

Randomized Algorithms: QuickSort and QuickSelect

Lecture 15
March 12, 2015

Part I

Slick analysis of QuickSort

A Slick Analysis of QuickSort

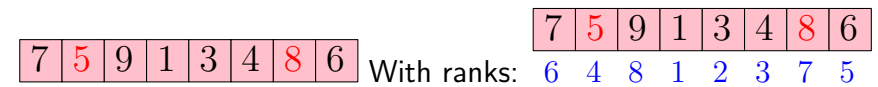
- 1 Let $Q(A)$ be number of comparisons done on input array A :
 - 1 R_{ij} : event that rank i element is compared with rank j element, for $1 \leq i < j \leq n$.
 - 2 X_{ij} is the indicator random variable for R_{ij} . That is, $X_{ij} = 1$ if rank i is compared with rank j element, otherwise 0 .
- 2 $Q(A) = \sum_{1 \leq i < j \leq n} X_{ij}$.
- 3 By linearity of expectation,

$$\begin{aligned}
 E[Q(A)] &= E\left[\sum_{1 \leq i < j \leq n} X_{ij}\right] = \sum_{1 \leq i < j \leq n} E[X_{ij}] \\
 &= \sum_{1 \leq i < j \leq n} \Pr[R_{ij}].
 \end{aligned}$$

A Slick Analysis of QuickSort

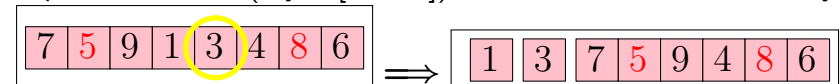
R_{ij} = rank i element is compared with rank j element.

Question: What is $\Pr[R_{ij}]$?



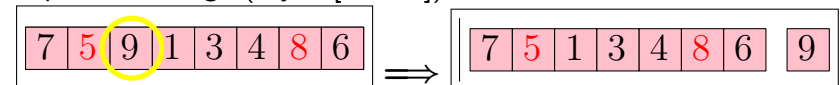
As such, probability of comparing **5** to **8** is $\Pr[R_{4,7}]$.

- 1 If pivot too small (say **3** [rank 2]). Partition and call recursively:



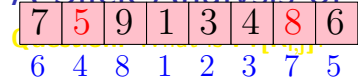
Decision if to compare **5** to **8** is moved to subproblem.

- 2 If pivot too large (say **9** [rank 8]):



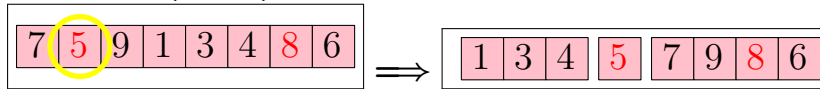
Decision if to compare **5** to **8** moved to subproblem.

A Slick Analysis of QuickSort

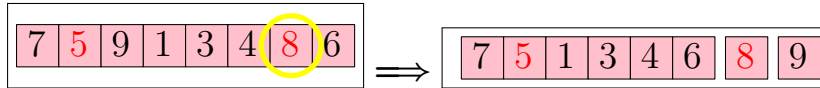


As such, probability of comparing 5 to 8 is $\Pr[R_{4,7}]$.

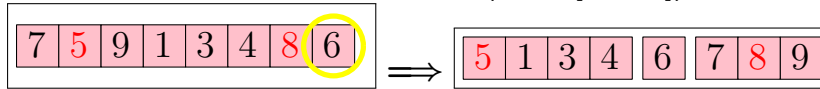
- 1 If pivot is 5 (rank 4). Bingo!



- 2 If pivot is 8 (rank 7). Bingo!



- 3 If pivot in between the two numbers (say 6 [rank 5]):



5 and 8 will never be compared to each other.

A Slick Analysis of QuickSort

Question: What is $\Pr[R_{i,j}]$?

Conclusion:

$R_{i,j}$ happens \iff :

i th or j th ranked element is the first pivot out of the elements of rank

$$i, i + 1, i + 2, \dots, j$$

How to analyze this? Thinking acrobatics!

- 1 Assign every element in array random priority (say in $[0, 1]$).
- 2 Choose pivot to be element with lowest priority in subproblem.
- 3 Equivalent to picking pivot uniformly at random (as QuickSort do).

A Slick Analysis of QuickSort

Question: What is $\Pr[R_{i,j}]$?

- 1 Choosing a pivot using priorities

- 1 Assign every element in array is a random priority (in $[0, 1]$).
- 2 pivot = the element with lowest priority in subproblem.

- 2 $\implies R_{i,j}$ happens if either i or j have lowest priority out of elements in rank $i \dots j$,

- 3 There are $k = j - i + 1$ relevant elements.

- 4 $\Pr[R_{i,j}] = \frac{2}{k} = \frac{2}{j - i + 1}$.

A Slick Analysis of QuickSort

Question: What is $\Pr[R_{ij}]$?

Lemma

$$\Pr[R_{ij}] = \frac{2}{j - i + 1}$$

Proof

- 1 $a_1, \dots, a_i, \dots, a_j, \dots, a_n$: elements of A in sorted order. Let $S = \{a_i, a_{i+1}, \dots, a_j\}$
- 2 **Observation:** If pivot is chosen outside S then all of S either in left or right recursive subproblem.
- 3 **Observation:** a_i and a_j separated when a pivot is chosen from S for the first time. Once separated never to meet again. $\implies a_i$ and a_j will not be compared.

A Slick Analysis of QuickSort

Continued...

Lemma

$$\Pr[R_{ij}] = \frac{2}{j-i+1}.$$

Proof.

- 1 Let $a_1, \dots, a_i, \dots, a_j, \dots, a_n$ be sort of A .
- 2 Let $S = \{a_i, a_{i+1}, \dots, a_j\}$
- 3 **Observation:** a_i is compared with $a_j \iff$ either a_i or a_j is chosen as a pivot from S at separation.
- 4 **Observation:** Given: Pivot chosen from S .
The probability that it is a_i or a_j is exactly $2/|S| = 2/(j-i+1)$ since the pivot is chosen uniformly at random from the array.

□

A Slick Analysis of QuickSort

Continued...

$$E[Q(A)] = \sum_{1 \leq i < j \leq n} E[X_{ij}] = \sum_{1 \leq i < j \leq n} \Pr[R_{ij}].$$

Lemma

$$\Pr[R_{ij}] = \frac{2}{j-i+1}.$$

$$\begin{aligned} E[Q(A)] &= \sum_{1 \leq i < j \leq n} \Pr[R_{ij}] = \sum_{1 \leq i < j \leq n} \frac{2}{j-i+1} \\ &= \sum_{i=1}^{n-1} \sum_{j=i+1}^n \frac{2}{j-i+1} \\ &\leq 2 \sum_{i=1}^{n-1} (H_{n-i+1} - 1) \leq 2 \sum_{1 \leq i < n} H_n \end{aligned}$$

Part II

Quick sort with high probability

Yet another analysis of QuickSort

You should never trust a man who has only one way to spell a word

- 1 Consider element e in the array.
- 2 S_1, S_2, \dots, S_k : subproblems e participates in during **QuickSort** execution:
- 3 **Definition**
 e is lucky in the j th iteration if $|S_j| \leq (3/4) |S_{j-1}|$.
- 4 **Key observation:** The event that e is lucky in j th iteration...
- 5 ... is independent of the event that e is lucky in k th iteration, (if $j \neq k$)
- 6 $X_j = 1 \iff e$ is lucky in the j th iteration.

Yet another analysis of QuickSort

Continued...

Claim

$$\Pr[X_j = 1] = 1/2.$$

Proof.

- 1 X_j determined by j recursive subproblem.
- 2 Subproblem has $n_{j-1} = |X_{j-1}|$ elements.
- 3 j th pivot rank $\in [n_{j-1}/4, (3/4)n_{j-1}] \implies e$ lucky in j th iter.
- 4 Prob. e is lucky $\geq |[n_{j-1}/4, (3/4)n_{j-1}]| / n_{j-1} = 1/2. \quad \square$

Observation

If $X_1 + X_2 + \dots + X_k = \lceil \log_{4/3} n \rceil$ then e subproblem is of size one. Done!

Yet another analysis of QuickSort

Continued...

Observation

Probability e participates in $\geq k = 40 \lceil \log_{4/3} n \rceil$ subproblems. Is equal to

$$\begin{aligned} \Pr[X_1 + X_2 + \dots + X_k \leq \lceil \log_{4/3} n \rceil] \\ \leq \Pr[X_1 + X_2 + \dots + X_k \leq k/4] \\ \leq 2 \cdot 0.68^{k/4} \leq 1/n^5. \end{aligned}$$

Conclusion

QuickSort takes $O(n \log n)$ time with high probability.

Randomized Quick Selection

Input Unsorted array A of n integers

Goal Find the j th smallest number in A (*rank j number*)

Randomized Quick Selection

- 1 Pick a pivot element *uniformly at random* from the array
- 2 Split array into 3 subarrays: those smaller than pivot, those larger than pivot, and the pivot itself.
- 3 Return pivot if rank of pivot is j .
- 4 Otherwise recurse on one of the arrays depending on j and their sizes.

Algorithm for Randomized Selection

Assume for simplicity that A has distinct elements.

QuickSelect(A, j):

```
Pick pivot  $x$  uniformly at random from  $A$ 
Partition  $A$  into  $A_{\text{less}}, x$ , and  $A_{\text{greater}}$  using  $x$  as pivot
if ( $|A_{\text{less}}| = j - 1$ ) then
    return  $x$ 
if ( $|A_{\text{less}}| \geq j$ ) then
    return QuickSelect( $A_{\text{less}}, j$ )
else
    return QuickSelect( $A_{\text{greater}}, j - |A_{\text{less}}| - 1$ )
```

QuickSelect analysis

- 1 S_1, S_2, \dots, S_k be the subproblems considered by the algorithm. Here $|S_1| = n$.
- 2 S_i would be **successful** if $|S_i| \leq (3/4)|S_{i-1}|$
- 3 Y_1 = number of recursive calls till first successful iteration. Clearly, total work till this happens is $O(Y_1 n)$.
- 4 n_i = size of the subproblem immediately after the $(i - 1)$ th successful iteration.
- 5 Y_i = number of recursive calls after the $(i - 1)$ th successful call, till the i th successful iteration.
- 6 Running time is $O(\sum_i n_i Y_i)$.

QuickSelect analysis

Example

S_i = subarray used in i th recursive call

$|S_i|$ = size of this subarray

Red indicates successful iteration.

Inst'	S_1	S_2	S_3	S_4	S_5	S_6	S_7	S_8	S_9
$ S_i $	100	70	60	50	40	30	25	5	2
Succ'	$Y_1 = 2$		$Y_2 = 4$			$Y_3 = 2$		$Y_4 = 1$	
$n_i =$	$n_1 = 100$		$n_2 = 60$			$n_3 = 25$		$n_4 = 2$	

- 1 All the subproblems after $(i - 1)$ th successful iteration till i th successful iteration have size $\leq n_i$.
- 2 Total work: $O(\sum_i n_i Y_i)$.

QuickSelect analysis

- 1 Total work: $O(\sum_i n_i Y_i)$.
- 2 $n_i \leq (3/4)n_{i-1} \leq (3/4)^{i-1}n$.
- 3 Y_i is a random variable with geometric distribution
Probability of $Y_i = k$ is $1/2^k$.
- 4 $E[Y_i] = 2$.
- 5 As such, expected work is proportional to

$$\begin{aligned}
 E\left[\sum_i n_i Y_i\right] &= \sum_i E[n_i Y_i] \leq \sum_i E[(3/4)^{i-1} n Y_i] \\
 &= n \sum_i (3/4)^{i-1} E[Y_i] = n \sum_{i=1}^{\infty} (3/4)^{i-1} 2 \leq 8n.
 \end{aligned}$$

QuickSelect analysis

Theorem

The expected running time of **QuickSelect** is $O(n)$.

QuickSelect analysis via recurrence

Analysis via Recurrence

- Given array A of size n let $Q(A)$ be number of comparisons of randomized selection on A for selecting rank j element.
- Note that $Q(A)$ is a random variable
- Let A_{less}^i and A_{greater}^i be the left and right arrays obtained if pivot is rank i element of A .
- Algorithm recurses on A_{less}^i if $j < i$ and recurses on A_{greater}^i if $j > i$ and terminates if $j = i$.

$$Q(A) = n + \sum_{i=1}^{j-1} \Pr[\text{pivot has rank } i] Q(A_{\text{greater}}^i) + \sum_{i=j+1}^n \Pr[\text{pivot has rank } i] Q(A_{\text{less}}^i)$$

Analyzing the Recurrence

As in **QuickSort** we obtain the following recurrence where $T(n)$ is the worst-case expected time.

$$T(n) \leq n + \frac{1}{n} \left(\sum_{i=1}^{j-1} T(n-i) + \sum_{i=j}^n T(i-1) \right).$$

Theorem

$$T(n) = O(n).$$

Proof.

(Guess and) Verify by induction (see next slide). □

Analyzing the recurrence

Theorem

$$T(n) = O(n).$$

Prove by induction that $T(n) \leq \alpha n$ for some constant $\alpha \geq 1$ to be fixed later.

Base case: $n = 1$, we have $T(1) = 0$ since no comparisons needed and hence $T(1) \leq \alpha$.

Induction step: Assume $T(k) \leq \alpha k$ for $1 \leq k < n$ and prove it for $T(n)$. We have by the recurrence:

$$\begin{aligned} T(n) &\leq n + \frac{1}{n} \left(\sum_{i=1}^{j-1} T(n-i) + \sum_{i=j}^n T(i-1) \right) \\ &\leq n + \frac{\alpha}{n} \left(\sum_{i=1}^{j-1} (n-i) + \sum_{i=j}^n (i-1) \right) \quad \text{by applying induction} \end{aligned}$$

Analyzing the recurrence

$$\begin{aligned} T(n) &\leq n + \frac{\alpha}{n} \left(\sum_{i=1}^{j-1} (n-i) + \sum_{i=j}^n (i-1) \right) \\ &\leq n + \frac{\alpha}{n} \left((j-1)(2n-j)/2 + (n-j+1)(n+j-2)/2 \right) \\ &\leq n + \frac{\alpha}{2n} (n^2 + 2nj - 2j^2 - 3n + 4j - 2) \\ &\quad \text{above expression maximized when } j = (n+1)/2: \text{ calculus} \\ &\leq n + \frac{\alpha}{2n} (3n^2/2 - n) \quad \text{substituting } (n+1)/2 \text{ for } j \\ &\leq n + 3\alpha n/4 \\ &\leq \alpha n \quad \text{for any constant } \alpha \geq 4 \end{aligned}$$

Comments on analyzing the recurrence

- 1 Algebra looks messy but intuition suggest that the median is the hardest case and hence can plug $j = n/2$ to simplify without calculus
- 2 Analyzing recurrences comes with practice and after a while one can see things more intuitively

John Von Neumann:

Young man, in mathematics you don't understand things. You just get used to them.