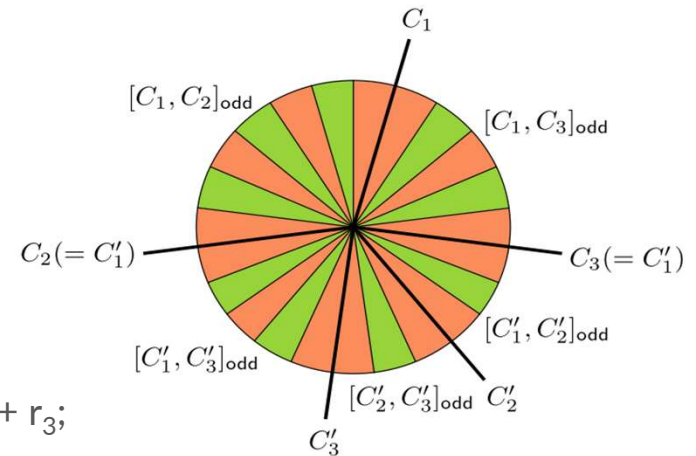
$$3/2 (r_1 + g_2 + r_3) + 2 (g_1 + g_2/2 + r_3) + ((g_1 + r_1)/2 + r_2 + g_3) \geq 2(r_1 + g_1 + r_2 + g_2 + r_3 + g_3) = 2.$$

Since we are having total weight = $1.5+2+1=4.5$, this mixed strategy has outcome = $2/4.5 = 4/9$

Second case: partial is hard.

We partition the partial pizza into 3 intervals again.

- (i) The fB-strategy associated to C_2 , which yields at least $r_1 + g_2 + r_3$;
- (ii) The mfB-strategy mfB2, which yields at least $g_1 + g_2 / 2 + r_3$;
- (iii) The fB-strategy associated to C'_2 plugged into the whole pizza, which yields at least $r'_1 + g'_2 + r'_3 + r_2 + g_3$;
- (iv) The strategy mfB1 for the partial pizza plugged into the whole pizza, which yields at least $g'_1 / 2 + g'_2 + r'_3 + r_2 + g_3$



Similarly we have:

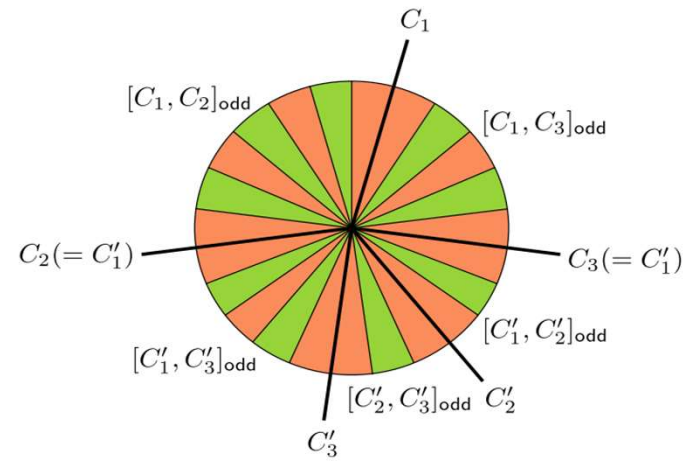
$$3/2 (r_1 + g_2 + r_3) + g_1 + g_2 / 2 + r_3 + (r'_1 + g'_2 + r'_3 + r_2 + g_3) + g'_1 / 2 + g'_2 + r'_3 + r_2 + g_3 .$$

Plug in $r_1 = g'_1 + r'_2 + r'_3$ and $g_1 = r'_1 + g'_2 + g'_3$

The above is greater or equal to 2, which again, would

Imply that this mixed strategy gives 4/9

Done! Thanks for watching!





Citations:

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