

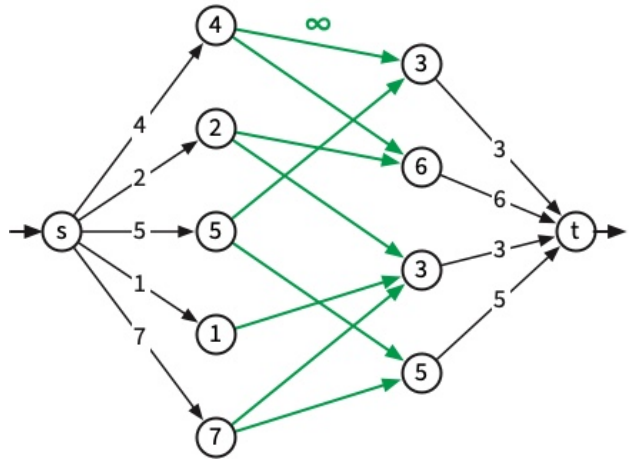
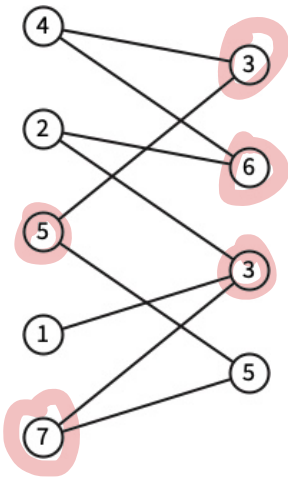
Next Tuesday — university holiday — GD VOTE!

HW 9 will be last graded HW

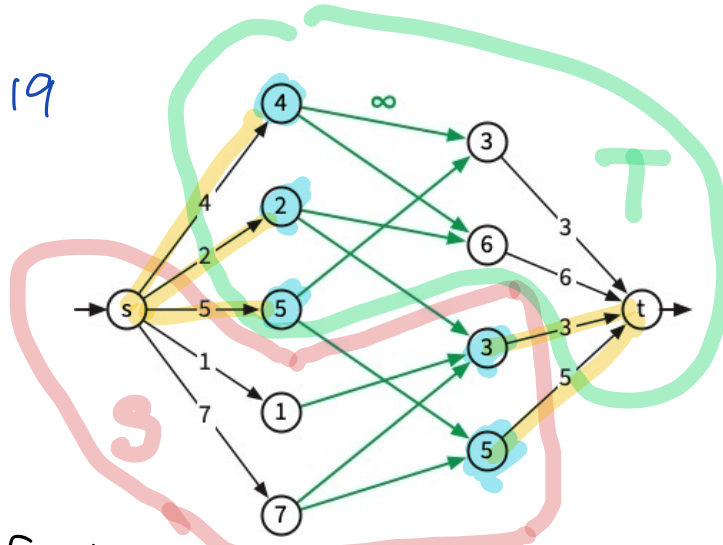
Bipartite Min Vertex Cover

Vertex cover = subset of vertices that touch every edge

NP-hard for non-bipartite graphs



$$||S, T|| = 4 + 2 + 5 + 3 + 5 = 19$$



Cuts with finite capacity
 $||S, T||$

Vertex covers

$$f(C) = \sum_{v \in C} f(v)$$

$L \cap T \rightarrow R \cap T$ $C = (L \cap T) \cup (R \cap S)$
 $L \cap S \rightarrow R \cap S$

$S = (L \cap C) \cup (R \cap C) \cup \{s\}$
 $T = V \setminus S$

Compute min cut $\|S, T\|$ in $[O(E)]$ time [Orlin]

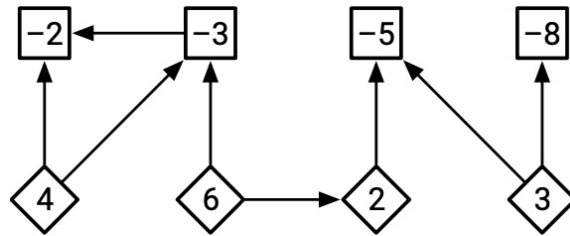
Project selection

Input is DAG
value $\phi(v)$ for every vertex v

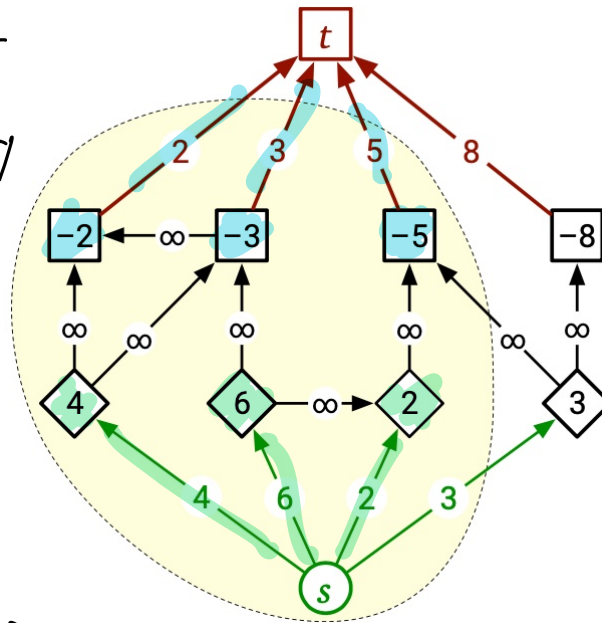
$V =$ projects
 $E =$ dependencies

$u \rightarrow v$
 u can only be done after v .

Output: $S \subseteq V$ max $\phi(S)$



Any (s, t) -cut S, T
with finite capacity
 \Downarrow
Valid selection
 $S \setminus \{s\}$



Claim: $\phi(S \setminus s) = P - \|S, T\|$ where $P = \sum_{\phi(v) > 0} \phi(v)$
 $P = \phi(S \setminus s) + \|S, T\|$

For any $X \subseteq V$: $\text{cost}(X) = \sum_{\substack{\phi(v) < 0 \\ v \in X}} -\phi(v) = \sum_{v \in X} c(v \rightarrow t)$

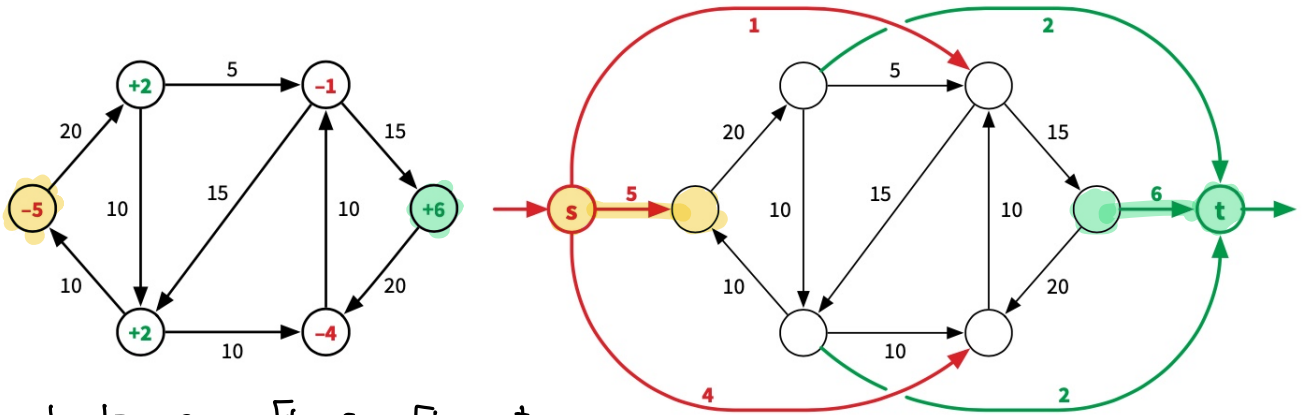
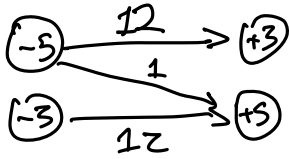
$\text{income}(X) = \sum_{\substack{v \in X \\ \phi(v) > 0}} \phi(v) = \sum_{v \in X} c(s \rightarrow v)$

$$\text{profit}(x) = \text{income}(x) - \text{cost}(x) = \sum_{v \in X} \Phi(v)$$

$$P = \text{income}(V) = \text{income}(S) + \text{income}(T)$$

$$\|S, T\| = \text{cost}(S) + \text{income}(T)$$

$$P - \|S, T\| = \text{income}(S) - \text{cost}(S)$$



$b(v)$ - balance = Flow in - Flow out

Is there a feasible flow?

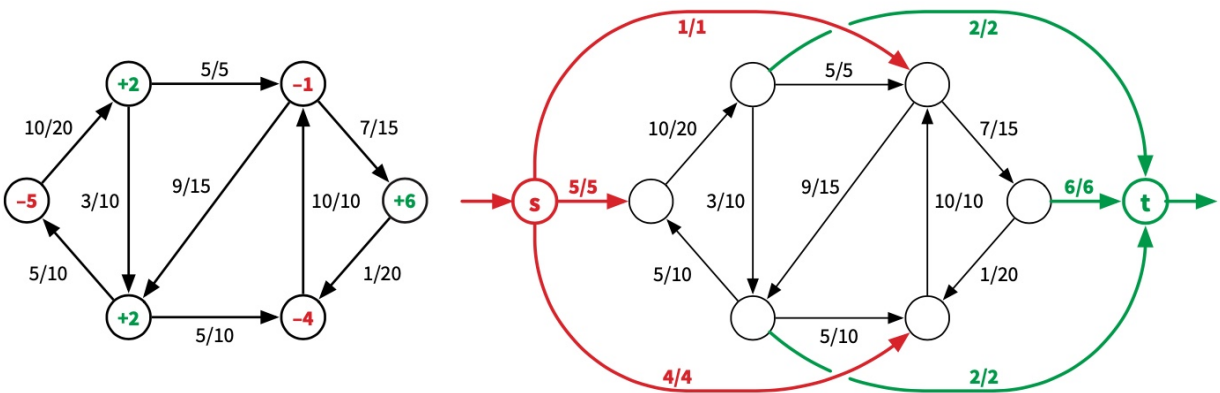
Necessary: $\sum_v b(v) = 0$

Feasible flow F in G

$\downarrow \uparrow$

Feasible flow F' in G'
with value

$\sum_{b(w) > 0} b(w)$ - saturates all edges from s



Maximum flow in network with nonzero balances

① Feasible flow $F \rightarrow$ max flow F' in G'

② Maximize it \rightarrow max flow F'' in G_F
return $F' + F''$

