

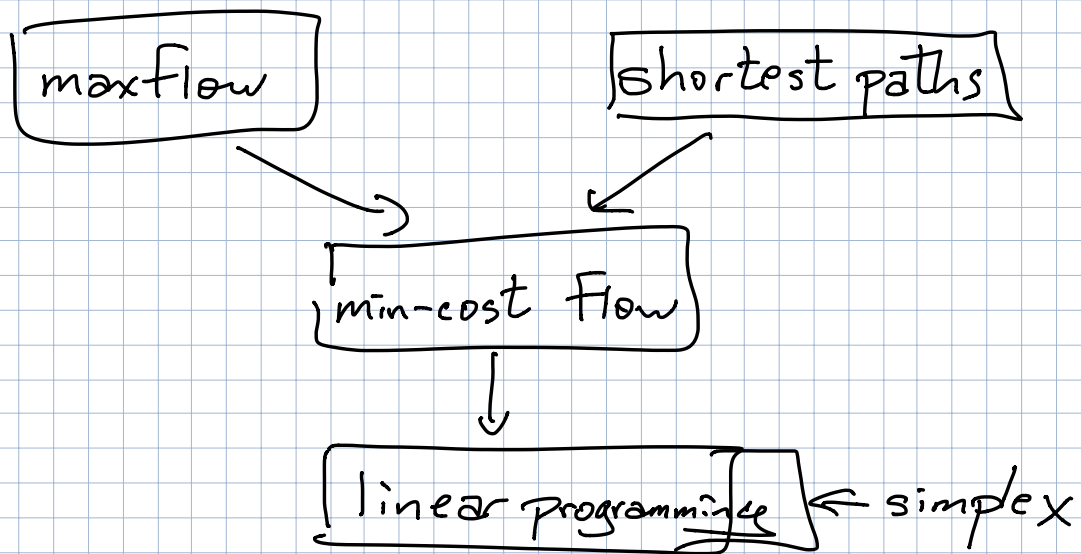
Midterm 2 - 1 week from today

no lecture / opt review session

Regrades back Thu. | HW 7 due today

Conflict exam Wed - register online

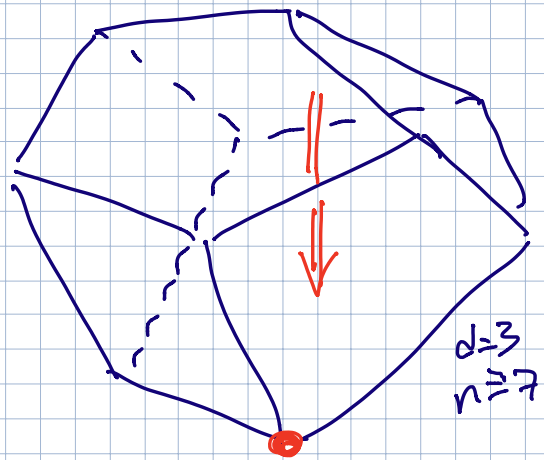
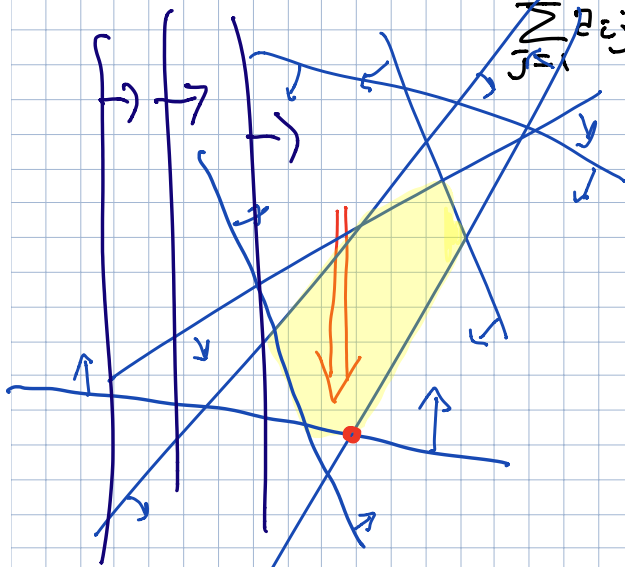
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# Linear Programming

$$\begin{aligned} & \max \sum_{j=1}^d c_j x_j && x_j \text{ --- variable} \\ & \text{subject to } \sum_{j=1}^d a_{ij} x_j \leq b_i && \text{for } i=1 \dots p \\ & \sum_{j=1}^d a_{ij} x_j \geq b_i && \text{for } i=p+1 \dots q \\ & \sum_{j=1}^d a_{ij} x_j = b_i && \text{for } i=q+1 \dots n \end{aligned}$$

$a_{ij}, b_i, c_j$  - constants



$$\begin{aligned} x & \in \mathbb{R}^d & c & \in \mathbb{R}^d \\ b & \in \mathbb{R}^n & A & \in \mathbb{R}^{n \times d} \end{aligned}$$

Canonical Form

$$\begin{aligned} & \max \quad c \cdot x \\ & \text{s.t.} \quad Ax \leq b \\ & \quad \quad x \geq 0 \end{aligned}$$

$$\begin{aligned} & \leftarrow \sum_j a_{ij} x_j \leq b_i \text{ for all } i \\ & \leftarrow x_1 \geq 0 \quad x_2 \geq 0 \quad \dots \quad x_d \geq 0 \end{aligned}$$

# Maximum flow

maximize  $\sum_w f_{s \rightarrow w} - \sum_n f_{u \rightarrow s}$

subj. to  $\sum_w f_{v \rightarrow w} - \sum_n f_{u \rightarrow v} = 0$  for all  $v \neq s, t$

A

$f_{u \rightarrow v} \leq C_{u \rightarrow v}$  for all  $u \rightarrow v$

$f_{u \rightarrow v} \geq 0$  for all  $u \rightarrow v$

# Minimum cut

maximize  
minimize  $\sum_{u \rightarrow v} -C_{u \rightarrow v} \cdot X_{u \rightarrow v}$

$X_{u \rightarrow v} \geq 0$  for all  $u \rightarrow v$

~~$X_{u \rightarrow v} \leq 1$  for all  $u \rightarrow v$~~

A

$X_{u \rightarrow v} + S_v - S_u \geq 0$  for all  $u \rightarrow v$

$S_s = 1$

$S_t = 0$

$X_{u \rightarrow v} = 1$  means  $u \in S$   
 $v \in T$

= 0 otherwise

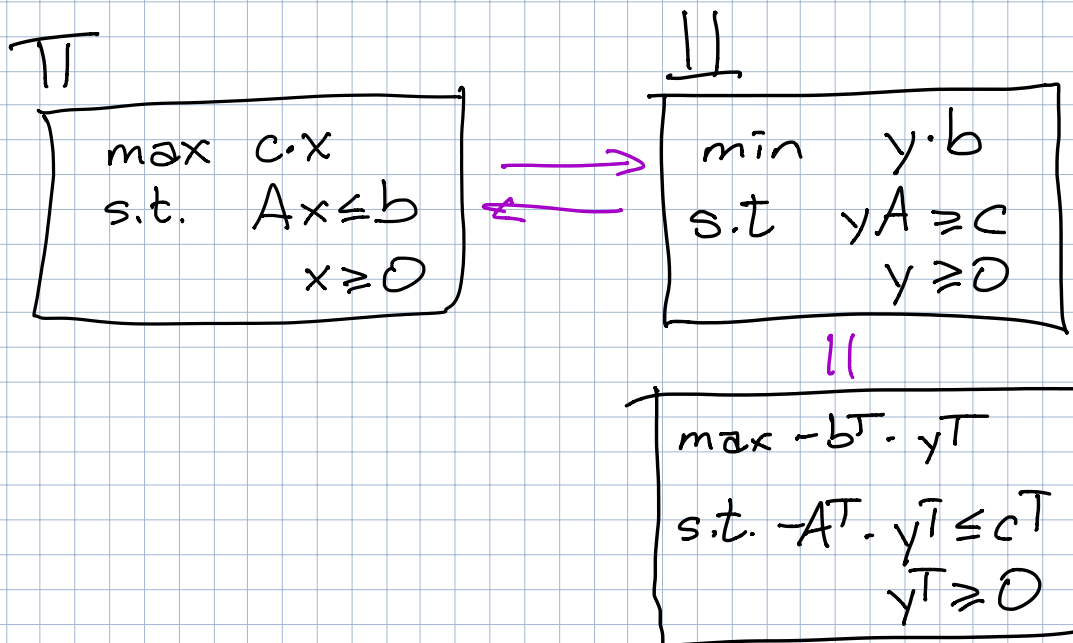
$S_v = 1$  if  $v \in S$   
0 if  $v \in T$

$0 \xrightarrow{\geq 0} 0$

$0 \xrightarrow{\geq -1} 1$

$1 \xrightarrow{\geq 0} 1$

$1 \xrightarrow{\geq 1} 0$



## Fundamental Theorem of LP!

$x^*$  is opt solution for  $\Pi$   
 $y^*$  is opt solution for  $\Pi$

iff  $c \cdot x^* = y^* A x^* = y^* b$

Weak FTLP:

$x$  is feasible for  $\Pi$  and  $y$  is feasible for  $\Pi$

then  $c x \leq \underline{y A x} \leq y b$

Proof:  $Ax \leq b \quad y \geq 0$

$$\begin{array}{ccc} \Downarrow & & \Downarrow \\ yAx & \leq & yb \end{array}$$



$$u \rightarrow v \quad \text{dist}(v) \leq \text{dist}(u) + l(u \rightarrow v)$$

$$\begin{array}{l} \text{maximize} \quad 4x_1 + x_2 + 3x_3 \\ \text{subj. to} \quad x_1 + 4x_2 \leq 2 \\ \quad \quad \quad 3x_1 - x_2 + x_3 \leq 4 \\ \quad \quad \quad x_1, x_2, x_3 \geq 0 \end{array}$$

$$\sigma^* = \text{opt value} \quad x = (1, 0, 0) \rightarrow \sigma^* \geq 4$$

$$x = (0, 0, 3) \rightarrow \sigma^* \geq 9$$

$$\sigma^* \leq ?$$

$$y_1, y_2 \geq 0$$

$$\begin{array}{l} \text{maximize} \quad 4x_1 + x_2 + 3x_3 \\ \text{subj. to} \quad y_1(x_1 + 4x_2) \leq 2y_1 \\ \quad \quad \quad y_2(3x_1 - x_2 + x_3) \leq 4y_2 \\ \quad \quad \quad x_1, x_2, x_3 \geq 0 \end{array}$$

$$y_1(x_1 + 4x_2) + y_2(3x_1 - x_2 + x_3) \leq 2y_1 + 4y_2$$

$$\underline{(y_1 + 3y_2)}x_1 + \underline{(4y_1 - y_2)}x_2 + \underline{y_2}x_3 \geq \underline{4}x_1 + \underline{1}x_2 + \underline{3}x_3$$

$$\begin{array}{l} \min \quad 2y_1 + 4y_2 \\ \text{s.t.} \quad y_1 + 3y_2 \geq 4 \\ \quad \quad 4y_1 - y_2 \geq 1 \\ \quad \quad y_2 \geq 3 \\ \quad \quad y_1, y_2 \geq 0 \end{array}$$