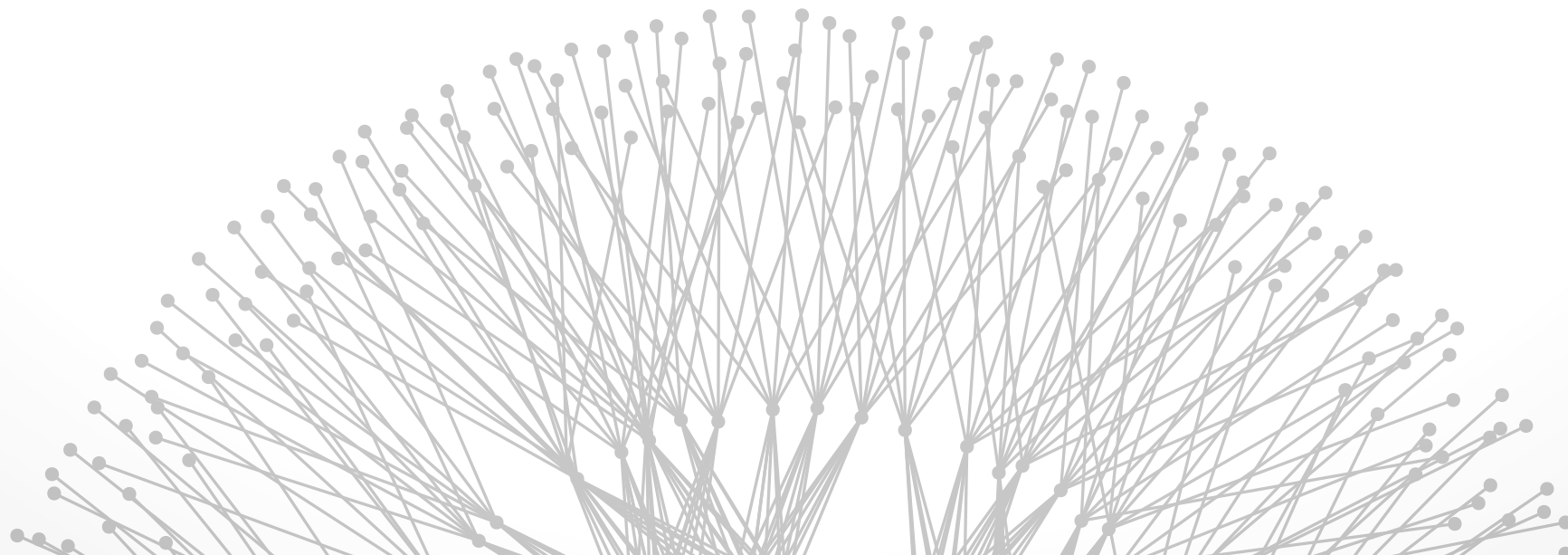


Denial of Service

Brighten Godfrey
CS 538 October 24 2013



but first,

Routing Security, continued

Not just malicious attackers



Many or most high-profile outages likely just configuration errors

Natural correspondence between attackers and bugs

- behavior unknown ahead of time
- defense is to limit and contain worst-case effects

What about a bug in the protocol?

- worst-case scenario: zero-day exploit on large fraction of routers across the entire Internet
- many are running the same software!

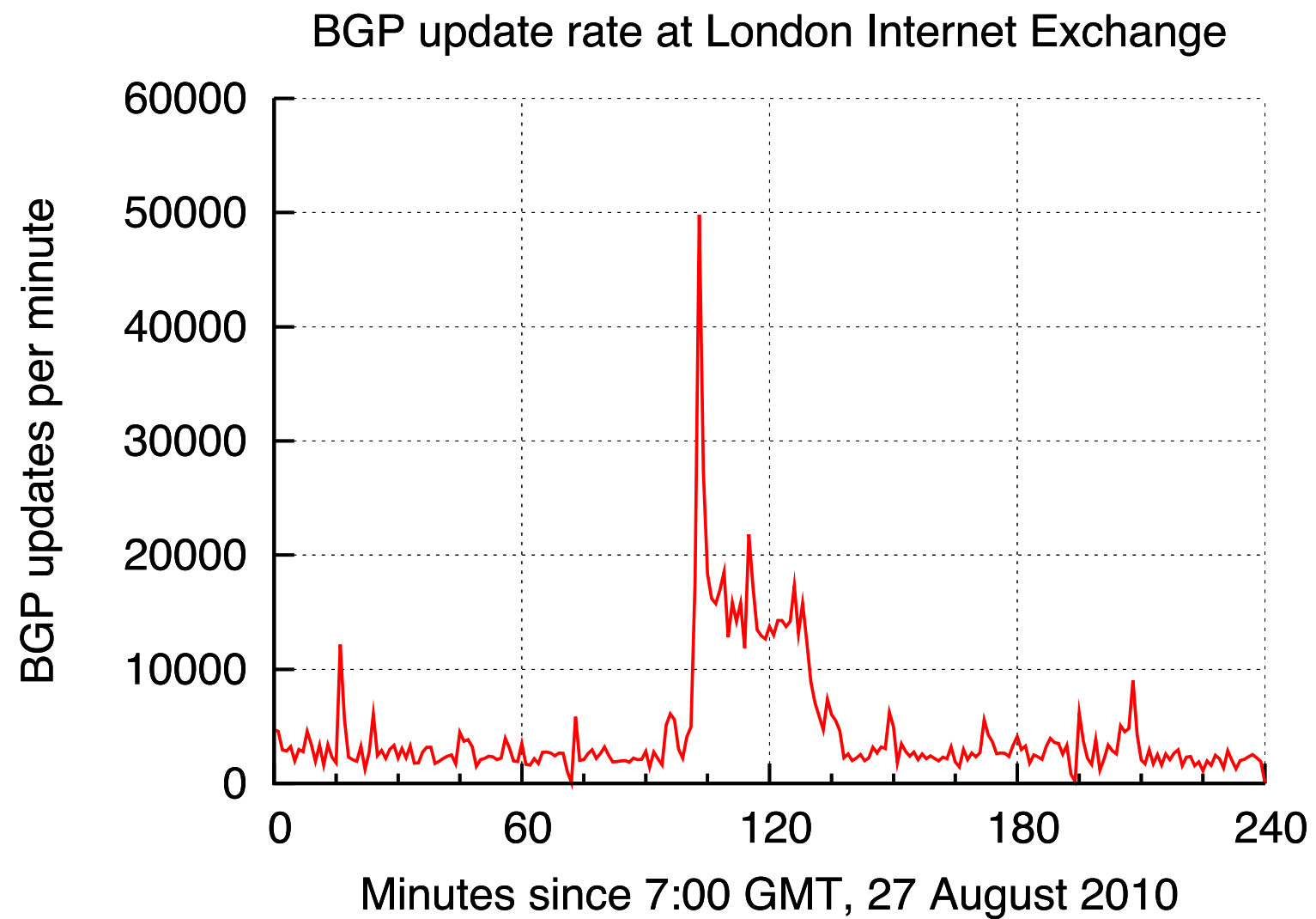
A (bad) day in the life of the Internet



About 1% of Internet destinations disrupted for about 30 minutes

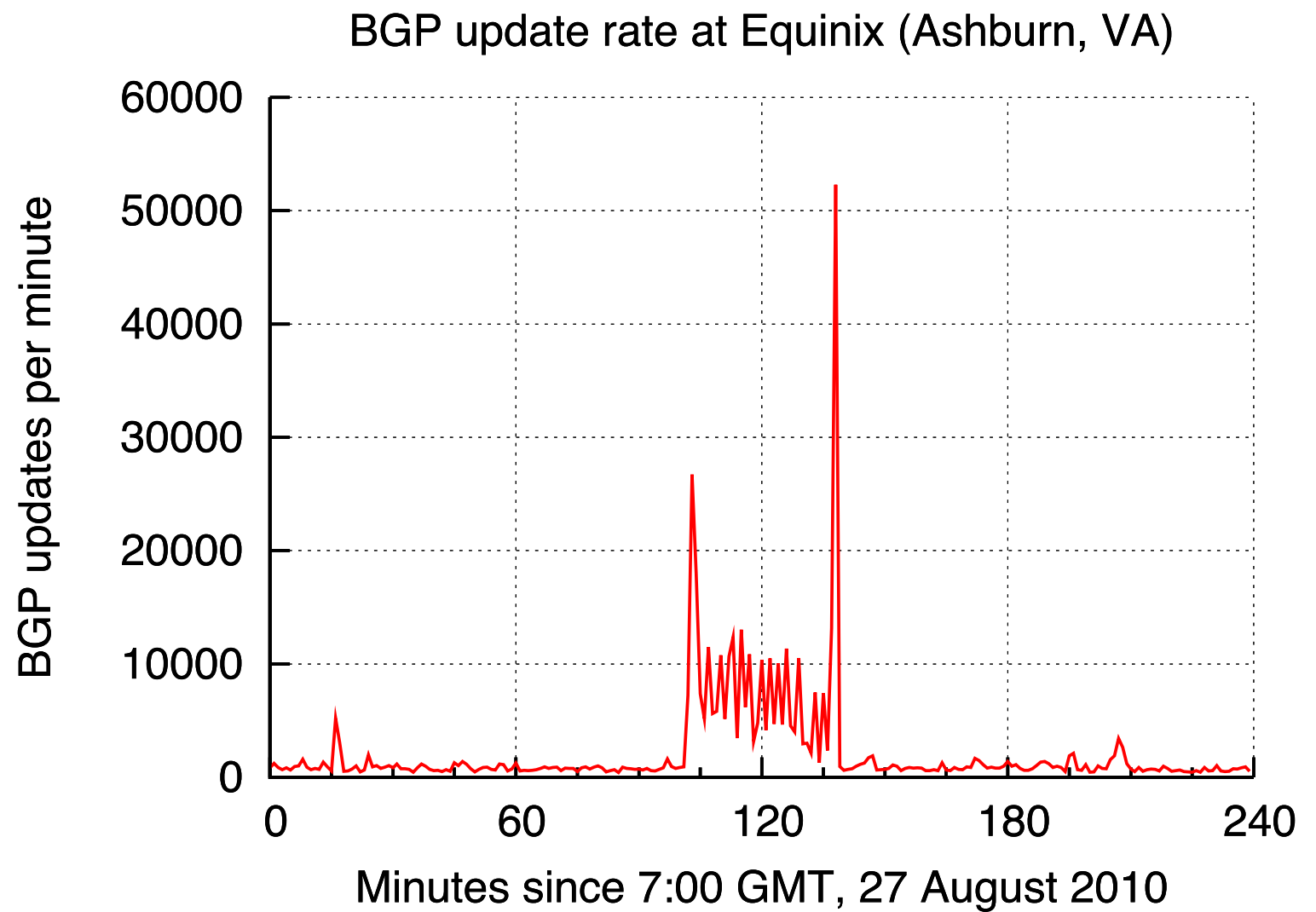
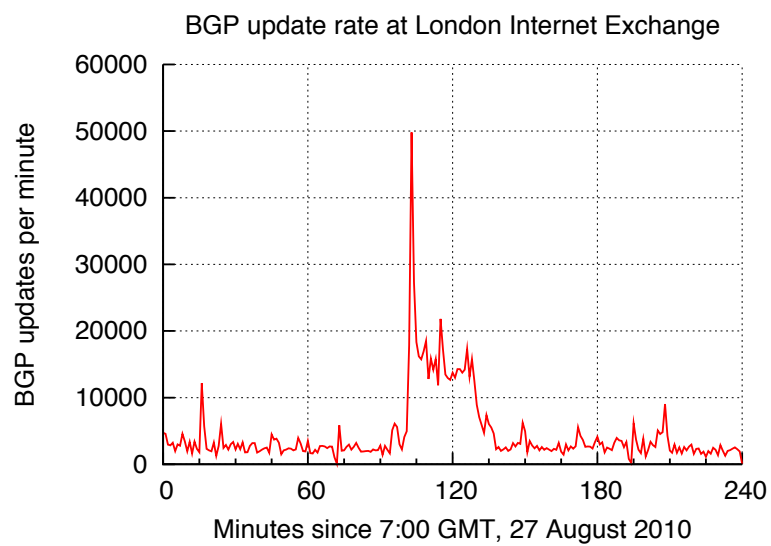
How did this happen?

Internet had a bad Friday



[Plots by Brighten based on raw update feeds from Route Views]

Internet had a bad Friday



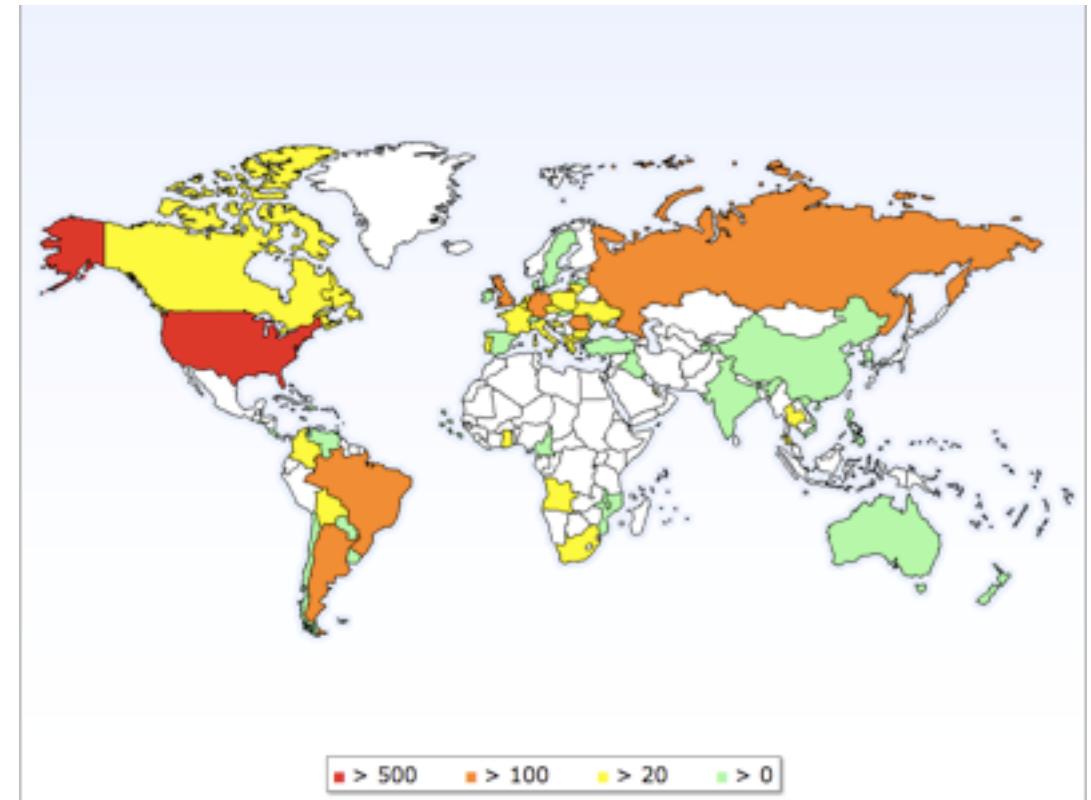
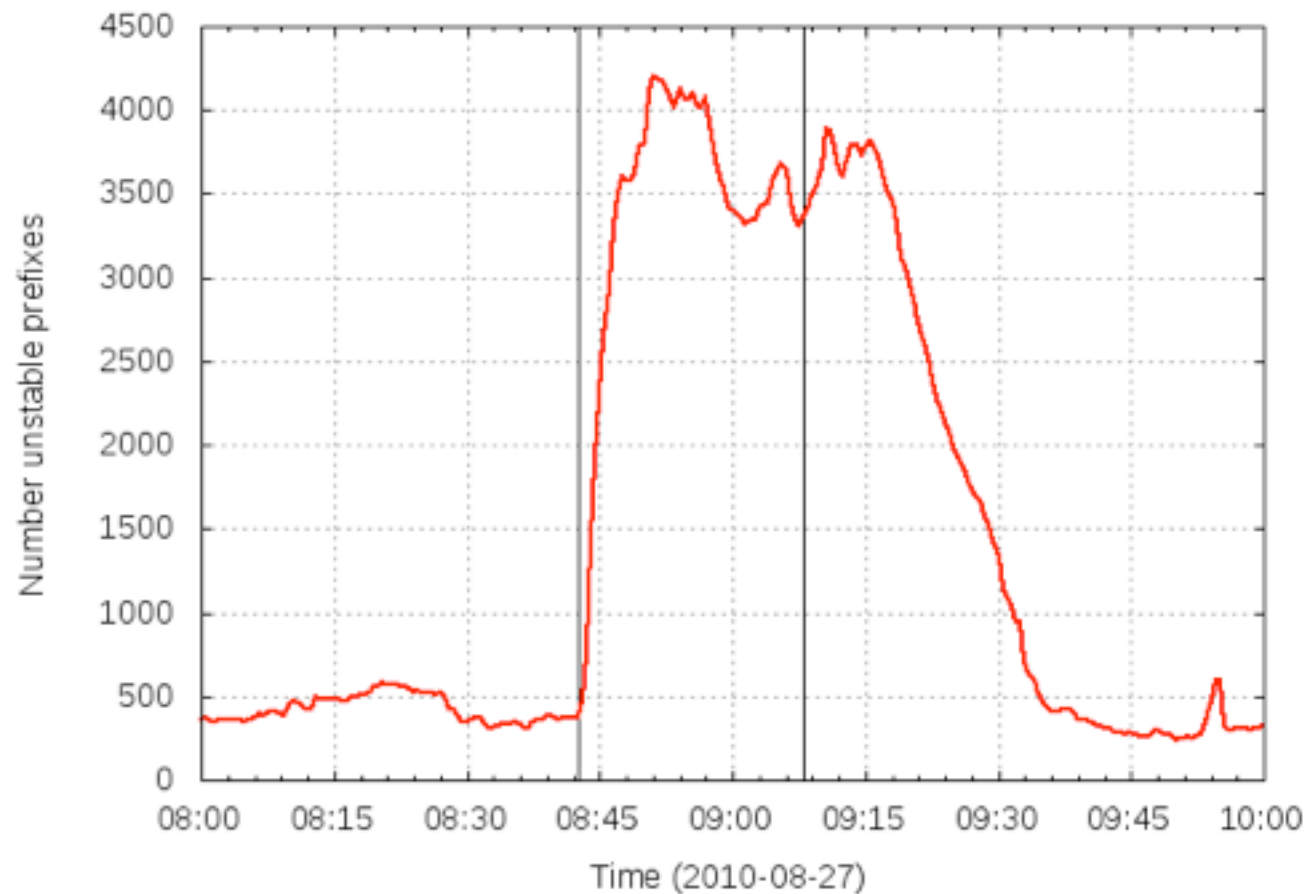
[Plots by Brighten based on raw update feeds from Route Views]

~1% of prefixes affected



[Earl Zmijewski, Renesys]

Unstable prefixes, 0800 to 1000 (UTC)



Brewing a storm



1. An unusual announcement
2. Propagation from router to router
3. Buggy software mangles announcement
4. `while(true)`
 1. Buggy router propagates announcement to neighbor
 2. BGP session dropped upon receipt of mangled message
 3. BGP session reestablished



Many unsavory BGP announcements can be contained, but this one wasn't

- Spread **geographically** because it was an entirely valid announcement
- Spread to **many prefixes** because BGP spec lets one bad announcement from a router affect all traffic to that router

Widespread correlated failures from similar software

Bugs and attacks can have similar effects and solutions

- Lucky in this case: bug triggered by researchers, not attackers!

Onward to Denial of Service

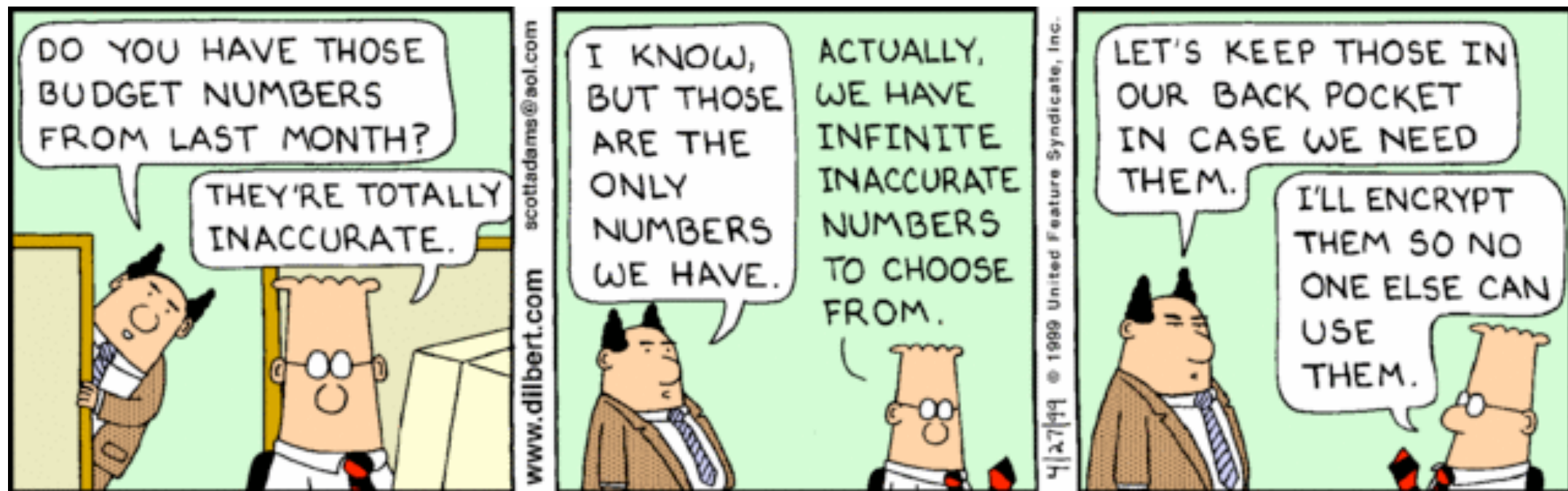
DoS in the real world



Source: Arbor Networks

Disclaimers:

- Survey of 130 network operators (mix of Tier 1, 2, 3; enterprise networks, etc.), not direct measurement
- Arbor sells network security solutions :-)



DDoS is frequent and can be big



Attack Frequency per Month

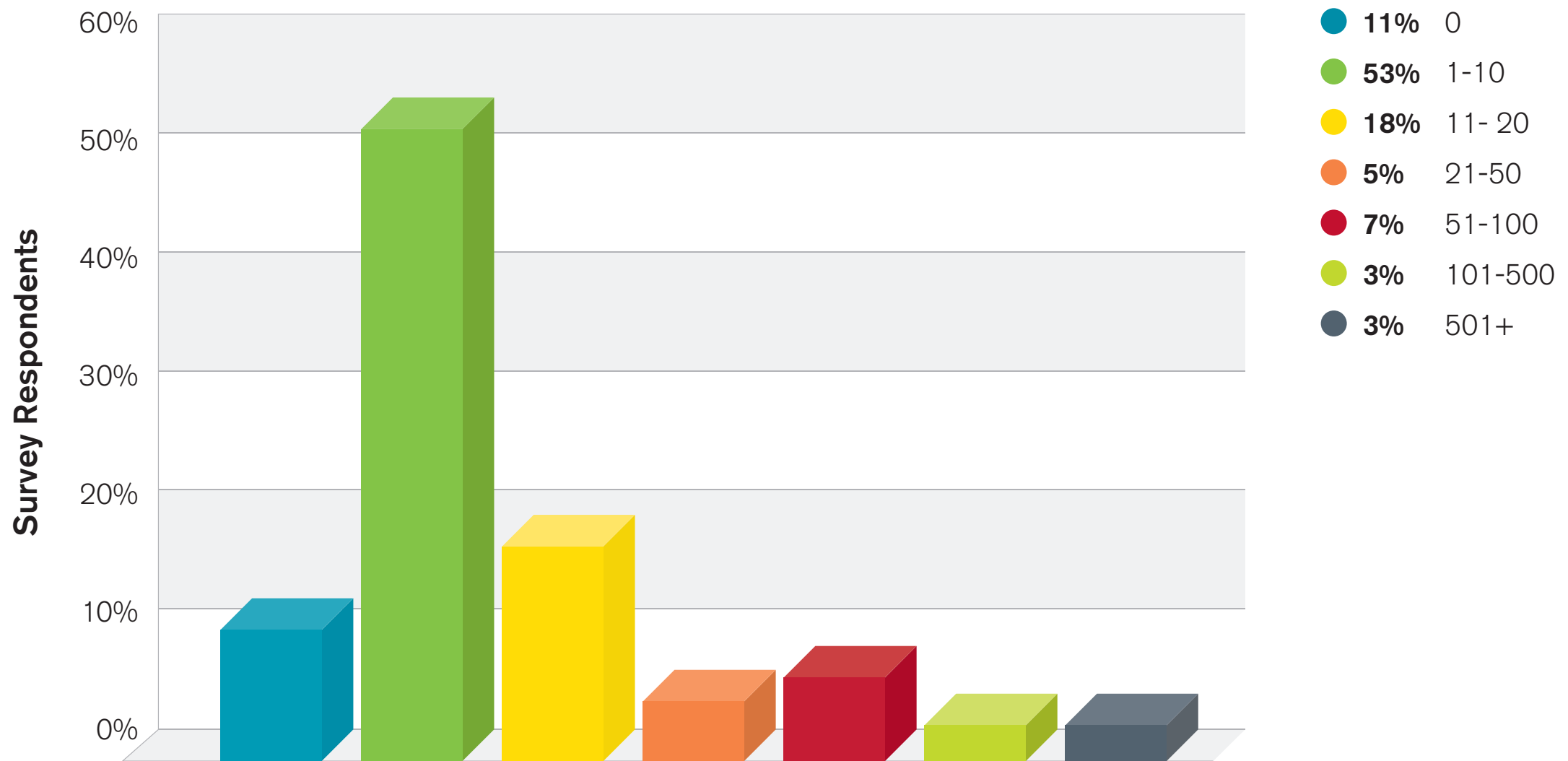


Figure 27 Source: Arbor Networks, Inc.

DDoS is frequent and can be big

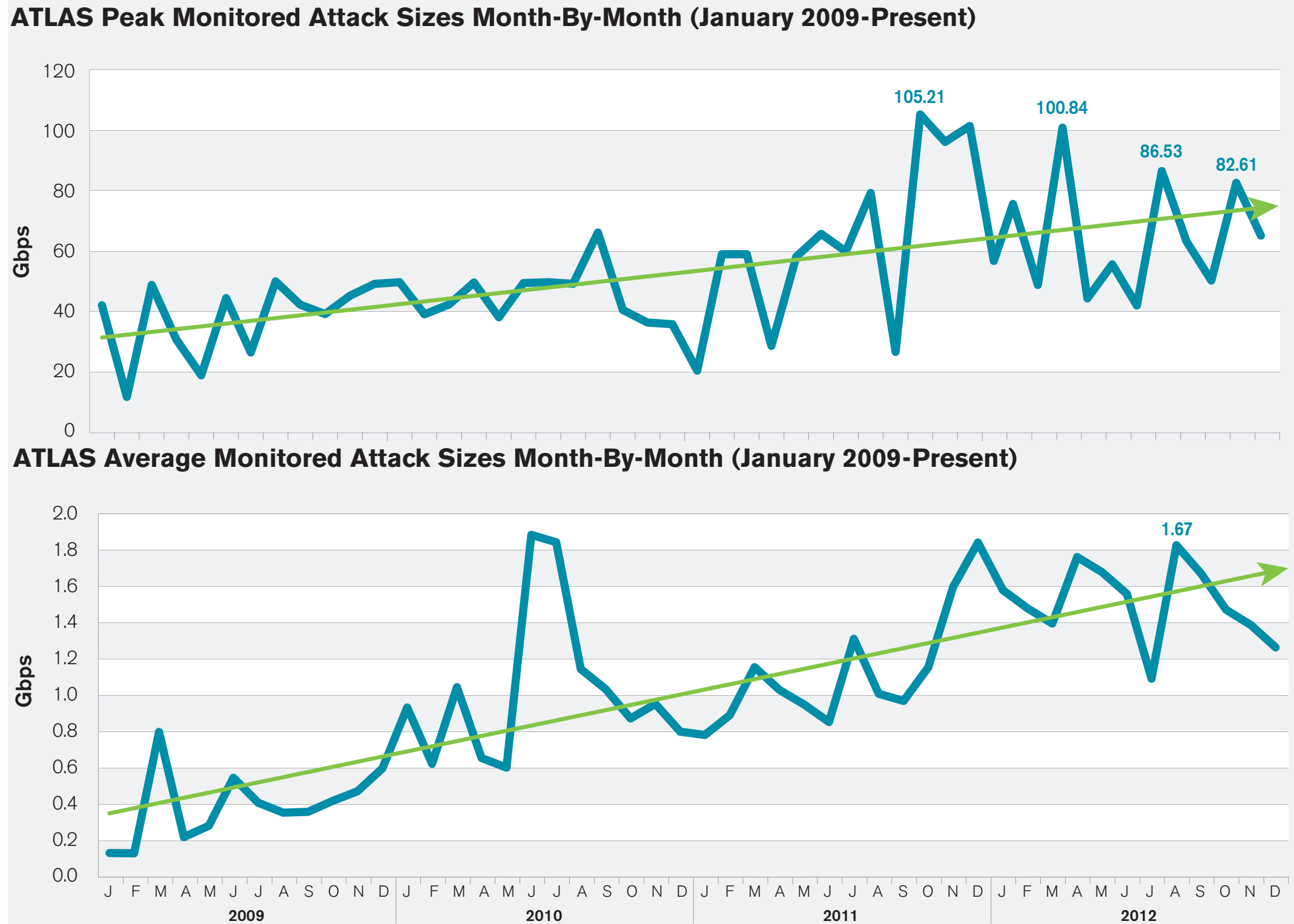


Figure 18 Source: Arbor Networks, Inc.

[Measurements from Arbor's ATLAS tool in 250 networks]

Many types of attacks



Targets of Application-Layer Attacks

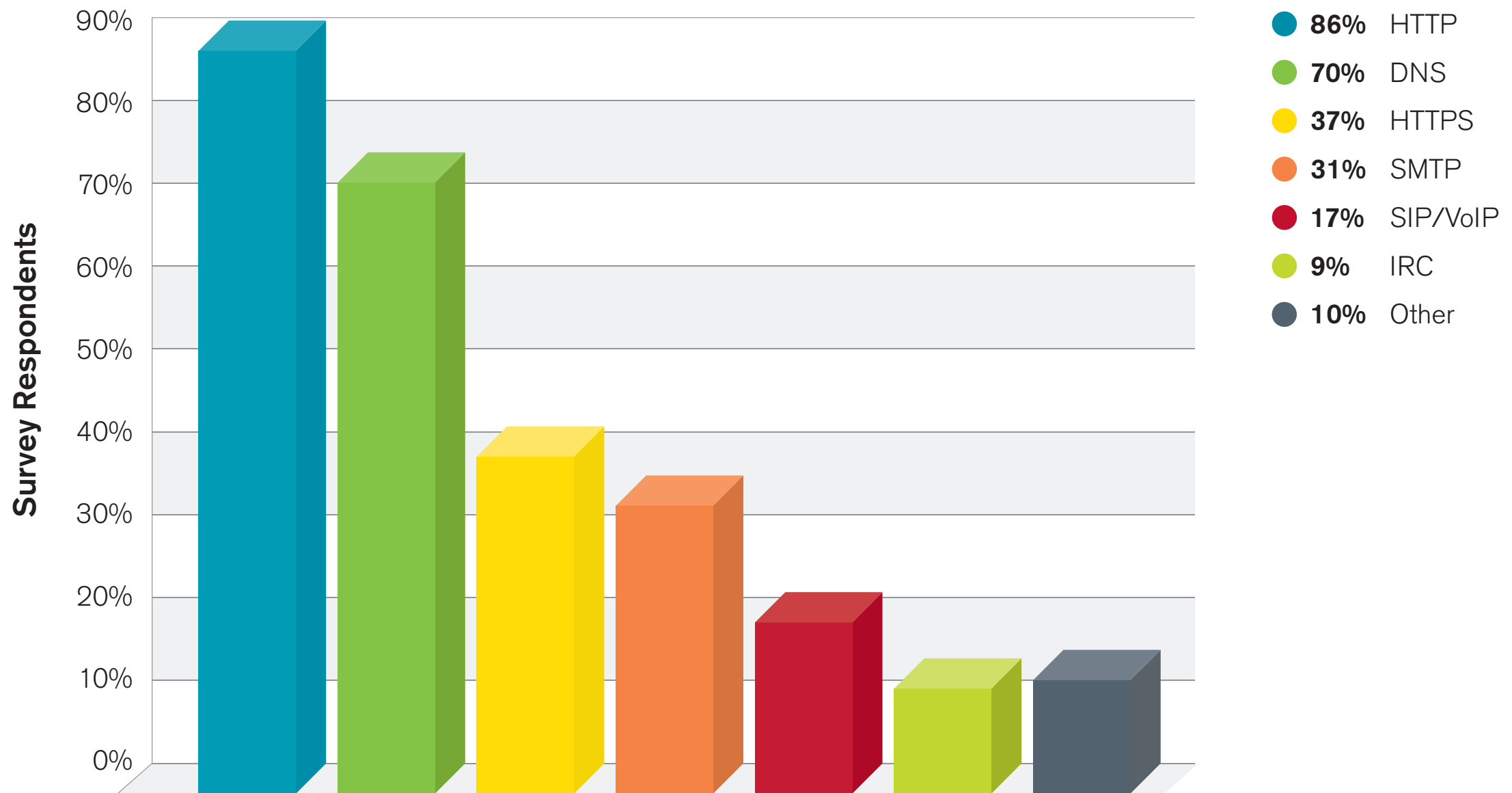


Figure 24 Source: Arbor Networks, Inc.

Motivation



Most Common Motivations Behind DDoS Attacks

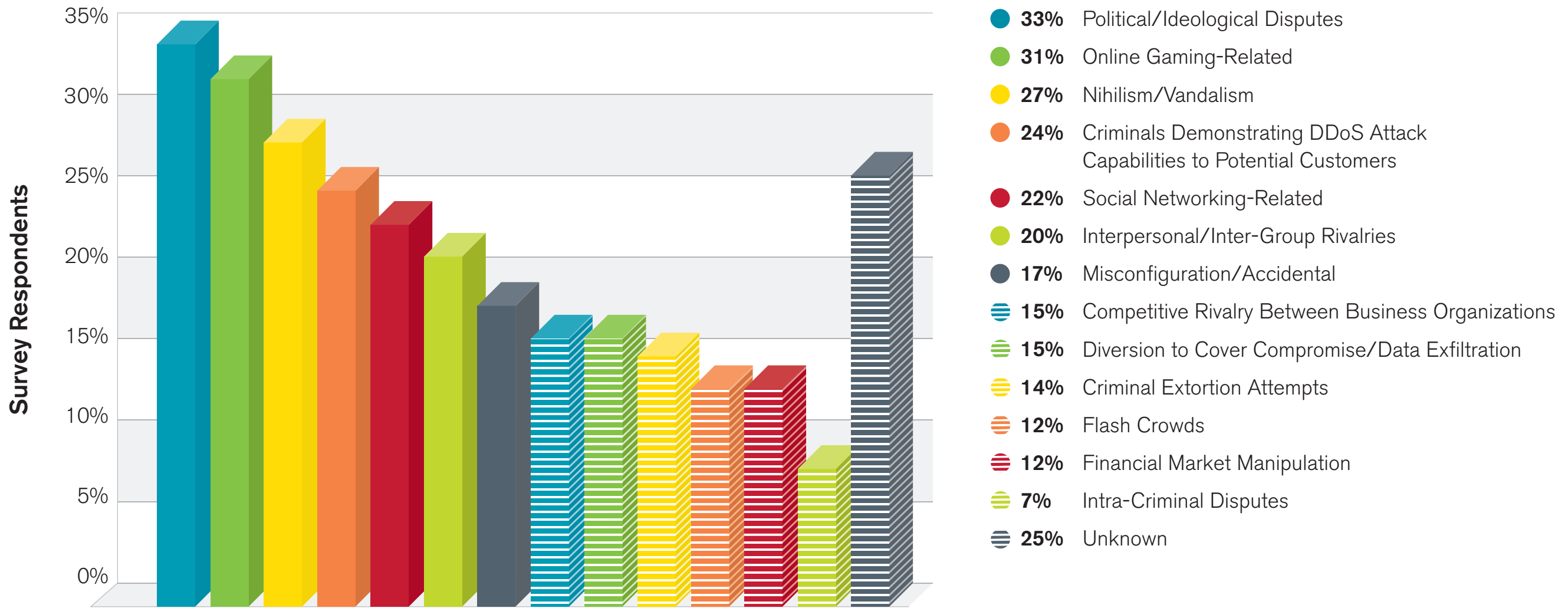


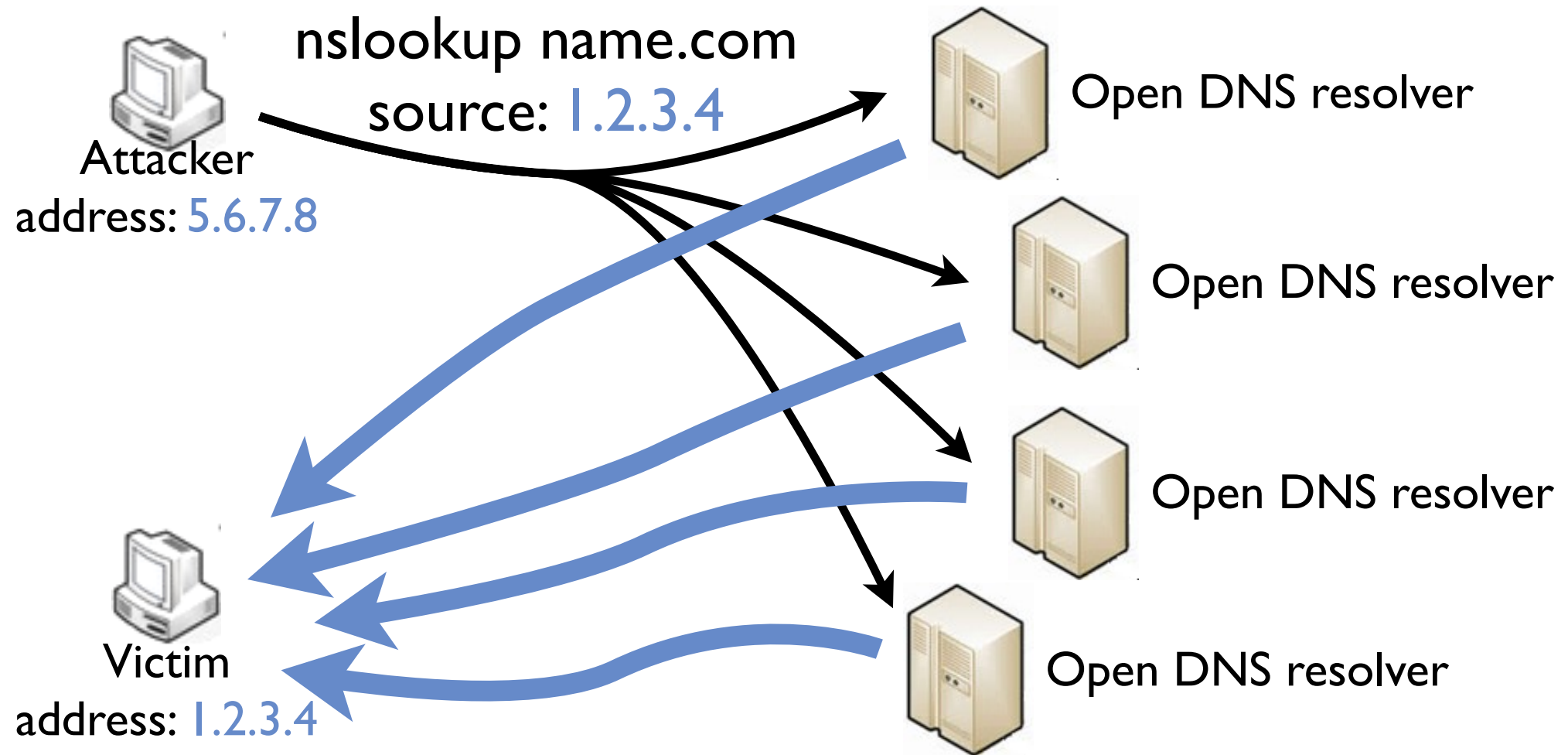
Figure 16 Source: Arbor Networks, Inc.

Comments on largest attacks



- “TCP/80 SYN flood toward Chinese online gaming (not gambling) site who was a DDoS mitigation customer of ours. Motivations unknown. Frequent **on-and-off waves** of attack traffic over several days, the largest of which topped out at **28.3 Mpps.**”
- “UDP port 22 small byte packets at high rate for less than 10 minutes, **overran firewalls** supposedly able to handle much higher pps rates.”
- “**UDP reflection/amplification attack**, primarily a mix of port 53 and 520 with some SYN and ICMP backscatter. Suspected attack motivation was retaliatory attack to something our users posted on a web forum (destination of the attack was a web proxy).”

DNS Reflection and Amplification



DNS Reflection and Amplification



Reflection

- via source address spoofing + lack of handshake in DNS
- hides source
- enables amplification by converting attacker's request traffic into response traffic

Amplification

- response is larger than request
- magnifies damage per unit of attacker work

DDoS defense



“destination-based remotely triggered blackholing”

Attack Mitigation Techniques

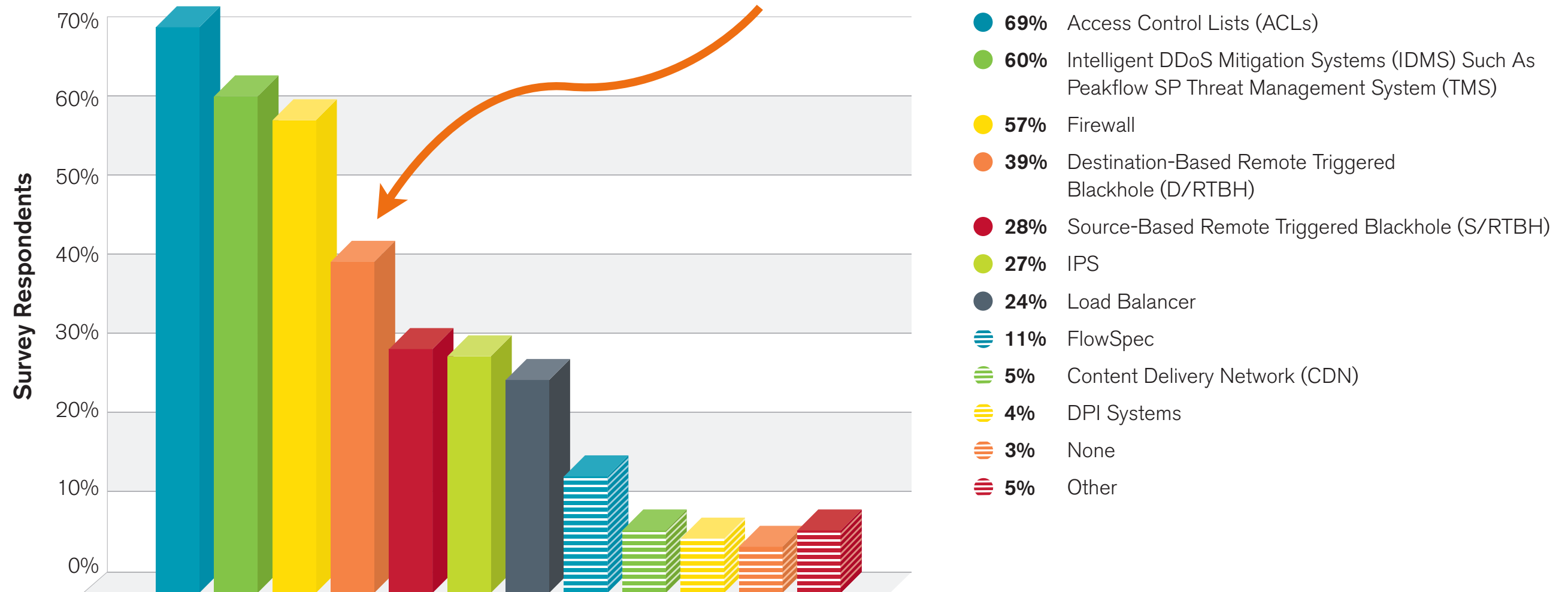


Figure 32 Source: Arbor Networks, Inc.

DoS in context



Most Significant Operational Threats Experienced

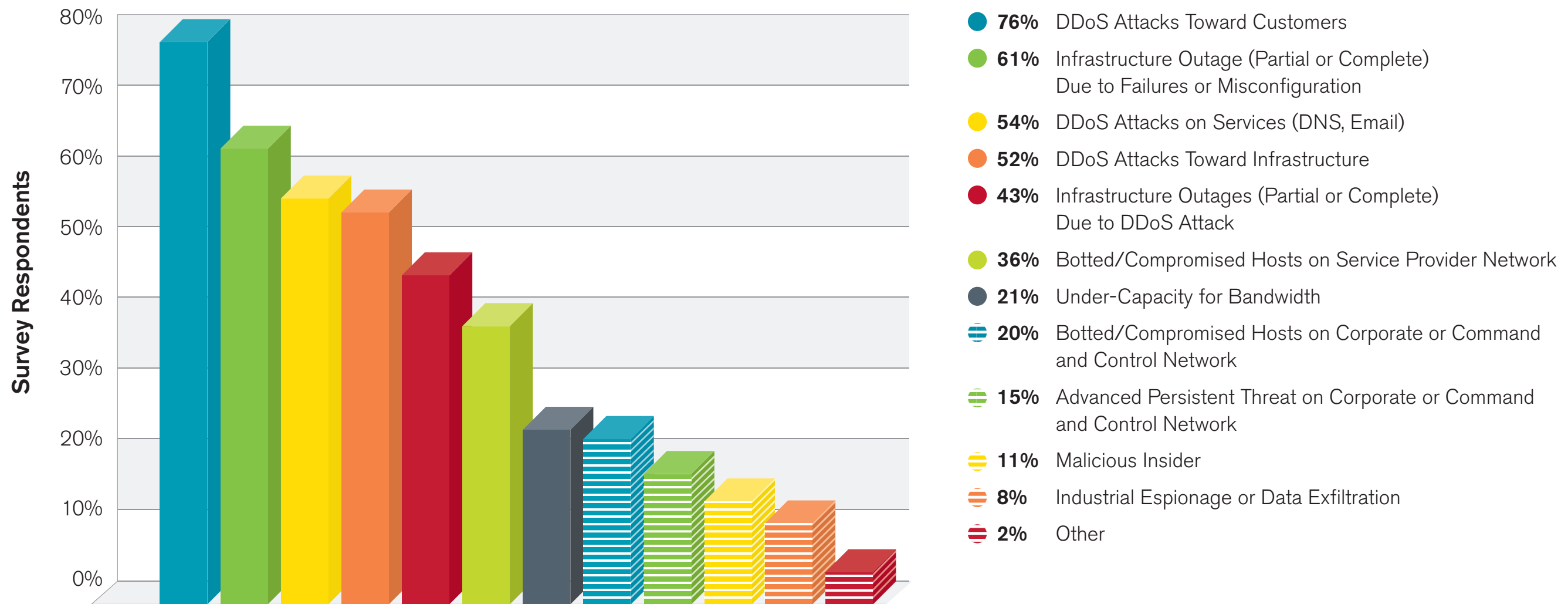
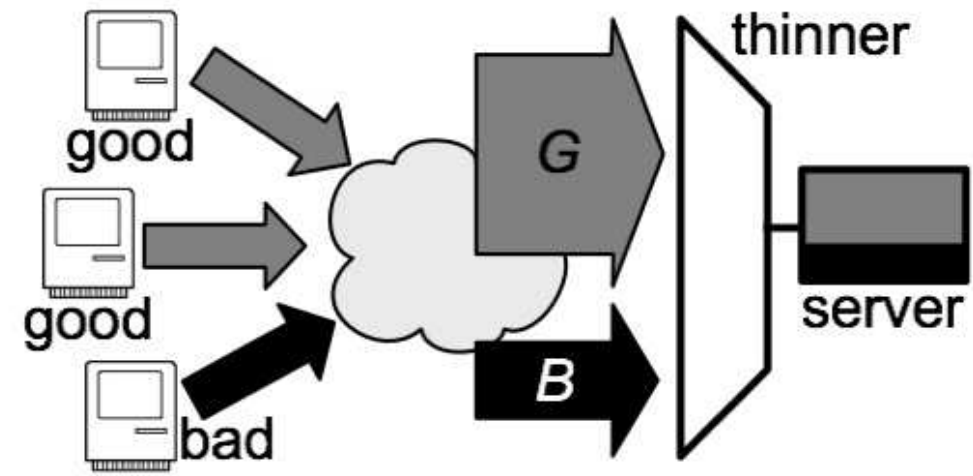
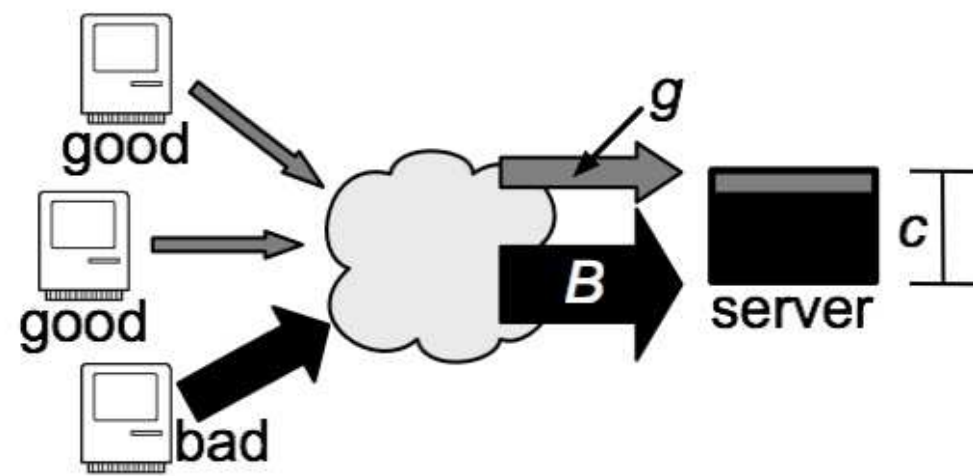


Figure 10 Source: Arbor Networks, Inc.

DDoS Defense by Offense

Walfish, Vutukuru, Balakrishnan, Karger, Shenker,
SIGCOMM 2006

Speak-up key idea



[Walfish et al.]

Mechanism

- Thinner guards access to server
- Runs “auction” for each service slot
- Whoever has sent most since last service gets service

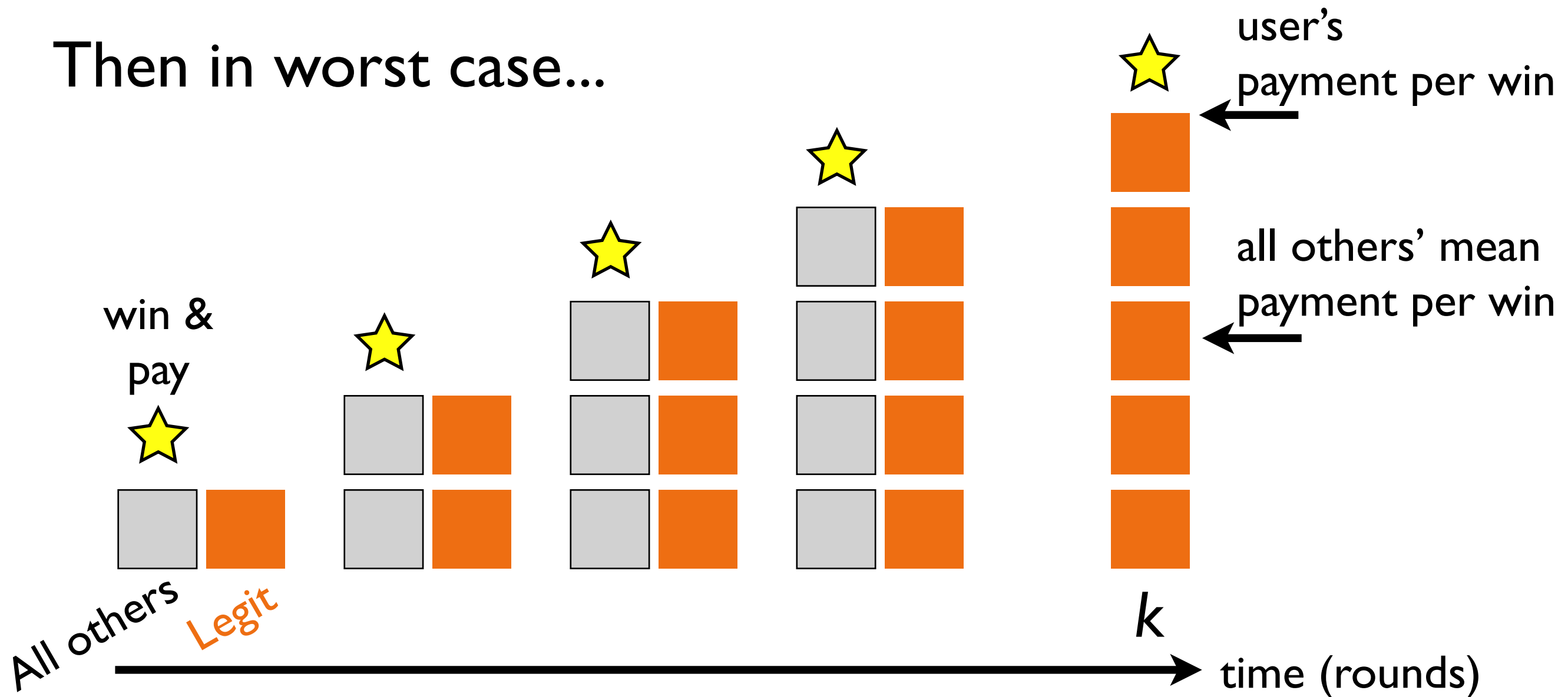
Bounding attacker damage



Assume

- Legitimate user sends one unit traffic per round
- Obtains one unit service after k rounds

Then in worst case...



Two ways to view the paper



If attackers are selfish, let good clients be selfish too

- At least they compete fairly
- Vastly improved situation

Charge clients a currency

- Reduces amplification due to HTTP request being much less work than response
- Here, currency = bandwidth
- This happens today under DoS attack, whether you like it or not
- Just need to inform legitimate clients about the situation

Clear drawback: Spends bandwidth



Extra charging

- Clients might be charged more [Yen Shine Low]
- Customers may thus switch to services that don't use Speak Up [Suna Kim]

Worse performance

- Could hurt other client apps' performance [Danny Xu]
- Could congest services co-located with server [Shayan Saeed, Alok Taigi]

Questioning Assumptions



Assumed: Good clients typically use small % of their bandwidth capacity

When is this false?

Questioning Assumptions



Assumed: Clients can all send about the same amount

What about users with low or costly bandwidth?

- Unfair to them [Anthony Lang, Nirupam Roy]
- **What would you do to fix it?**



What if multiple bots pretend to be one? [Zhenhuan]

- They get 2x the service
- But that's just a name change (2x1 vs. 1x2)
- “Ironically, taxing clients is easier than identifying them”



Week after next: midterm project presentations

- Be ready by Tuesday of that week
- 5 minute presentation, 5 minute questions
 - What **problem** are you solving?
 - Why has **past work** not addressed the problem?
 - What is your **approach** for solving it?
 - What are your **preliminary results** & progress?