CS576 Topics in Automated Deduction

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Slides based in part on slides by Tobias Nipkow March 11, 2015

Records in Isabelle/HOL

• Records in HOL are basically tuples, but...

```
record ('a)graph_sig =
Vertices :: "'a set"
Edges :: "('a \times 'a) set"
definition one :: "(unit) graph_sig" where
"one \equiv ( Vertices = \{()\}, Edges = \{((),())\} )"
```

Components accessed by field

```
lemma "()∈ Vertices one"
by (simp add: one_def)
```

• Unlike functional programming languages, position of fields matters

```
\{ Edges = \{((),())\}, Vertices = \{()\} \} "
causes an "Error in record input"
```

Record Field Update

Records support field update

```
definition no_edge :: "(unit) graph_sig" where
"no_edge = one() \overline{E}dges := {} ) "
lemma "no_edge = ( Vertices = \{()\}, Edges = \{\} )"
by (simp add: no_edge_def one_def)
```

Record Field Update

Record update is functional

```
definition no_edge :: "(unit) graph_sig" where
"no_edge = one(Edges := {})"
lemma "no_edge = (Vertices = {()}, Edges = {} )"
by (simp add: no_edge_def one_def)
\texttt{lemma "no\_edge} \neq \texttt{one"}
by (simp add: no_edge_def one_def)
lemma "(Vertices no_edge = Vertices one) \wedge
       (Edges no_edge = {})"
by (simp add: no_edge_def one_def)
```

record type representations

• Every record type many be given by field syntax:

```
lemma "(one::unit graph_sig) =
       (one::(Vertices::unit set, Edges::(unit × unit) set))
by (rule refl)
```

- Every record is extensible
 - Every record rec_ty defines a polymorphic type ('a) rec_ty_scheme; type rec_ty same as (unit) rec_ty_scheme

term "one::(unit,unit) graph_sig_scheme"

• Every record type has a "hidden field" more

New Record Types From Old

• New record types may be created by extending existing record types with addtional fields:

```
record ('a,'b) labeled_graph_sig = "'a graph_sig" +
Label :: "('a \times 'a) \Rightarrow ('b) option"
definition two :: "(Vertices :: nat set,
                 Edges :: (nat \times nat) set,
                 Label :: nat \times nat \Rightarrow bool option) where
"two ≡
 ( Vertices = \{1,2\}, Edges = \{(1,2)\},
   Label = (\lambda (m,n)). (if (m,n) = (1,2) then Some True else No.
 two :: "(nat, bool) labeled_graph_sig"
```

Record Polymorphism

- The _scheme types allow for (weak) record polymorphism
- Also referrd to as record subtyping
- Generally, input to functions should use _scheme type instead of strict record type

```
definition is_graph :: "('a,'b) graph_sig_scheme ⇒ bool" when
"is_graph G \equiv (\forall e \in Edges G. {fst e, snd e} \subseteq Vertices G)"
lemma shows "is_graph one"
by (simp add: one_def is_graph_def)
lemma shows "is_graph two"
by (simp add: two_def is_graph_def)
```

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