

CS477 Formal Software Development Methods

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How to build a Reduced Ordered BDD

- Let u, l, h be bdd nodes, i be a variable index
 - $\text{var}(u)$ gives the variable at u
 - $\text{low}(u)$ gives node pointed to by false out-edge of u
 - $\text{high}(u)$ gives node pointed to by true out-edge of u
- Assume to tables: T and H
- $T : u \rightarrow (i, l, h)$
 - init T , update $(T, (u, (i, l, h)))$,
- $H : (i, l, h) \rightarrow u$ - Inverse of T
 - init H , member $(H, (i, l, h))$, lookup $(H, (i, l, h))$, insert $(H, ((i, l, h), u))$

How to build a Reduced Ordered BDD

- For efficiency, use memoizing function
- function $\text{Mk}[T, H](i, l, h) =$

```

if l = h then return l
else if member (H, (i, l, h))
  then return (lookup (H, (i, l, h)))
else update (T, (u, (i, l, h)));
  insert(H, ((i, l, h), u));
  return u
end

```

How to build a Reduced Ordered BDD


- Main function for building ROBDD
- p proposition; $x_1 > \dots > x_n$ variables in p
- function $\text{Build}[T, H](p, i) =$

```


if i = 0 then if p = False then False else True
else let l = Build[T, H](p[False/xi], (i-1)) in
  let h = Build[T, H](p[True/xi], (i-1)) in
  let u = Mk[T, H](i, l, h) in
  return u
end

```

$(A \wedge B) \vee (\text{not } C)$ Variables: $C > B > A$

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$(A \wedge B) \vee (\text{not } C)$ 

$((A \wedge B) \vee (\text{not } C))[\text{False}/C] =$
True

$((A \wedge B) \vee (\text{not } C))[\text{True}/C]$
 $= (A \wedge B)$

