CS477 Formal Software Development Methods

Elsa L Gunter 2112 SC, UIUC egunter@illinois.edu http://courses.engr.illinois.edu/cs477

Slides based in part on previous lectures by Mahesh Vishwanathan, and by Gul Agha

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Getting Started with Isabelle

- Choice
 - Use Isabelle on EWS
 - Install on your machine
 - Both
- On EWS
 - Assuming you are running an X client, log in to EWS:

ssh -Y <netid>@remlnx.ews.illinois.edu

- -Y used to forward X packets securely
- To start Isabelle with jedit /class/cs477/bin/isabelle jedit
- Older versions of Isabelle used emacs and ProofGeneral
- Will assume jedit here

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My First Theory File

File name: my_theory.thy Contents:

theory my_theory imports Main begin

thm impI

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lemma trivial: "A \longrightarrow A"
apply (rule impI)
apply assumption
done (* of lemma *)
```

```
thm trivial
```

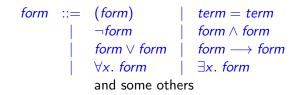
end (* of theory file *)

Overview of Isabelle/HOL

- HOL = Higher-Order Logic
- HOL = Types + Lambda Calculus + Logic
- HOL has
 - datatypes
 - recursive functions
 - logical operators (\land , \lor , \neg , \longrightarrow , \forall , \exists , ...)
- Contains propositional logic, first-order logic
- HOL is very similar to a functional programming language
- Higher-order = functions are values, too!
- We'll start with propositional and first order logic

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• Syntax (in decreasing priority):



• Scope of quantifiers: as far to the right as possible

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- $\neg A \land B \lor C \equiv ((\neg A) \land B) \lor C$
- $A \wedge B = C \equiv A \wedge (B = C)$
- $\forall x. P x \land Q x \equiv \forall x. (P x \land Q x)$
- $\forall x. \exists y. P \times y \land Q x \equiv \forall x. (\exists y. (P \times y \land Q x))$

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General schema:

lemma name: "..."
apply (...)
:
done

First ... theorem statement (...) are *proof methods*

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- "completes" any proof (by giving up, and accepting it)
- Suitable for top-down development of theories:
- Assume lemmas first, prove them later.

Only allowed for interactive proof!

- Distinct from HOL syntax
- Contains HOL syntax within it
- Also the same as HOL need to not confuse them

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Syntax:

theory MyThimports $ImpTh_1 \ldots ImpTh_n$ begin declarations, definitions, theorems, proofs, ... end

- MyTh: name of theory being built. Must live in file MyTh.thy.
- ImpTh_i: name of imported theories. Importing is transitive.

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Implication: \implies (==>)

For separating premises and conclusion of theorems / rules

Equality: \equiv (==) For definitions

Universal Quantifier: Λ (!!)

Usually inserted and removed by Isabelle automatically

Do not use inside HOL formulae

abbreviates

 $[|A_1; \ldots; A_n|] \Longrightarrow B$ $A_1 \Longrightarrow \ldots \Longrightarrow A_n \Longrightarrow B$ and means the rule (or potential rule):

 $\frac{A_1;\ldots;A_n}{B}$

;
$$\approx$$
 "and"

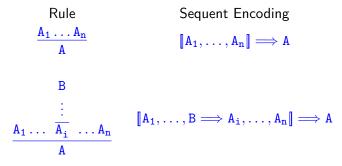
Note: A theorem is a rule; a rule is a theorem.

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- 1. $\Lambda x_1 \dots x_m$. $[|A_1; \dots; A_n|] \Longrightarrow B$
- $x_1 \dots x_m$ Local constants (fixed variables) $A_1 \dots A_n$ Local assumptionsBActual (sub)goal

• Isabelle uses Natural Deduction proofs

- Uses (modified) sequent encoding
- Rule notation:



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For each logical operator \oplus , have two kinds of rules:

Introduction: How can I prove $A \oplus B$?

 $\frac{?}{A \oplus B}$

Elimination: What can I prove using $A \oplus B$?

 $\frac{\ldots A \oplus B \ldots}{?}$

 $\frac{A_1 \dots A_n}{A}$

Introduction rule:

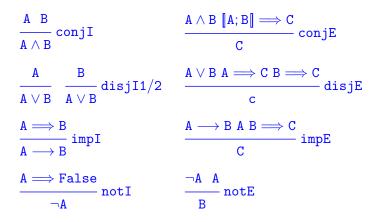
To prove A it suffices to prove $A_1 \dots A_n$.

Elimination rule:

If we know A_1 and we want to prove Ait suffices to prove $A_2 \dots A_n$

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Natural Deduction for Propositional Logic



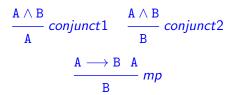
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Natural Deduction for Propositional Logic

$$\frac{A \Longrightarrow B \quad B \Longrightarrow A}{A = B} \quad iffI \qquad \frac{A = B \quad A}{B} \quad iffD1$$
$$\frac{A = B \quad B}{A} \quad iffD2$$

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Compare to elimination rules:

 $\frac{A \wedge B \ \|A; B\| \Longrightarrow C}{C} \ conjE \quad \frac{A \longrightarrow B \ A \ B \Longrightarrow C}{C} \ impE$

$$\frac{\cancel{A} \Longrightarrow \texttt{False}}{\cancel{A}} \textit{ccontr} \quad \frac{\cancel{A} \Longrightarrow \cancel{A}}{\cancel{A}} \textit{classical}$$

- ccontr and classical are not derivable from the Natural Deduction rules.
- They make the logic *"classical"*, i.e. *"non-constructive* or *"non-intuitionistic"*.

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Proof by Assumption

$$\frac{\mathtt{A}_1 \ldots \mathtt{A}_i \ldots \mathtt{A}_n}{\mathtt{A}_i}$$

- Proof method: assumption
- Use:

Proves:

apply assumption $[\![A_1;\ldots;A_n]\!] \Longrightarrow A$

by unifying A with one of the A_i

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Rule Application: The Rough Idea

Applying rule $[A_1; \ldots; A_n] \implies A$ to subgoal *C*:

- Unify A and C
- Replace C with n new subgoals: $A'_1 \ldots A'_n$

Backwards reduction, like in Prolog Example: rule: $[?P; ?Q] \implies ?P \land ?Q$ subgoal: 1. $A \land B$

Result: 1. A2. B

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Rule Application: More Complete Idea

Applying rule $[A_1; \ldots; A_n] \implies A$ to subgoal *C*:

- Unify A and C with (meta)-substitution σ
- Specialize goal to $\sigma(C)$
- Replace C with n new subgoals: $\sigma(A_1) \ldots \sigma(A_n)$

Note: schematic variables in C treated as existential variables Does there exist value for ?X in C that makes C true? (Still not the whole story)

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rule Application

Rule: $[A_1; \ldots; A_n] \Longrightarrow A$

Subgoal: 1. $[B_1; ...; B_m] \Longrightarrow C$

Substitution: $\sigma(A) \equiv \sigma(C)$

New subgoals:

Proves: Command:

1.
$$[\sigma(B_1); ...; \sigma(B_m)] \Longrightarrow \sigma(A_1)$$

:
n. $[\sigma(B_1); ...; \sigma(B_m)] \Longrightarrow \sigma(A_n)$
 $\sigma(B_1); ...; \sigma(B_m)] \Longrightarrow \sigma(C)$
apply (rule)

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apply (erule < elim-rule>)

Like rule but also

- Unifies first premise of rule with an assumption
- Eliminates that assumption instead of conclusion

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| Rule: | $\llbracket ?P \land ?Q; \llbracket ?P; ?Q \rrbracket \Longrightarrow ?R \rrbracket \Longrightarrow ?R$ |
|--------------|---|
| Subgoal: | $1. \ [\![X; A \land B; Y]\!] \Longrightarrow Z$ |
| Unification: | $P \land Q \equiv A \land B$ and $R \equiv Z$ |
| New subgoal: | $1. \ \llbracket X; Y \rrbracket \Longrightarrow \llbracket A; B \rrbracket \Longrightarrow Z$ |
| Same as: | $1.[\![X;Y;A;B]\!] \Longrightarrow Z$ |

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