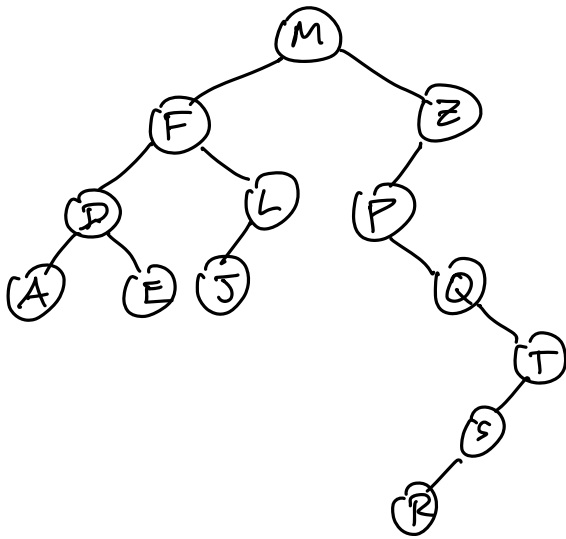


# Binary Search Trees



Balanced:  $O(\log n)$  search time

Worst case:  $O(n)$

Find(x) ← Pred(x)  
                    ← Succ(x)

Insert(x)  
Delete(x)

$T_L, T_R \leftarrow \text{Split}(x)$

Join( $T_L, T_R$ )

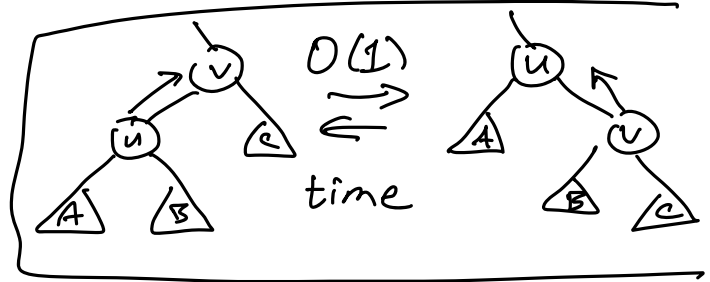
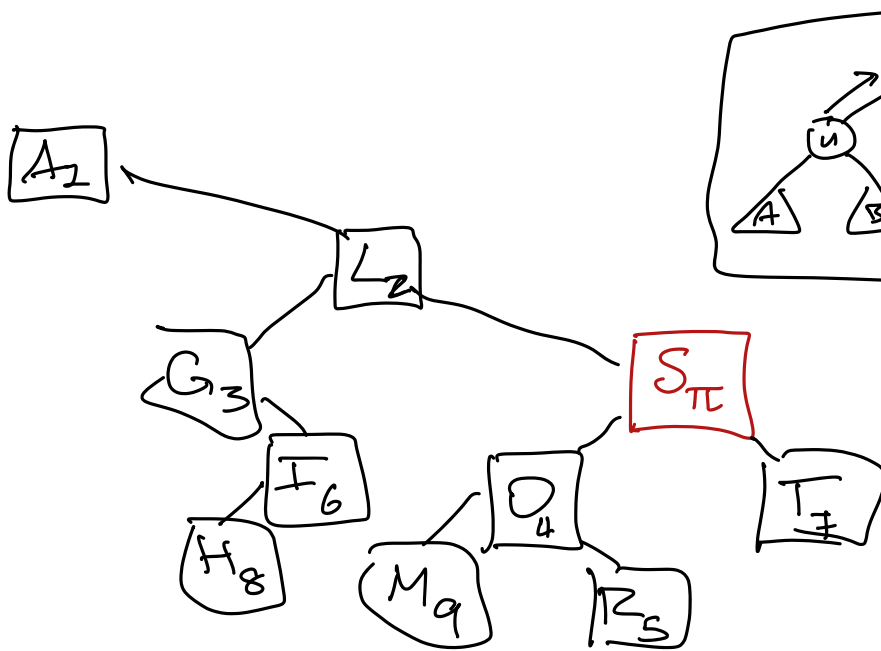
ideally:  
 $O(\log n)$   
time

Treap — Aragon + Seidel '90

$O(\log n)$  exp. time

Every node has a search key and a priority

← Binary search tree  
← min-heap



Insert( $k, p$ )

New node( $k, p$ )

Insert looking only at keys

Bubble up new node looking only at priorities

All operations:

time =  $O(\text{depth of some node})$

Random priorities  $\Rightarrow$

For any node  $v$

$E(\text{depth}(v)) = O(\log n)$

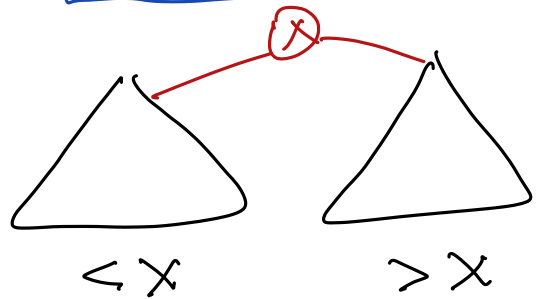
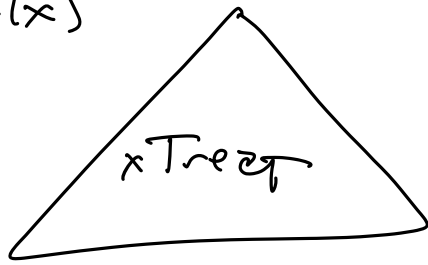
Delete:

Set priority  $\leftarrow \infty$   
Rotate down (promoting child with smaller priority)  
Chop off leaf

Stronger:

$E[\text{max depth}(v)] = O(\log n)$

Split(x)



set priority(x) ← -∞  
bubble up to root

WLOG: keys are integers 1..n

$$E[\text{depth}(k)] = E[\# \text{proper ancestors of } k]$$

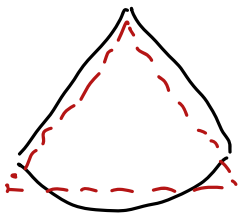
$$= \sum_{i=1}^n \Pr[i \uparrow k]$$

$i \uparrow k = i$  is proper ancestor of  $k$

Lemma: For all  $i < k$ ;  
 $i \uparrow k \iff \text{priority}(i) = \min\{\text{priority}(j) \mid i \leq j \leq k\}$

Lemma  $\Rightarrow$   $\Pr[1 \uparrow n] = \frac{1}{n}$   
 $\Pr[i \uparrow k] = \frac{1}{k-i+1}$

$$E[\text{depth}(k)] = \sum_{i=1}^k \Pr[i \uparrow k]$$



$$= \sum_{i=1}^{k-1} \Pr[i \uparrow k] + \sum_{i=k+1}^n \Pr[i \uparrow k]$$

$$= \underbrace{\sum_{i=1}^{k-1} \frac{1}{k-i+1}}_{j=k-i+1} + \underbrace{\sum_{i=k+1}^n \frac{1}{i-k+1}}_{j=i-k+1}$$

$$= \sum_{j=2}^k \frac{1}{j}$$

$$+ \sum_{j=2}^{n-k+1} \frac{1}{j}$$

$$= H_k - 1 + H_{n-k+1} - 1$$

$$< 2 \ln n - 2$$

Lemma: For all  $i < k$ ,

$$i \uparrow k \iff \text{priority}(i) = \min \{ \text{priority}(j) \mid i \leq j \leq k \}$$

Proof:

Five cases:

•  $\text{root} < i \rightarrow$  by IH

•  $\text{root} = i \rightarrow i \uparrow k$        $\text{priority}(i)$  is min

•  $i < \text{root} < k \rightarrow i \not\uparrow k$        $\text{priority}(i)$  is not min

•  $\text{root} = k$

•  $\text{root} > k \rightarrow$  by IH

□

Treap = BST defined by random priorities

Treap = BST obtained by inserting keys in random order

Treap = recursion tree for randomized quicksort

RANDOM QUICKSORT (A):

randomly permute A

for each  $x$  in A  
 Insert  $x$  into BST

inorder (BST)