

## CS477 Formal Software Development Methods

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Slides based in part on previous lectures by Mahesh Vishwanathan, and  
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February 26, 2014

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## $\alpha$ -equivalence

- $\psi \stackrel{\alpha}{\equiv} \psi$
- If  $\psi_1 \stackrel{\alpha}{\equiv} \psi_2$  then  $\psi_2 \stackrel{\alpha}{\equiv} \psi$ .
- If  $\psi_1 \stackrel{\alpha}{\equiv} \psi_2$  and  $\psi_2 \stackrel{\alpha}{\equiv} \psi_3$  then  $\psi_1 \stackrel{\alpha}{\equiv} \psi_3$
- If  $x \notin \text{fv}(\psi)$  and  $y \notin \text{fv}(\psi)$  then  $\psi \stackrel{\alpha}{\equiv} \psi[x \leftrightarrow y]$ .
- If  $\psi_i \stackrel{\alpha}{\equiv} \psi'_i$  for  $i = 1, 2$  then
  - $(\psi_1) \stackrel{\alpha}{\equiv} (\psi'_1) \quad \neg\psi_1 \stackrel{\alpha}{\equiv} \neg\psi'_1$
  - $\psi_1 \otimes \psi_2 \stackrel{\alpha}{\equiv} \psi'_1 \otimes \psi'_2$  for  $\otimes \in \{\wedge, \vee, \Rightarrow, \Leftrightarrow\}$
  - $\mathcal{Q} z. \psi_1 \stackrel{\alpha}{\equiv} \mathcal{Q} z. \psi'_1$  for  $\mathcal{Q} \in \{\forall, \exists\}$

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## $\alpha$ -equivalence: Example

$$(x > 3 \wedge (\exists y. (\forall z. z \geq (y - x)) \vee (z \geq y))) \stackrel{\alpha}{\equiv} (x > 3 \wedge (\exists w. (\forall z. z \geq (w - x)) \vee (z \geq w)))$$

$$(x > 3 \wedge (\exists y. (\forall z. z \geq (y - x)) \vee (z \geq y))) \stackrel{\alpha}{\equiv} (x > 3 \wedge (\exists w. (\forall y. y \geq (w - x)) \vee (z \geq w)))$$

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## Proof Rules

Will give Sequent version of Natural Deduction rules  
All rules from Propositional Logic included

$$\frac{\Gamma \vdash \psi'[t/x]}{\Gamma \vdash \exists x. \psi'} \text{Ex I}$$

provided  $\psi \stackrel{\alpha}{\equiv} \psi'$

$$\frac{\Gamma \vdash \exists x. \psi \quad \Gamma \cup \{\psi[y/x]\} \vdash \varphi}{\Gamma \vdash \varphi} \text{Ex E}$$

provided  
 $y \notin \text{fv}(\varphi) \cup (\text{fv}(\psi) \setminus \{x\}) \cup \bigcup_{\psi' \in \Gamma} \text{fv}(\psi')$

$$\frac{\Gamma \vdash \psi[y/x]}{\Gamma \vdash \forall x. \psi} \text{All I}$$

provided  
 $y \notin (\text{fv}(\psi) \setminus \{x\}) \cup \bigcup_{\psi' \in \Gamma} \text{fv}(\psi')$

$$\frac{\Gamma \vdash \forall x. \psi \quad \Gamma \cup \{\psi[t/x]\} \vdash \varphi}{\Gamma \vdash \varphi} \text{All E}$$

provided  $\psi \stackrel{\alpha}{\equiv} \psi'$

## Example

Show

$$\frac{}{\{\} \vdash (\exists x. \forall y. x \leq y) \Rightarrow (\forall x. \exists y. y \leq x)}$$

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## Example

Show

$$\frac{\{\exists x. \forall y. x \leq y\} \vdash \forall x. \exists y. y \leq x}{\{\} \vdash (\exists x. \forall y. x \leq y) \Rightarrow (\forall x. \exists y. y \leq x)} \text{Imp I}$$

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## Example

Show

$$\begin{array}{c}
 \frac{\text{Hyp}}{\left\{ \exists x. \forall y. x \leq y; \forall y. z \leq y \right\} \vdash \forall y. z \leq y} \\
 \frac{\text{Hyp}}{\left\{ \exists x. \forall y. x \leq y; \forall y. z \leq y; z \leq x \right\} \vdash z \leq x} \text{ Hyp} \\
 \frac{\left\{ \exists x. \forall y. x \leq y; \forall y. z \leq y; z \leq x \right\} \vdash z \leq x}{\left\{ \exists x. \forall y. x \leq y; \forall y. z \leq y; z \leq x \right\} \vdash \exists y. y \leq x} \text{ Ex I} \\
 \frac{\left\{ \exists x. \forall y. x \leq y; \forall y. z \leq y; z \leq x \right\} \vdash \exists y. y \leq x}{\left\{ \exists x. \forall y. x \leq y; \forall y. z \leq y \right\} \vdash \exists y. y \leq x} \text{ All E} \\
 \frac{\left\{ \exists x. \forall y. x \leq y \right\} \vdash \exists x. \forall y. x \leq y \quad \left\{ \exists x. \forall y. x \leq y; \forall y. z \leq y \right\} \vdash \exists y. y \leq x}{\left\{ \exists x. \forall y. x \leq y \right\} \vdash \exists y. y \leq x} \text{ Hyp Ex E} \\
 \frac{\left\{ \exists x. \forall y. x \leq y \right\} \vdash \exists y. y \leq x}{\left\{ \exists x. \forall y. x \leq y \right\} \vdash \exists y. y \leq x} \text{ All I} \\
 \frac{\left\{ \exists x. \forall y. x \leq y \right\} \vdash \exists y. y \leq x}{\left\{ \right\} \vdash (\exists x. \forall y. x \leq y) \Rightarrow (\forall x. \exists y. y \leq x)} \text{ Imp I}
 \end{array}$$

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## Example of Failure

Let's try to show 1

$$\left\{ \right\} \vdash (\forall x. \exists y. y \leq x) \Rightarrow (\exists x. \forall y. x \leq y)$$

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## Example of Failure

Let's try to show 2

$$\frac{\left\{ \forall x. \exists y. y \leq x \right\} \vdash \exists x. \forall y. x \leq y}{\left\{ \right\} \vdash (\forall x. \exists y. y \leq x) \Rightarrow (\exists x. \forall y. x \leq y)} \text{ Imp I}$$

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## Example of Failure

Let's try to show 3

$$\frac{\left\{ \forall x. \exists y. y \leq x \right\} \vdash \forall y. z \leq y \quad \left\{ \forall x. \exists y. y \leq x \right\} \vdash \exists x. \forall y. x \leq y}{\left\{ \right\} \vdash (\forall x. \exists y. y \leq x) \Rightarrow (\exists x. \forall y. x \leq y)} \text{ Ex I Imp I}$$

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## Example of Failure

Let's try to show 4

$$\begin{array}{c}
 \frac{\left\{ \forall x. \exists y. y \leq x \right\} \vdash z \leq x}{\left\{ \forall x. \exists y. y \leq x \right\} \vdash \forall y. z \leq y} \text{ All I} \\
 \frac{\left\{ \forall x. \exists y. y \leq x \right\} \vdash \forall y. z \leq y}{\left\{ \forall x. \exists y. y \leq x \right\} \vdash \exists x. \forall y. x \leq y} \text{ Ex I} \\
 \frac{\left\{ \forall x. \exists y. y \leq x \right\} \vdash \exists x. \forall y. x \leq y}{\left\{ \right\} \vdash (\forall x. \exists y. y \leq x) \Rightarrow (\exists x. \forall y. x \leq y)} \text{ Imp I}
 \end{array}$$

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## Example of Failure

Let's try to show 5

$$\begin{array}{c}
 \frac{\left\{ \forall x. \exists y. y \leq x \right\} \vdash \forall x. \exists y. y \leq x}{\left\{ \forall x. \exists y. y \leq x; \exists y. y \leq x \right\} \vdash z \leq x} \\
 \frac{\left\{ \forall x. \exists y. y \leq x; \exists y. y \leq x \right\} \vdash z \leq x}{\left\{ \forall x. \exists y. y \leq x \right\} \vdash z \leq x} \text{ All E} \\
 \frac{\left\{ \forall x. \exists y. y \leq x \right\} \vdash z \leq x}{\left\{ \forall x. \exists y. y \leq x \right\} \vdash \forall y. z \leq y} \text{ All I} \\
 \frac{\left\{ \forall x. \exists y. y \leq x \right\} \vdash \forall y. z \leq y}{\left\{ \forall x. \exists y. y \leq x \right\} \vdash \exists x. \forall y. x \leq y} \text{ Ex I} \\
 \frac{\left\{ \forall x. \exists y. y \leq x \right\} \vdash \exists x. \forall y. x \leq y}{\left\{ \right\} \vdash (\forall x. \exists y. y \leq x) \Rightarrow (\exists x. \forall y. x \leq y)} \text{ Imp I}
 \end{array}$$

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## Example of Failure

Let's try to show

6

$$\frac{\{\forall x. \exists y. y \leq x\} \vdash \forall x. \exists y. y \leq x \quad \text{Hyp}}{\frac{\{\forall x. \exists y. y \leq x; \exists y. y \leq x\} \vdash z \leq x \quad \text{All E}}{\frac{\{\forall x. \exists y. y \leq x\} \vdash z \leq x \quad \text{All I}}{\frac{\{\forall x. \exists y. y \leq x\} \vdash \forall y. z \leq y \quad \text{Ex I}}{\frac{\{\forall x. \exists y. y \leq x\} \vdash \exists x. \forall y. x \leq y \quad \text{Imp I}}{\{\} \vdash (\forall x. \exists y. y \leq x) \Rightarrow (\exists x. \forall y. x \leq y)}}}}$$

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## Example of Failure

Let's try to show

7

$$\frac{\frac{\frac{\frac{\frac{\frac{\frac{\text{Something } \{ \text{Lab. } \exists y. y \leq x; } \exists y. y \leq x; z \leq x\} \vdash z \leq x \quad \text{Hyp}}{\frac{\{\forall x. \exists y. y \leq x\} \vdash \forall x. \exists y. y \leq x \quad \text{Ex E}}{\frac{\{\forall x. \exists y. y \leq x; \exists y. y \leq x; z \leq x\} \vdash z \leq x \quad \text{All E}}{\frac{\{\forall x. \exists y. y \leq x\} \vdash z \leq x \quad \text{All I}}{\frac{\{\forall x. \exists y. y \leq x\} \vdash \forall y. z \leq y \quad \text{Ex I}}{\frac{\{\forall x. \exists y. y \leq x\} \vdash \exists x. \forall y. x \leq y \quad \text{Imp I}}{\{\} \vdash (\forall x. \exists y. y \leq x) \Rightarrow (\exists x. \forall y. x \leq y)}}}}}}}}}}}$$

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## Example of Failure

Let's try to show

8

$$\frac{\{\forall x. \exists y. y \leq x\} \vdash \forall x. \exists y. y \leq x \quad \text{Hyp}}{\frac{\frac{\frac{\frac{\frac{\frac{\text{Something } \{ \text{Lab. } \exists y. y \leq x; } \exists y. y \leq x; z \leq x\} \vdash z \leq x \quad \text{Hyp}}{\{\forall x. \exists y. y \leq x; \exists y. y \leq x; z \leq x\} \vdash z \leq x \quad \text{Ex E}}{\frac{\{\forall x. \exists y. y \leq x\} \vdash z \leq x \quad \text{All E}}{\frac{\{\forall x. \exists y. y \leq x\} \vdash \forall y. z \leq y \quad \text{All I}}{\frac{\{\forall x. \exists y. y \leq x\} \vdash \exists x. \forall y. x \leq y \quad \text{Ex I}}{\frac{\{\} \vdash (\forall x. \exists y. y \leq x) \Rightarrow (\exists x. \forall y. x \leq y) \quad \text{Imp I}}{\{\} \vdash (\forall x. \exists y. y \leq x) \Rightarrow (\exists x. \forall y. x \leq y)}}}}}}}}}}$$

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## Example of Failure

Let's try to show

9

$$\frac{\{\forall x. \exists y. y \leq x\} \vdash \forall x. \exists y. y \leq x \quad \text{Hyp}}{\frac{\frac{\frac{\frac{\frac{\frac{\text{Something } \{ \text{Lab. } \exists y. y \leq x; } \exists y. y \leq x; z \leq x\} \vdash z \leq x \quad \text{Hyp}}{\{\forall x. \exists y. y \leq x; \exists y. y \leq x; z \leq x\} \vdash z \leq x \quad \text{Ex E}}{\frac{\{\forall x. \exists y. y \leq x\} \vdash z \leq x \quad \text{All E}}{\frac{\{\forall x. \exists y. y \leq x\} \vdash \forall y. z \leq y \quad \text{All I}}{\frac{\{\forall x. \exists y. y \leq x\} \vdash \exists x. \forall y. x \leq y \quad \text{Ex I}}{\frac{\{\} \vdash (\forall x. \exists y. y \leq x) \Rightarrow (\exists x. \forall y. x \leq y) \quad \text{Imp I}}{\{\} \vdash (\forall x. \exists y. y \leq x) \Rightarrow (\exists x. \forall y. x \leq y)}}}}}}}}}}$$

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## Example of Failure

Let's try to show

10

$$\frac{\{\forall x. \exists y. y \leq x\} \vdash \forall x. \exists y. y \leq x \quad \text{Hyp}}{\frac{\frac{\frac{\frac{\frac{\frac{\text{Something } \{ \text{Lab. } \exists y. y \leq x; } \exists y. y \leq x; z \leq x\} \vdash z \leq x \quad \text{Hyp}}{\{\forall x. \exists y. y \leq x; \exists y. y \leq x; z \leq x\} \vdash z \leq x \quad \text{Ex E}}{\frac{\{\forall x. \exists y. y \leq x\} \vdash z \leq x \quad \text{All E}}{\frac{\{\forall x. \exists y. y \leq x\} \vdash \forall y. z \leq y \quad \text{All I}}{\frac{\{\forall x. \exists y. y \leq x\} \vdash \exists x. \forall y. x \leq y \quad \text{Ex I}}{\frac{\{\} \vdash (\forall x. \exists y. y \leq x) \Rightarrow (\exists x. \forall y. x \leq y) \quad \text{Imp I}}{\{\} \vdash (\forall x. \exists y. y \leq x) \Rightarrow (\exists x. \forall y. x \leq y)}}}}}}}}}}$$

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## Floyd-Hoare Logic

- Also called **Axiomatic Semantics**
- Based on formal logic (first order predicate calculus)
- Logical system built from **axioms** and **inference rules**
- Mainly suited to simple imperative programming languages
- Ideas applicable quite broadly

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## Floyd-Hoare Logic

- Used to formally prove a property (**post-condition**) of the **state** (the values of the program variables) after the execution of program, assuming another property (**pre-condition**) of the state holds before execution

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## Floyd-Hoare Logic

- Goal: Derive statements of form

$$\{P\} C \{Q\}$$

- $P, Q$  logical statements about state,  $P$  precondition,  $Q$  postcondition,  $C$  program

- Example:

$$\{x = 1\} x := x + 1 \{x = 2\}$$

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## Floyd-Hoare Logic

- Approach:** For each type of language statement, give an axiom or inference rule stating how to derive assertions of form

$$\{P\} C \{Q\}$$

where  $C$  is a statement of that type

- Compose axioms and inference rules to build proofs for complex programs

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## Partial vs Total Correctness

- An expression  $\{P\} C \{Q\}$  is a **partial correctness** statement
- For **total correctness** must also prove that  $C$  terminates (i.e. doesn't run forever)
  - Written:  $[P] C [Q]$
- Will only consider partial correctness here

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## Simple Imperative Language

- We will give rules for simple imperative language

$$\begin{aligned} \langle \text{command} \rangle ::= & \langle \text{variable} \rangle := \langle \text{term} \rangle \\ | & \langle \text{command} \rangle ; \dots ; \langle \text{command} \rangle \\ | & \text{if } \langle \text{statement} \rangle \text{ then } \langle \text{command} \rangle \text{ else } \langle \text{command} \rangle \\ | & \text{while } \langle \text{statement} \rangle \text{ do } \langle \text{command} \rangle \end{aligned}$$

- Could add more features, like for-loops

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## Substitution

- Notation:  $P[e/v]$  (sometimes  $P[v \rightarrow e]$ )
- Meaning: Replace every  $v$  in  $P$  by  $e$
- Example:

$$(x + 2)[y - 1/x] = ((y - 1) + 2)$$

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## The Assignment Rule

$$\overline{\{P[e/x]\}} \ x := e \ \{P\}$$

Example:

$$\overline{\{ \quad ? \quad \}} \ x := y \ \{ x = 2 \}$$

## The Assignment Rule

$$\overline{\{P[e/x]\}} \ x := e \ \{P\}$$

Example:

$$\overline{\{ \square = 2 \}} \ x := y \ \{ \square = 2 \}$$

## The Assignment Rule

$$\overline{\{P[e/x]\}} \ x := e \ \{P\}$$

Example:

$$\overline{\{ \square = 2 \}} \ x := y \ \{ \square = 2 \}$$

## The Assignment Rule

$$\overline{\{P[e/x]\}} \ x := e \ \{P\}$$

Examples:

$$\overline{\{y = 2\}} \ x := y \ \{x = 2\}$$

$$\overline{\{y = 2\}} \ x := 2 \ \{y = x\}$$

$$\overline{\{x + 1 = n + 1\}} \ x := x + 1 \ \{x = n + 1\}$$

$$\overline{\{2 = 2\}} \ x := 2 \ \{x = 2\}$$

## The Assignment Rule – Your Turn

- What is the weakest precondition of

$$x := x + y \ \{x + y = wx\}?$$

$$\{ \quad \quad \quad ? \quad \quad \quad \} \\ x := x + y \\ \{x + y = wx\}$$

## The Assignment Rule – Your Turn

- What is the weakest precondition of

$$x := x + y \ \{x + y = wx\}?$$

$$\{ (x + y) + y = w(x + y) \} \\ x := x + y \\ \{x + y = wx\}$$

## Precondition Strengthening

$$\frac{(P \Rightarrow P') \{P'\} C \{Q\}}{\{P\} C \{Q\}}$$

- Meaning: If we can show that  $P$  implies  $P'$  (i.e.  $(P \Rightarrow P')$ ) and we can show that  $\{P\} C \{Q\}$ , then we know that  $\{P\} C \{Q\}$
- $P$  is **stronger** than  $P'$  means  $P \Rightarrow P'$

## Precondition Strengthening

- Examples:

$$\frac{x = 3 \Rightarrow x < 7 \quad \{x < 7\} \ x := x + 3 \ \{x < 10\}}{\{x = 3\} \ x := x + 3 \ \{x < 10\}}$$

$$\frac{\text{True} \Rightarrow (2 = 2) \quad \{2 = 2\} \ x := 2 \ \{x = 2\}}{\{\text{True}\} \ x := 2 \ \{x = 2\}}$$

$$\frac{x = n \Rightarrow x + 1 = n + 1 \quad \{x + 1 = n + 1\} \ x := x + 1 \ \{x = n + 1\}}{\{x = n\} \ x := x + 1 \ \{x = n + 1\}}$$

## Which Inferences Are Correct?

$$\frac{\{x > 0 \wedge x < 5\} \ x := x * x \ \{x < 25\}}{\{x = 3\} \ x := x * x \ \{x < 25\}}$$

$$\frac{\{x = 3\} \ x := x * x \ \{x < 25\}}{\{x > 0 \wedge x < 5\} \ x := x * x \ \{x < 25\}}$$

$$\frac{\{x * x < 25\} \ x := x * x \ \{x < 25\}}{\{x > 0 \wedge x < 5\} \ x := x * x \ \{x < 25\}}$$

## Which Inferences Are Correct?

$$\frac{\{x > 0 \wedge x < 5\} \ x := x * x \ \{x < 25\}}{\{x = 3\} \ x := x * x \ \{x < 25\}} \text{ YES}$$

$$\frac{\{x = 3\} \ x := x * x \ \{x < 25\}}{\{x > 0 \wedge x < 5\} \ x := x * x \ \{x < 25\}}$$

$$\frac{\{x * x < 25\} \ x := x * x \ \{x < 25\}}{\{x > 0 \wedge x < 5\} \ x := x * x \ \{x < 25\}}$$

## Which Inferences Are Correct?

$$\frac{\{x > 0 \wedge x < 5\} \ x := x * x \ \{x < 25\}}{\{x = 3\} \ x := x * x \ \{x < 25\}} \text{ YES}$$

$$\frac{\{x = 3\} \ x := x * x \ \{x < 25\}}{\{x > 0 \wedge x < 5\} \ x := x * x \ \{x < 25\}} \text{ NO}$$

$$\frac{\{x * x < 25\} \ x := x * x \ \{x < 25\}}{\{x > 0 \wedge x < 5\} \ x := x * x \ \{x < 25\}}$$

## Which Inferences Are Correct?

$$\frac{\{x > 0 \wedge x < 5\} \ x := x * x \ \{x < 25\}}{\{x = 3\} \ x := x * x \ \{x < 25\}} \text{ YES}$$

$$\frac{\{x = 3\} \ x := x * x \ \{x < 25\}}{\{x > 0 \wedge x < 5\} \ x := x * x \ \{x < 25\}} \text{ NO}$$

$$\frac{\{x * x < 25\} \ x := x * x \ \{x < 25\}}{\{x > 0 \wedge x < 5\} \ x := x * x \ \{x < 25\}} \text{ YES}$$