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TOWARDS ENERGY-EFFICIENT DATABASE CLUSTER DESIGN

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MAIN IDEA

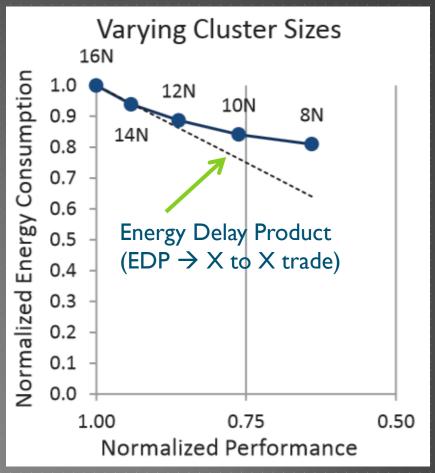
Study the trade-offs between performance and energy consumption in a cluster to identify bottlenecks and propose a model that considers these bottlenecks and other query parameters to predict the performance and energy consumption of the cluster. With the findings provide cluster design principles.

INTRODUCTION

- Energy growing cost of operational cost
- ► CHALLENGES to increasing energy efficiency
 - Inherent scaling inefficiency
 - Choosing energy-efficient hardware
- Architectural design space of energy efficient database clusters

VARYING CLUSTER SIZES

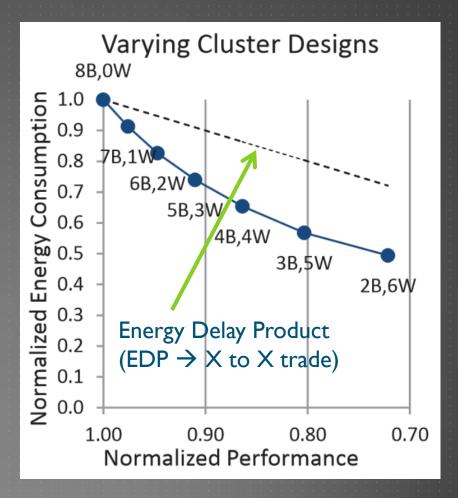
- Trade performance for reduced energy consumption
- Results
 - ► Sub-linear speedup
 - Performance
- Data points above EDP curve
- ▶ Eg: 10 N
 - ▶ 24% penalty in performance
 - ► 16% decrease in energy



Vertica – running TPC-H Q12

VARYING CLUSTER DESIGNS

- Performing parallel hash joinP-store
- Heterogeneous cluster design
- Wimpy scan and filter data and send to Beefy
- Results
 - Data points below EDP curve
 - Greater energy savings for less performance penalty

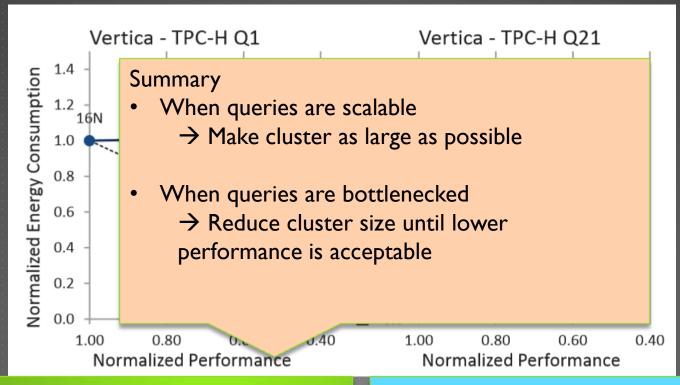


8 node cluster of Beefy (B) and Wimpy (W)

CONTRIBUTIONS

- Explore trade-offs in performance versus energy efficiency
 - ▶ Vertica, P-store, HadoopDB
- Identify bottlenecks for performance and energy efficiency
- Build a model which predicts performance and energy efficiency for various cluster configurations
- Illustrate interesting cluster design points
- Provide guiding principles for energy-efficient data processing
- Seed future research in this area

PERFORMANCE VERSUS ENERGY EFFICIENCY



- Simple aggregation, no joins
- Performance scales linearly
- Constant energy consumption
- Hence, add as many nodes as possible

- 94.5% of query on local machines
- Performance scales linearly
- Constant energy consumption
- Hence, add as many nodes as possible
- Refer (I)

BOTTLENECKS

- Hardware (network and disk)
 - ▶ Repartitioning → internode communication
 - ▶ Node waits for data from network
- Algorithmic (broadcast)
 - ► Broadcast takes same time regardless of number of nodes
 - ► Eg: 16 N \rightarrow each node receives 15m/16, 32 N \rightarrow each node receives 31m/32
- Data skew
 - Part of future work

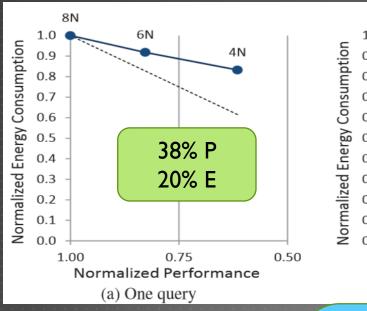
BUILDING A MODEL

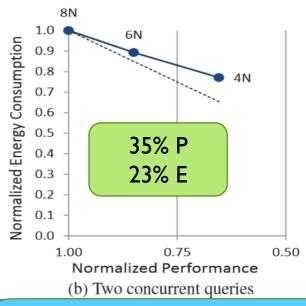
- ▶ P-store
 - \triangleright Custom built parallel engine \rightarrow scan, project, select, hash join, network exchange
- Explore performance bottlenecks affecting energy efficiency
 - Hash join query
 - Cluster V configuration

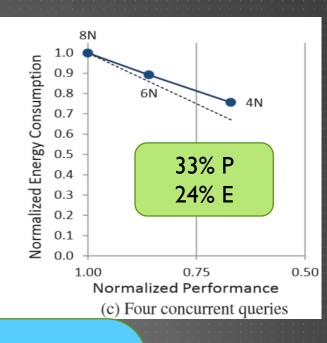
DBMS	Vertica	RAM	48GB
# nodes	16	Disks	8x300GB
	10		
TPC-H size	1TB (scale 1000)	Network	1Gb/s
CPU	Intel X5550 2 sockets	SysPower	$130.03C^{0.2369}$
			C = CPU utilization

TPC-H Q3 hash join

EXPERIMENT #1



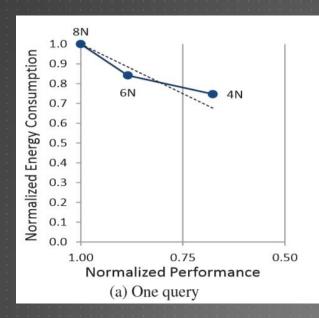


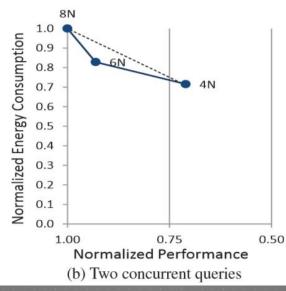


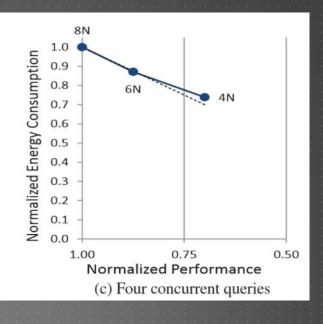
- Poor performance scalability
- Energy savings increases as concurrency increases (data points closer to EDP)
- Reason → CPU utilization does not scale because of network bottleneck

EXPERIMENT #2

TPC-H Q3 broadcast join

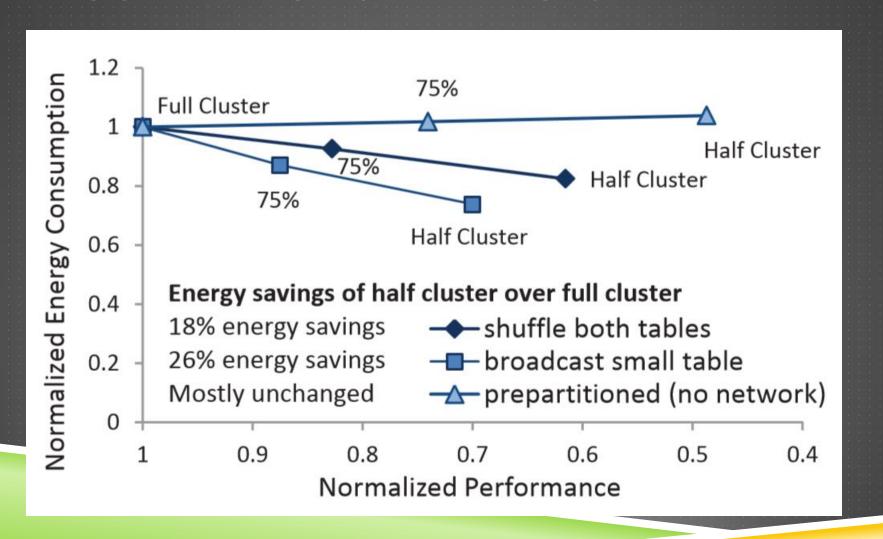






- Energy savings increases as concurrency increases (data points closer to EDP)
- Suffers non-linear scalability
- Reason → Broadcast does not scale

SUMMARY OF NETWORK AND ALGORITHMIC BOTTLENECKS



ENERGY EFFICIENCY OF INDIVIDUAL NODES

System	CPU (cores/threads)	RAM	Idle Power (W)
Workstation A	i7 920 (4/8)	12GB	93W
Workstation B	Xeon (4/4)	24GB	69W
Desktop	Atom (2/4)	4GB	28W
Laptop A	Core 2 Duo (2/2)	4GB	12W (screen off)
Laptop B	i7 620m (2/4)	8GB	11W (screen off)

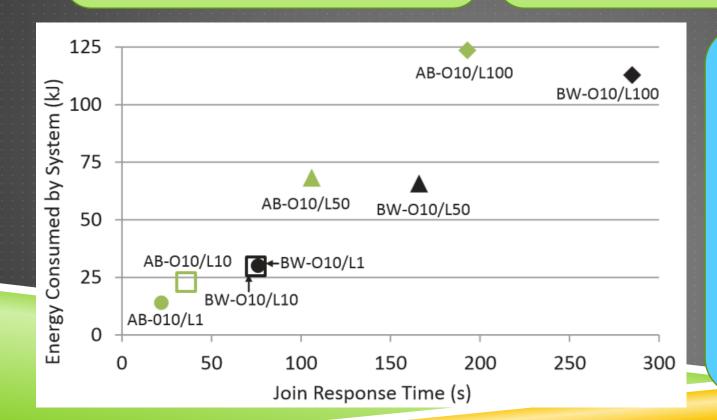
- In-memory workload
- Hash join



CLUSTER DESIGN POINTS

- Beefy →4 HP ProLiant servers with quad core Nehalem Xeon processors
- Each node has 32 GB of memory
- Average power I54 W

- 2 Beefy/2 Wimpy → 2 Beefy nodes
 + 2 Laptop Bs with i7 processor
- Wimpy node has 8 GB of memory
- Average laptop power 37W



- 1%, 10%, 50%, 100% selectivity on LINEITEM
- I 0% selectivity on ORDERS
- Beefy nodes build hash tables, Wimpy scan and filter
- Heterogeneous Execution

MODELING P-STORE AND BOTTLENECKS

- Understand the nature of query parameters and scalability bottleneck
- Predict the performance and energy consumption of various ways to execute a hash join
- Parameters

- Hash join
- I. Build phase
- 2. Probe phase

```
T_{prb} Probe phase time (s)
T_{bld} Build phase time (s)
                                     E_{prb} Probe phase energy (J)
E_{bld} Build phase energy (J)
 N_B # Beefy nodes
                                      N_W # Wimpy nodes
M_B Beefy memory size (MB)
                                     M_W | Wimpy memory size (MB)
    I|Disk bandwidth (MB/s)
                                         L Network bandwidth (MB/s)
 Bld|Hash join build table size (MB)|| Prb|Hash join probe table size (MB)
S_{bld} Build table predicate selectivity ||S_{prb}|| Probe table predicate selectivity
          R_{Wbld} Rate at which a Wimpy node builds its hash table (MB/s)
           R_{Bbld} Rate at which a Beefy node builds its hash table (MB/s)
          U_{Wbld} Wimpy node CPU bandwidth during the build phase
           U_{Bbld} Beefy node CPU bandwidth during the build phase
          R_{Wprb} Rate at which the Wimpy node probes its hash table (MB/s)
          R_{Bprb} Rate at which the Beefy node probes its hash table (MB/s)
          U_{Wprb} | Wimpy node CPU bandwidth during the probe phase
          U_{Bprb} Beefy node CPU bandwidth during the probe phase
     C_B = 5037 | Maximum CPU bandwidth of a Beefy node (MB/s)
    C_W = 1129 Maximum CPU bandwidth of a Wimpy node (MB/s)
     G_B = 0.25 Beefy CPU utilization constants for P-store
     G_W = 0.13 Wimpy CPU utilization constants for P-store
 f_B(c) = 130.03 \times (100c)^{0.2369} (c=CPU util.) Beefy node power model
f_W(c) = 10.994 \times (100c)^{0.2875} (c=CPU util.) Wimpy node power model
  H = M_W \ge (Bld * Bld_{sel})/(N_B + N_W) Wimpy can build the hash table
```

HOMOGENEOUS EXECUTION

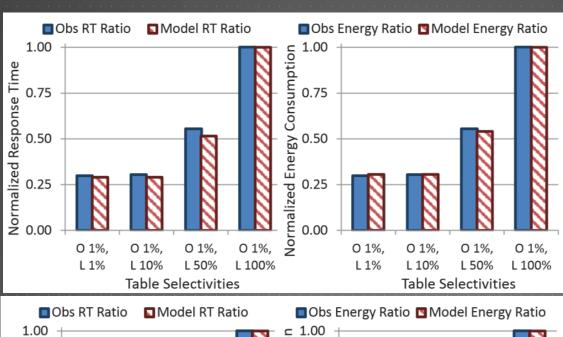
$$T_{prb} = \frac{Prb \times S_{prb}}{(N_B R_{Bprb}) + (N_W R_{Wprb})}$$

$$E_{prb} = T_{prb} \times (N_B f_B (G_B + \frac{U_{Bprb}}{C_B}) + N_W f_W (G_W + \frac{U_{Wprb}}{C_W}))$$

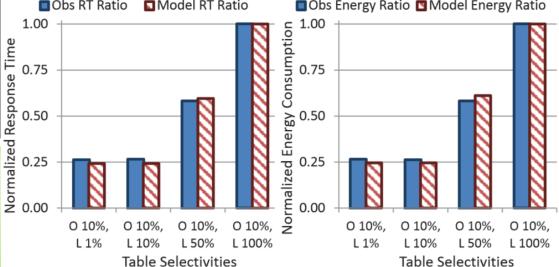
Probe phase response time and cluster energy consumption

T_{bld} Build phase time (s)	T_{prb} Probe phase time (s)			
$ E_{bld} $ Build phase energy (J)	$ E_{prb} $ Probe phase energy (J)			
N_B # Beefy nodes	N_W # Wimpy nodes			
M_B Beefy memory size (MB)	M_W Wimpy memory size (MB)			
$I \mid \text{Disk bandwidth (MB/s)}$	L Network bandwidth (MB/s)			
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MODEL VALIDATION

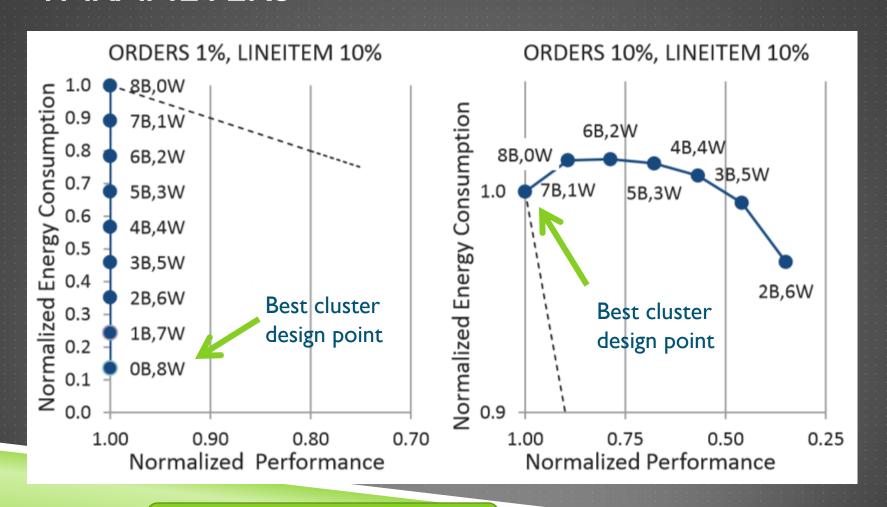


- Homogeneous Execution
- Error < 5%



- Heterogeneous Execution
- Error < 10%

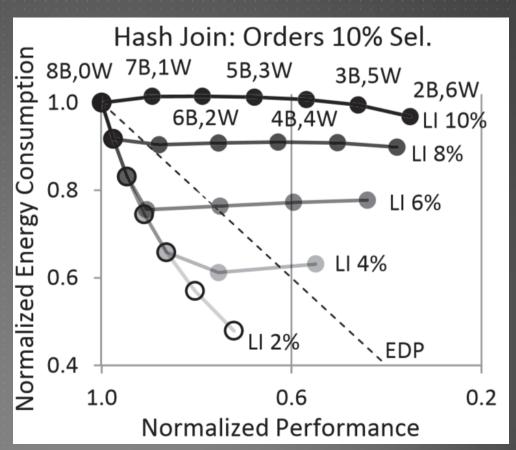
EXPLORING QUERY AND CLUSTER PARAMETERS



P-store hash join performance and energy efficiency

EXPLORING QUERY AND CLUSTER PARAMETERS

P-store hash join performance and energy efficiency LINEITEM (2-10% selectivity)



CLUSTER DESIGN PRINCIPLES (SUMMARY OF THE PAPER)

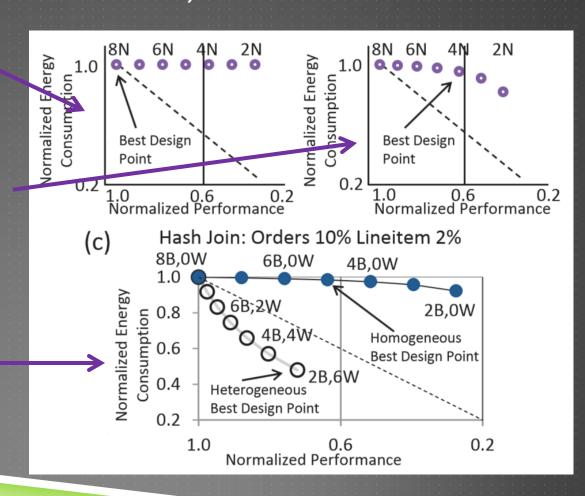
Query is highly scalable

Query is not scalable

→ Reduce performance
to meet required target
(SLAs)

Query is not scalable

→Homogeneous/
Heterogeneous cluster



DISCUSSION

- ▶ Modeling a system. How easy/difficult?
- ▶ Only single queries used. Acknowledged to include more workloads.
- Max cluster size is 16N
- The break-even time of installing new clusters has not been discussed
 - ► Homegeneous ← → Heterogenous
- Dynamic configuration of workloads/servers

Thank you!!