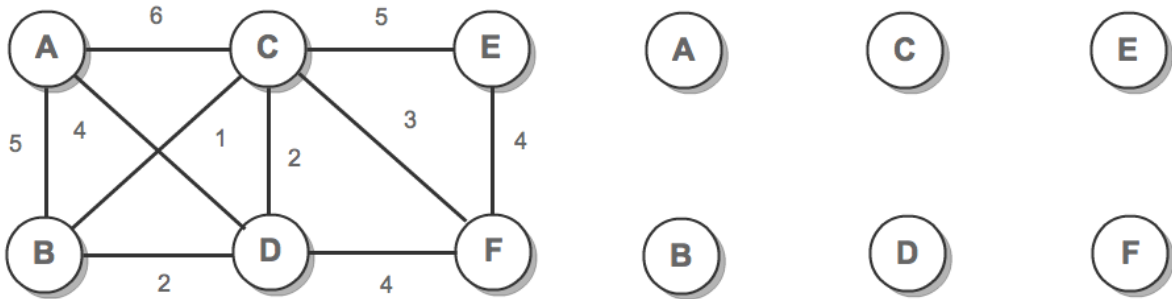


# CS 473: Algorithms, Fall 2010

## HBS 6

**Problem 1. [Minimum Spanning Tree]**



- Draw the edges in the Minimum Spanning Tree for the following graph.
- Given  $G$  and MST  $T$ , suppose you decrease the weight of an edge  $e$  not in  $T$ . Give an algorithm to recompute the MST in  $O(n)$  time.

**Problem 2. [Stock Picking]**

You have a group of investor friends who are looking at  $n$  consecutive days of a given stock at some point in the past. The days are numbered.  $i = 1, 2, \dots, n$ . For each day  $i$ , they have a price  $p(i)$  per share for the stock on that day.

For certain (possibly large) values of  $k$ , they want to study what they call *k-shot strategies*. A *k-shot strategy* is a collection of  $m$  pairs of days  $(b_1, s_1), \dots, (b_m, s_m)$ , where  $0 \leq m \leq k$  and

$$1 \leq b_1 < s_1 < b_2 < s_2 \cdots < b_m < s_m \leq n$$

We view these as a set of up to  $k$  nonoverlapping intervals, during each of which the investors buy 1,000 shares of the stock (on day  $b_i$  and then sell it (on day  $s_i$ . The *return* of a given *k-shot strategy* is simply the profit obtained from the  $m$  buy-sell transactions, namely,

$$1000 \cdot \sum_{i=1}^m p(s_i) - p(b_i)$$

- Design an efficient algorithm that determines, given the sequence of prices, the *k-shot strategy* with the maximum possible return. Since  $k$  may be relatively large, your running time should be polynomial in both  $n$  and  $k$ .

- Now, modify your algorithm to only use  $O(n)$  space.

**Problem 3. [Weighted Scheduling]**

We have  $n$  jobs  $J_1, J_2, \dots, J_n$  which we need to schedule on a machine. Each job  $J_i$  has a processing time  $t_i$  and a weight  $w_i$ . A schedule for the machine is an ordering of the jobs. Given a schedule, let  $C_i$  denote the finishing time of job  $J_i$ . For example, if job  $J_j$  is the first job in the schedule, its finishing time  $C_j$  is equal to  $t_j$ ; if job  $J_j$  follows job  $J_i$  in the schedule, its finishing time  $C_j$  is equal to  $C_i + t_j$ . The weighted completion time of the schedule is  $\sum_{i=1}^n w_i C_i$ .

- Given an efficient algorithm that finds a minimum weighted schedule when  $w_i = 1$  for all  $i$ .
- Give an efficient algorithm that finds a schedule with minimum weighted completion time given arbitrary weights.