CS576 Topics in Automated Deduction

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Records in Isabelle/HOL

Records in HOL are basically tuples, but...

```
record ('a)graph_sig =
Vertices :: "'a set"
Edges :: "('a × 'a) set"

definition one :: "(unit) graph_sig" where
"one = { Vertices = {()}, Edges = {((),())} } "
```

Components accessed by field

```
lemma "()\in Vertices one" by (simp add: one_def)
```

Unlike functional programming languages, position of fields matters

```
\emptyset Edges = \{((),())\}, Vertices = \{()\} \emptyset" causes an "Error in record input"
```

Record Field Update

Records support field update

```
definition no_edge :: "(unit) graph_sig" where
"no_edge = one() Edges := {} )"

lemma "no_edge = () Vertices = {()}, Edges = {} )"

by (simp add: no_edge_def one_def)
```

Record Field Update

Record update is functional

```
definition no_edge :: "(unit) graph_sig" where
"no_edge = one(Edges := {})"
lemma "no_edge = (Vertices = {()}, Edges = {} )"
by (simp add: no_edge_def one_def)
lemma "no_edge \neq one"
by (simp add: no_edge_def one_def)
lemma "(Vertices no_edge = Vertices one) ∧
       (Edges no_edge = {})"
by (simp add: no_edge_def one_def)
```

record type representations

Every record type many be given by field syntax:

- Every record is extensible
 - Every record rec_ty defines a polymorphic type
 ('a) rec_ty_scheme;
 type rec_ty same as (unit) rec_ty_scheme

```
term "one::(unit,unit) graph_sig_scheme"
```

Every record type has a "hidden field" more

New Record Types From Old

 New record types may be created by extending existing record types with additional fields:

```
record ('a, 'b) labeled_graph_sig = "'a graph_sig" +
Label :: "('a \times 'a) \Rightarrow ('b) option"
definition two :: "(Vertices :: nat set,
                  Edges :: (nat \times nat) set,
                  Label :: nat \times nat \Rightarrow bool option)" where
"two =
 \{ Vertices = \{1,2\}, Edges = \{(1,2)\}, \}
   Label = (\lambda \text{ (m,n)}). (if (\text{m,n}) = (1,2) then Some True else No
constants
 two :: "(nat, bool) labeled_graph_sig"
```

Record Polymorphism

- The _scheme types allow for (weak) record polymorphism
- Also referrd to as record subtyping

by (simp add: two_def is_graph_def)

 Generally, input to functions should use _scheme type instead of strict record type

```
definition is_graph :: "('a,'b) graph_sig_scheme ⇒ bool" when
"is_graph G ≡ (∀ e ∈ Edges G. {fst e, snd e} ⊆ Vertices G)"

lemma shows "is_graph one"
by (simp add: one_def is_graph_def)

lemma shows "is_graph two"
```