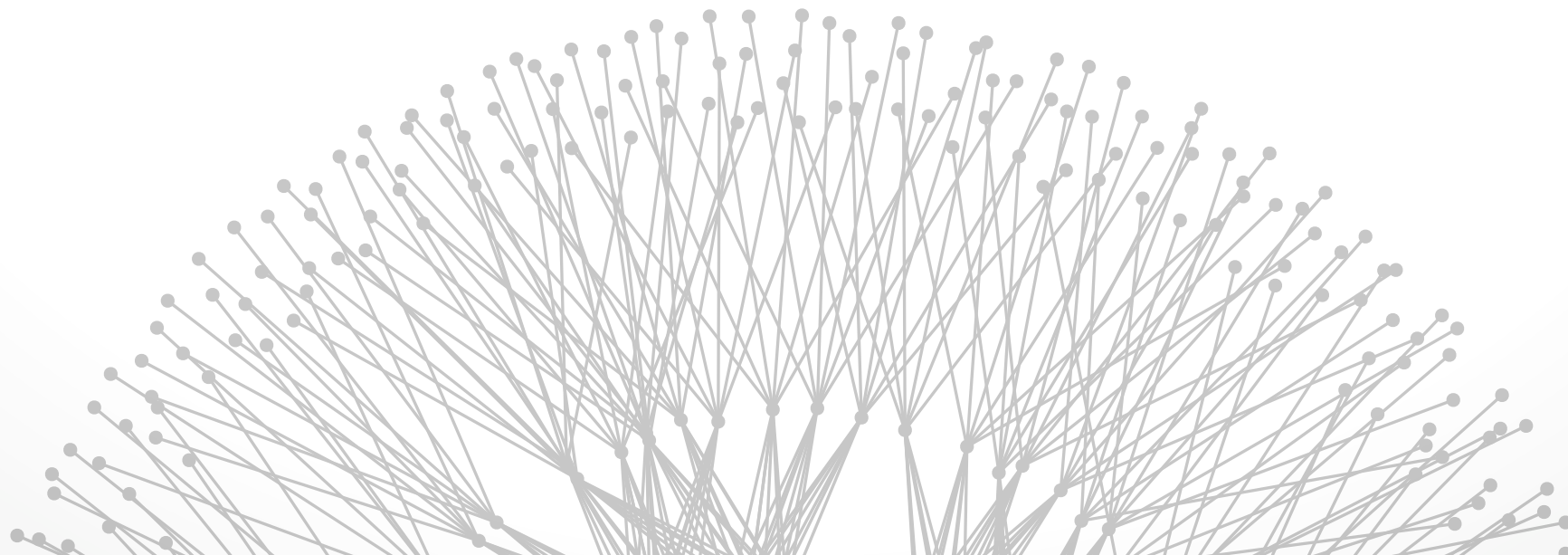


# Interdomain Routing and Connectivity

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CS 538 October 1 2013





Choosing paths along which messages will travel from source to destination.

# Problems for intradomain routing



Distributed path finding

Optimize link utilization (traffic engineering)

React to dynamics

High reliability even with failures

Scale

# Problems for interdomain routing



All of intradomain's problems

Bigger scale

Multiple parties

- No central control
- Conflicting interests
- Greater volume and diversity of attacks

Harder to change architecture

- Intradomain evolution: RIP, ISIS, OSPF, MPLS, OpenFlow, ...
- Interdomain: BGP.



## BGP: Border Gateway Protocol

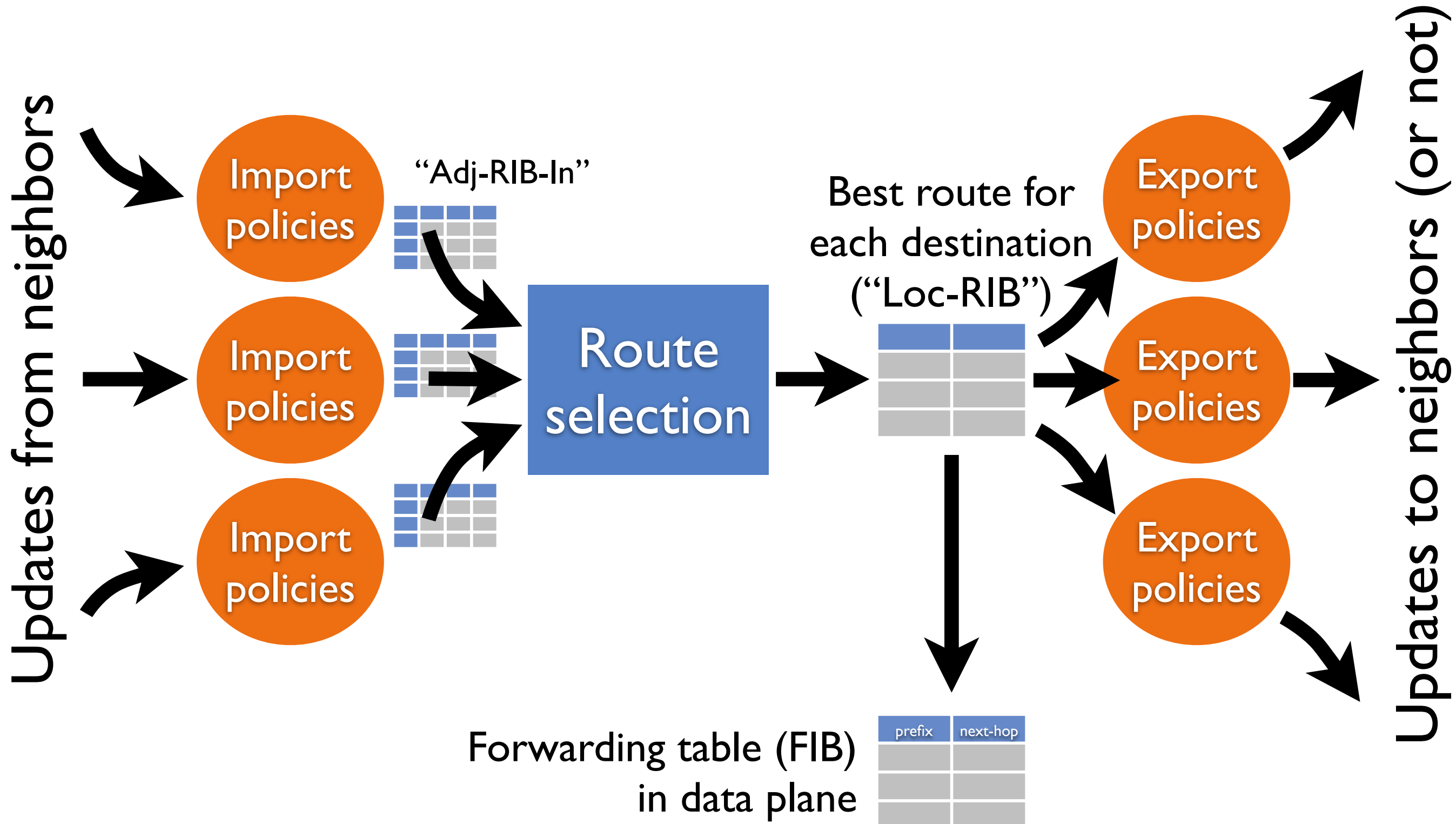
### Distance vector variant

- Send incremental changes, not whole vector
- Path vector: Remember path instead of distance

### Why path vector?

- Avoid DV's transient loops; but more importantly...
- **Policy support:** can pick any path offered by neighbors, not necessarily the shortest (Link State cannot)
- **Privacy support:** path choice policy is applied locally, not announced globally
  - Q: How much privacy is there?


# BGP: The picture at one router



# Route selection process



Import policies



Step	Attribute	Controlled by local or neighbor AS?
1.	Highest LocalPref	local
2.	Lowest AS path length	neighbor
3.	Lowest origin type	neither
4.	Lowest MED	neighbor
5.	eBGP-learned over iBGP-learned	neither
6.	Lowest IGP cost to border router	local
7.	Lowest router ID (to break ties)	neither

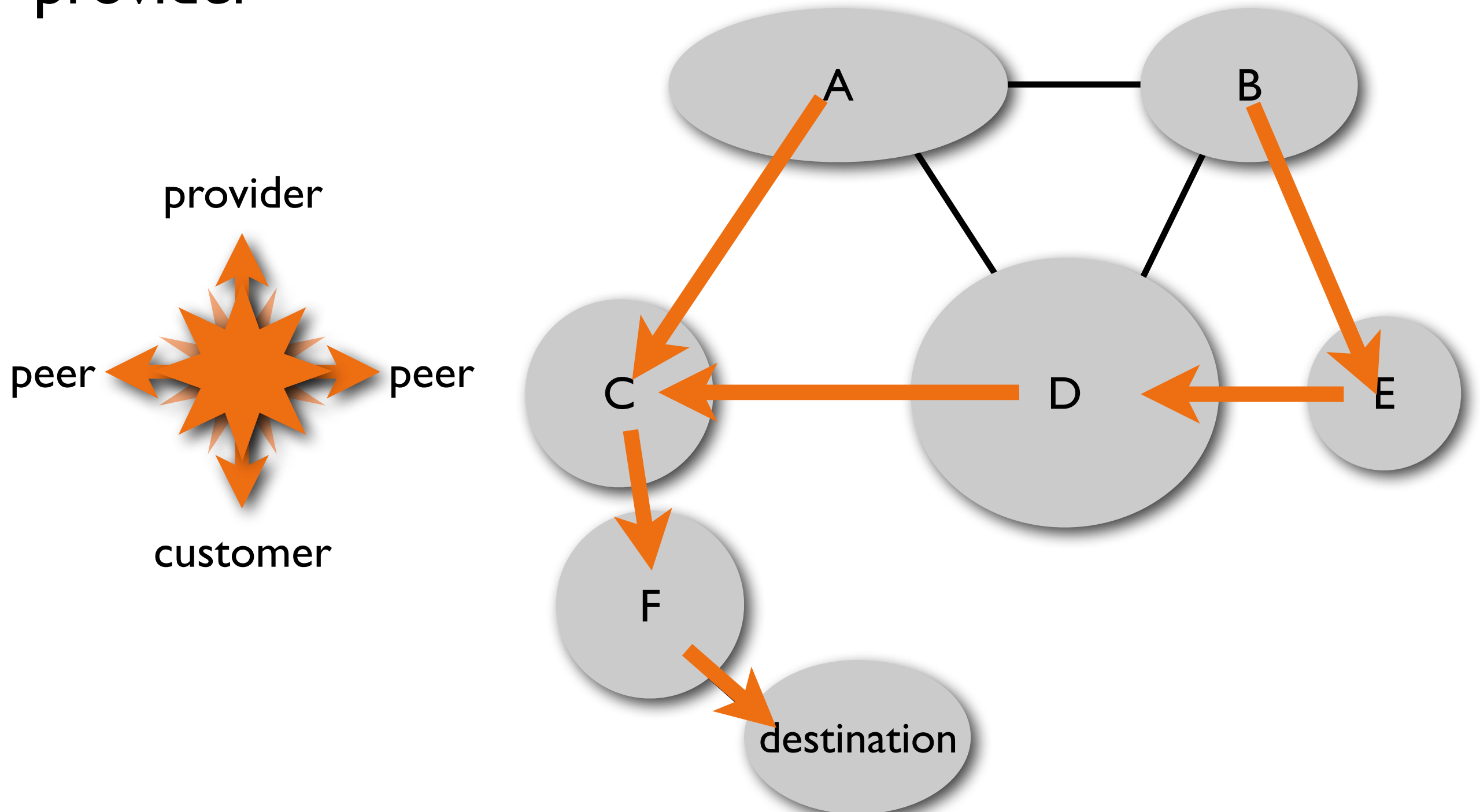
[Caesar, Rexford, IEEE Network Magazine, 2005]

This process is extended in many real implementations.

# Common policies



Route selection: prefer customer over peer over provider

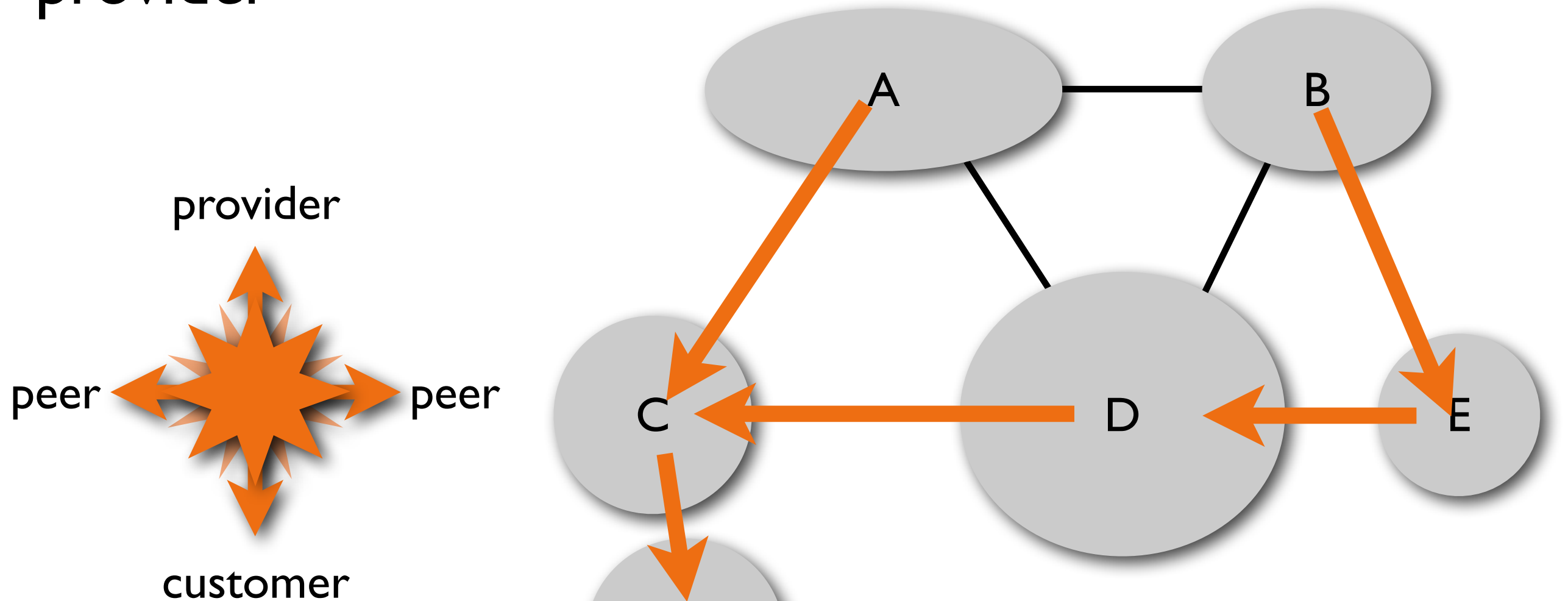




# Common policies



Route selection: prefer customer over peer over provider



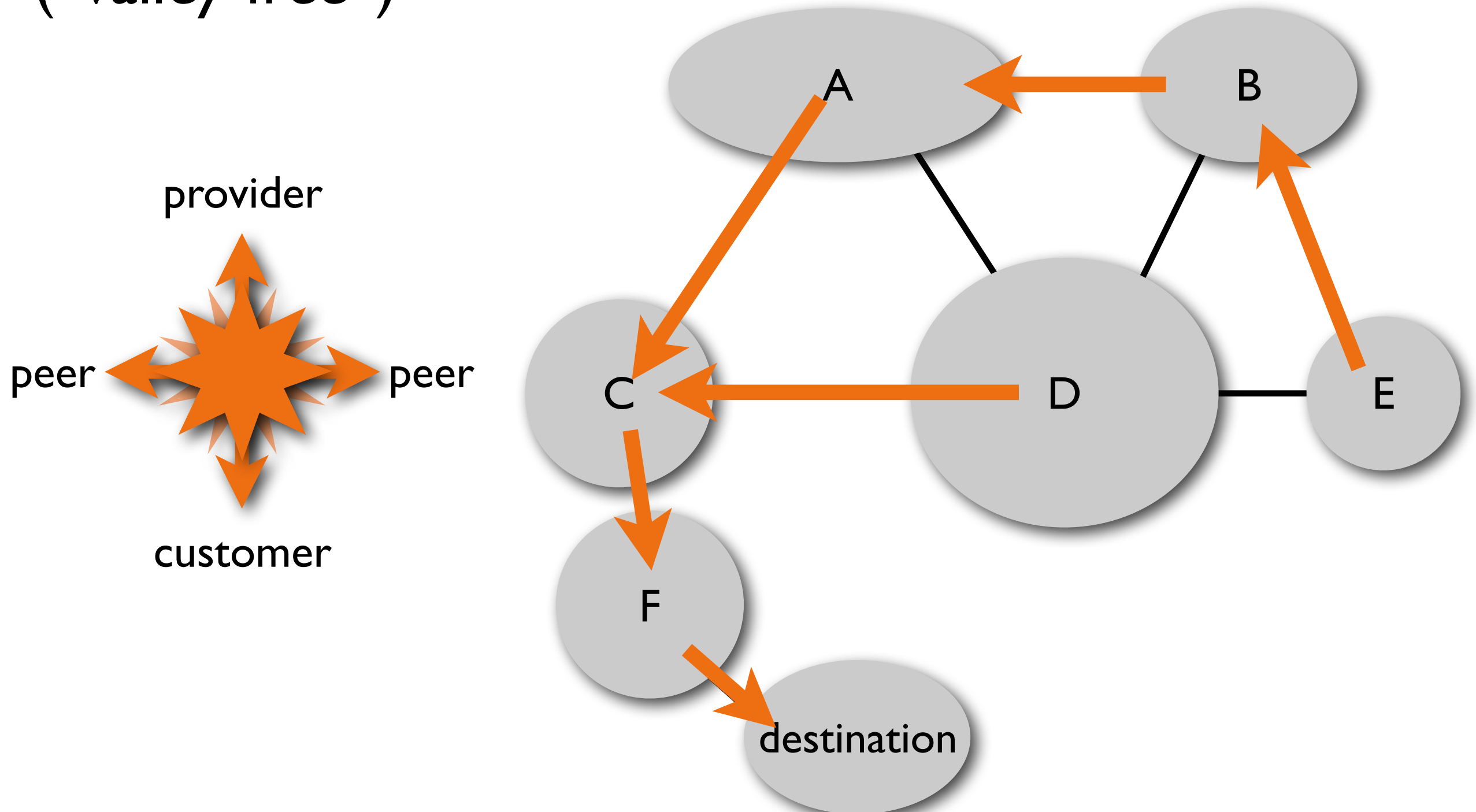
But ... What's wrong with this picture?

Falsely assumed all routes are exported

# Common policies



Route export (most common): to/from customer only  
("valley-free")



# Paper discussion



How does BGP inbound traffic engineering fit with TeXCP? Are they solving the same problem?

How can ISPs perform interdomain outbound TE?

*Peng Zhang*  
*netid : pzhang29*

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*CS 538 Fall 2013/9/16*  
*Assignment 1*

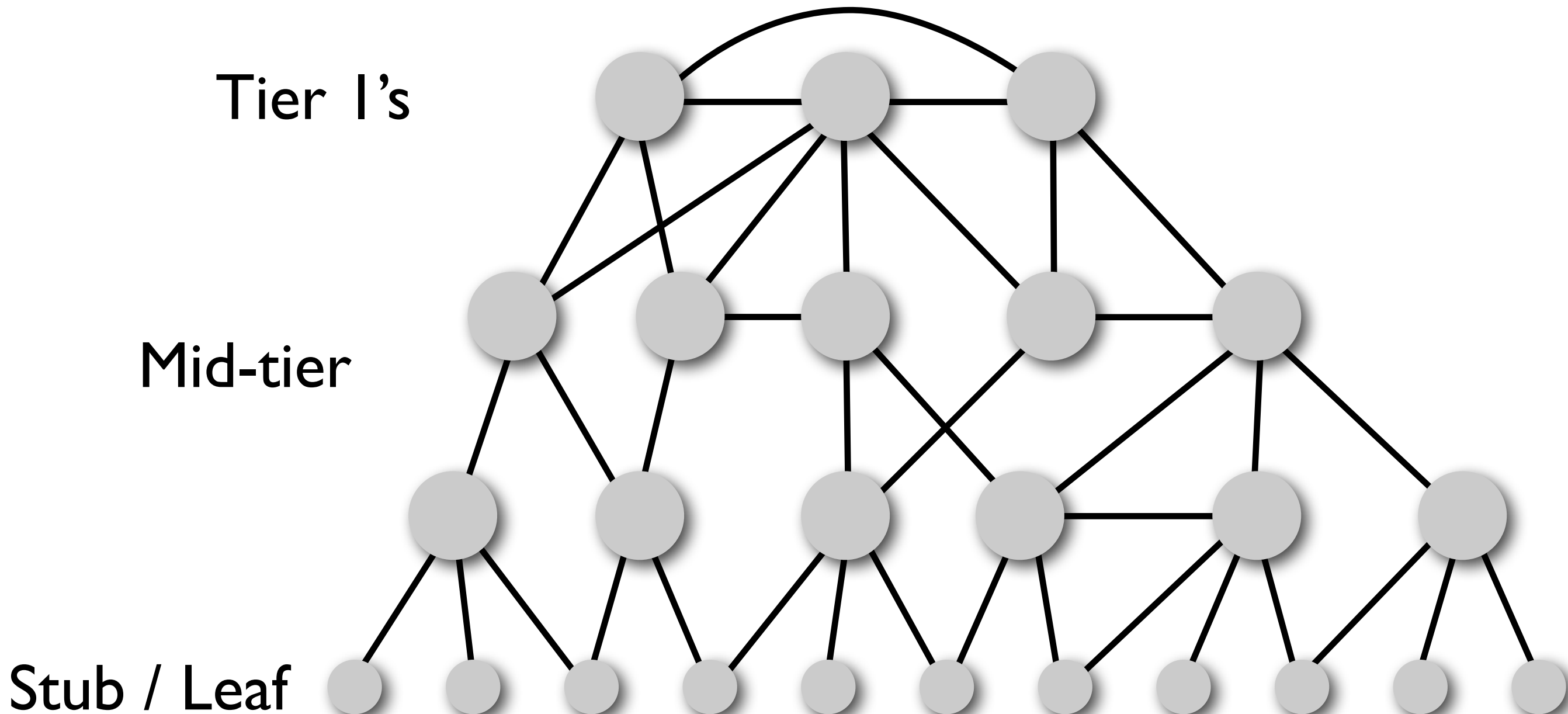
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2#1. The sequence of ISPs (AS numbers and/or business names) from the last step.

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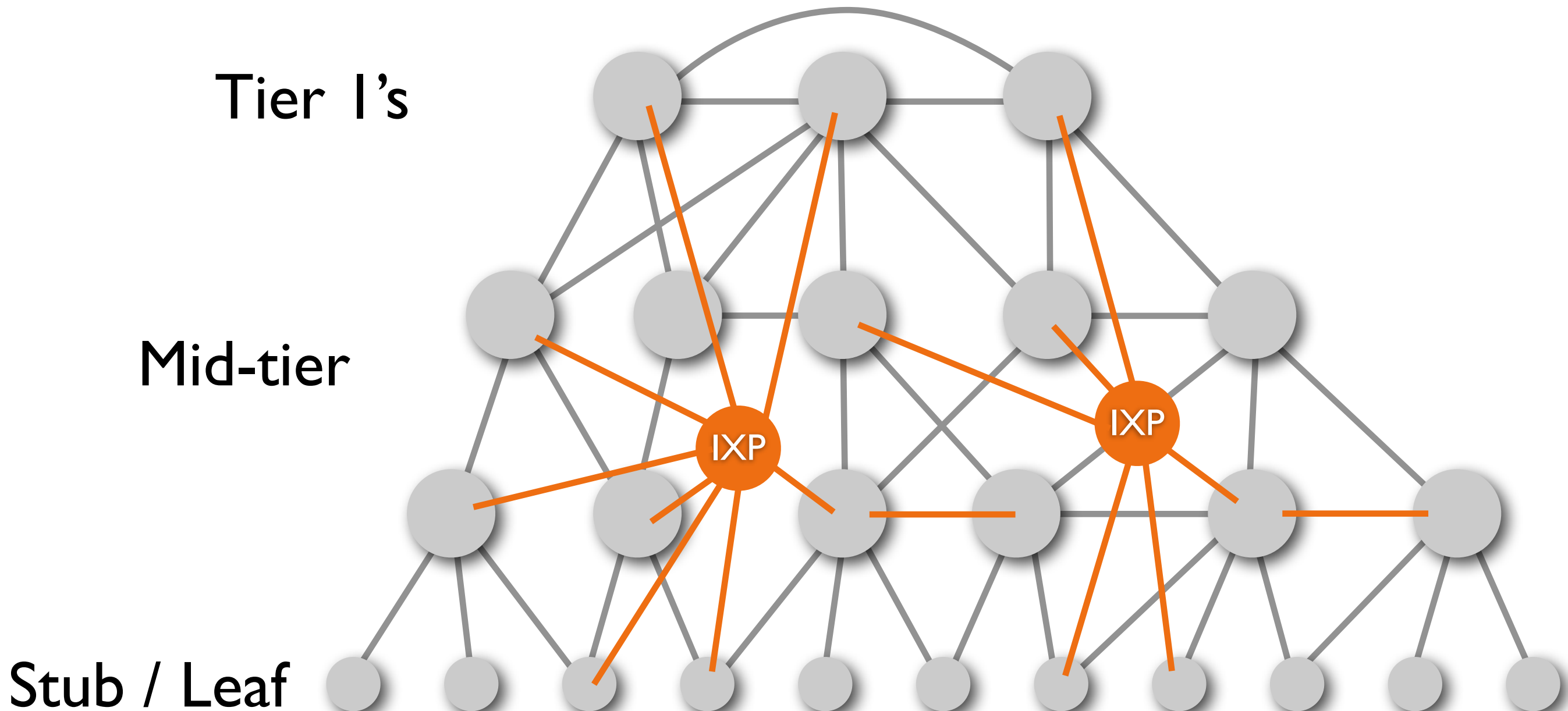
13030 11537 40387 38  
31500 174 40387 38  
8928 7132 40387 40387 40387 40387 38  
1299 174 40387 38  
5413 1299 174 40387 38  
6067 174 40387 38  
8426 3549 11537 40387 38  
19151 11537 40387 38  
6939 11537 40387 38

# Interconnection: Traditional view



Hierarchical, limited peering at lower tiers

# Interconnection: Modern view



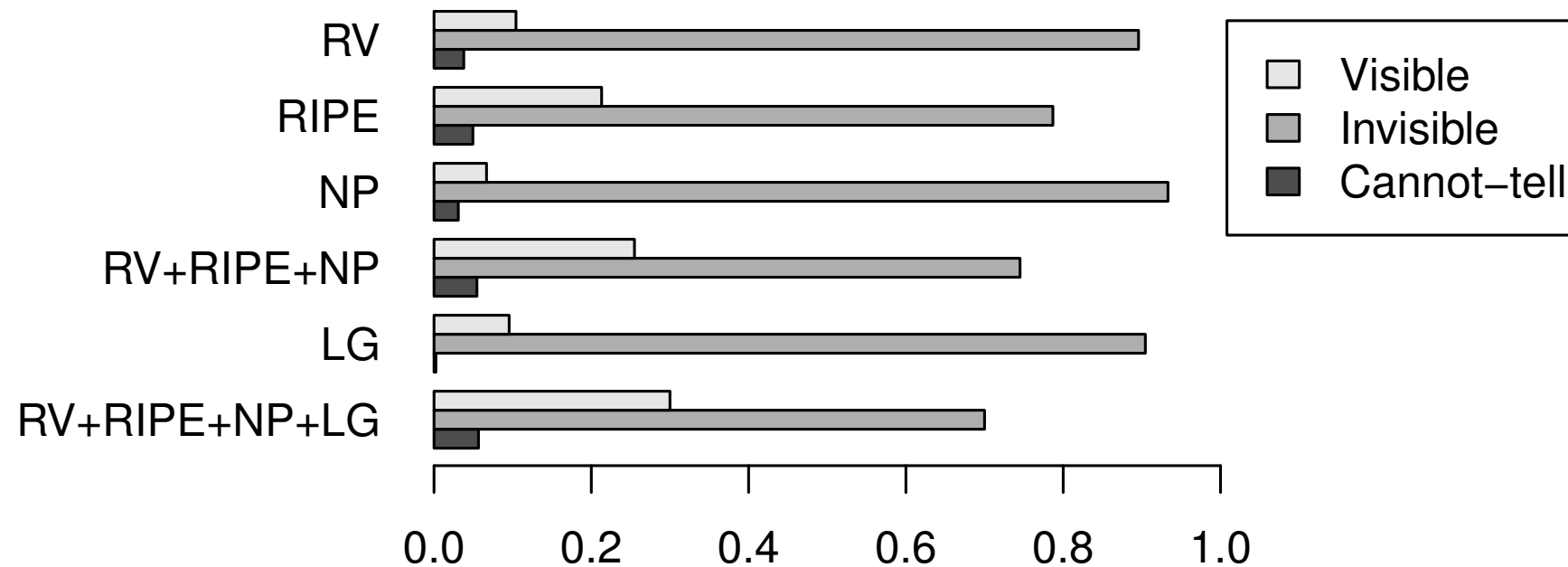
Significant and increasing peering at lower tiers



## Significant peering

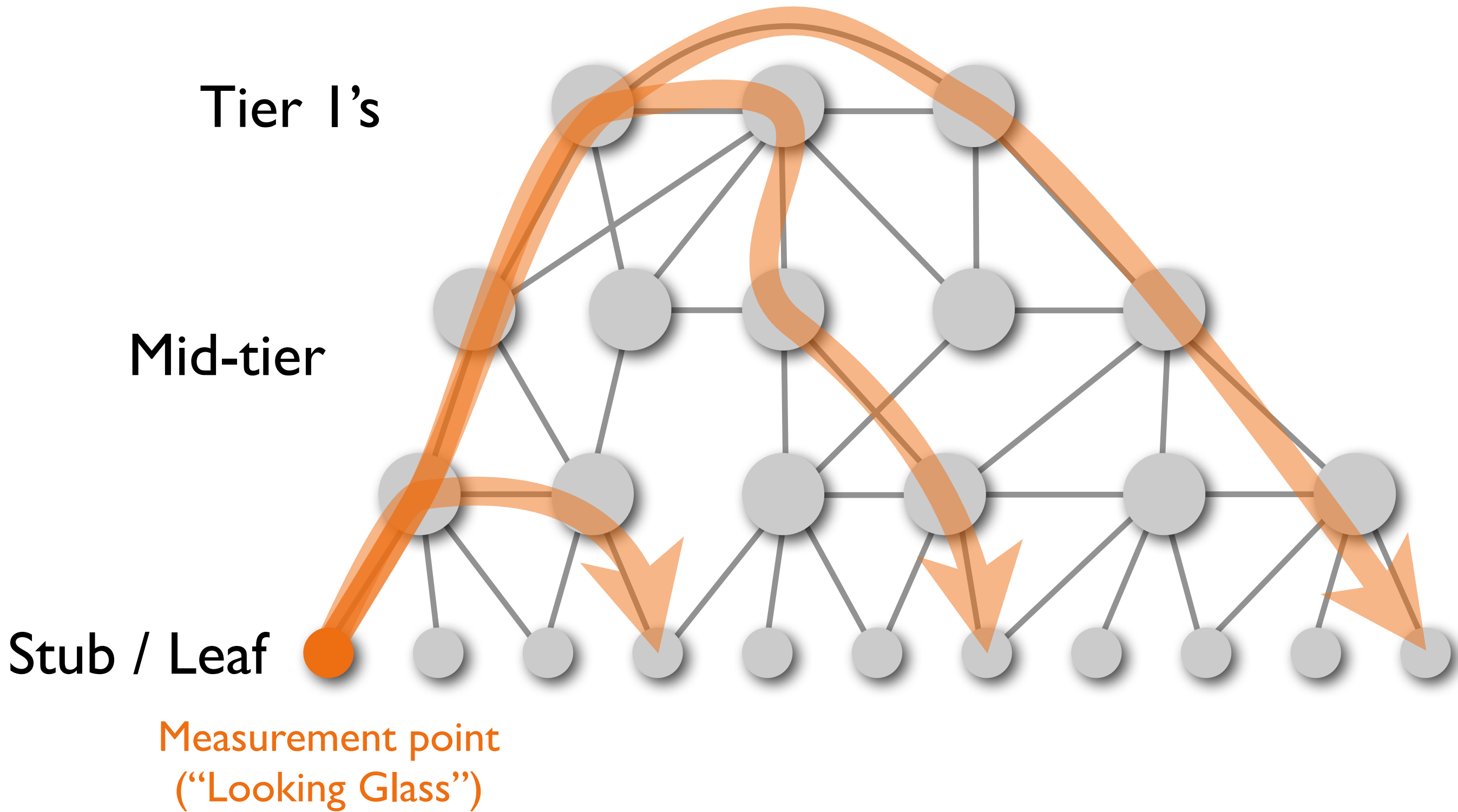
- Estimated 200,000 peerings just in Europe
- More than 2x as many as non-peering links!

## These peerings missed in past measurements



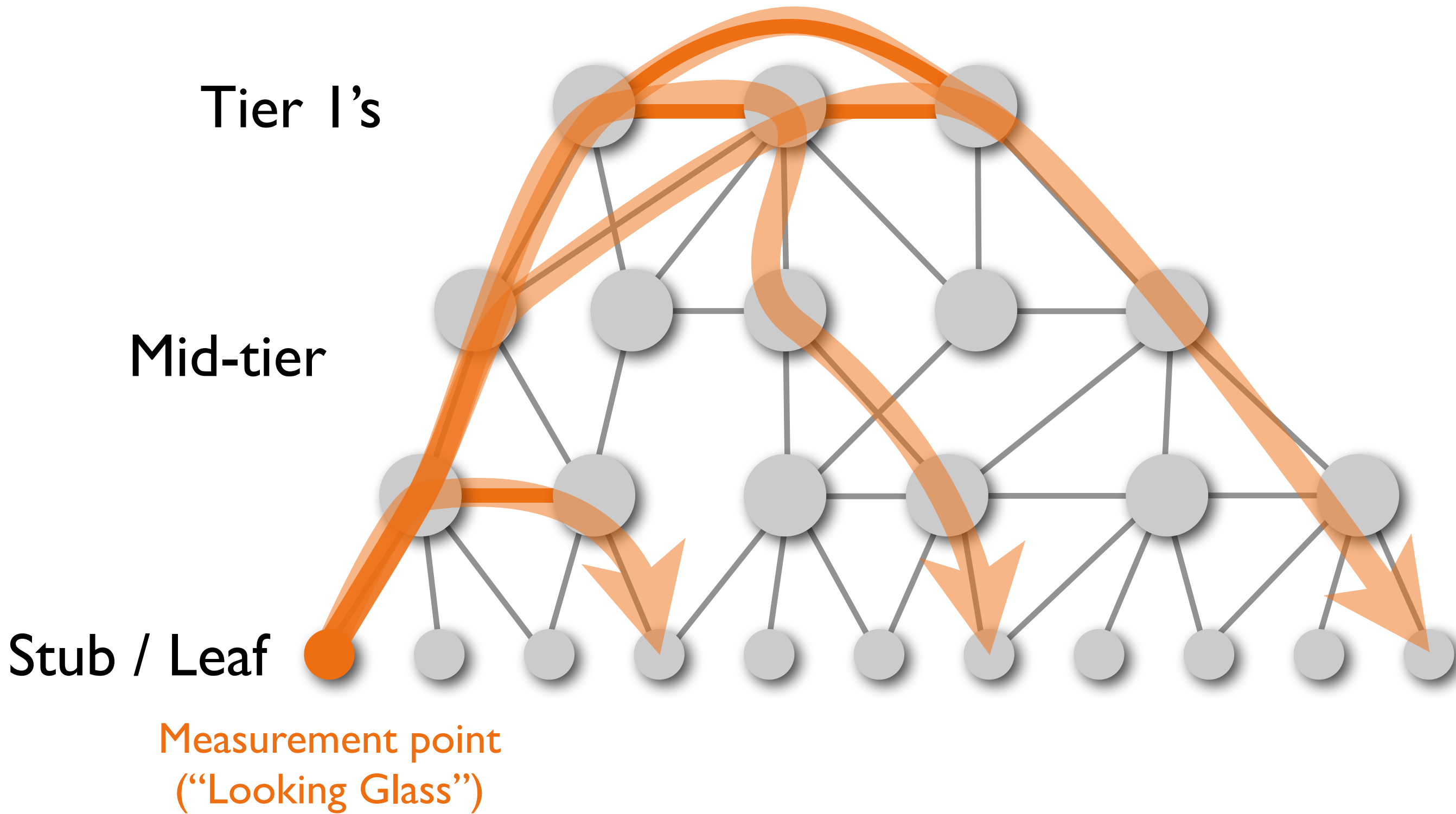
**Figure 2: Peering links and visibility in control/data plane (normalized by number of detected P-P links).**

# Why measurements miss so much



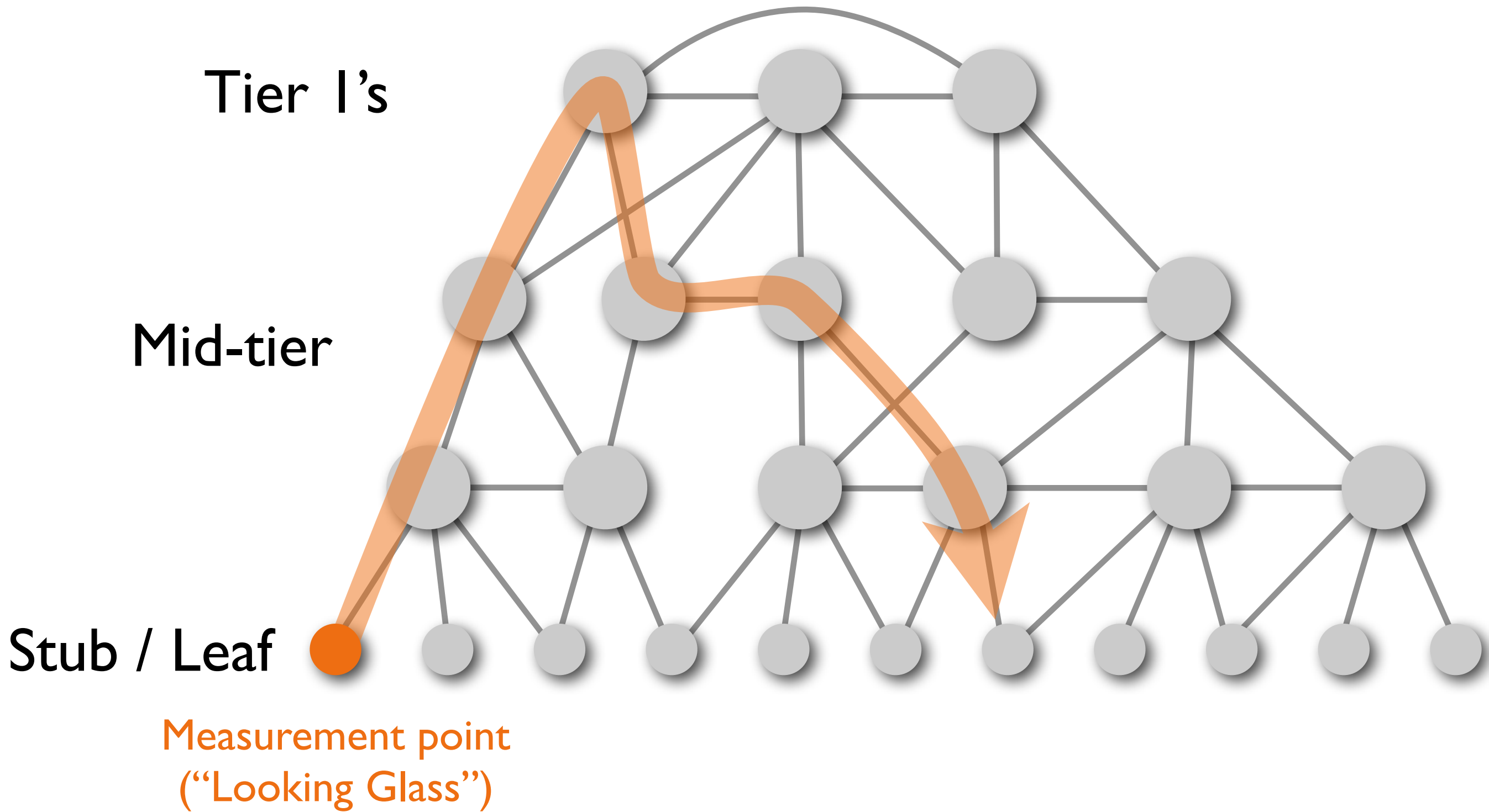


# Why measurements miss so much

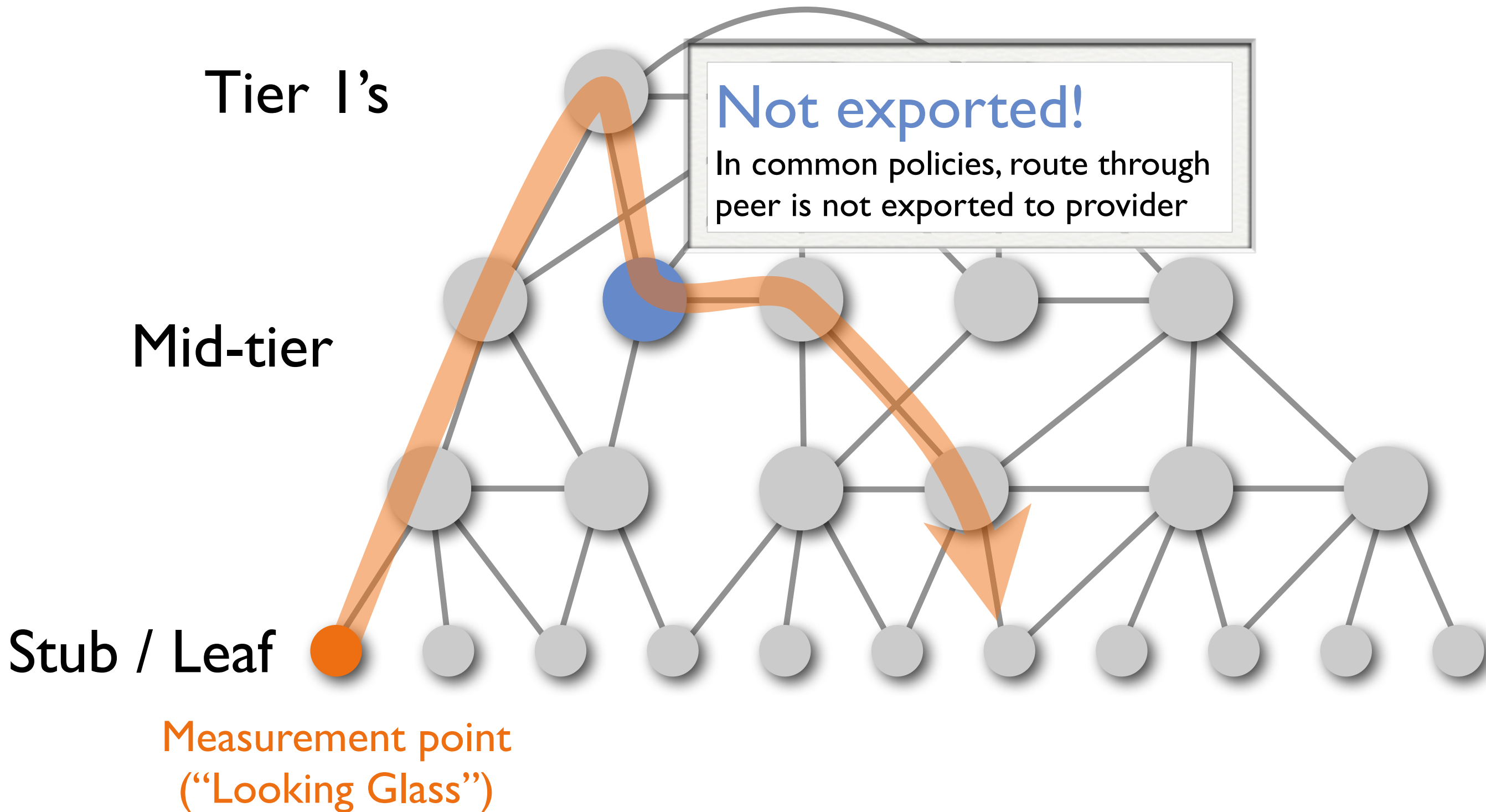




# Why measurements miss so much



# Why measurements miss so much



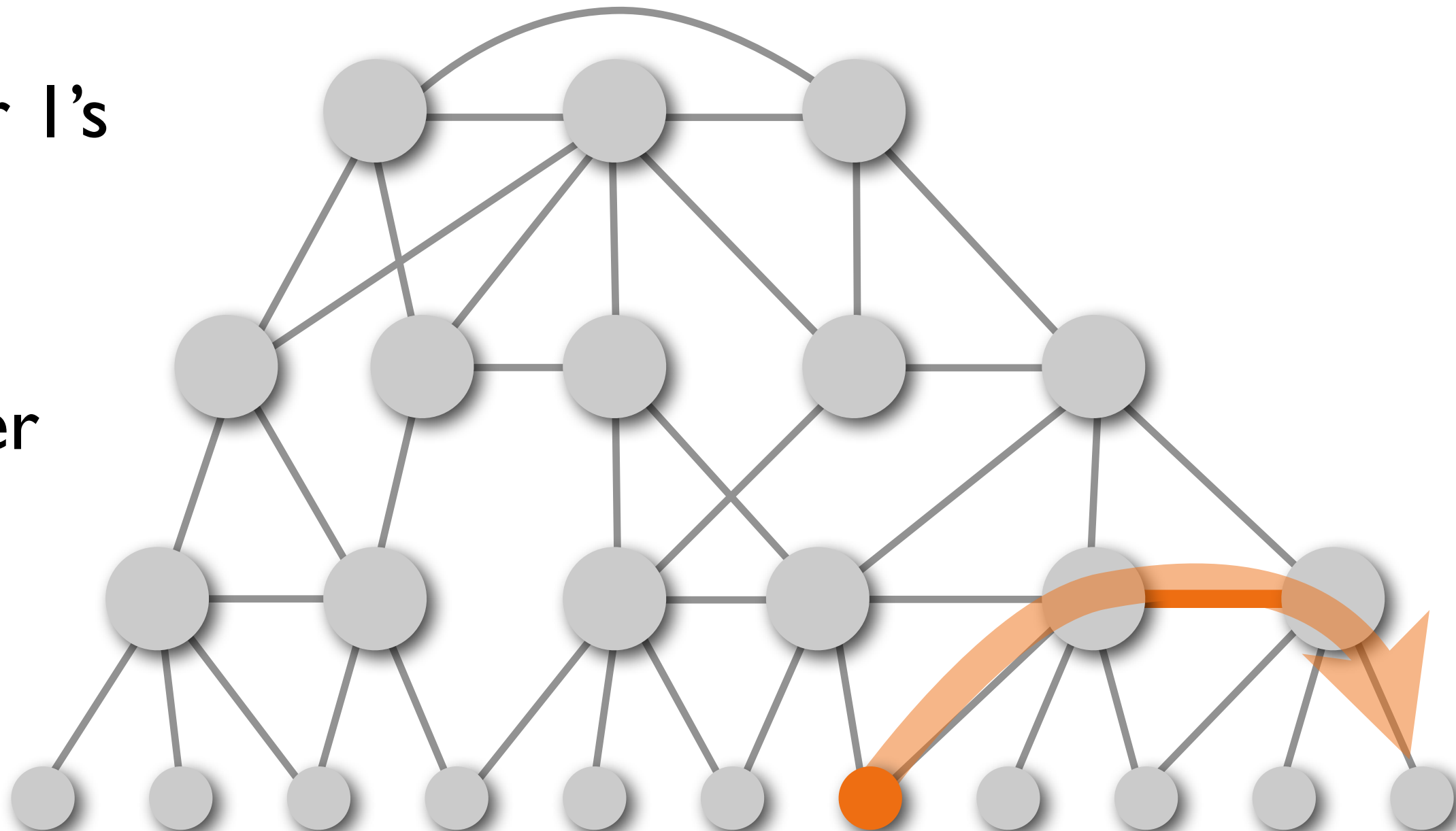
# Why measurements miss so much



Tier 1's

Mid-tier

Stub / Leaf



Measurement point  
("Looking Glass")

To see peer-peer link, *both*  
probe source & dest. must  
be in localized area

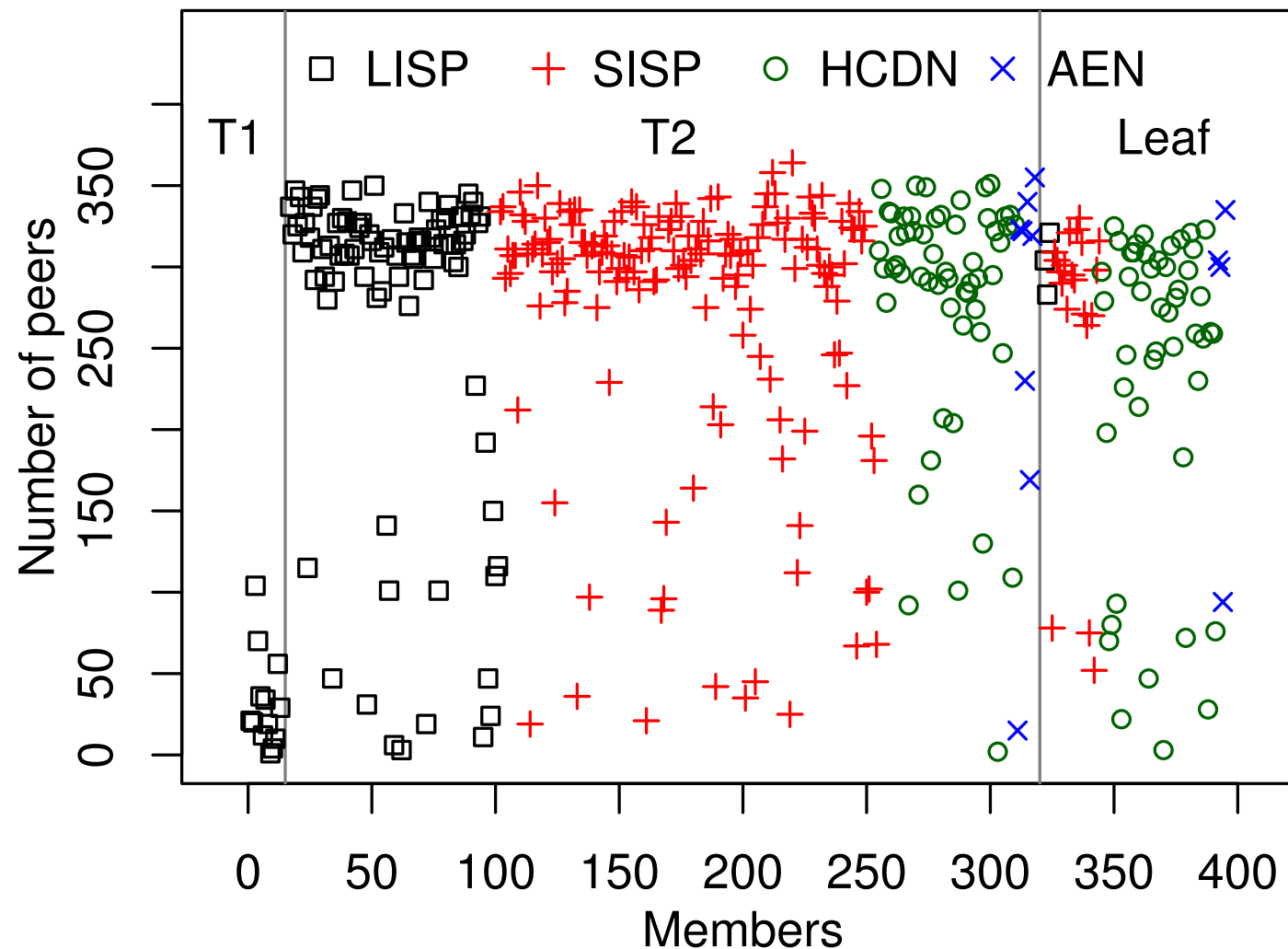
# Paper discussion



What's the purpose of an IXP?

- “Metcalf’s law”: value of net is  $O(n^2)$  when  $n$  participants

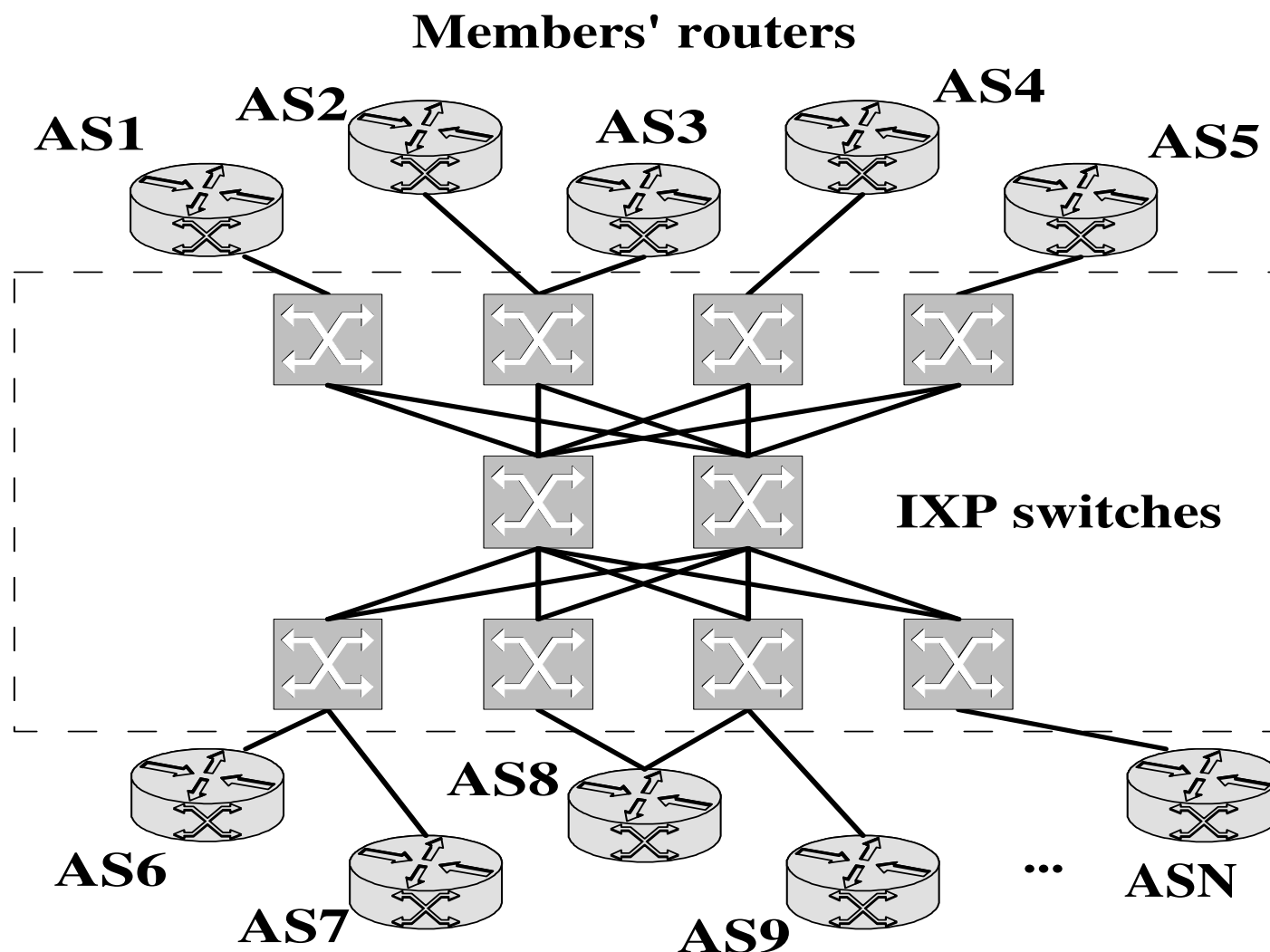
Why don't top-tier ISPs peer much at the IXP?



# Paper discussion



How might **router-level interconnection** differ from **AS-level peering**? Would this paper's conclusions be the same for router-level?



# physical links = 27

# potential peerings = 45



Similarly ... suppose we treat the IXP as an AS “in the middle” of each member AS-to-AS connection

Now how many links are there?

- 396 total members of this IXP, so 396 links
- vs. 50,000 reported in the paper!
- $O(n^2)$  inflation factor for  $n$  member ASes

This suggests interesting measurement projects:

- If you care about only the router level, what fraction of the links are observable?
- If you treat the IXP as an AS “in the middle”, what fraction of the links are observable?

# What's to come



## Project proposals

- Comments back by Thursday

## Next: Part Two of the course: Grand Challenges

- scalability
- complexity: SDN
- reliability
- selfishness
- security & privacy