

# Parallel Numerical Algorithms

## Chapter 11 – QR Factorization

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CS 554 / CSE 512

# Outline

- 1 QR Factorization
- 2 Householder Transformations
- 3 Givens Rotations

## QR Factorization

- For given  $m \times n$  matrix  $A$ , with  $m > n$ , **QR factorization** has form

$$A = Q \begin{bmatrix} R \\ O \end{bmatrix}$$

where matrix  $Q$  is  $m \times m$  and orthogonal, and  $R$  is  $n \times n$  and upper triangular

- Can be used to solve linear systems, least squares problems, etc.
- As with Gaussian elimination, zeros are introduced successively into matrix  $A$ , eventually reaching upper triangular form, but using orthogonal transformations instead of elementary eliminators

## Householder Transformations

- Householder transformation** has form

$$H = I - 2 \frac{vv^T}{v^T v}$$

where  $v$  is nonzero vector

- From definition,  $H = H^T = H^{-1}$ , so  $H$  is both orthogonal and symmetric
- For given vector  $a$ , choose  $v$  so that

$$Ha = \begin{bmatrix} \alpha \\ 0 \\ \vdots \\ 0 \end{bmatrix} = \alpha \begin{bmatrix} 1 \\ 0 \\ \vdots \\ 0 \end{bmatrix} = \alpha e_1$$

## Householder QR Factorization

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for k = 1 to n
     $\alpha_k = -\text{sign}(a_{kk}) \sqrt{a_{kk}^2 + \dots + a_{mk}^2}$ 
     $v_k = [0 \ \dots \ 0 \ a_{kk} \ \dots \ a_{mk}]^T - \alpha_k e_k$ 
     $\beta_k = v_k^T v_k$ 
    if  $\beta_k = 0$  then
        continue with next k
    for j = k to n
         $\gamma_j = v_k^T a_j$ 
         $a_j = a_j - (2\gamma_j / \beta_k) v_k$ 
    end
end
    
```

## Methods for QR Factorization

- Householder transformations (elementary reflectors)
- Givens transformations (plane rotations)
- Gram-Schmidt orthogonalization

## Householder Transformations

- Substituting into formula for  $H$ , we see that we can take

$$v = a - \alpha e_1$$

and to preserve norm we must have  $\alpha = \pm \|a\|_2$ , with sign chosen to avoid cancellation

## Parallel Householder QR

- Householder QR factorization is similar to Gaussian elimination for LU factorization
- Forming Householder vector  $v_k$  is analogous to computing multipliers in Gaussian elimination
- Subsequent updating of remaining unreduced portion of matrix is also analogous to Gaussian elimination
- Thus, parallel implementation is similar to parallel LU, but with Householder vectors broadcast horizontally instead of multipliers
- For this reason, we will not go into details

## Givens Rotations

- **Givens rotation** operates on pair of rows to introduce single zero
- For given 2-vector  $\mathbf{a} = [a_1 \ a_2]^T$ , if

$$c = \frac{a_1}{\sqrt{a_1^2 + a_2^2}}, \quad s = \frac{a_2}{\sqrt{a_1^2 + a_2^2}}$$

then

$$\mathbf{G}\mathbf{a} = \begin{bmatrix} c & s \\ -s & c \end{bmatrix} \begin{bmatrix} a_1 \\ a_2 \end{bmatrix} = \begin{bmatrix} \alpha \\ 0 \end{bmatrix}$$

- Scalars  $c$  and  $s$  are cosine and sine of angle of rotation, and  $c^2 + s^2 = 1$ , so  $\mathbf{G}$  is orthogonal

## Parallel Givens QR Factorization

- With 1-D partitioning of  $\mathbf{A}$  by columns, parallel implementation of Givens QR factorization is similar to parallel Householder QR factorization, with cosines and sines broadcast horizontally and each task updating its portion of relevant rows
- With 1-D partitioning of  $\mathbf{A}$  by rows, broadcast of cosines and sines is unnecessary, but there is no parallelism unless multiple pairs of rows are processed simultaneously
- Fortunately, it is possible to process multiple pairs of rows simultaneously without interfering with each other

## Parallel Givens QR Factorization

- Communication cost is high, but can be reduced by having each task initially reduce its entire local set of rows to upper triangular form, which requires no communication
- Then, in subsequent phase, task pairs cooperate in annihilating additional entries using one row from each of two tasks, exchanging data as necessary
- Various strategies can be used for combining results of first phase, depending on underlying network topology
- With hypercube, for example, final upper triangular form can be reached in  $\log p$  combining steps

## References

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## Givens QR Factorization

- Givens rotations can be systematically applied to successive pairs of rows of matrix  $\mathbf{A}$  to zero entire strict lower triangle
- Subdiagonal entries of matrix can be annihilated in various possible orderings (but once introduced, zeros should be preserved)
- Each rotation must be applied to all entries in relevant pair of rows, not just entries determining  $c$  and  $s$
- Once upper triangular form is reached, product of rotations,  $\mathbf{Q}$ , is orthogonal, so we have QR factorization of  $\mathbf{A}$

## Parallel Givens QR Factorization

- Stage at which each subdiagonal entry can be annihilated is shown here for  $8 \times 8$  example

$$\begin{bmatrix} \times & & & & & & & \\ 7 & \times & & & & & & \\ 6 & 8 & \times & & & & & \\ 5 & 7 & 9 & \times & & & & \\ 4 & 6 & 8 & 10 & \times & & & \\ 3 & 5 & 7 & 9 & 11 & \times & & \\ 2 & 4 & 6 & 8 & 10 & 12 & \times & \\ 1 & 3 & 5 & 7 & 9 & 11 & 13 & \times \end{bmatrix}$$

- Maximum parallelism is  $n/2$  at stage  $n - 1$  for  $n \times n$  matrix

## Parallel Givens QR Factorization

- With 2-D partitioning of  $\mathbf{A}$ , parallel implementation combines features of 1-D column and 1-D row algorithms
- In particular, sets of rows can be processed simultaneously to annihilate multiple entries, but updating of rows requires horizontal broadcast of cosines and sines

## References

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