

Admin: Hw 1 - due 9pm

Hw 2 - out today, due next Tue

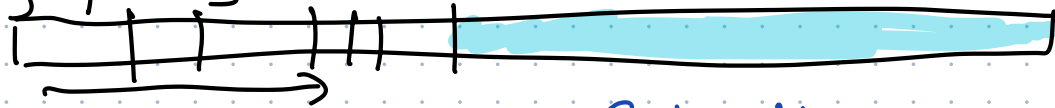
Sources and collaborators

Dynamic Programming

What are the subproblems?

inputs to

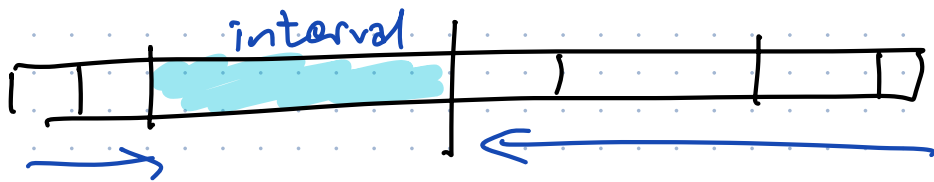
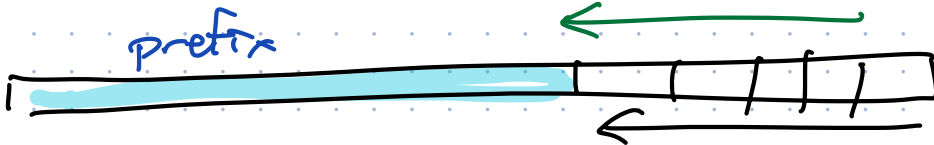
string splitting



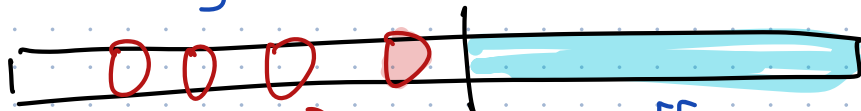
Decisions: left to right word bdy's

Subproblem: suffix

memoization



Longest Increasing Subsequence:



decisions

suffix + one element

memo

Longest Common <sup>Increasing</sup> Subsequence

A [ALXANDER | ZHAM | ILTON]

B [ALMOND | ALLERGY]

two suffixes + fence post

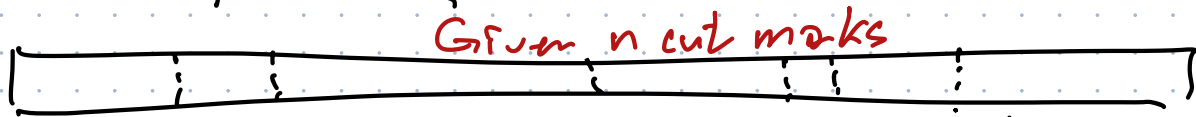
Longest Palindrome Subsequence



MADAMIMADAM

P -> e | a | o P a

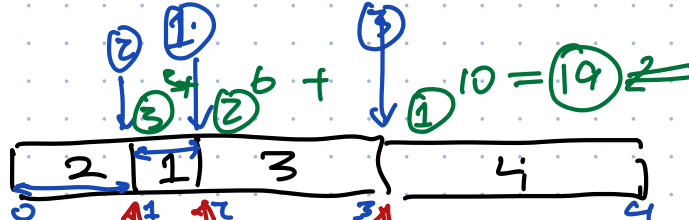
# Woodcutter's problem



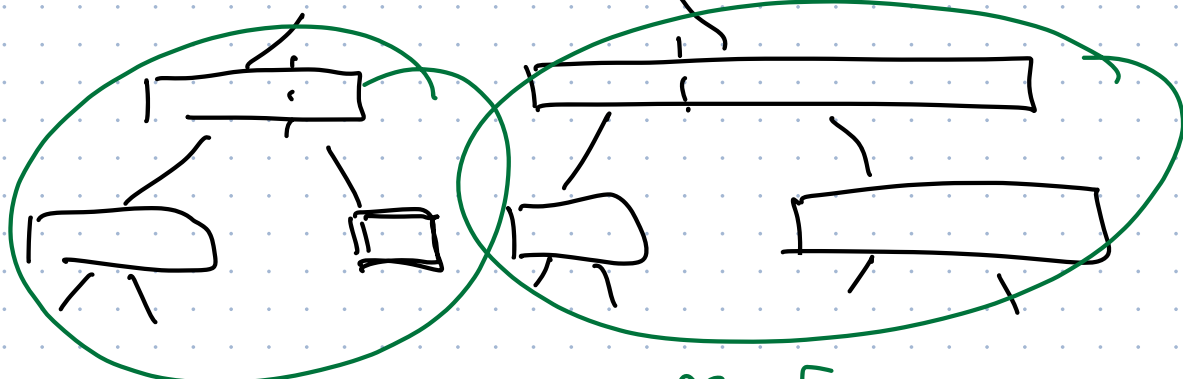
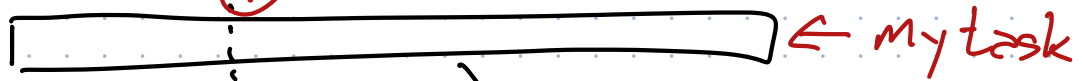
Sawmill can cut any  $L$ -foot board once for  $\$L$   
 Find order of cuts with min total cost

$$3 + 10 + 7 = 20$$

$$L = [2, 1, 3, 4]$$



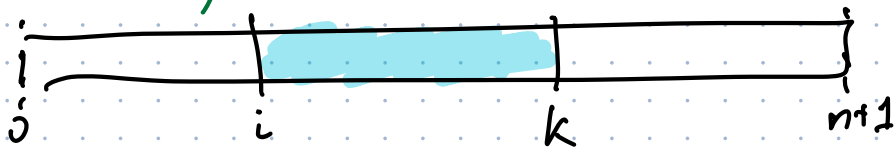
$$10 + 8 + 7 = 25$$



rec fairy

rec fairy

subproblems = intervals

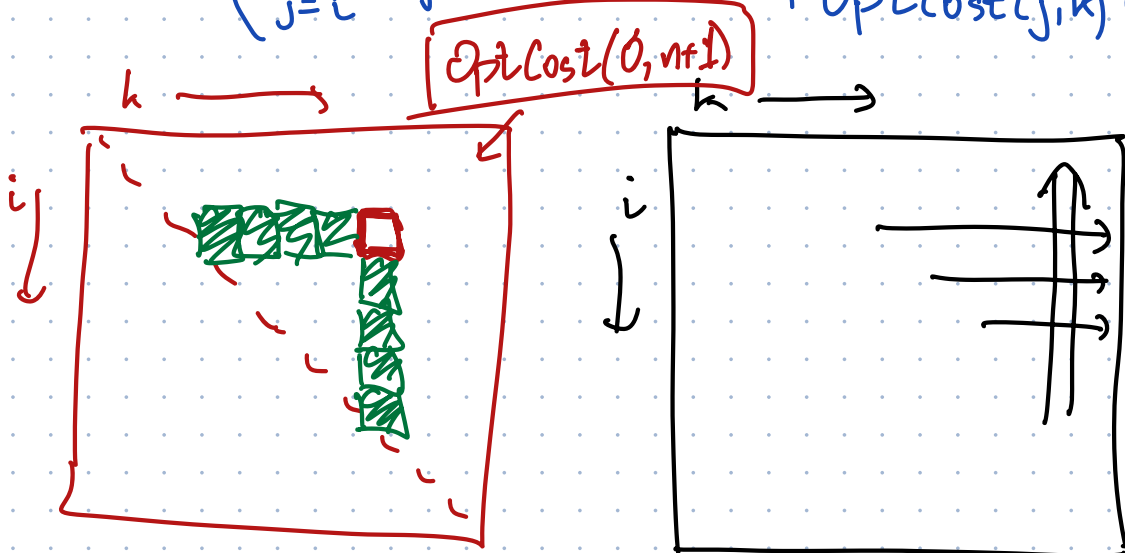


$OptCost(i, k) = \text{min. cost to cut up segment from cut } i \text{ to cut } k.$

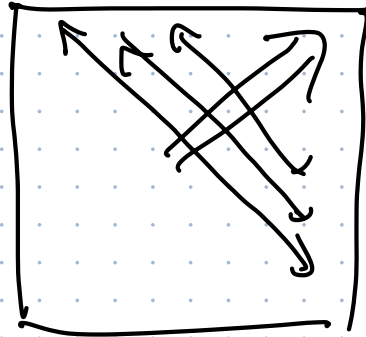
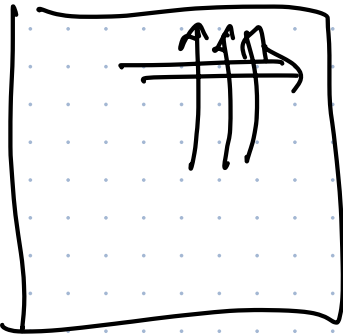
Input:  $L[0..n]$  = array of board lengths

$L[i]$  = length between cut  $i$  and cut  $i+1$

$$\text{OptCost}(i, k) = \begin{cases} 0 & \text{if } k = i+1 \\ \sum_{j=i}^{k-1} L[i, j] + \min \{ \text{OptCost}(i, j) + \text{OptCost}(j, k) \} & \{ i < j < k \} \end{cases}$$



$O(n^2)$  subprobs  $\times$   $O(n)$  time each =  $O(n^3)$  time



We can do better!  $O(n^2)$  time Knuth-Yao  
 $O(n \log n)$