Compression, Information and Entropy – Huffman's coding

Lecture 25
December 1, 2015

25.1: Huffman coding

- \bullet Σ : alphabet.
- e binary code: assigns a string of 0s and 1s to each character in the alphabet.
- \odot each symbol in input = a codeword over some other alphabet.
- ullet Useful for transmitting messages over a wire: only 0/1
- receiver gets a binary stream of bits...
- o ... decode the message sent.
- prefix code: reading a prefix of the input binary string uniquely match it to a code word.
- Ontinuing to decipher the rest of the stream.
- binary/prefix code is prefix-free if no code is a prefix of any other.
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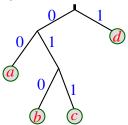
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- Morse code is binary+prefix code but not prefix-free.
- \bigcirc ... code for S (\cdots) includes the code for E (\cdot) as a prefix.
- Prefix codes are binary trees...
- ...characters in leafs, code word is path from root
- o prefix treestree!prefix tree or code trees.
- Open Decoding / Encoding is easy.

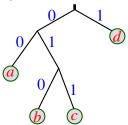
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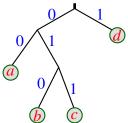
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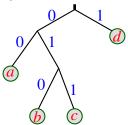
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- Encoding: given frequency table: $f[1 \dots n]$.
- ode(i): binary string for ith character. len(s): length (in bits) of binary string s.
- Compute tree T that minimizes

$$cost(\mathfrak{I}) = \sum_{i=1}^{n} f[i] * len(code(i)), \tag{1}$$

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Frequency table for...

"A tale of two cities" by Dickens

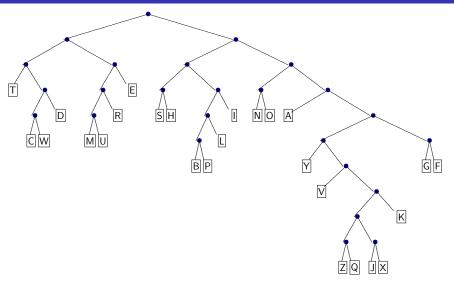
\ n	16,492	'1'	61	'C'	13,896	'Q'	667
, ,	130,376	'2'	10	'D'	28,041	'R'	37,187
'!'	955	'3'	12	'Ε'	74,809	'S'	37,575
411 1	5,681	'4'	10	'F'	13,559	'T'	54,024
'\$'	2	'5'	14	'G'	12,530	'U'	16,726
'%'	1	'6'	11	'H'	38,961	'V'	5,199
477	1,174	'7'	13	'1'	41,005	'W'	14,113
'('	151	'8'	13	'J'	710	'X'	724
()'	151	'9'	14	'K'	4,782	'Y'	12,177
·*¹	70	' :'	267	'L'	22,030	'Z'	215
, ,	13,276	٠.,	1,108	'M'	15,298	-	182
' _'	2,430	'?'	913	'N'	42,380	,,,	93
. ,	6,769	'A'	48,165	'O'	46,499	'@'	2
'0'	20	'B'	8,414	'P'	9,957	'/'	26

Computed prefix codes...

char	frequency	code	char	freq	code
'A'	48165	1110	'N'	42380	1100
'B'	8414	101000	'O'	46499	1101
'C'	13896	00100	'P'	9957	101001
'D'	28041	0011	'Q'	667	1111011001
'E'	74809	011	'R'	37187	0101
'F'	13559	111111	'S'	37575	1000
'G'	12530	111110	'T'	54024	000
'H'	38961	1001	'U'	16726	01001
'1'	41005	1011	'V'	5199	1111010
'J'	710	1111011010	'W'	14113	00101
'K'	4782	11110111	'X'	724	1111011011
L'	22030	10101	'Y'	12177	111100
'M'	15298	01000	ʻZ'	215	1111011000

The Huffman tree generating the code

Build only on A-Z for clarity.



- two trees for some disjoint parts of the alphabet...
- Merge into larger tree by creating a new node and hanging the trees from this common node.



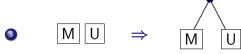


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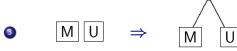


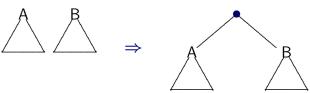
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- 1 take two least frequent characters in frequency table...
- 2 ... merge them into a tree, and put the root of merged tree back into table.
- ...instead of the two old trees.
- Algorithm stops when there is a single tree.
- Intuition: infrequent characters participate in a large number of merges. Long code words.
- Algorithm is due to David Huffman (1952).
- Resulting code is best one can do.
- Huffman coding: building block used by numerous other compression algorithms.

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Building optimal prefix code trees

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- **①** \mathfrak{I} : optimal code tree (prefix free!).
- \bigcirc Then \Im is a full binary tree.
- $oldsymbol{0}$... every node of $oldsymbol{\mathfrak{T}}$ has either $oldsymbol{0}$ or $oldsymbol{2}$ children
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- ② New code tree is better compressor: $cost(\mathfrak{I}) = \sum_{i=1}^{n} f[i] * len(code(i)).$
- $oldsymbol{u}$: leaf u with maximum depth d in $oldsymbol{\mathfrak{T}}$. Consider parent $v=\overline{\mathrm{p}}(u)$.
- \bigcirc \Longrightarrow v: has two children, both leafs

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Infrequent characters are stuck together...

Lemma

x, y: two least frequent characters (breaking ties arbitrarily).

 \exists optimal code tree in which x and y are siblings.

- **①** Claim: \exists optimal code s.t. x and y are siblings + deepest.
- ② \mathfrak{I} : optimal code tree with depth d.
- $exttt{ iny By lemma... } \mathfrak{T}$ has two leafs at depth d that are siblings,
- 1 If not x and y, but some other characters α and β .
- $lacksquare{5}{3}$ \mathfrak{I}' : swap $oldsymbol{x}$ and $oldsymbol{\alpha}$.
- $\mathbf{0} \cos t(\mathfrak{T}') = \cos t(\mathfrak{T}) (f[\alpha] f[x])\Delta.$
- $oldsymbol{x}$: one of the two least frequent characters. ...but $oldsymbol{lpha}$ is not.
- @ Swapping x and α does not increase cost.
- $exttt{0}$ $exttt{0}$: optimal code tree, swapping x and lpha does not decrease cost.
- ① T' is also an optimal code tree
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- Claim: \exists optimal code s.t. x and y are siblings + deepest.
- T: optimal code tree with depth d.
- **3** By lemma... $\mathfrak T$ has two leafs at depth d that are siblings,
- If not x and y, but some other characters α and β .
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Huffman's codes are optimal

Theorem

Huffman codes are optimal prefix-free binary codes.

- If message has 1 or 2 diff characters, then theorem easy.
- ullet Assume f[1] and f[2] are the two smallest.
- Let f[n+1] = f[1] + f[2].
- lacksquare lemma \Longrightarrow \exists opt. code tree $\mathcal{T}_{\mathrm{opt}}$ for f[1..n]
- \bigcirc Remove 1 and 2 from \Im_{opt} .
- $\mathfrak{I}'_{\mathrm{opt}}$: Remaining tree has $3,\ldots,n$ as leafs and "special" character n+1 (i.e., parent 1,2 in $\mathfrak{I}_{\mathrm{opt}}$)

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Average length of codewords...

Special case

Lemma

```
1, \ldots, n: symbols.
```

Assume, for $i = 1, \ldots, n$:

- $oldsymbol{0} p_i = 1/2^{l_i}$: probability for the ith symbol
- $\mathbf{0} \ l_i \geq \mathbf{0}$: integer.

Then, in Huffman coding for this input, the code for i is of length l_i .

- induction of the Huffman algorithm.
- ② n=2: claim holds since there are only two characters with probability 1/2.
- $exttt{3}$ Let i and j be the two characters with lowest probability.
- lacksquare Must be $p_i=p_j$ (otherwise, $\sum_k p_k
 eq 1$).
- ullet Huffman's tree merges this two letters, into a single "character" that have probability $2p_i$.
- New "character" has encoding of length l_i-1 , by induction (on remaining n-1 symbols).
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