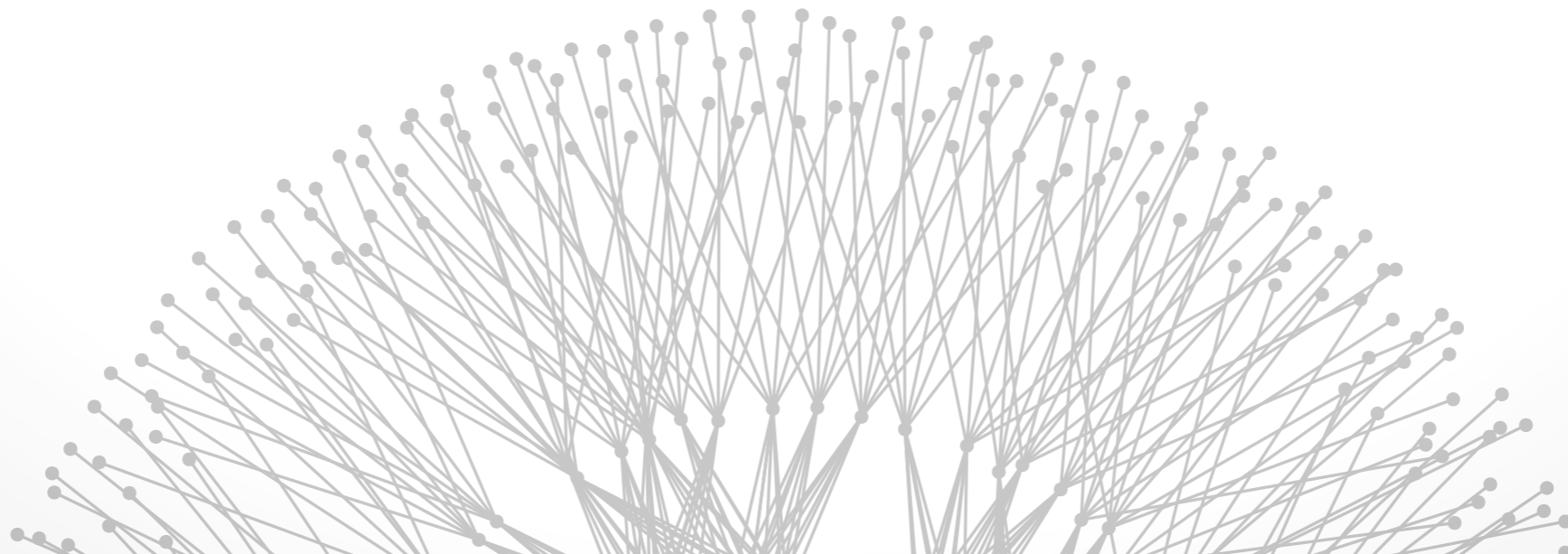


# Forwarding Architecture

Brighten Godfrey  
CS 538 February 14 2018



# Building a fast router

# Partridge: 50 Gb/sec router



## A fast IP router

- well, fast at the time...

Good exhibition of the guts of a router and problems to be solved in router architecture

# Routing vs. forwarding



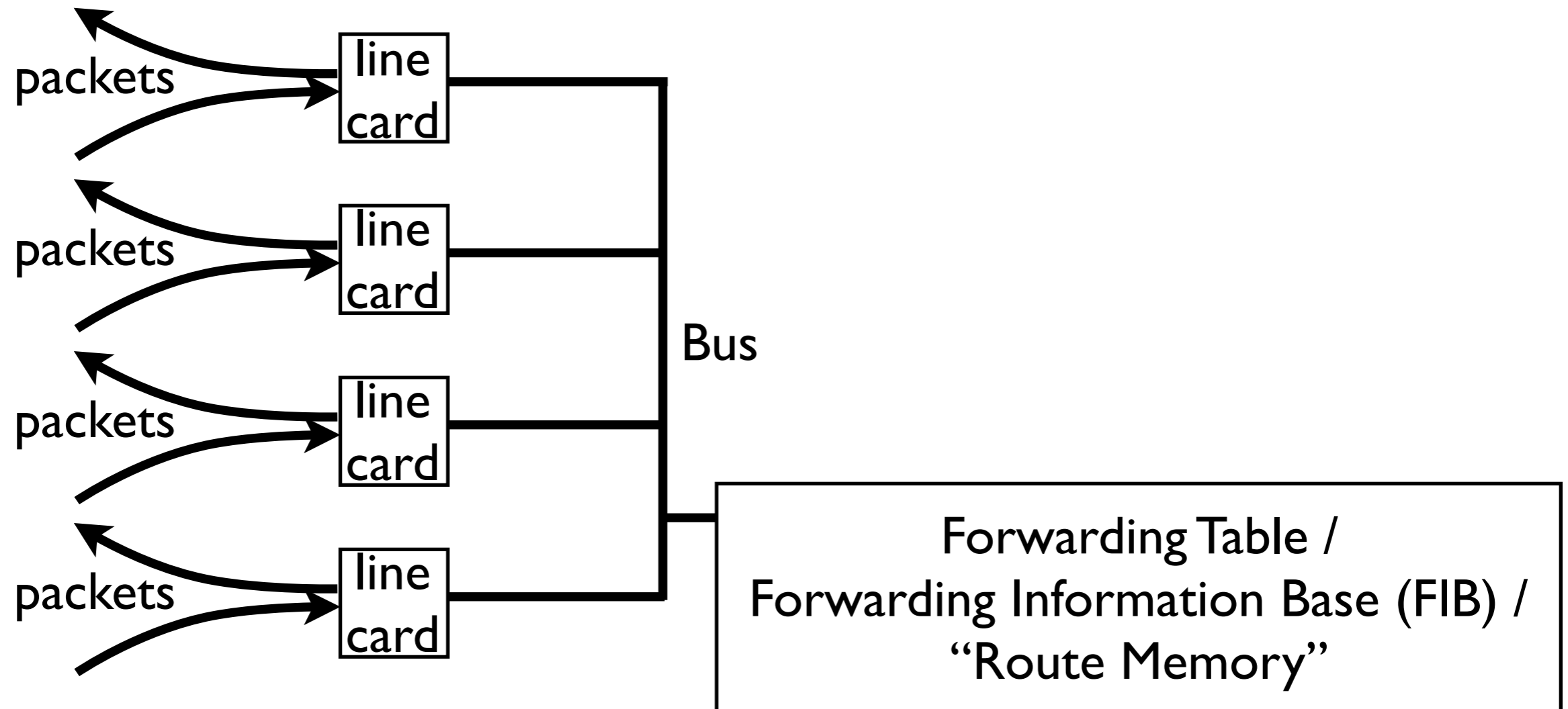
## Control plane ← Upcoming lectures

- Decides how data should flow through the network
- Uses operator configs & distributed routing protocols
- Can use commodity hardware (often Linux today)
- Output: Forwarding table
  - Contains instructions for what to do with each packet based on its header
  - Relatively simple data structure (lookup table, tree, ...)

## Data plane ← Today

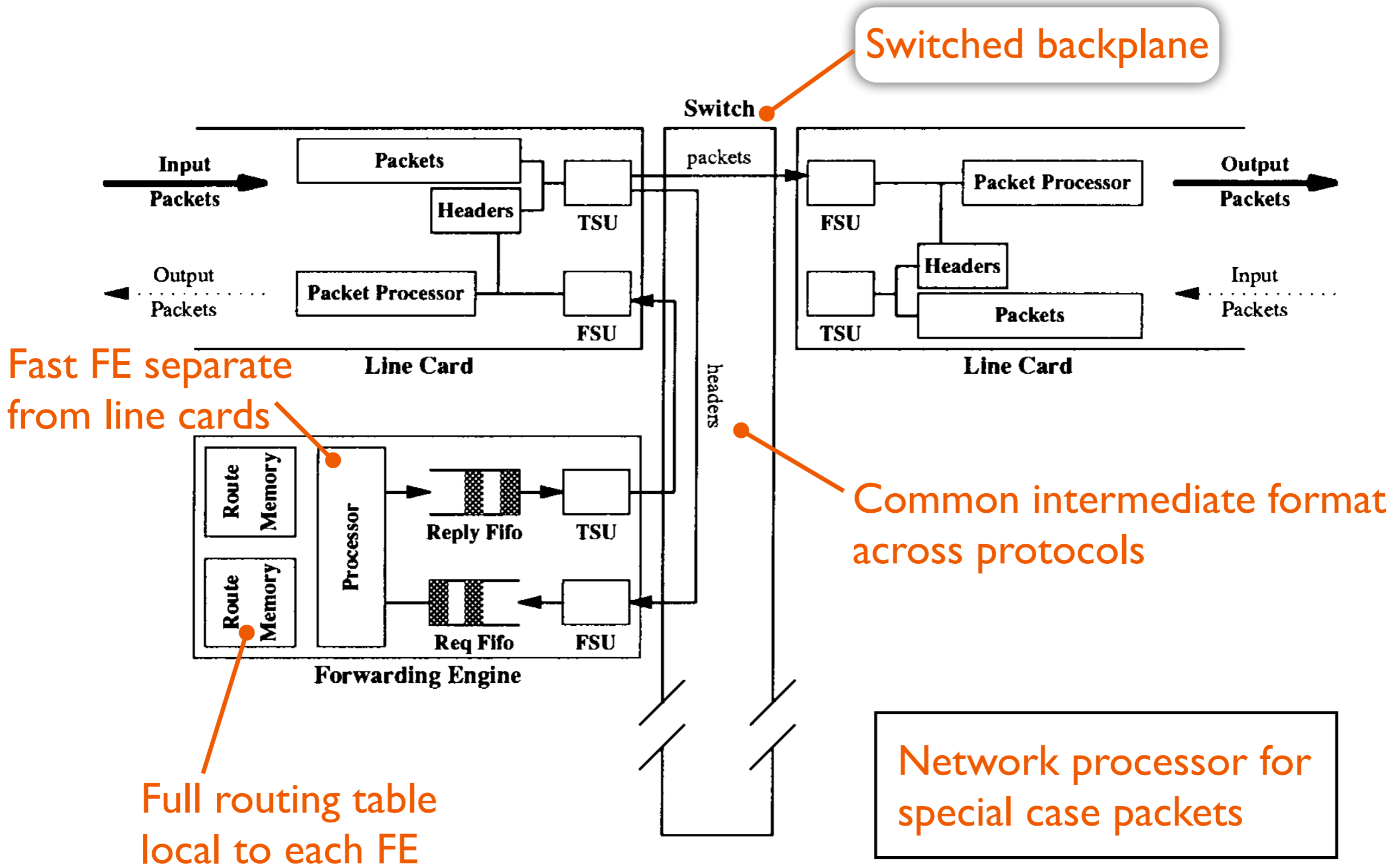
- Input: Forwarding table
- Forwards data through the network
- Typically specialized ASIC hardware for fast forwarding

# Simplistic router diagram



What's wrong with this picture?

# Inside the router



# Switching fabric



## Operates in epochs

- 128 bytes sent by each line card to next-hop line card
- Each line card can send to only one other card, and can receive from only one other card

... to outputs

Inputs ready to send...

	1	2	3	4	5	6	7
1	0	0	1	0	1	0	0
2	1	1	1	1	0	0	0
3	0	1	1	1	0	0	1
4	1	0	1	0	0	1	0
5	0	0	0	0	0	0	0
6	0	0	0	0	1	0	1
7	0	1	1	0	0	1	0

Allocator assigns inputs to outputs & tells line cards

What fundamental problem is being solved?

# Maximum bipartite matching



Maximize number of matched input-output pairs, such that each input & output only matched once

50 Gbit/s router uses approximate solution

Inputs ready to send...

... to outputs

	1	2	3	4	5	6	7
1	0	0	1	0	1	0	0
2	1	1	1	1	0	0	0
3	0	1	1	1	0	0	1
4	1	0	1	0	0	1	0
5	0	0	0	0	0	0	0
6	0	0	0	0	1	0	1
7	0	1	1	0	0	1	0



# Beyond the 50 Gb/sec router

# Many more problems to solve



Buffering (how big?)

Queueing (where? what order?)

FIB memory (how fast, large, and expensive?)

FIB algorithms (what data structure? how many memory accesses?)

Specialized functionality (how to do it efficiently?)



Modern routers are jacks-of-all-trades:

- Load balance across links
- Access control
- Filtering attacks
- Quality of service
- Accounting, traffic metering
- Virtual private networks
- Protocol support: IPv4, IPv6, MPLS, ethernet, ...
- ...

Can we make forwarding **flexible, extensible**?

# Data plane flexibility over the ages



## 1990s: Software routers

- Zebra, later Quagga, Click, Xorp

## 1997: Label switching (MPLS)

- Set up explicit paths for classes of packets

## 1999: Active networks

- Packet header carries (pointer to) program code

## 2008: Software Defined Networks

- Open interface to data plane, programmed by software controller

# Data plane flexibility over the ages

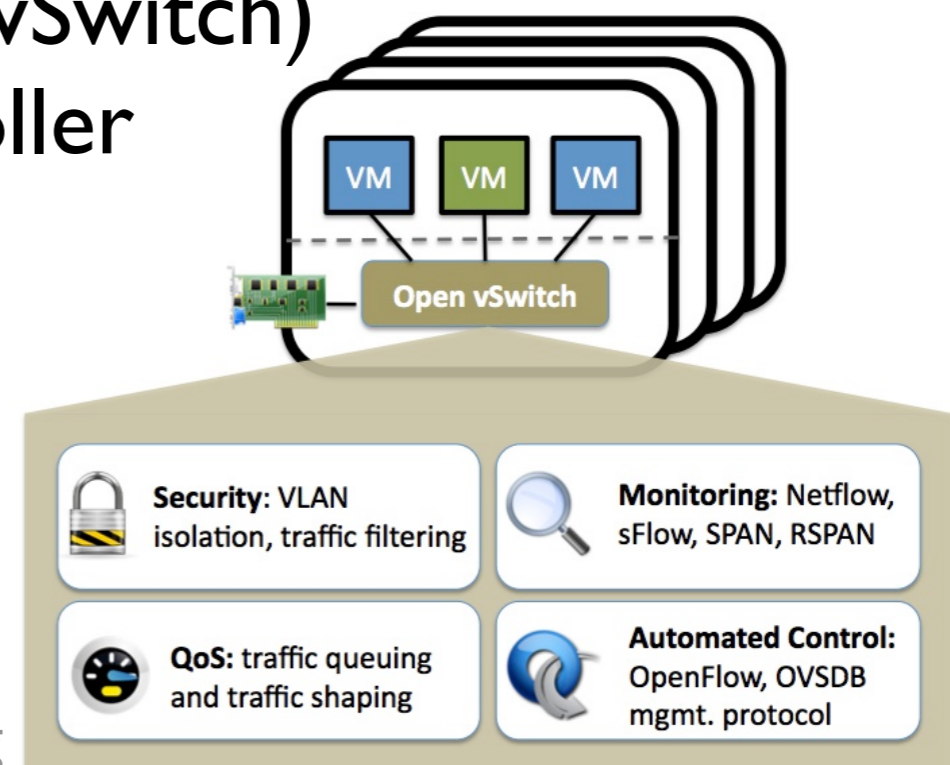


## 2008: Software Defined Networks

- Open interface to data plane, programmed by software controller

## 2009: Network virtualization

- Data plane complexity moved to edge devices: software in servers (Hypervisor, Open vSwitch)
- Managed by SDN-style controller





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## 2014: Next-gen programmable switches

# Scaling software routers



## Hardware routers

- Fast
- Specific functionality
- Result: many physical devices (routers, firewalls, intrusion detection, ...)

## Software routers

- Slow
- Extensible

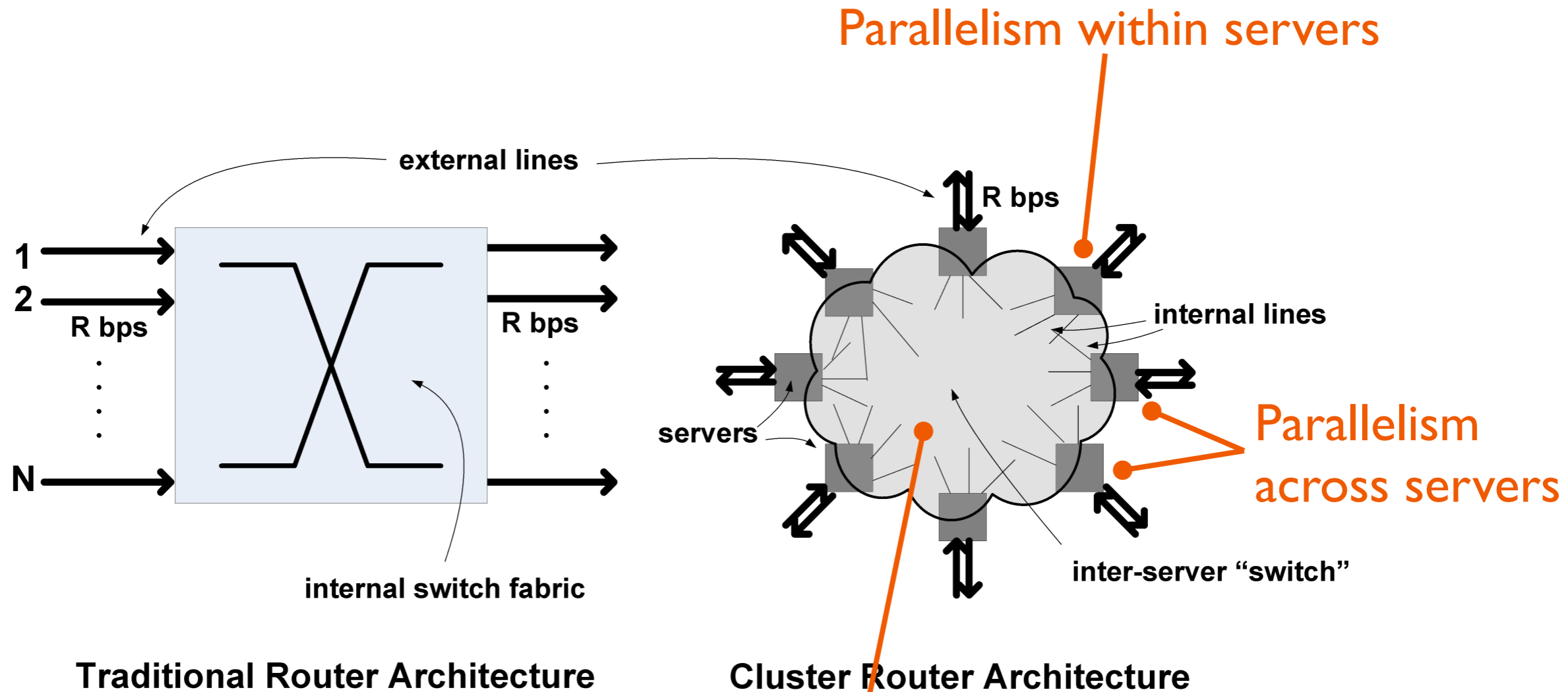
Can we get the best of both worlds?



# RouteBricks approach



[Dobrescu, Egi, Argyraki, Chun, Fall, Iannaccone, Knies, Manesh, Ratnasamy, NSDI 2009]



High bandwidth switching fabric built from commodity hardware

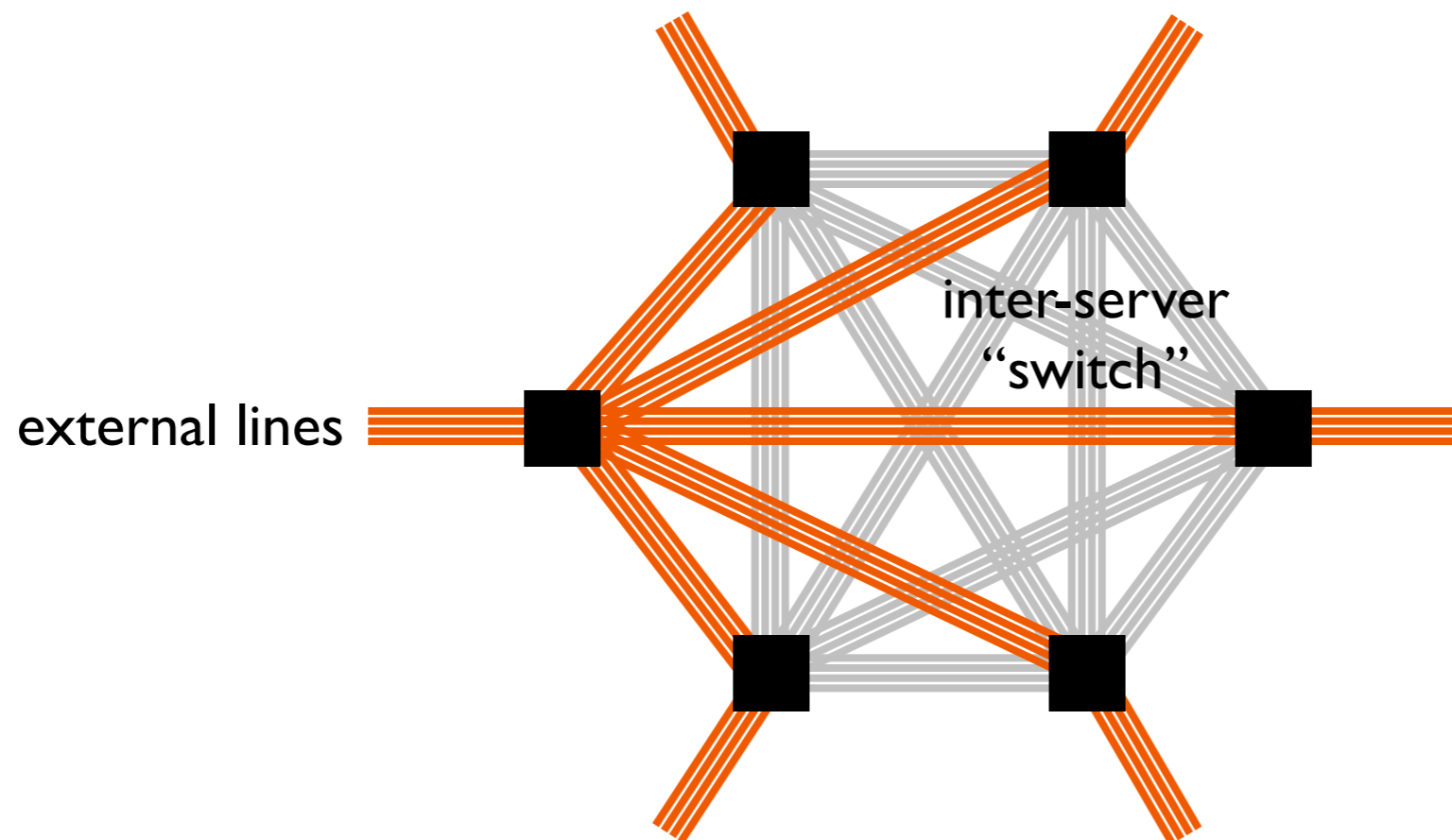
# Switching fabric challenges



**Handle any traffic pattern:** for example, all input traffic at a server might go to any one output server

**Low degree (ports):** we're using commodity hardware

Naïve approach:



Useless: might  
as well just use  
one server!

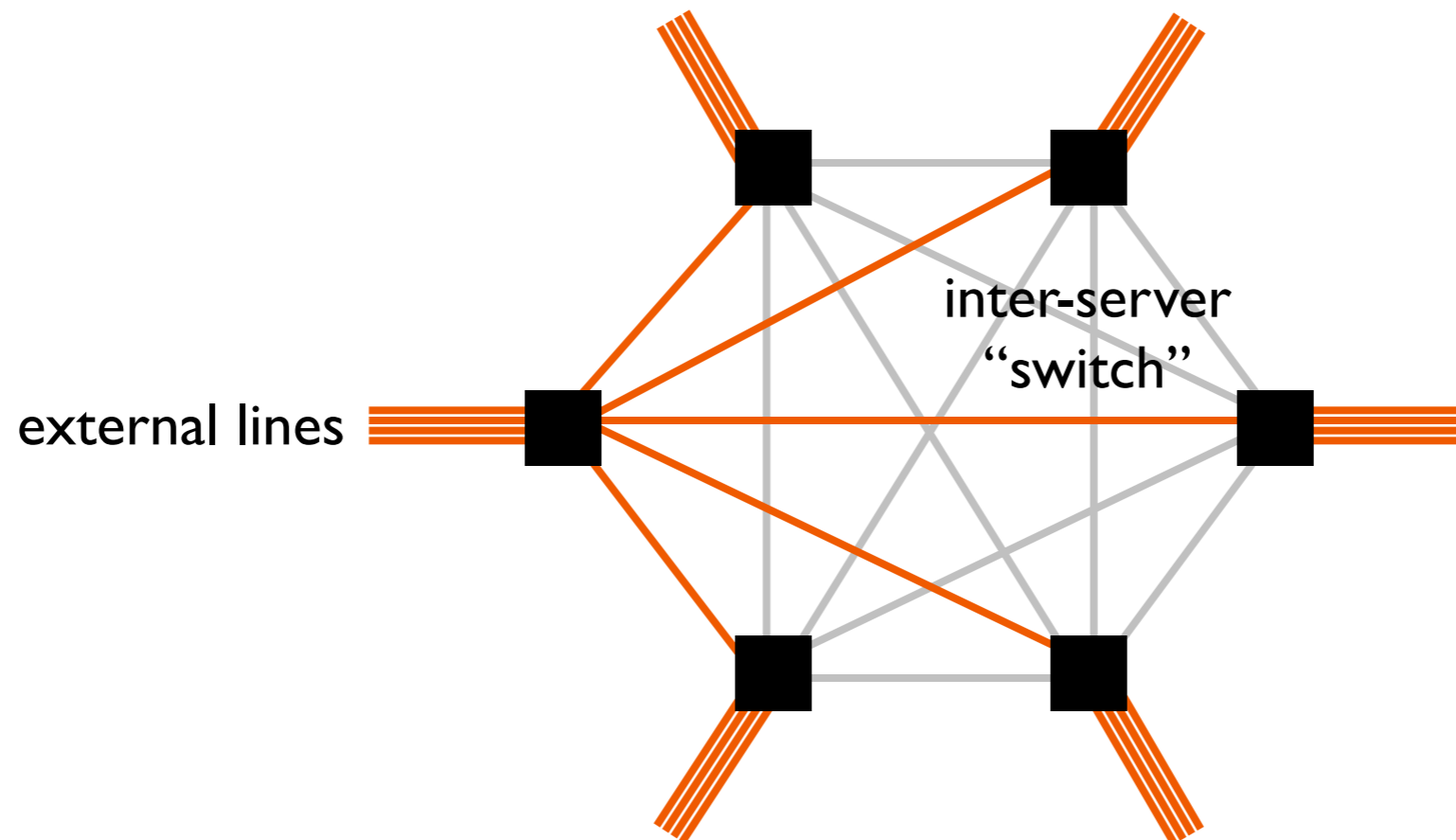
# Low degree solution



Just one link out for each link in

Total out b/w enough, but doesn't go where we need

Solution (**Valiant load balancing**): send packet to



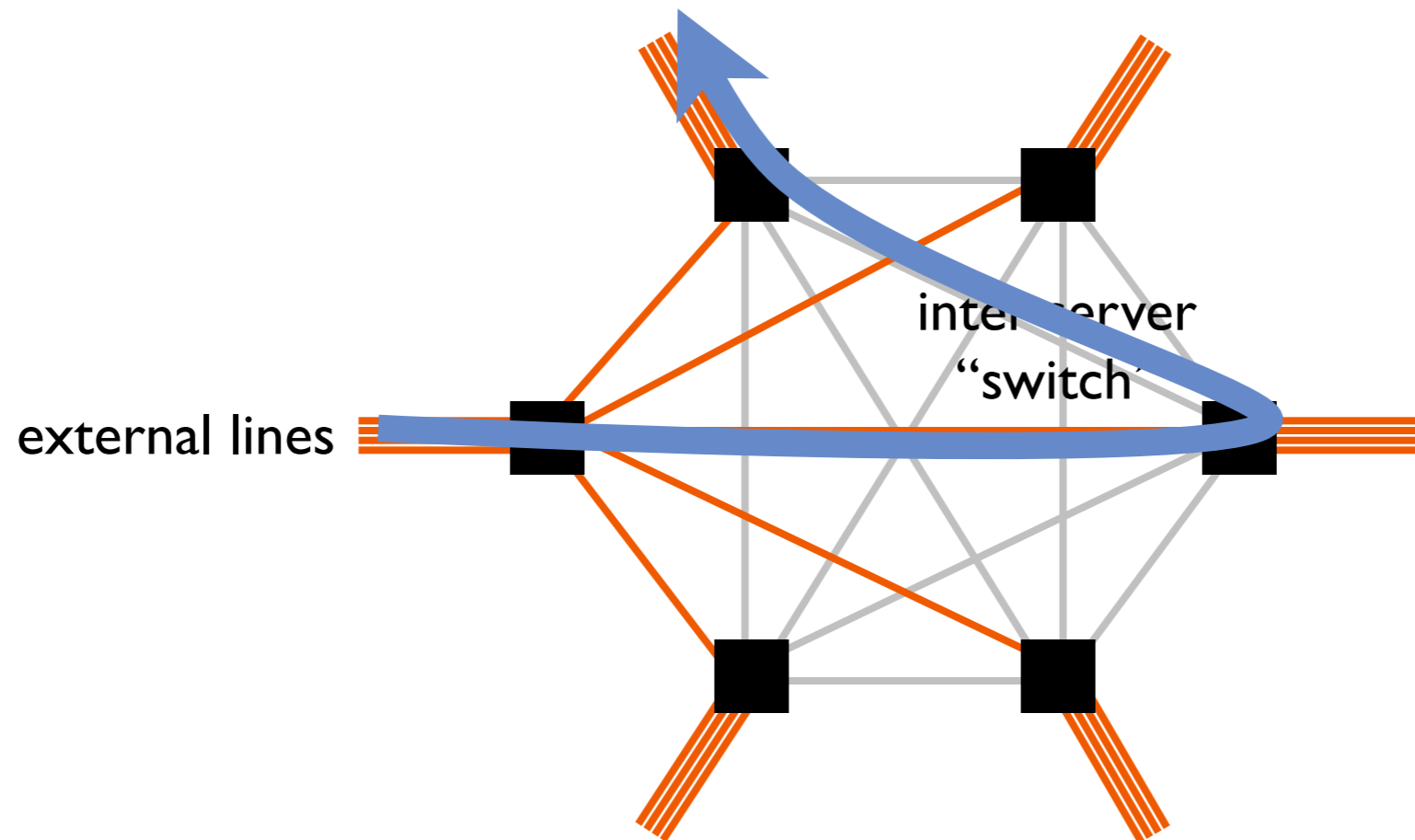
# Low degree solution



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# VLB guarantees & questions



Guaranteed to nearly full throughput for any traffic demands!

- “nearly” = 2x. Why?
- So, switch fabric needs to be 2x as fast as external links to provide guarantees

Why does sending to a random intermediate node work?



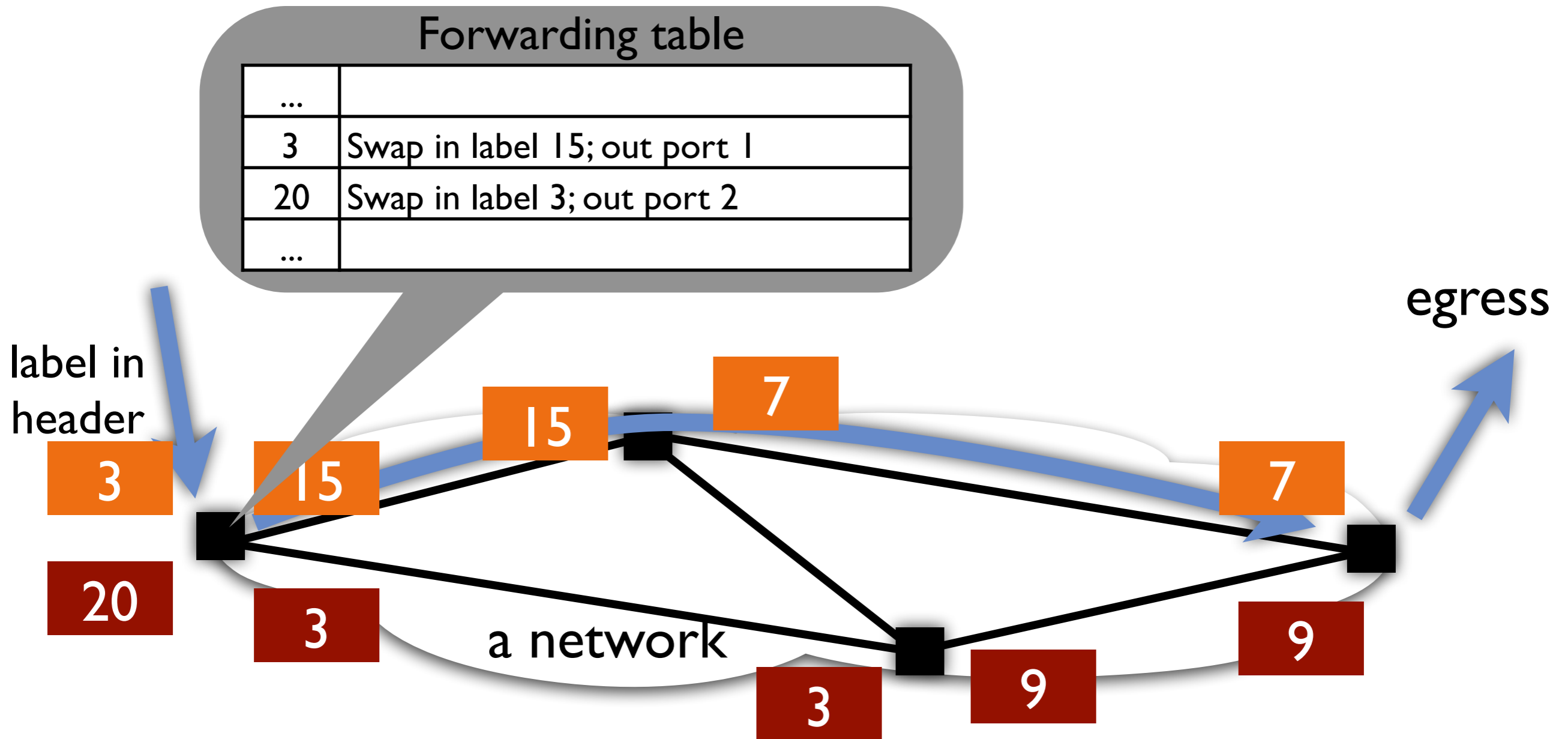
Software routers are not replacing big iron routers

However, increasing use of software in the fast path

- ‘Middleboxes’ / Virtual Network Functions: load balancing, logging, firewalls, intrusion detection, ...
- Software switches in software-defined data centers
- Solutions from multiple vendors to speed software processing of data flows (e.g. 10 Gbps)

**MPLS**

# MPLS design



Why is this more flexible than shortest path routing?

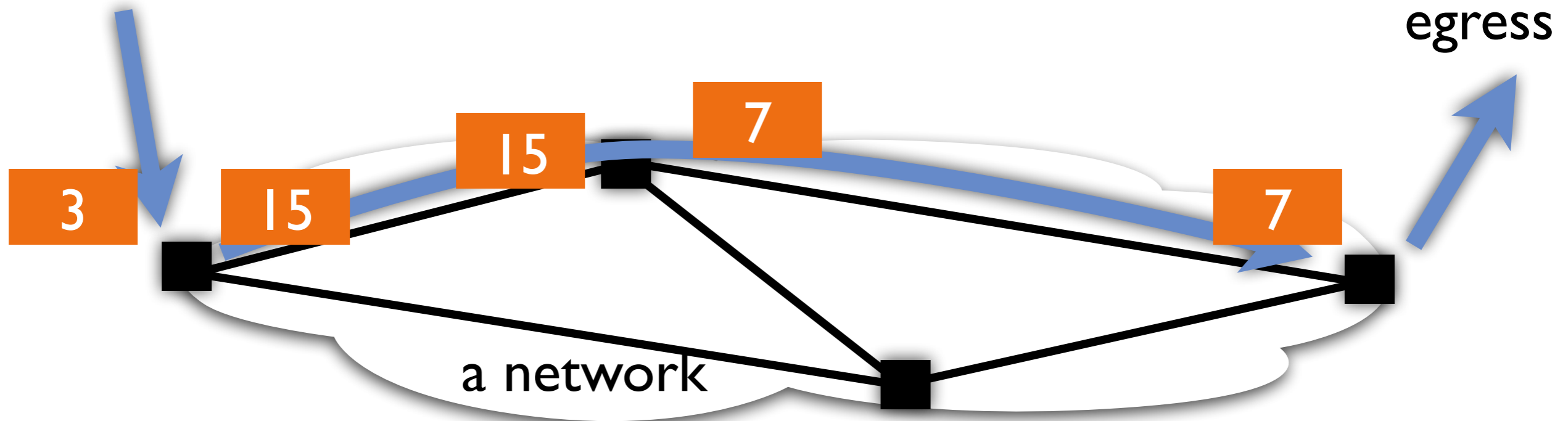


# MPLS design



Ingress:

Traffic classification, label packets (“forwarding equivalence class”)



Control plane constructs label-switched paths and coordinates labels

Can also stack labels = concatenate paths



## In the design doc

- High performance forwarding
- Minimal forwarding requirements, so can interface well with many types of media such as ATM
- Flexible control of traffic routing

## What matters today?

**Flexibility.** Widely used to achieve:

- Virtual Private Network (VPN) service along dedicated paths between enterprise sites
- Control backup paths with MPLS Fast ReRoute
- Traffic engineering (load balancing)



## Next Monday: Intradomain routing

- Papers will be posted by midnight tonight