

# CS473: Algorithms

<https://courses.engr.illinois.edu/cs473>

Fibonacci #s:

$$F_0 = 0$$

$$F_1 = 1$$

$$F_n = F_{n-1} + F_{n-2}$$

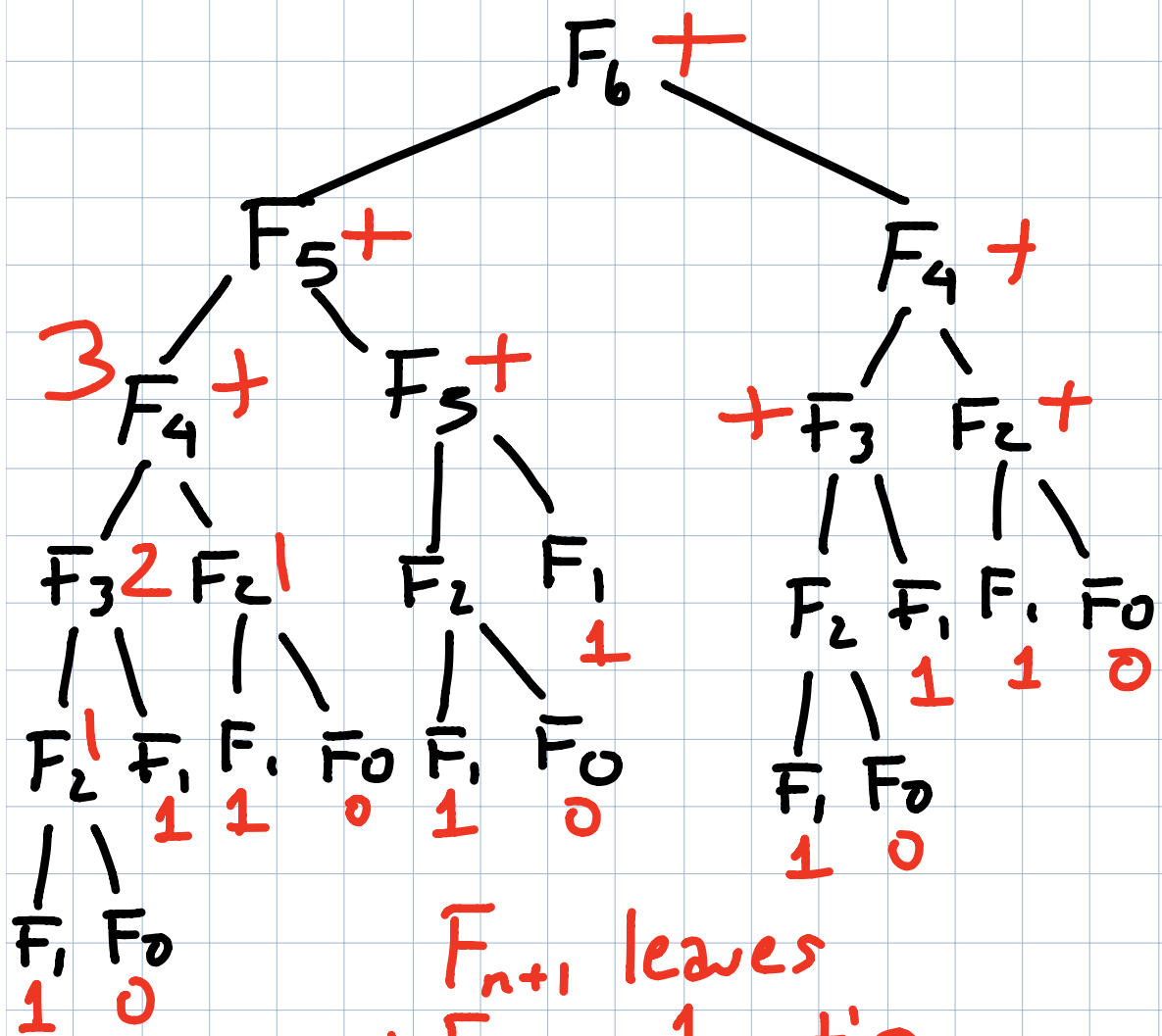
RECFIBO(n):

if ( $n < 2$ )

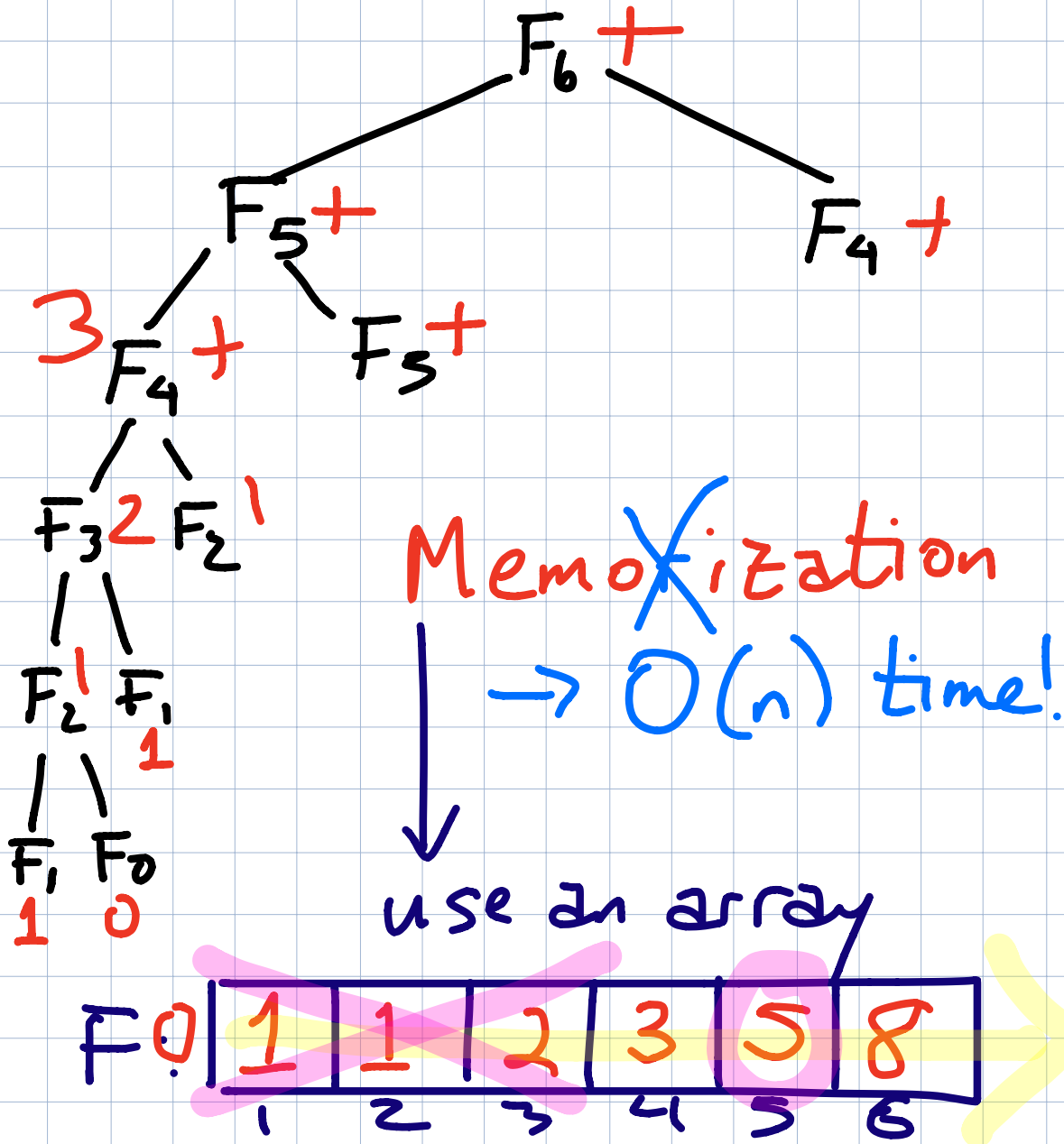
return  $n$

else

return RECFIBO( $n - 1$ ) + RECFIBO( $n - 2$ )



$F_{n+1}$  leaves  
 $\rightarrow F_{n+1} - 1$  '+'s  
 $\rightarrow O(F_n)$  time



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MEMFIBO(n):
  if (n < 2)
    return n
  else
    if F[n] is undefined
      F[n] ← MEMFIBO(n - 1) + MEMFIBO(n - 2)
    return F[n]

```

ITERFIBO(n):

$F[0] \leftarrow 0$

$F[1] \leftarrow 1$

for  $i \leftarrow 2$  to  $n$

$F[i] \leftarrow F[i-1] + F[i-2]$

return  $F[n]$

## Dynamic Programming

1. Recurrence
2. Memoize
  - Subproblems
  - Data Structure
3.  $O_n$  Purpose
  - Dependencies
  - Eval. Order
4. Time, Space

## 5. Optimize!!

ITERFIBO2(n):

prev  $\leftarrow$  1

curr  $\leftarrow$  0

for  $i \leftarrow 1$  to  $n$

    next  $\leftarrow$  curr + prev

    prev  $\leftarrow$  curr

    curr  $\leftarrow$  next

return curr

$$\begin{bmatrix} 0 & 1 \\ 1 & 1 \end{bmatrix} \begin{bmatrix} \text{prev} \\ \text{curr} \end{bmatrix} = \begin{bmatrix} \text{curr} \\ \text{next} \end{bmatrix}$$

$$\begin{bmatrix} 0 & 1 \\ 1 & 1 \end{bmatrix}^n \begin{bmatrix} 1 \\ 0 \end{bmatrix} = \begin{bmatrix} F_{n-1} \\ F_n \end{bmatrix}$$

Repeated Squaring  $\rightarrow O(\log n)$  X's

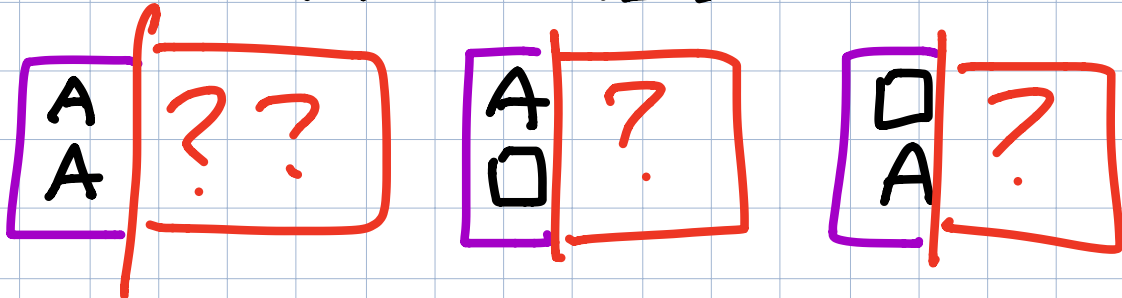
# Edit Distance

FOOD  
MOOD  
MOND  
MONEY  
MONEY

Min #  
insertions  
deletions  
replacements  
to change A to B



ALGORITHMS  
ALTRUISTIC



$$\text{Edit}(A[1..n], B[1..m]) = \min \left\{ [A[i] \neq B[j]] + \text{Edit}(A[2..n], B[1..m]) \right\}$$

Subproblem:

ALGORZI  
ALTRUIST

$$\text{Edit}(A[1..m], B[1..n]) = \min \left\{ \begin{array}{l} \text{Edit}(A[1..m-1], B[1..n]) + 1 \\ \text{Edit}(A[1..m], B[1..n-1]) + 1 \\ \text{Edit}(A[1..m-1], B[1..n-1]) + [A[m] \neq B[n]] \end{array} \right\}$$

$$Edit(i, j) = \begin{cases} i & \text{if } j = 0 \\ j & \text{if } i = 0 \\ \min \left\{ \begin{array}{l} Edit(i-1, j) + 1, \\ Edit(i, j-1) + 1, \\ Edit(i-1, j-1) + [A[i] \neq B[j]] \end{array} \right\} & \text{otherwise} \end{cases}$$



EDITDISTANCE( $A[1..m], B[1..n]$ ):

for  $j \leftarrow 1$  to  $n$

$Edit[0, j] \leftarrow j$

for  $i \leftarrow 1$  to  $m$

$Edit[i, 0] \leftarrow i$

    for  $j \leftarrow 1$  to  $n$

        if  $A[i] = B[j]$

$Edit[i, j] \leftarrow \min \{Edit[i-1, j] + 1, Edit[i, j-1] + 1, Edit[i-1, j-1]\}$

        else

$Edit[i, j] \leftarrow \min \{Edit[i-1, j] + 1, Edit[i, j-1] + 1, Edit[i-1, j-1] + 1\}$

return  $Edit[m, n]$

• ALGORITHM

•  
A  
J  
R  
S  
H  
T  
H  
C

	A	L	G	O	R	I	T	H	M	
	0	→1	→2	→3	→4	→5	→6	→7	→8	→9
A	↓ 1	<b>0</b> →1→2→3→4→5→6→7→8								
L	↓ 2	↓ 1	<b>0</b> →1→2→3→4→5→6→7							
T	↓ 3	↓ 2	↓ 1	↓ 1	↓ 2	↓ 3	↓ 4	↓ <b>4</b>	↓ 5	↓ 6
R	↓ 4	↓ 3	↓ 2	↓ 2	↓ 2	↓ <b>2</b>	↓ 3	↓ 4	↓ 5	↓ 6
U	↓ 5	↓ 4	↓ 3	↓ 3	↓ 3	↓ 3	↓ 3	↓ 3	↓ 4	↓ 5
I	↓ 6	↓ 5	↓ 4	↓ 4	↓ 4	↓ 4	↓ 4	↓ <b>3</b>	↓ 4	↓ 5
S	↓ 7	↓ 6	↓ 5	↓ 5	↓ 5	↓ 5	↓ 5	↓ 4	↓ 4	↓ 5
T	↓ 8	↓ 7	↓ 6	↓ 6	↓ 6	↓ 6	↓ 6	↓ 5	↓ <b>4</b>	↓ 5
I	↓ 9	↓ 8	↓ 7	↓ 7	↓ 7	↓ 7	↓ 7	↓ <b>6</b>	↓ 5	↓ 5
C	↓ 10	↓ 9	↓ 8	↓ 8	↓ 8	↓ 8	↓ 8	↓ 7	↓ 6	↓ 6