Lecture 17
Other Solution Concepts and Game Models

CS580

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Some slides are borrowed from V. Conitzer's presentations.

So far

- Normal-form games
  - Multiple rational players, single shot, simultaneous move
- Bayesian Games (Incomplete Information)
- Nash equilibrium
  - Existence
  - Computation in two-player games.

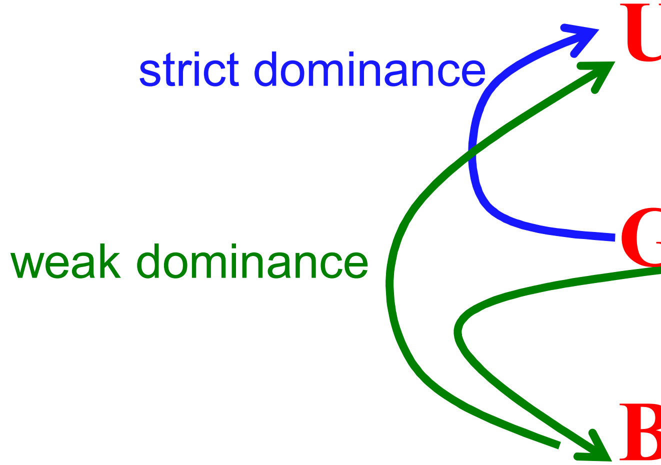
Today:

- Issues with NE
  - Multiplicity
  - Selection: How players decide/reach any particular NE
- Possible Solutions
  - Dominance: Dominant Strategy equilibria
  - Arbitrator/Mediator: Correlated equilibria, Coarse-correlated equilibria
  - Communication/Contract: Stackelberg equilibria, Nash bargaining
- Other Games
  - Extensive-form Games

Dominance

- Strict dominance: For a player, move s strictly dominates t if no matter what others play, s gives her better payoff than t
  - for all s_{-i}, $u_i(s, s_{-i}) > u_i(t, s_{-i})$
- s weakly dominates t if
  - for all s_{-i}, $u_i(s, s_{-i}) \geq u_i(t, s_{-i})$; and
  - for some s_{-i}, $u_i(s, s_{-i}) > u_i(t, s_{-i})$

-i = "the player(s) other than i"



	L	M	R
U	0, 0	1, -1	1, -1
G	-1, 1	0, 0	-1, 1
B	-1, 1	1, -1	0, 0

Dominant Strategy Equilibrium

Playing move s_i is best for agent i, no matter what others play.

- For each player i, there is a (strategy) move s_i that (weakly) dominates all other moves.
  - for all i, $s_i', s_{-i},\ u_i(s_i, s_{-i}) \geq u_i(s_i', s_{-i})$;

Example?

Prisoner's Dilemma

- Pair of criminals has been caught
- They have two choices: {confess, don't confess}

	confess	don't confess
confess	-5, -5	0, -6
don't confess	-6, 0	-1, -1

"Should I buy an SUV?"

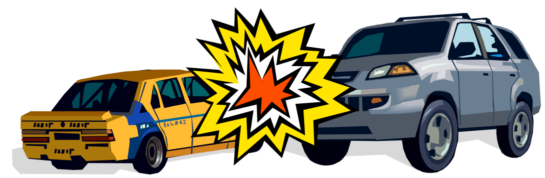
purchasing cost

cost: 5

cost: 3

accident cost

cost: 5  cost: 5

cost: 8  cost: 2

cost: 5  cost: 5

		
	-10, -10	-7, -11
	-11, -7	-8, -8

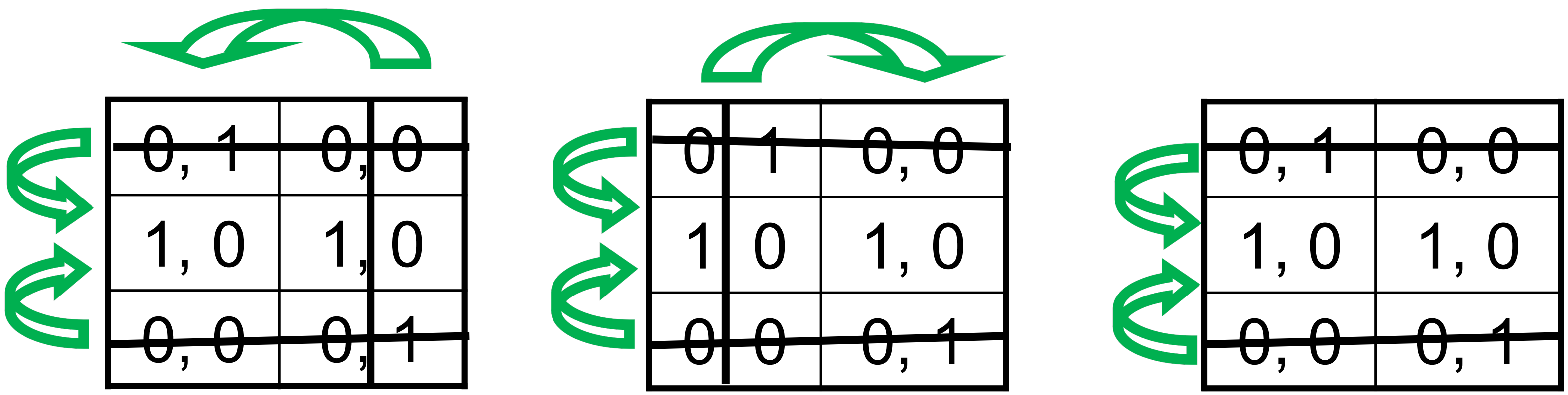
Dominance by Mixed strategies

- Example of dominance by a mixed strategy:

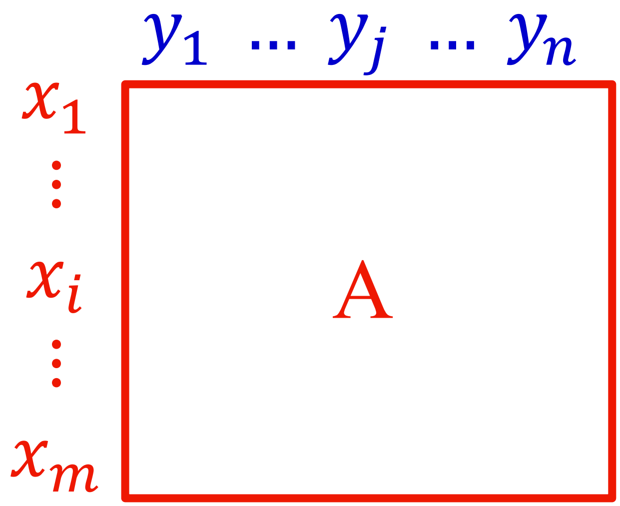
1/2	3, 1	0, 0
1/2	0, 0	3, 2
	1, 0	1, 1

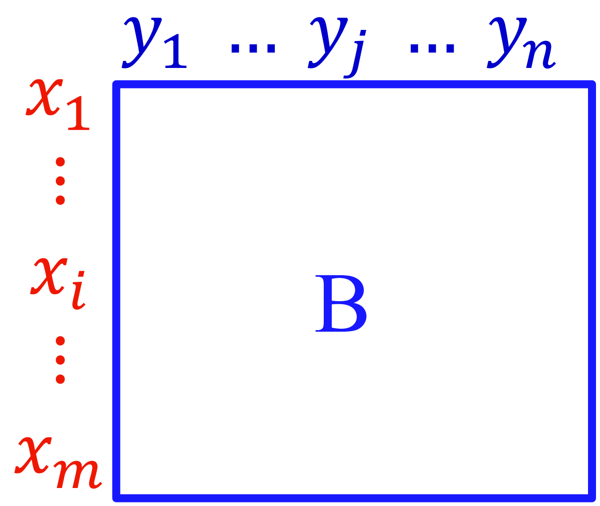
Iterated dominance: path (in)dependence

Iterated weak dominance is path-dependent: sequence of eliminations may determine which solution we get (if any)
(whether or not dominance by mixed strategies allowed)



Iterated strict dominance is path-independent: elimination process will always terminate at the same point
(whether or not dominance by mixed strategies allowed)





NE: $x^T A y \geq x'^T A y, \ \forall x'$ $\quad$ $x^T B y \geq x^T B y', \ \forall y'$

No one plays dominated strategies.

Why?

What if they can discuss beforehand?

Players: {Alice, Bob}

Two options: {Football, Tennis}

		2/3 F	1/3 T
1/3	F	1 2 0.5	0 0
2/3	T	0 0	2 1 0.5

At Mixed NE
both get 2/3 < 1

Instead they agree on ½(F, T), ½(T, F)
Payoffs are (1.5, 1.5) Fair!

Needs a common coin toss!

Correlated Equilibrium – (CE)
(Aumann'74)

- **Mediator** declares a joint distribution P over $S = \times_i S_i$
- Tosses a coin, chooses $s = (s_1, \dots, s_n) \sim P$.
- Suggests s_i to player i **in private**

- P is at equilibrium if each player wants to follow the suggestion when others do.
  - $U_i\left(s_i, P_{(s_i,.)}\right) \geq U_i\left(s_i', P_{(s_i,.)}\right), \forall s_i' \in S_1$

CE for 2-Player Case

- **Mediator** declares a joint distribution $P = \begin{bmatrix} p_{11} & \cdots & p_{1n} \\ \vdots & \vdots & \vdots \\ p_{m1} & \cdots & p_{mn} \end{bmatrix}$
- Tosses a coin, chooses $(i, j) \sim P$.
- Suggests i to Alice, j to Bob, in private.

- P is a CE if each player wants to follow the suggestion, when the other does.

Given Alice is suggested i, she knows Bob is suggested $j \sim P(i,.)$

$$\langle A(i,.), P(i,.)\rangle \geq \langle A(i',.), P(i,.)\rangle \quad : \forall i' \in S_1$$

$$\langle B(.,j), P(.,j)\rangle \geq \langle B(.,j'), P(.,j)\rangle \quad : \forall j' \in S_2$$

Players: {Alice, Bob}

Two options: {Football, Shopping}

	F	S
F	1 2 (0.5)	0 0
S	0 0	2 1 (0.5)

Instead they agree on ½(F, S), ½(S, F)

Payoffs are (1.5, 1.5)   Fair!

CE!

Prisoner's Dilemma

	C	NC
C	-5, -5 (1)	0, -6 (0)
NC	-6, 0 (0)	-1, -1 (0)

NC is dominated

Rock-Paper-Scissors (Aumann)

	R	P	S
R	0, 0 (0)	0, 1 (1/6)	1, 0 (1/6)
P	1, 0 (1/6)	0, 0 (0)	0, 1 (1/6)
S	0, 1 (1/6)	1, 0 (1/6)	0, 0 (0)

When Alice is suggested R

Bob must be following $P_{(R,.)} \sim (0,1/6,1/6)$

Following the suggestion gives her 1/6

While P gives 0, and S gives 1/6.

Computation: Linear Feasibility Problem

Game (A, B). Find, joint distribution $P = \begin{bmatrix} p_{11} & \cdots & p_{1n} \\ \vdots & \vdots & \vdots \\ p_{m1} & \cdots & p_{mn} \end{bmatrix}$

s.t.
$$\sum_j A_{ij} p_{ij} \geq \sum_j A_{i'j} p_{ij} \quad \forall i, i' \in S_1$$
$$\sum_i B_{ij} p_{ij} \geq \sum_i B_{ij'} p_{ij} \quad \forall j, j' \in S_2$$
$$\sum_{ij} p_{ij} = 1; \quad p_{ij} \geq 0, \ \forall (i, j)$$

N-player game: Find distribution P over $S = \times_{i=1}^{N} S_i$

s.t. $U_i\left(s_i, P_{(s_i,.)}\right) \geq U_i\left(s_i', P_{(s_i,.)}\right), \ \forall s_i, s_i' \in S_i$

$$\sum_{s \in S} P(s) = 1$$

↑

$\sum_{s_{-i} \in S_{-i}} U_i(s_i, s_{-i}) P(s_i, s_{-i})$ Linear in P variables!

Computation: Linear Feasibility Problem

N-player game: Find distribution P over $S = \times_{i=1}^{N} S_i$

s.t. $U_i\left(s_i, P_{(i,.)}\right) \geq U_i\left(s_i', P_{(s_i,.)}\right),\ \forall s_i, s_i' \in S_i$

$\uparrow$ $\sum_{s \in S} P(s) = 1$

$\sum_{s_{-i} \in S_{-i}} U_i(s_i, s_{-i}) P(s_i, s_{-i})$ Linear in P variables!

Can optimize any convex function as well!

Coarse-Correlated Equilibrium

- After mediator declares P, each player opts in or out.
- Mediator tosses a coin, and chooses s ~ P.
- If player i opted in, then the mediator suggests her s_i in private, and she has to obey.
- If she opted out, then (knowing nothing about s) plays a fixed strategy $t \in S_i$
- At equilibrium, each player wants to opt in, if others are.

$$U_i(P) \geq U_i(t, P_{-i}),\ \forall t \in S_i$$

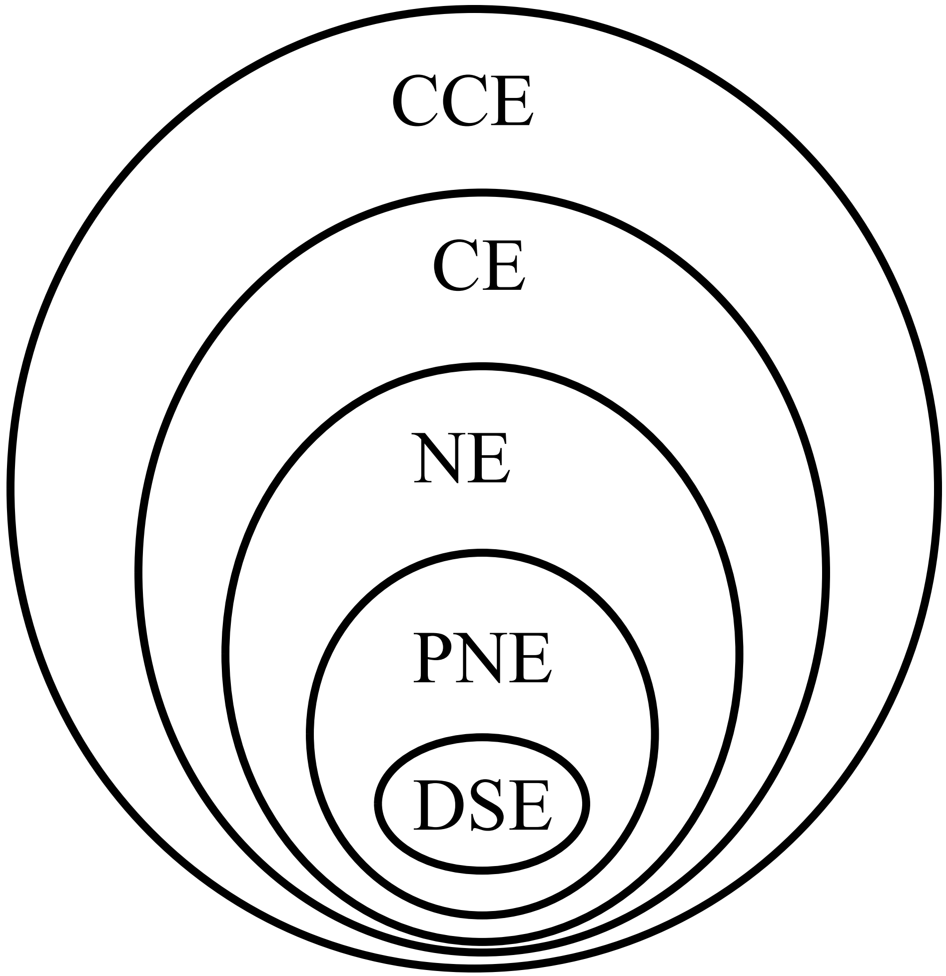
Where P_{-i} is joint distribution of all players except i.

Importance of (Coarse) CE

- Natural dynamics quickly arrive at approximation of such equilibria.
  - No-regret, Multiplicative Weight Update (MWU)

- Poly-time computable in the size of the game.
  - Can optimize a convex function too.

Show the following

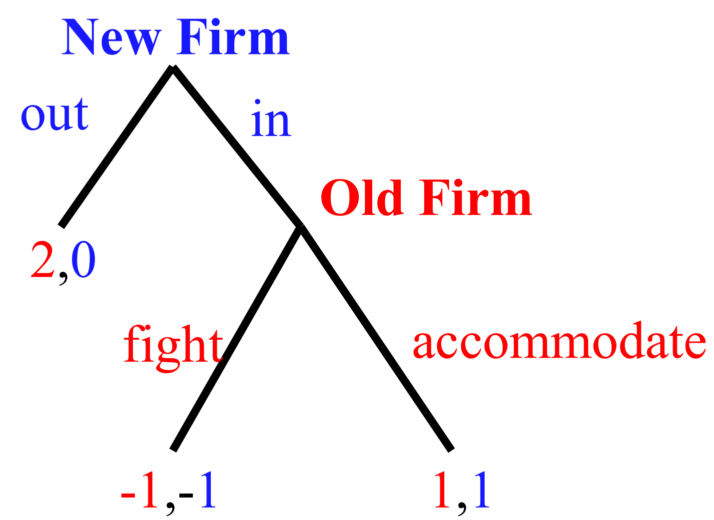


Extensive-form Game

- Players move one after another
  - Chess, Poker, etc.
  - Tree representation.

Strategy of a player:
What to play at each of its node.

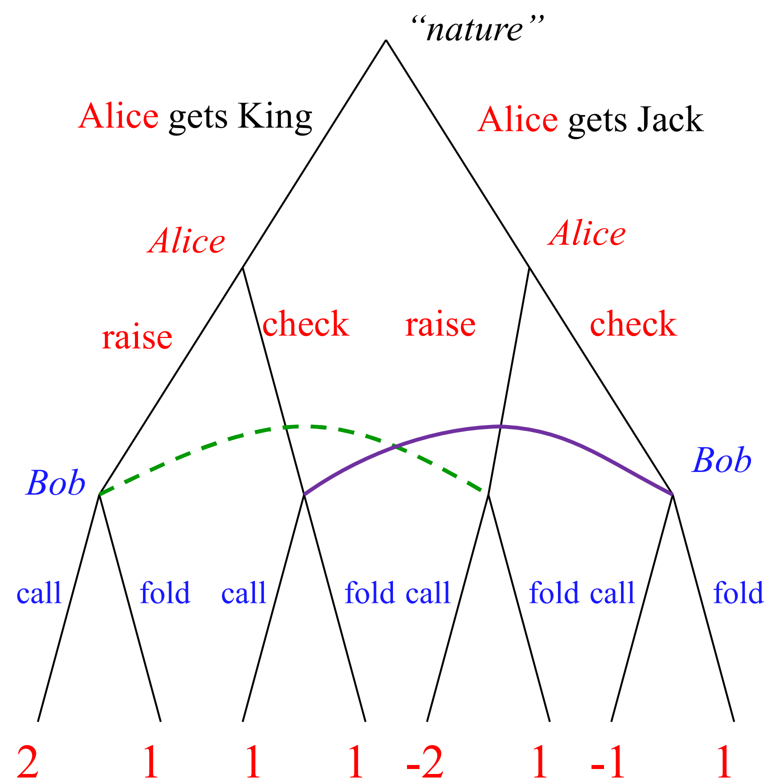
	I	O
F	-1, -1	2, 0
A	1, 1	2, 0



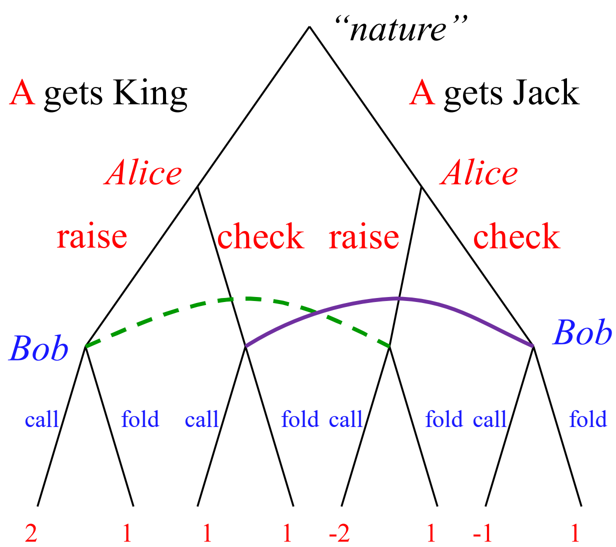
Entry game

A poker-like game

- Both players put 1 chip in the pot
- Alice gets a card (King is a winning card, Jack a losing card)
- Alice decides to raise (add one to the pot) or check
- Bob decides to call (match) or fold (Alice wins)
- If Bob called, he adds one to the pot. Alice's card determines pot winner.



Poker-like game in normal form

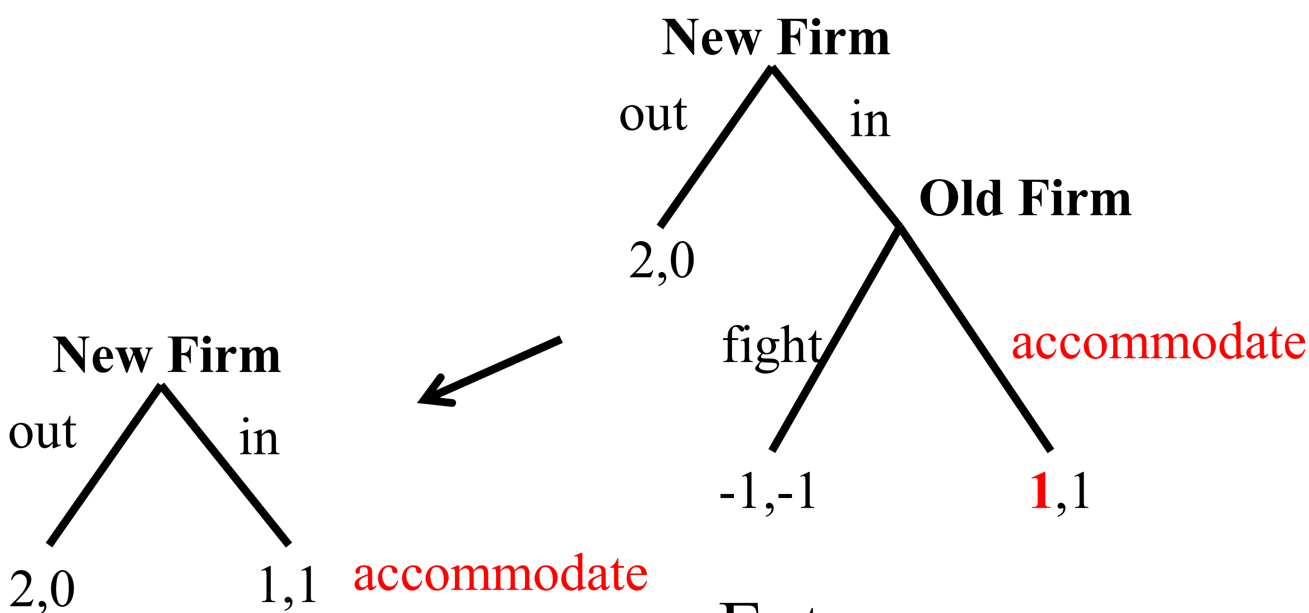


	cc	cf	fc	ff
rr	0, 0	0, 0	1, -1	1, -1
rc	.5, -.5	1.5, -1.5	0, 0	1, -1
cr	-.5, .5	-.5, .5	1, -1	1, -1
cc	0, 0	1, -1	0, 0	1, -1

Can be exponentially big!

Sub-Game Perfect Equilibrium

- Every sub-tree is at equilibrium
- Computation when perfect information (no nature/chance move): Backward induction



Entry game

Sub-Game Perfect Equilibrium

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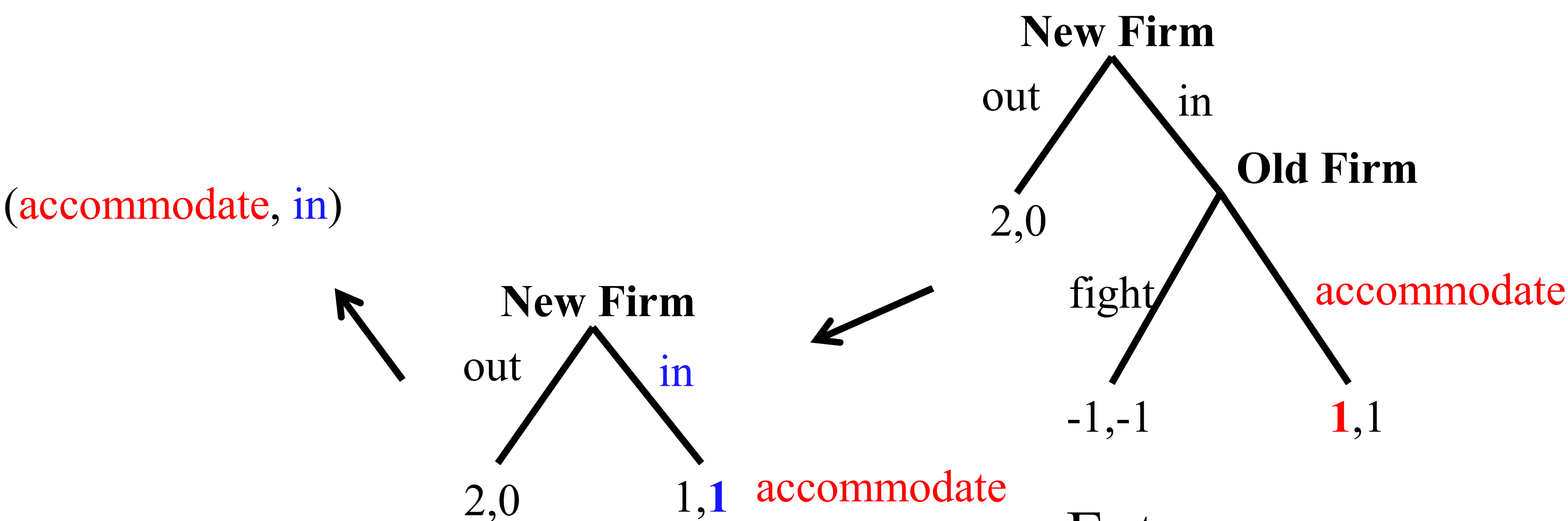


Entry game

Corr. Eq. in Extensive form Game

- How to define?
  - CE in its normal-form representation.
- Is it computable?
  - Recall: exponential blow up in size.
- Can there be other notions?

See "Extensive-Form Correlated Equilibrium: Definition and Computational Complexity" by von Stengel and Forges, 2008.

Commitment
(Stackelberg strategies)

Commitment

Unique Nash equilibrium (iterated strict dominance solution) → (top-left cell)

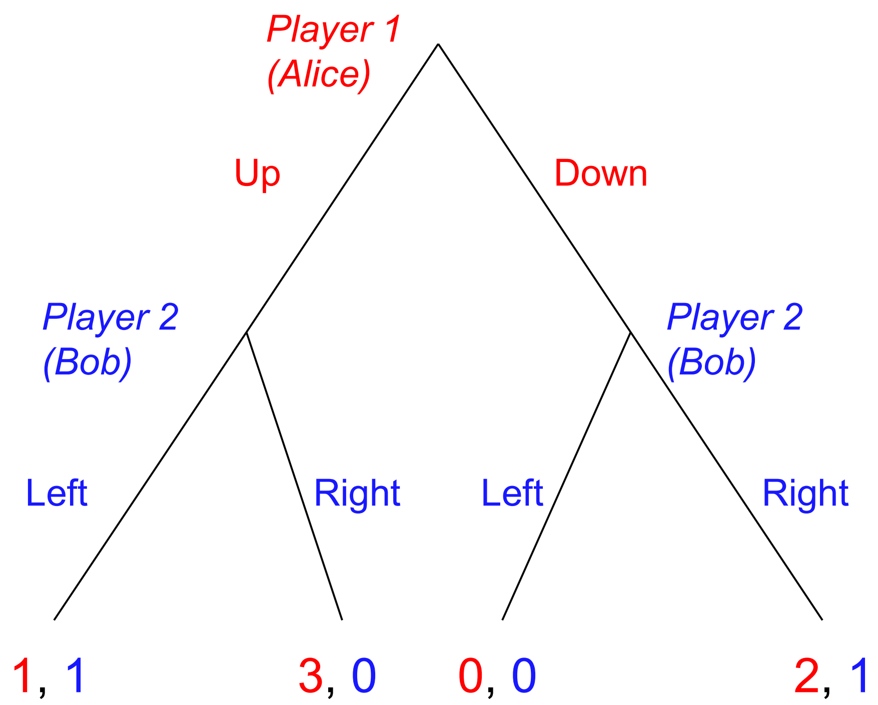
1, 1	3, 0
0, 0	2, 1

von Stackelberg

- Suppose the game is played as follows:
  - Alice **commits** to playing one of the rows,
  - Bob observes the commitment and then chooses a column
- Optimal strategy for Alice: commit to Down

Commitment: an extensive-form game

For the case of committing to a pure strategy:



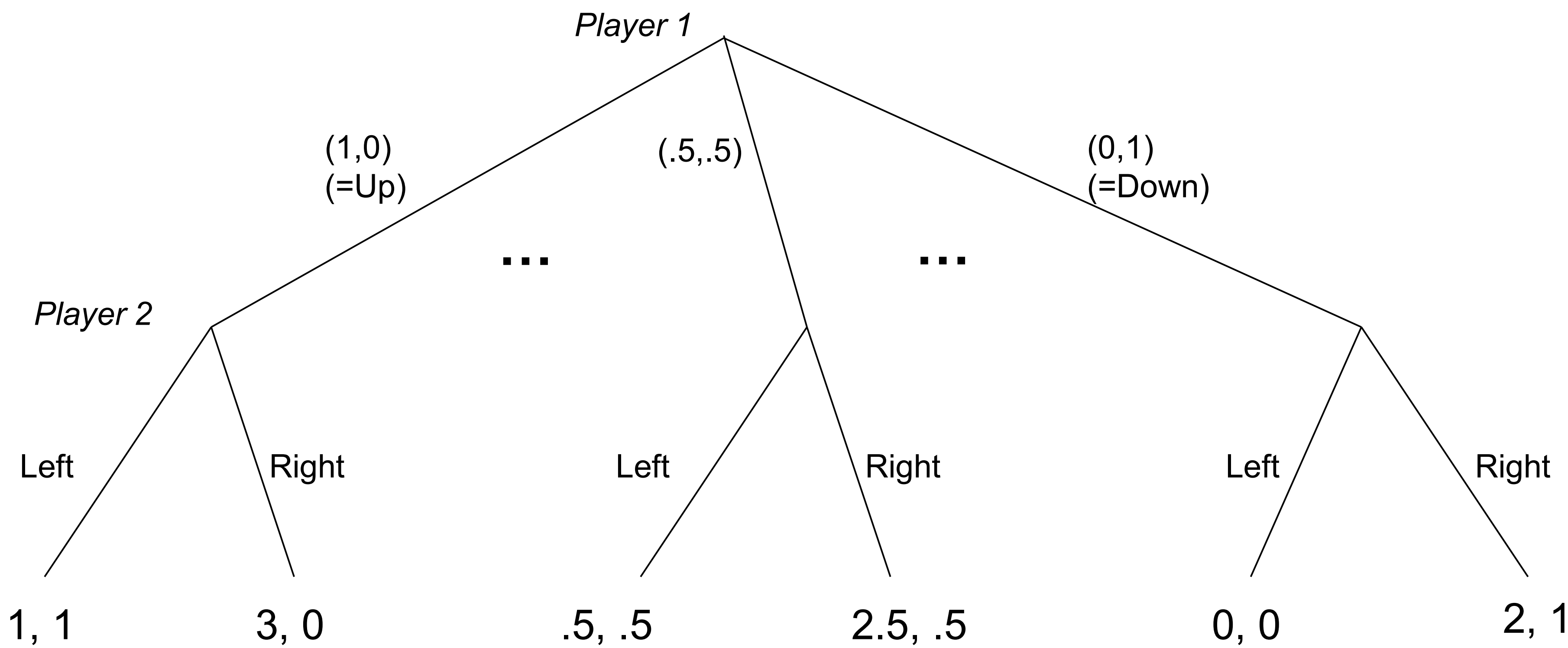
Commitment to mixed strategies

	0	1
.49	1, 1	3, 0
.51	0, 0	2, 1

Also called a Stackelberg (mixed) strategy

Commitment: an extensive-form game

- … for the case of committing to a mixed strategy:



- Economist: Just an extensive-form game, nothing new here
- Computer scientist: Infinite-size game! Representation matters

Computing the optimal mixed strategy to commit to

[Conitzer & Sandholm EC'06]

- Player 1 (Alice) is a leader.
- Separate LP for Bob's move (column) $j^* \in S_2$:

maximize $\sum_i x_i A_{ij^*}$ — Alice's utility when Bob plays j^*

subject to $\forall j, \; \left(x^T B\right)_{j^*} \geq \left(x^T B\right)_j$ — Playing j^* is best for Bob

$x \geq 0, \; \sum_i x_i = 1$ — x is a probability distribution

Among soln. of all the LPs,
pick the one that gives max utility.

On the game we saw before

x_1	1, 1	3, 0
x_2	0, 0	2, 1

$$\text{maximize } 1x_1 + 0\, x_2$$

$$\text{subject to}$$

$$1\, x_1 + 0\, x_2 \geq 0\, x_1 + 1\, x_2$$

$$x_1 + x_2 = 1$$

$$x_1 \geq 0,\, x_2 \geq 0$$

$$\text{maximize } 3\, x_1 + 2\, x_2$$

$$\text{subject to}$$

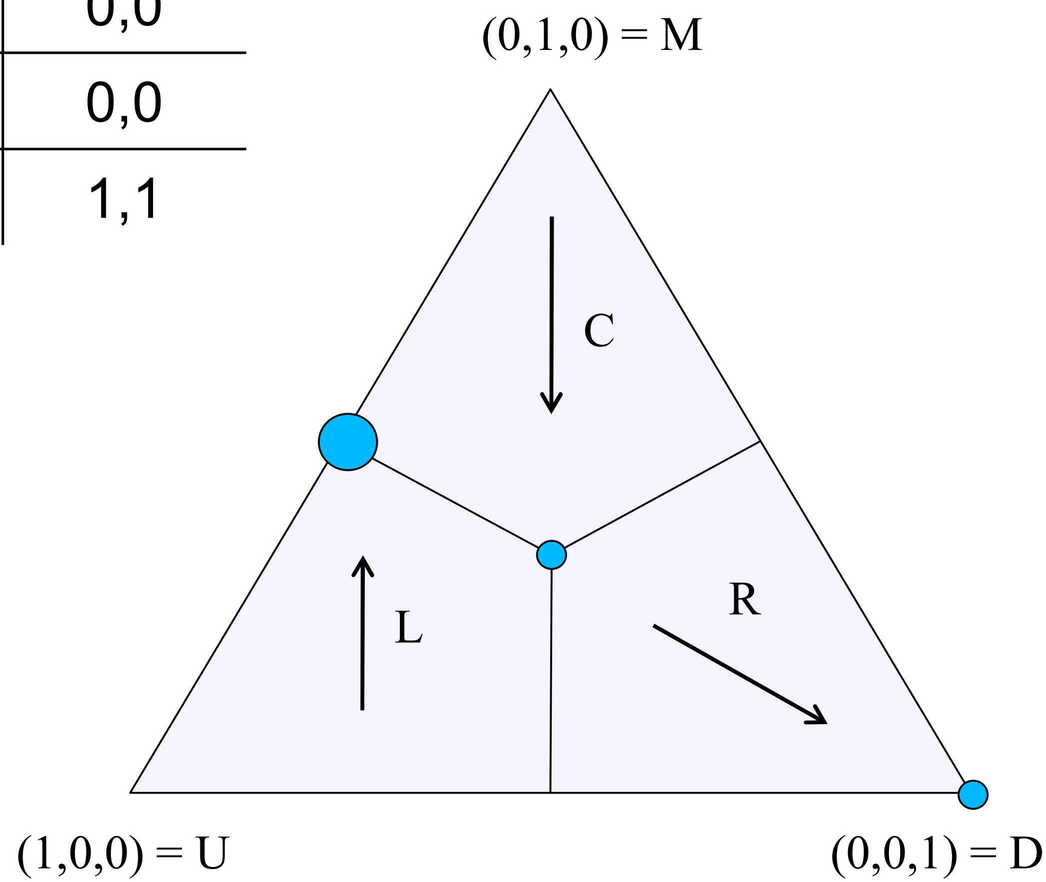
$$0\, x_1 + 1\, x_2 \geq 1\, x_1 + 0\, x_2$$

$$x_1 + x_2 = 1$$

$$x_1 \geq 0,\, x_2 \geq 0$$

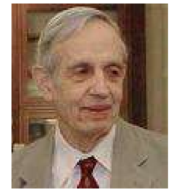
Visualization

	L	C	R
U	0,1	1,0	0,0
M	4,0	0,1	0,0
D	0,0	1,0	1,1

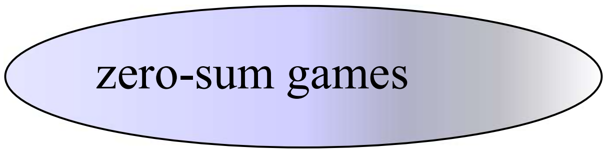


Generalizing beyond zero-sum games

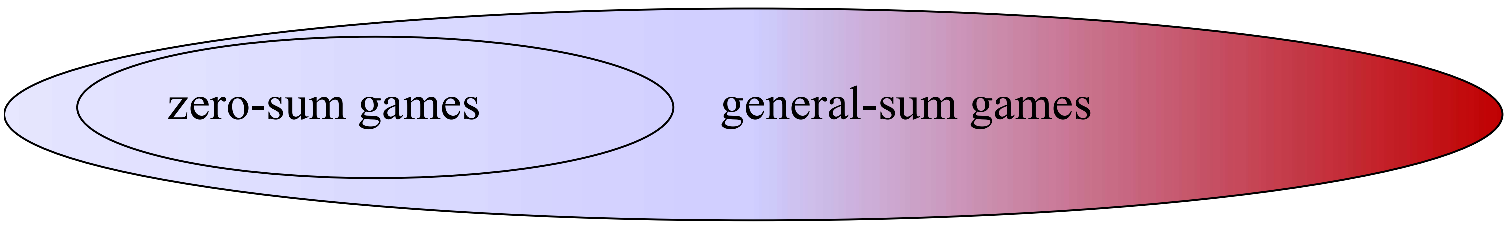
Minimax, Nash, Stackelberg all agree in zero-sum games



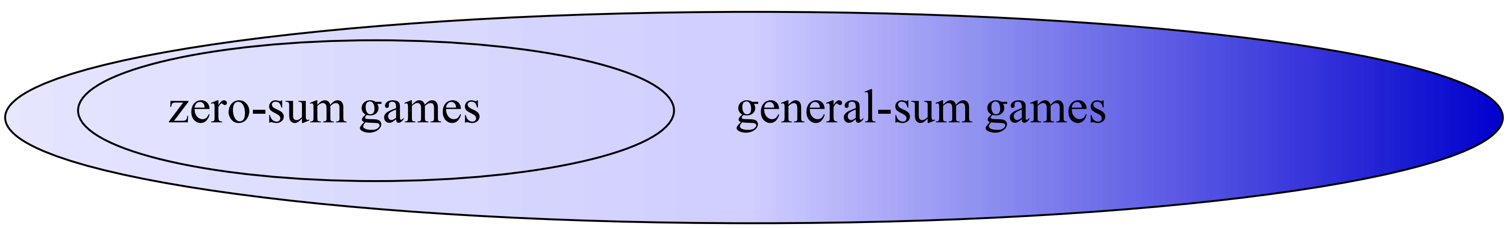
0, 0	-1, 1
-1, 1	0, 0



minimax strategies



Nash equilibrium



Stackelberg mixed strategies

Other nice properties of commitment to mixed strategies

- No equilibrium selection problem

0, 0	-1, 1
1, -1	-5, -5

- Leader's payoff at least as good as any Nash eq. or even correlated eq. (von Stengel & Zamir [GEB '10])

$\geq$

Nash Bargaining

Nash Bargaining: Dividing Utilities

Two agents: 1, 2
Outside option utilities: c_1, c_2
Feasible set of Utilities: $U \subseteq R^2$ (convex),

$$(c_1, c_2) \in U$$

Goal: define a bargaining function $\mathrm{f}(c_1, c_2, U) \in U$
Satisfying certain good properties

Nash Bargaining: Axioms

Two agents: 1, 2
Outside option with utilities: c_1, c_2
Feasible set of Utilities: $U \subseteq R^2$ (convex), $(c_1, c_2) \in U$

Goal: f$(c_1, c_2, U) \in U$ that is

1. Scale free
2. Symmetric
3. Pareto Optimal
4. Independent of Irrelevant Alternatives (IIA)
5. Individually Rational

Nash Bargaining: Theorem

Two agents: 1, 2
Outside option with utilities: c_1, c_2
Feasible set of Utilities: $U \subseteq R^2$ (convex), $(c_1, c_2) \in U$
Goal: f$(c_1, c_2, U) \in U$ that is

1. Scale free
2. Symmetric
3. Pareto Optimal
4. Independent of Irrelevant Alternatives (IIA)
5. Individually Rational

Theorem (Nash'50). f satisfies the 5 axioms if and only if, $f(c_1, c_2, U)$ is

$$\begin{aligned} &argmax\ (u_1 - c_1)(u_2 - c_2) \\ &\text{s.t.} \quad (u_1, u_2) \in U \end{aligned}$$

Nash Bargaining: Theorem

Theorem (Nash'50). f satisfies the 5 axioms if and only if, $f(c_1, c_2, U)$ is

$$\begin{aligned} &argmax\ (u_1 - c_1)(u_2 - c_2) \\ &\text{s.t.} \qquad (u_1, u_2) \in U \end{aligned}$$

Proof. ($\Leftarrow$)

1. Scale free

2. Symmetric

3. Pareto Optimal

4. Independent of Irrelevant Alternatives (IIA)

5. Individually Rational