Incentivizing Truthful Feedback on Crowdsourcing Platforms

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Introduction

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 - Goal is to rate the quality of a product.
- Challenge is that users who report their reviews may not be truthful.
- This paper proposes a rewarding mechanism to the users so that they are incentivized to provide truthful feedback.

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- For each type x of plumber, the distribution of observations is $p_v(x)$.
- Each agent the arrival time of one or more plumbers, and they review the plumbers accordingly.
- We want to estimate $p_y(x)$ for every plumber type from the users' feedback.
 - Therefore, we want to incentivize workers to be truthful.

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- Y_{wt} denotes the response provided by the wth worker for the tth task
- v_w (y) is rewarded to worker w for a response y.

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- One way to ensure that agents are truthful is to induce a Bayes-Nash equilibrium.
- A rewarding mechanism induces a Bayes-Nash equilibrium if the strategy profile (response of workers) { Y_{wt} : w ∈ W, t ∈ T_w } satisfies the inequality:

```
E[v_w (\{Y_{wt} : w \in W, t \in T_w\})] \ge E[v_w (\{Y'_{wt} : w \in W, t \in T_w\} \cup \{Y_{w't} : w' \in W, w' \ne w, t \in T_{w'}\})]
```

for each alternative strategy $\{Y'_{wt} : t \in T_w\} \neq \{Y_{wt} : t \in T_w\}$ and all workers w.

Intuitive Explanation

- The Bayes-Nash Equilibrium says that if all the other workers adhere to the strategy profile $\{Y_{wt} : w \in W, t \in T_w \}$ then it is beneficial to worker w to also adhere to $\{Y_{wt} : w \in W, t \in T_w \}$.
- In our case, we would want the truthful strategy to be a Bayes-Nash Equilibrium.
- Further, if the inequality is strict, then it is called a strict Bayes-Nash Equilibrium.

Informational Requirements and Assumptions

- Assumptions and Public Knowledge
 - Generating model (P_X, p) is separated: For every $y \neq y'$, $E_{X \sim P_X} \Big[p_y(X) p_{y'}(X) \Big] < E_{X \sim P_X} \Big[p_y^2(X) \Big] E_{X \sim P_X} \Big[p_y^2(X) \Big]$.

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 - α -separation quantifies the gap:

$$\alpha + E_{X \sim P_X} [p_y(X)p_{y'}(X)] < E_{X \sim P_X} [p_y^2(X)]E_{X \sim P_X} [p_{y'}^2(X)].$$

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- Workers and host do not know the generating model (P_X, p)
 - Know that the model is separated.
- Workers and host assume that everyone is statistically identical.

Square Root Agreement Rule

- W_t = set of workers that respond to task t.
- T_w = set of tasks that worker w submitted a response to.
 - $T = AII \text{ tasks}; |T_w| \leq T \forall w.$
- Y_{wt} = Response of worker w to task t.

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- Repeat the following for every worker $w \in [n]$:
 - Let T_w be the set of tasks answered by worker w.
 - For every task $t \notin T_w$ **not** answered by this worker, select two different workers i, j who replied to task t.
 - Compute their #agreements on every possible observation y:

$$A(y) = \sum_{t \notin T_{w}} 1\{Y_{it} = Y_{jt} = y\}$$

• Scale appropriately to define a popularity index $I_w(y)$ for every observation y:

$$I_{w}(y) = \frac{1}{T - |T_{w}|} \left(1 + \sum_{t \notin T_{w}} 1\{Y_{it} = Y_{jt} = y\} \right).$$

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• For a task t that worker w answered (say y), reward the worker only if she agrees with an arbitrarily selected co-worker w':

$$v_w(Y_w = s) = \frac{1\{Y_{w'} = s\}}{\sqrt{I_w(y)}}, s \in \mathcal{Y}$$

inversely proportional to the popularity of the agreement.

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$$E[1\{Y_{w'} = s \mid Y_w = y\}] = \begin{cases} P(Y_{w'} = y \mid Y_w = y) : s = y \text{ (truthful)} \\ P(Y_{w'} = y' \mid Y_w = y) : s = y' \neq y \end{cases}.$$

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- Agents are statistically identical; If I observe y, it is likely that others observe y.
 - This incentivizes truthful behavior, based on an agent's observation.
- Issue: Workers may find it more likely for others to submit a highly popular observation.

Intuition 2: Importance of Popularity Index

• Popularity index, asymptotic:

$$I_{w}(y) = \frac{1}{T - |T_{w}|} \left(1 + \sum_{t \notin T_{w}} 1\{Y_{i} = Y_{j} = y\} \right)$$

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Repels workers from submitting an answer based on its popularity.

Result: SRA induces a Bayes-Nash Eq.

• Suppose all other agents
$$\neq w$$
 are truthful.
$$E[v_w(s)] = \begin{cases} \frac{P(Y_{w'} = y | Y_w = y)}{I_w(y)} : s = y \text{ (truthful)} \\ \frac{P(Y_{w'} = y' | Y_w = y)}{I_w(y')} : s = y' \neq y \end{cases}$$

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• Truthful reward ≥ False reward if

$$P(Y_{w} = Y_{w'} = y)P(Y_{w} = Y_{w'} = y') \ge P(Y_{w'} = y', Y_{w} = y)$$

$$\Leftrightarrow \sum_{x} P_{X}(x)p_{y}^{2}(x) \sum_{x} P_{X}(x)p_{y'}^{2}(x) \ge \left(\sum_{x} P_{X}(x)p_{y}(x)p_{y'}(x)\right)^{2}$$

for every $y' \neq y$.

Result Cont'd: Strict Bayes Nash

$$P(Y_w = Y_{w'} = y)P(Y_w = Y_{w'} = y') > P(Y_{w'} = y', Y_w = y)$$

- SRA is strictly Bayes-Nash incentive compatible in expectation if model is separable for some $\alpha > 0$.
 - Difference in reward for submitting true vs. false response is lower bounded by a value

$$\omega = \omega(T, \alpha, |\mathcal{Y}|, n)$$

when *T* is sufficiently large.

Extensions Overview

• Ex-ante collaboration: Workers may agree to submit a strategy $\sigma_w(s|y) \in \Delta(\mathcal{Y})$

before the game starts.

- Symmetric strategy: $\sigma_w = \sigma \ \forall w$.
 - Everyone submits s = 1 regardless of observation y.
- [Theorem] Truthful responses strictly maximize rewards over all symmetric strategies when #tasks $\rightarrow \infty$.

Collusion and Equilibrium

- Workers are not allowed to communicate throughout the game.
- It may be possible that workers communicate/collaborate beforehand to reduce effort while maintaining high reward.
- One such implication is a symmetric strategy profile, where all workers agree on a fixed modification $y \mapsto q(y)$.
 - For example, we want to get rid of cases when submitting trivial answers can achieve a high payoff. Otherwise, workers can receive high reward with low effort.

Collusion and Equilibrium

ullet Uninformativeness of a symmetric strategy profile q is define as

$$\Omega(q) = \frac{1}{|\mathcal{Y}|(|\mathcal{Y}| - 1)} \sum_{y,y'} \sum_{y'':y'' \neq y'} \sqrt{q_y(y')q_y(y'')}$$

and say that a strategy q is ω -uninformative if $\Omega(q) \geq \omega$.

- $\Omega(q) = 0 \Leftrightarrow q$ is fully-informative: q(y) has disjoint supports.
- $\Omega(q)$ is maximized if the report q is chosen independently of the true answer.
- Given enough tasks, a fully-informative strategy profile maximizes reward. In other words, SRA is strongly truthful across symmetric equilibria, asymptotically.

Strong Truthfulness Over All Equilibria

- Crowdsourcing host can choose how to assign tasks.
- Suppose $\frac{wN}{n} < M$ (fix to our notation).
- Asymmetric strategies that may arise (Sec. 5.3).
- State theorem 4.

Limitations and Our Extension

- Because workers are assumed to be statistically identical, a majority vote is the most accurate estimate of the underlying truth.
- When workers are heterogeneous, e.g. follow the Dawid-Skene model $p_w = P(Y_{wt} = y^*) \ \forall t$,

then SRA fails to incentivize truthful behavior.

- Most algorithms that efficiently aggregate worker responses are designed for the Dawid-Skene model, and it is of interest how to incentivize truthful behavior under such settings.
- Mechanisms designed for heterogeneous workers require extraneous reports, i.e. is not minimal.
- We are currently designing a mechanism that incentivizes truthful response when workers are heterogeneous.

Summary

- SRA incentivizes truthful responses without requiring extraneous reports.
- Under SRA, an honest response maximizes workers' rewards over symmetric equilibria.