

Expectation 1/20/22

Ω : Sample space
 Atomic event / elementary events
 $E \subseteq \Omega$
 Event: Dice is even $= \{2, 4, 6\}$
 $P: \Omega \rightarrow [0, 1]$
 - $P[\Omega] = 1$
 - $\forall X, Y \subseteq \Omega, X \cap Y = \emptyset \Rightarrow P[X \cup Y] = P[X] + P[Y]$
 $P: E \rightarrow [0, 1]$
 $E \subseteq \Omega$: σ -algebra.

Union bound

$E_1, E_2, \dots, E_m \subseteq \Omega$
 $P[\cup E_i] \leq \sum P[E_i]$

Random Variable RV

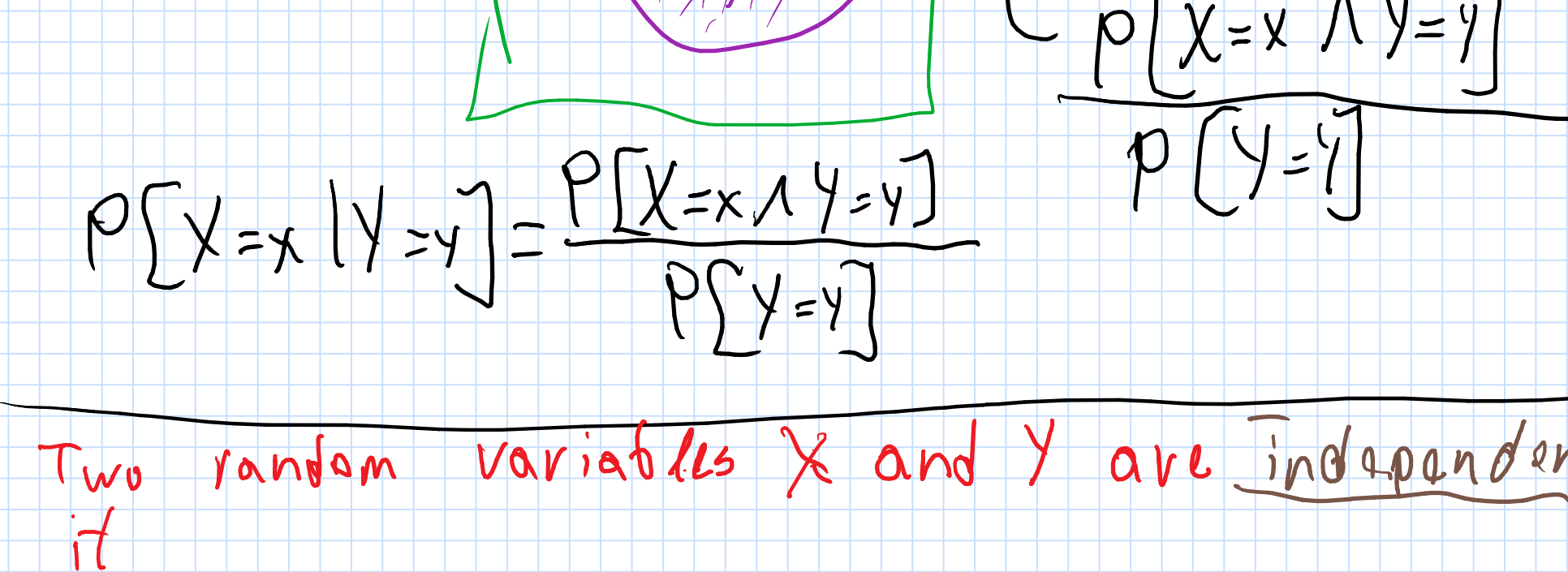
$f: \Omega \rightarrow T$
 $X \in \mathbb{Z}$
 $i \in \mathbb{Z} P[X=i]$

Expectation of a RV X is
 $\mu_X = E[X] = \sum_{x \in \text{all values that } X \text{ might get}} x \cdot P[X=x]$ - average value of X

Conditional probability

X, Y : Two RV
 $P[X=x | Y=y] =$ probability of X to be x given that Y is y .
 $X=1 \Leftrightarrow$ the dice is even **Indicator variables**
 $= 0$ otherwise
 $Y=1 \Leftrightarrow$ The dice ≤ 3
 $= 0$ otherwise

$P[X=1 | Y=1] = P[d=2,4,6 | d=1,2,3] = \frac{1}{3}$
 $P[X=1] = \frac{1}{2}$



Two random variables X and Y are independent if
 $P[X=x | Y=y] = P[X=x]$
 $\Leftrightarrow P[X=x \wedge Y=y] = P[X=x] P[Y=y]$

Expectation

Lemma $\forall X, Y$ RV we have
 $E[X+Y] = E[X] + E[Y]$.

Linearity of expectation

E is a linear operator.
 $E[X+Y] = E[X] + E[Y]$
 $E[c \cdot X] = c E[X]$.
 c constant

Lemma
 $\forall X, Y$ RV that are independent
 $E[XY] = E[X] E[Y]$.

Variance

$\mu = E[X]$

-1	0	+1	-10	0	+10
$\frac{1}{2}$		$\frac{1}{2}$	$\frac{1}{2}$		$\frac{1}{2}$

$V[X] = E[(X-\mu)^2]$
 $= E[X^2 - 2XM + M^2]$
 $= E[X^2] - 2E[X]M + M^2$
 $= E[X^2] - 2M^2 + M^2 = E[X^2] - M^2$
 Variance = second moment.

Bernoulli distribution

Flip coin with prob p for heads: Get 1 otherwise 0.
 $X \sim \text{Bernoulli}(p)$
 $E[X] = 0 \cdot P[X=0] + 1 \cdot P[X=1] = 0 \cdot (1-p) + 1 \cdot p = p$.
 $V[X] = pq$.

Binomial distribution

$b(k; n, p) = P[\text{getting } k \text{ heads}] = \binom{n}{k} p^k (1-p)^{n-k}$
 $\binom{n}{k}$ = # flips prob in the k trials

$Y = X_1 + X_2 + \dots + X_n$, X_1, \dots, X_n are Bernoulli(p) independent
 $E[Y] = np$
 $V[Y] = V[\sum X_i] = \sum V[X_i] = npq$ $q = (1-p)$.
 because X_i 's are independent.

Geometric dist

$X \sim \text{Geom}(p)$ $X \in \{1, 2, 3, \dots\}$
 $E[X] = \frac{1}{p}$
 $V[X] = \frac{1-p}{p^2}$

Approximating k-SAT

kSAT: a formula that is CNF
 $\wedge(\dots)$
 every clause is the \vee of exactly k literals, all literals in a clause use different variables.
 F: Given kSAT formula assume it uses n variables and it has m clauses
 max kSAT
 Compute an assignment that satisfy the largest number of clauses.

Theorem 2SAT is NP-Hard.

PCP Theorem

Given a 3SAT formula with m clauses unless $P=NP$ one can not in polynomial time compute a satisfying assignment that satisfies more than $\frac{7}{8}m$ of the clauses.

F: 3SAT, m clauses
 $C = (x \vee y \vee z)$
 Pick a random assignment $x \in \{0, 1\}$ with equal prob
 $P[C \text{ is satisfied}] = P[(x \vee y \vee z) = 1] = 1 - \frac{P[x=0]P[y=0]P[z=0]}{1} = 1 - \frac{1}{8}$
 $P[(x \vee y \vee z) = 0] = P[x=0]P[y=0]P[z=0] = \frac{1}{2} \cdot \frac{1}{2} \cdot \frac{1}{2} = \frac{1}{8}$

$P[x \vee y \vee z = 0] = P[x=0 \wedge y=0 \wedge z=0]$
 $= P[x=0] P[y=0] P[z=0]$.

$F = C_1 \wedge C_2 \wedge \dots \wedge C_m$
 $P[C_i = 1 \text{ (its satisfied)}] = \frac{7}{8}$

$X_i = 1 \Leftrightarrow C_i$ is satisfied
 $= 0$ otherwise
 X_i is an indicator variable.

$Y = \# \text{ of satisfied clauses} = \sum_{i=1}^m X_i$
 $E[Y] =$ how many clauses are satisfied on average
 $= E[\sum_{i=1}^m X_i] = \sum_{i=1}^m E[X_i] = \sum_{i=1}^m \frac{7}{8} = \frac{7}{8} m$.
by linearity of expectations

$E[X_i] = 0 \cdot P[X_i=0] + 1 \cdot P[X_i=1]$
 $= P[X_i=1] = \frac{7}{8}$

Markov inequality

X : RV with real values
 $P[X \geq t] \leq \frac{E[X]}{t}$, $t \geq E[X]$
 proof

kSAT \neq with m clauses
 $Z = m - Y = \# \text{ of clauses not satisfied by the assignment}$
 $E[Z] = m - \frac{7}{8}m = \frac{1}{8}m$
 $\epsilon \in (0, 1)$ $\epsilon = 0.1$
 $P[Z \geq (1+\epsilon) \frac{1}{8}m] = P[Z \geq (1+\epsilon) E[Z]] \leq \frac{E[Z]}{(1+\epsilon) E[Z]}$
 $= \frac{1}{1+\epsilon} \leq 1 - \frac{\epsilon}{2}$
 $\leq (1 - \frac{\epsilon}{2})^t \rightarrow 0$
 $t = O(\frac{1}{\epsilon^2})$