

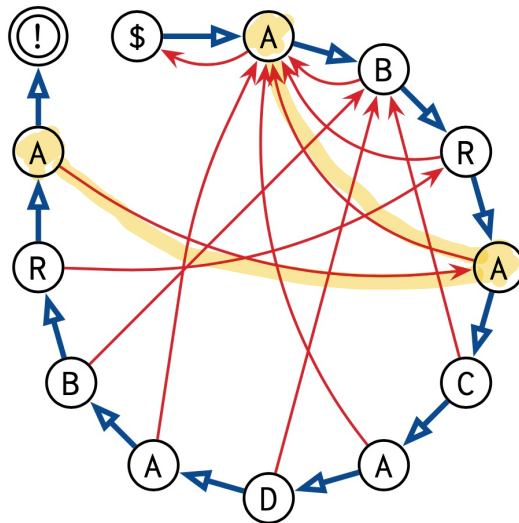
Once we've matched  $T[i]$  to  $P[j]$

we don't need to explicitly compare  $T[i]$  to anything else

The next reasonable shift is the smallest value of  $s$  such that  $T[s..i-1]$  is a shorter prefix of  $P$ .

A B R A Y

$P[i]$	A	B	R	A	C	A	D	A	B	R	A
$fail[i]$	0	1	1	1	2	1	2	1	2	3	4



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KNUTHMORRISPRATT( $T[1..n], P[1..m]$ ):
  fail[1..m] ← COMPUTEFAILURE( $P$ )
  j ← 1
  for i ← 1 to n
    while j > 0 and  $T[i] \neq P[j]$ 
      j ← fail[j]
    if j = m    <<Found it!>>
      return i - m + 1
    j ← j + 1
  return NONE

```

ABRACADABRA  
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ABRABRACADABRA  
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Running time =  $O(n)$

$\leq 2n$  comparisons

Two types of comparisons:

$T[i] = P[j] \rightarrow i++, j++ \leq n$  times

$T[i] \neq P[j] \rightarrow i \text{ same, } j \downarrow \leq n$  times

$P[1..fail[j]-1]$  is the longest proper prefix of  $P[1..j-1]$  that is also a suffix of  $T[1..i-1]$ .

$P[1..fail[j]-1]$  is the longest proper prefix of  $P[1..j-1]$  that is also a suffix of  $P[1..j-1]$ .

border of string is proper prefix that is also suffix

$P[i]$	A	B	R	A	C	A	D	A	B	R	A
fail[i]	0	1	1	1	2	1	2	<del>2</del>	<del>2</del>	<del>3</del>	<del>4</del>

$O(n^3)$  time brute force  $\uparrow$   $O$   
 $O(n^2)$  time DP

KNUTHMORRISPRATT( $T[1..n], P[1..m]$ ):

```

fail[1..m] ← COMPUTEFAILURE(P)
j ← 1
for i ← 1 to n
  while j > 0 and T[i] ≠ P[j]
    j ← fail[j]
  if j = m    <<Found it!>>
    return i - m + 1
  j ← j + 1
return NONE

```

COMPUTEFAILURE( $P[1..m]$ ):

```

j ← 0
for i ← 1 to m
  fail[i] ← j (*)
  while j > 0 and P[i] ≠ P[j]
    j ← fail[j]
  j ← j + 1
return fail[1..m]

```

$fail[j] = P[i]$   
 $fail[i] = fail[j]$   
 else  
 $fail[i] = j$

ABACABADABACABA | BABACABADABACABA

A border of P is either  
 - the longest border of P or  
 - a border of a border of P

ABAABABABAABC

AAAAAAAAAAAAAAAA

$$fail[i] = \begin{cases} 0 & \text{if } i = 0, \\ \max_{c \geq 1} \{ fail^c[i-1] + 1 \mid P[i-1] = P[fail^c[i-1]] \} & \text{otherwise.} \end{cases}$$

KMP makes  $O(n)$  comparisons (given fail)

not  $O(1)$  at each position in T

$O(n)$  at each value of i

Optimized fail:  $O(\log m)$  failures at each i.

$j \leftarrow 0, i \leftarrow 1$ $fail[i] \leftarrow j$	$\$^j$ $A^i$ B R A C A D A B R X ... $0$
$j \leftarrow j + 1, i \leftarrow i + 1$ $fail[i] \leftarrow j$	$\$$ $A^j$ $B^i$ R A C A D A B R X ... $0$ $1$
$j \leftarrow fail[j]$	$\$^j$ A $B^i$ R A C A D A B R X ...
$j \leftarrow j + 1, i \leftarrow i + 1$ $fail[i] \leftarrow j$	$\$$ $A^j$ B $R^i$ A C A D A B R X ... $0$ $1$ $1$
$j \leftarrow fail[j]$	$\$^j$ A B $R^i$ A C A D A B R X ...
$j \leftarrow j + 1, i \leftarrow i + 1$ $fail[i] \leftarrow j$	$\$$ $A^j$ B R $A^i$ C A D A B R X ... $0$ $1$ $1$ $1$
$j \leftarrow j + 1, i \leftarrow i + 1$ $fail[i] \leftarrow j$	$\$$ A $B^j$ R A $C^i$ A D A B R X ... $0$ $1$ $1$ $1$ $2$
$j \leftarrow fail[j]$	$\$$ $A^j$ B R A $C^i$ A D A B R X ...
$j \leftarrow fail[j]$	$\$^j$ A B R A $C^i$ A D A B R X ...
$j \leftarrow j + 1, i \leftarrow i + 1$ $fail[i] \leftarrow j$	$\$$ $A^j$ B R A C $A^i$ D A B R X ... $0$ $1$ $1$ $1$ $2$ $1$ $1$
$j \leftarrow j + 1, i \leftarrow i + 1$ $fail[i] \leftarrow j$	$\$$ A $B^j$ R A C A $D^i$ A B R X ... $0$ $1$ $1$ $1$ $2$ $1$ $2$
$j \leftarrow fail[j]$	$\$$ $A^j$ B R A C A $D^i$ A B R X ...
$j \leftarrow fail[j]$	$\$^j$ A B R A C A $D^i$ A B R X ...
$j \leftarrow j + 1, i \leftarrow i + 1$ $fail[i] \leftarrow j$	$\$$ $A^j$ B R A C A D $A^i$ B R X ... $0$ $1$ $1$ $1$ $2$ $1$ $2$ $1$
$j \leftarrow j + 1, i \leftarrow i + 1$ $fail[i] \leftarrow j$	$\$$ A $B^j$ R A C A D A $B^i$ R X ... $0$ $1$ $1$ $1$ $2$ $1$ $2$ $1$ $2$
$j \leftarrow fail[j]$	$\$$ A B R A C A D A B $R^i$ X ...
$j \leftarrow fail[j]$	$\$^j$ A B R A C A D A B $R^i$ X ...
$j \leftarrow j + 1, i \leftarrow i + 1$ $fail[i] \leftarrow j$	$\$$ A B R $A^j$ C A D A B R $X^i$ ... $0$ $1$ $1$ $1$ $2$ $1$ $2$ $1$ $2$ $3$ $4$
$j \leftarrow fail[j]$	$\$$ $A^j$ B R A C A D A B R $X^i$ ...
$j \leftarrow fail[j]$	$\$^j$ A B R A C A D A B R $X^i$ ...