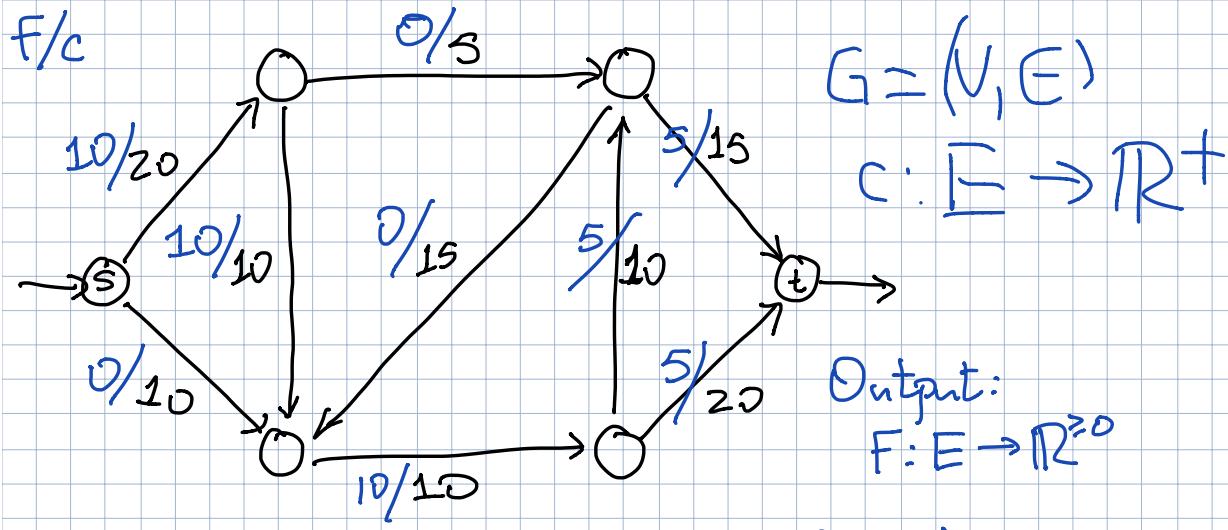


Harris and Ross 1950 (1999)

## Maximum Flow

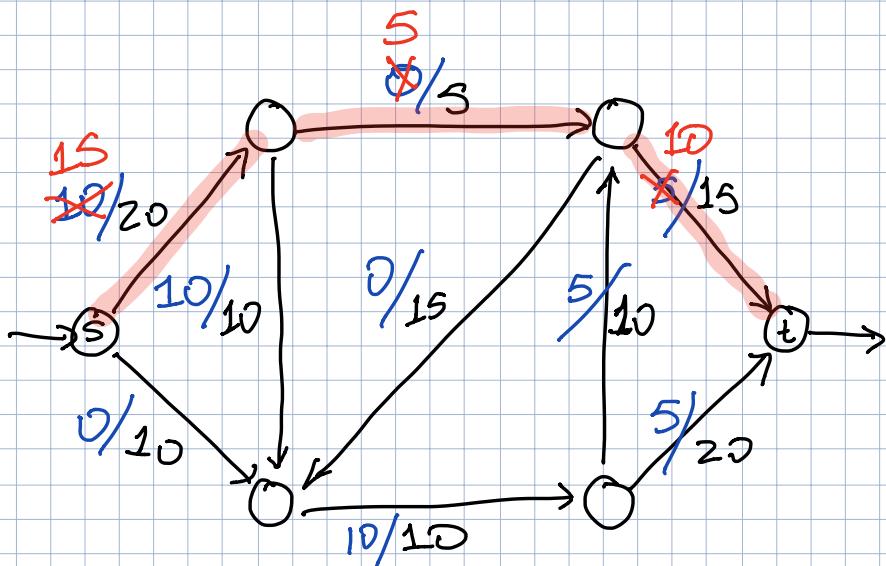
# Minimum Cut



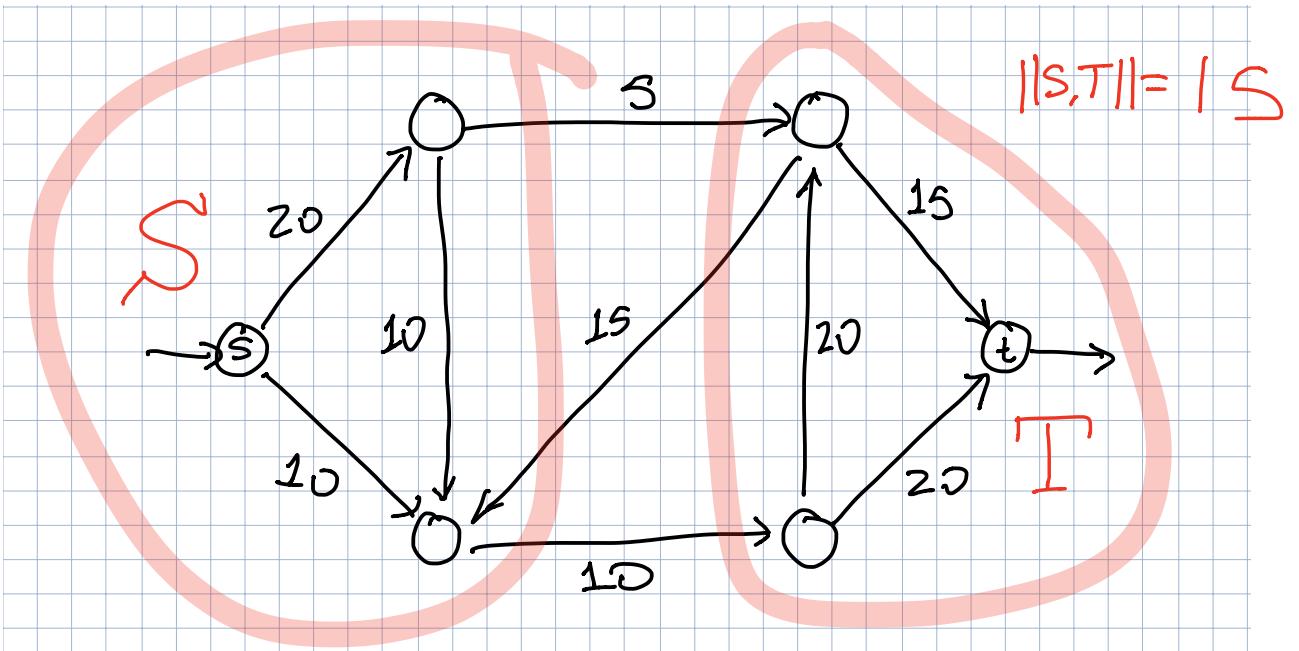
capacity:  $0 \leq F(u \rightarrow v) \leq c(u \rightarrow v)$

conservation balance  $\sum_u F(u \rightarrow v) = \sum_w F(v \rightarrow w)$  for all  $v \neq s, t$

maximize  $|F| = \sum_w F(s \rightarrow w) - \sum_u F(u \rightarrow s)$



If  $|F| = |(S, T)|$ , then  $F$  is max flow  
 $S, T$  is min cut



$(s,t)$ -cut = partition of vertices into  
 $S \ni s$  and  $T \ni t$

$$\text{capacity } ||S, T|| := \sum_{u \in S} \sum_{v \in T} c(u \rightarrow v)$$

Lemma: Let  $F$  be any flow  
 $S, T$  be any cut  
Then  $|F| \leq |S, T|$ .

Proof:

$$|F| = \sum_w F(s \rightarrow w) - \sum_u F(u \rightarrow s)$$

$|S, T| \leftarrow |F|$   
def  $|F|$

$$= \sum_{v \in S} \left( \sum_w F(v \rightarrow w) - \sum_u F(u \rightarrow v) \right)$$

balance

$$= \sum_{v \in S} \left( \sum_{w \in T} F(v \rightarrow w) - \sum_{u \in T} F(u \rightarrow v) \right)$$

remove dups ±

$$\leq \sum_{v \in S} \sum_{w \in T} F(v \rightarrow w)$$

$f \geq 0$

$$\leq \sum_{v \in S} \sum_{w \in T} c(v \rightarrow w)$$

$f \leq c$

$$= |S, T|$$

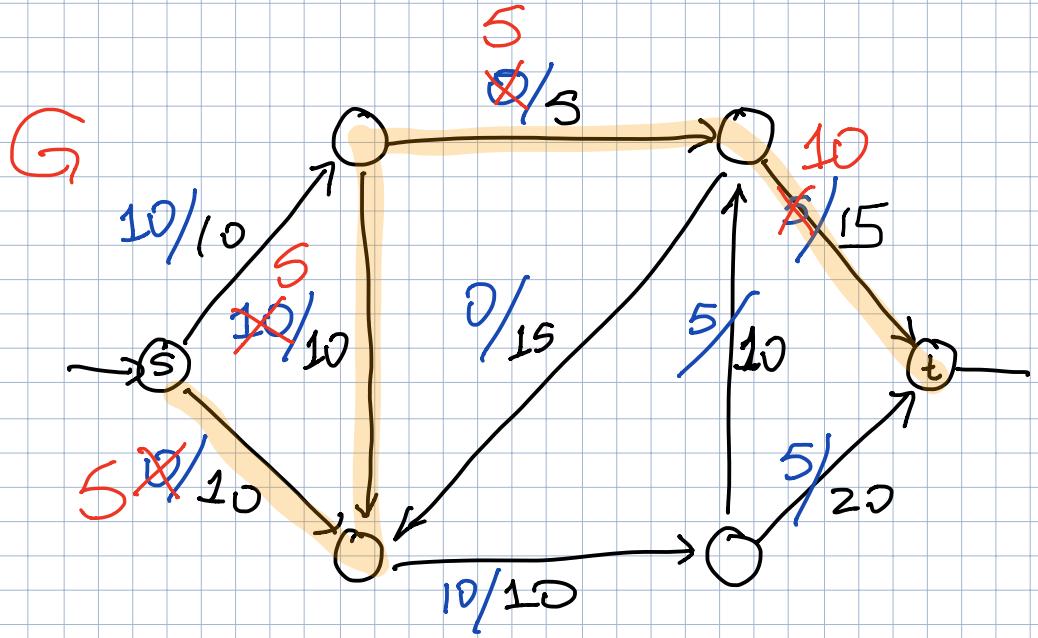
def  $|S, T|$

Corollary:

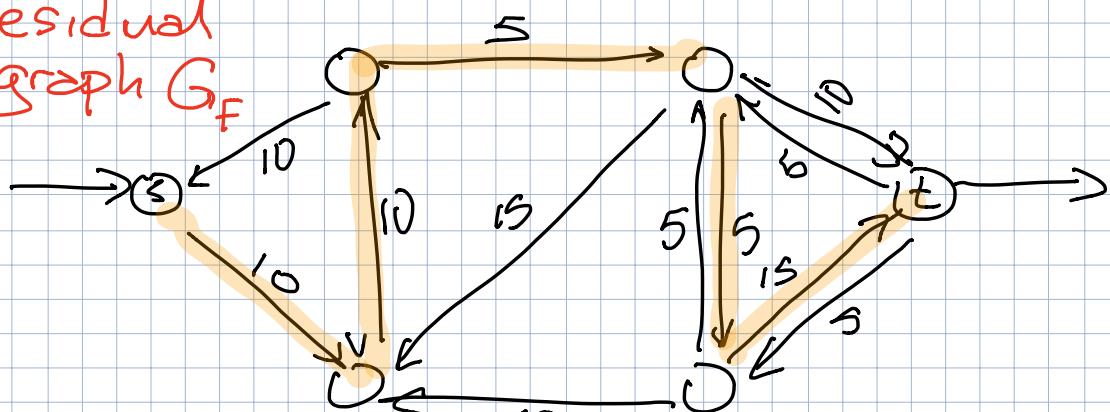
$$f(u \rightarrow v) = \begin{cases} 0 & \text{if } u \notin T \vee v \in S \\ c(u \rightarrow v) & \text{if } u \in S \wedge v \in T \end{cases}$$

IFF  $|F| = |S, T|$

$F$  is max     $S, T$  is min cut



Residual graph  $G_F$



$$C_f(u \rightarrow v) = \begin{cases} c(u \rightarrow v) - f(u \rightarrow v) & \text{if } u \rightarrow v \in E \\ f(v \rightarrow u) & \text{if } v \rightarrow u \in E \end{cases}$$

IF  $G_F$  has a path from  $s$  to  $t$ :  $O(V+E)$  time

$f$  is not maxflow

Let  $F = \min C_f(u \rightarrow v)$  along path

$f' \leftarrow f + F \cdot \text{path}$

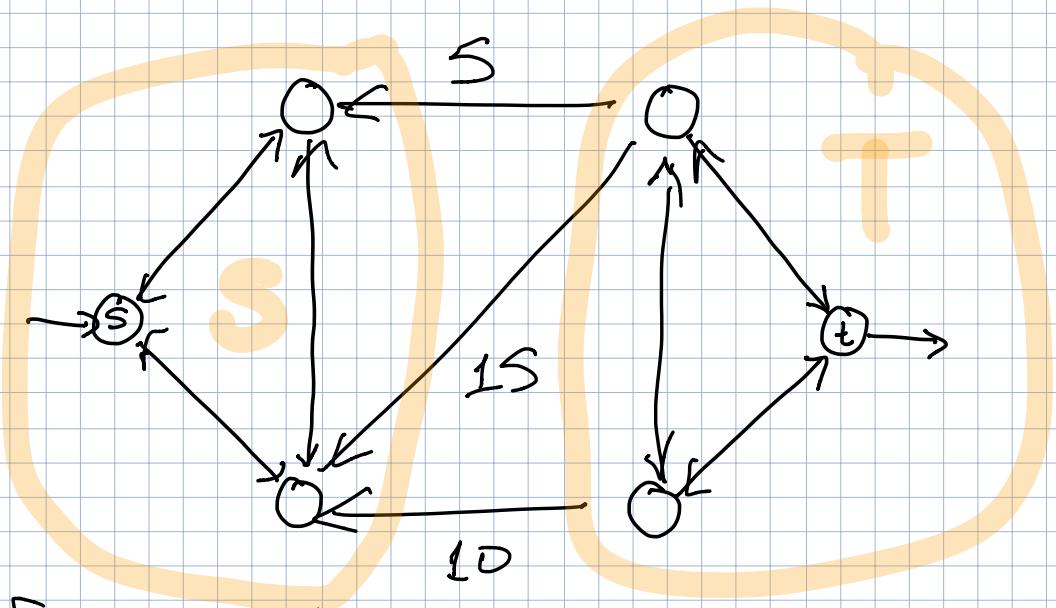
## Ford Fulkerson

$f \leftarrow 0$

while  $G_f$  has path from  $s$  to  $t$

augment  $f$  along that path

return  $f$



if no  $s \rightarrow t$  path

$$S = \text{Reach}(s) \quad T = V \setminus S$$

All edges  $S \rightarrow T$  full

$T \rightarrow S$  empty

$\Rightarrow f$  is maxflow and  $S, T$  are min cut

① IF all capacities are ints

Every iter increases  $|F|$  by  $\geq 1$

$$\Rightarrow \text{#iters} \leq |F^*|$$

↑  
output

in worst case,

exponential.

② If capacities are real #s

