

CS440/ECE448: Intro to Artificial Intelligence

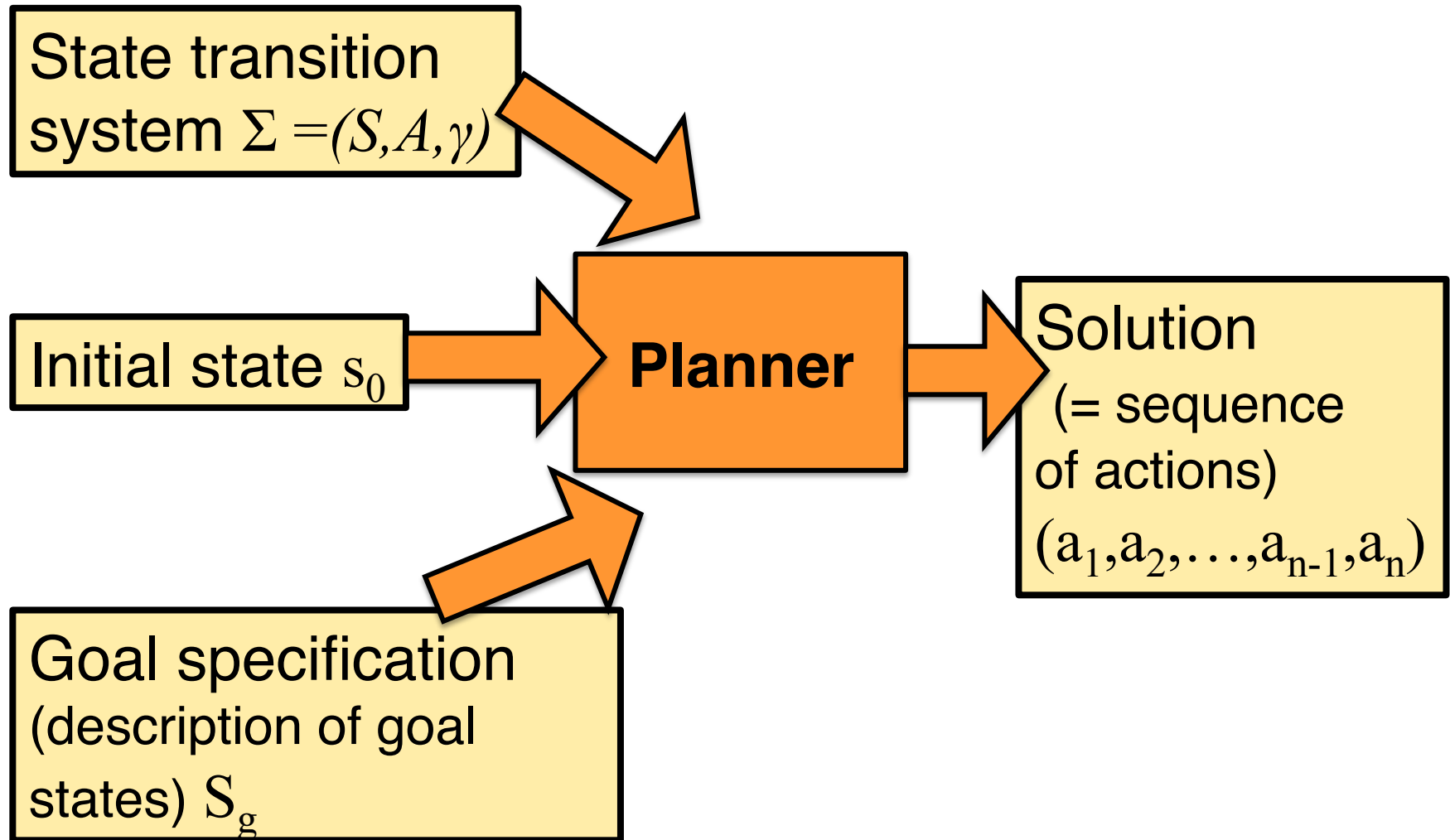
Lecture 12:

Planning algorithms

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Classical Planning



Review:

representations for planning

Situation Calculus

Specify fluents

Add-set

Persist-set

By default fluents
are deleted

Strips

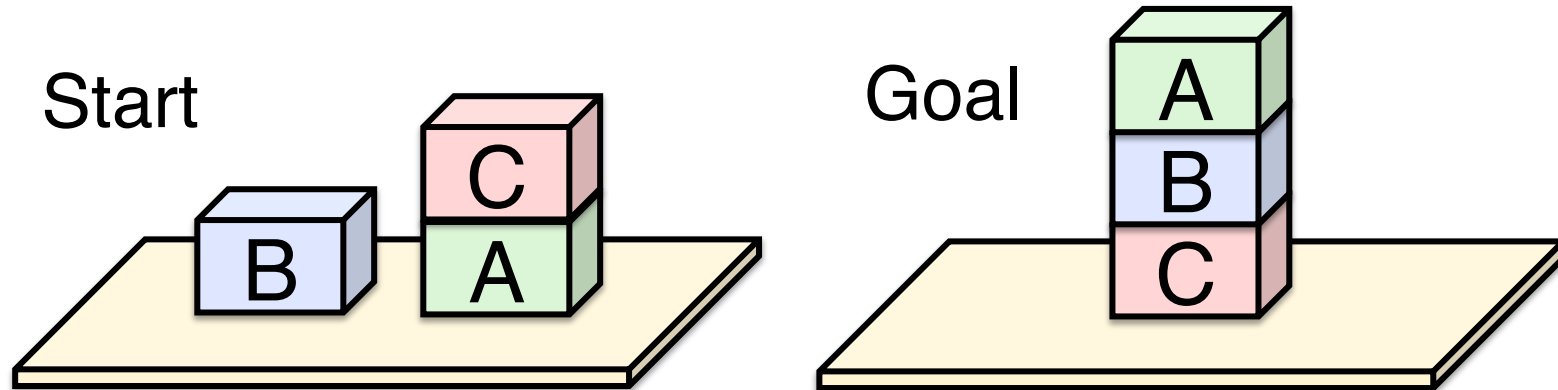
Specify fluents

Add-set

Delete-set

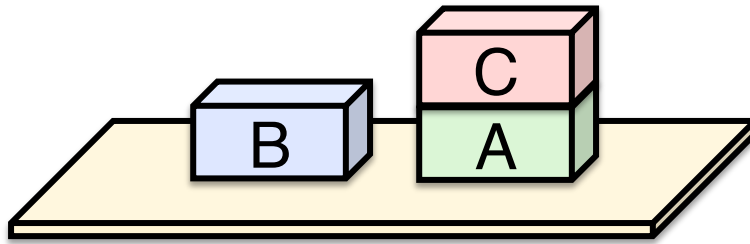
By default fluents persist

Sussman anomaly

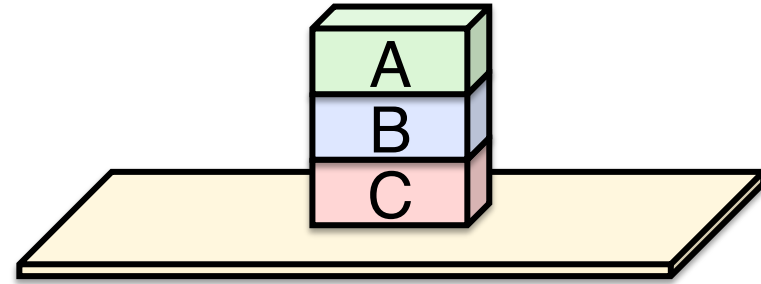


Start: $On(C,A)$

Goal: $On(A,B) \wedge On(B,C)$

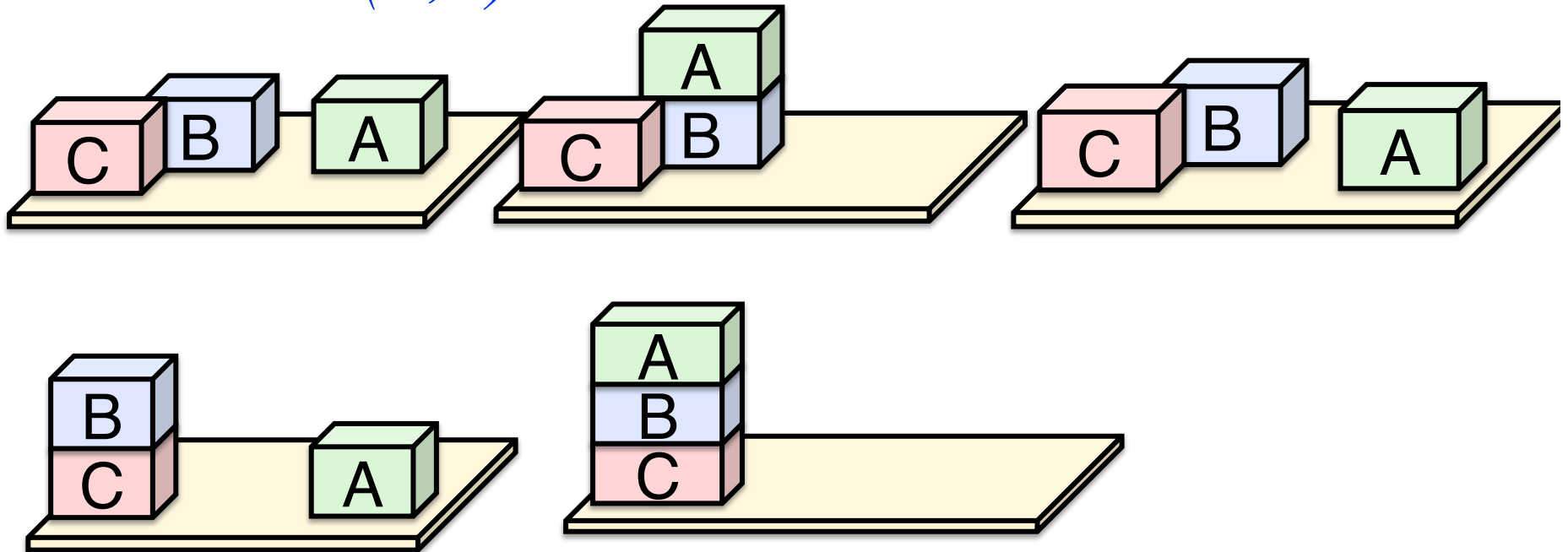


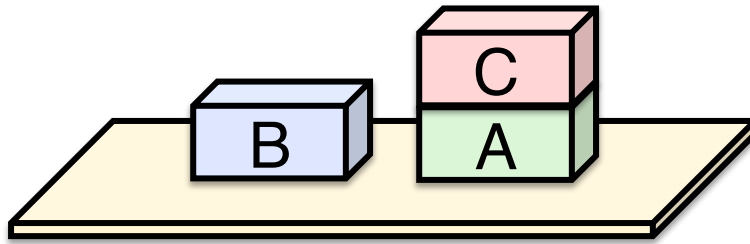
Start: $On(C,A)$



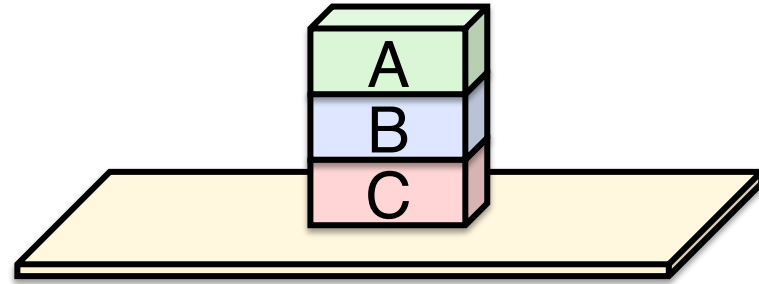
Goal: $On(A,B) \wedge On(B,C)$

Solve $On(A,B)$ first:



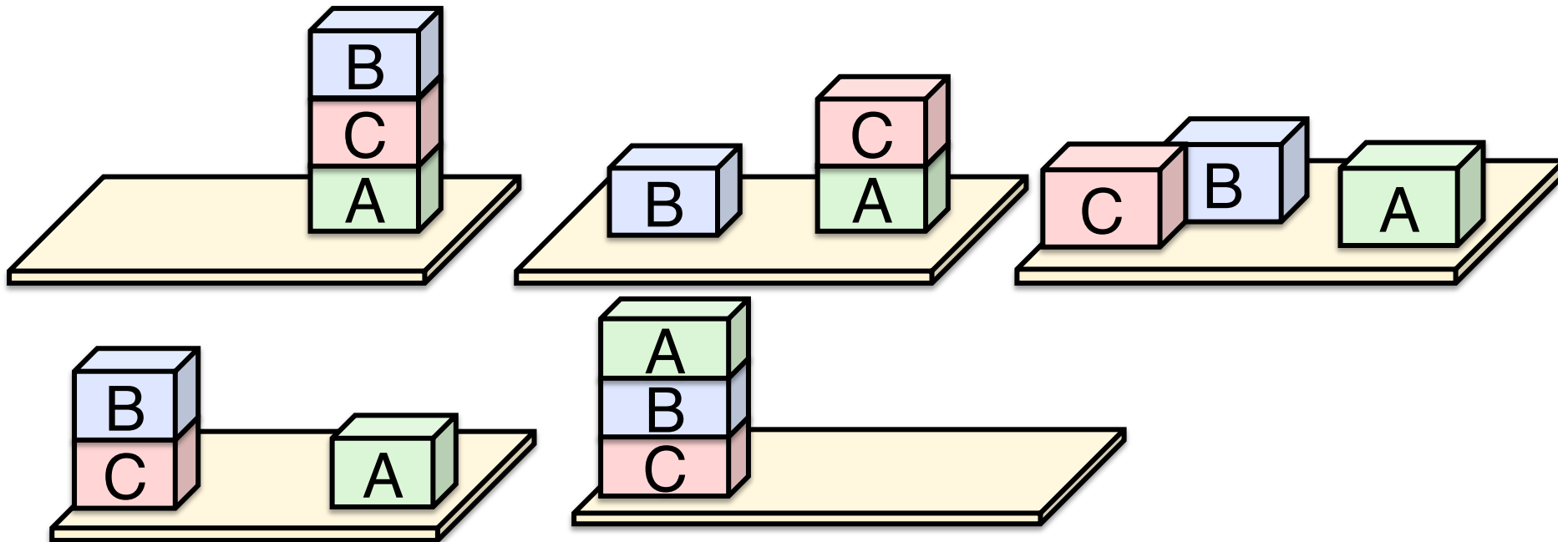


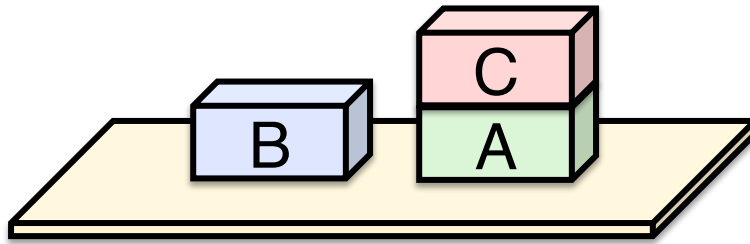
Start: $On(C,A)$



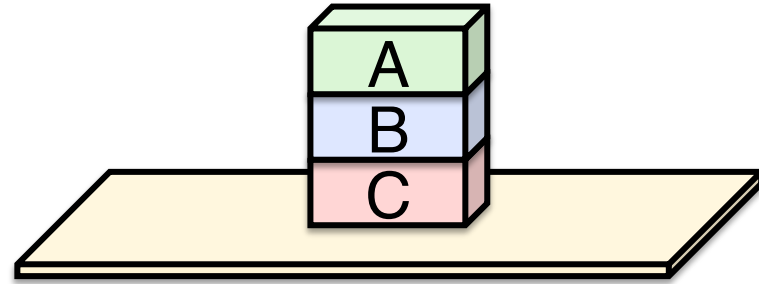
Goal: $On(A,B) \wedge On(B,C)$

Solve $On(B,C)$ first:



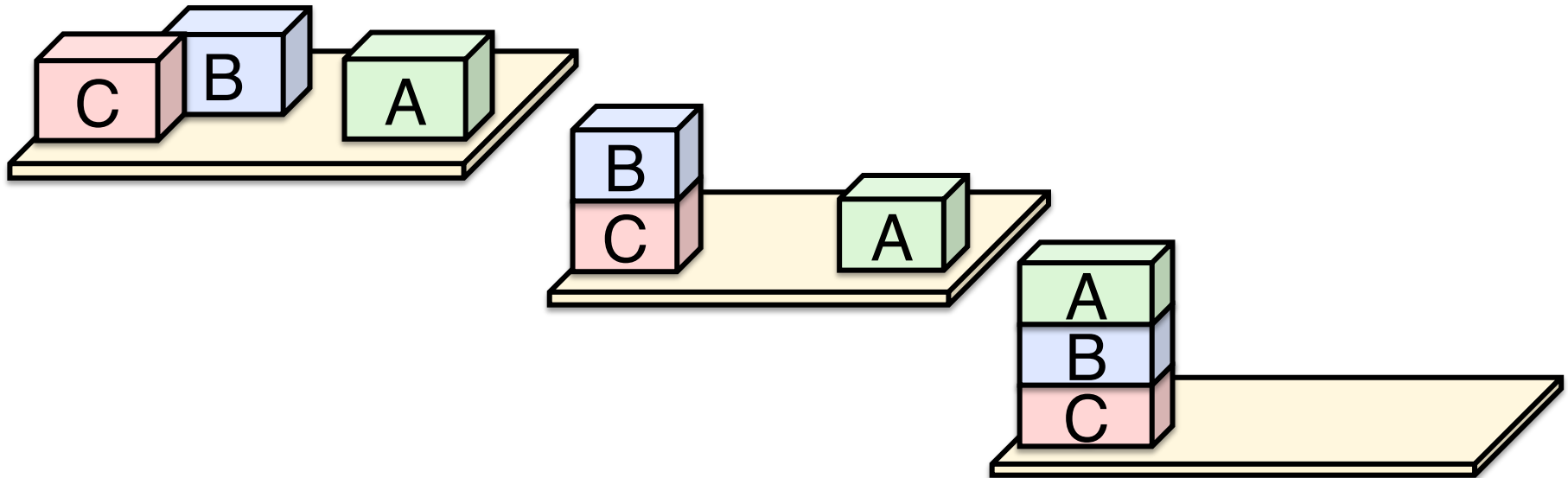


Start: $On(C,A)$



Goal: $On(A,B) \wedge On(B,C)$

Most efficient solution requires **interleaved planning**:



Planning algorithms

State space search (DFS, BFS, etc.)

Nodes = states; edges = actions;

Heuristics (make search more efficient)

Compute $h()$ using relaxed version of the problem

Plan space search (refinement of partial plans)

Nodes = partial plans; edges: fix flaws in plan

SATplan (encode plan in propositional logic)

Solution = *true* variables in a model for the plan

Graphplan (reduce search space to planning graph)

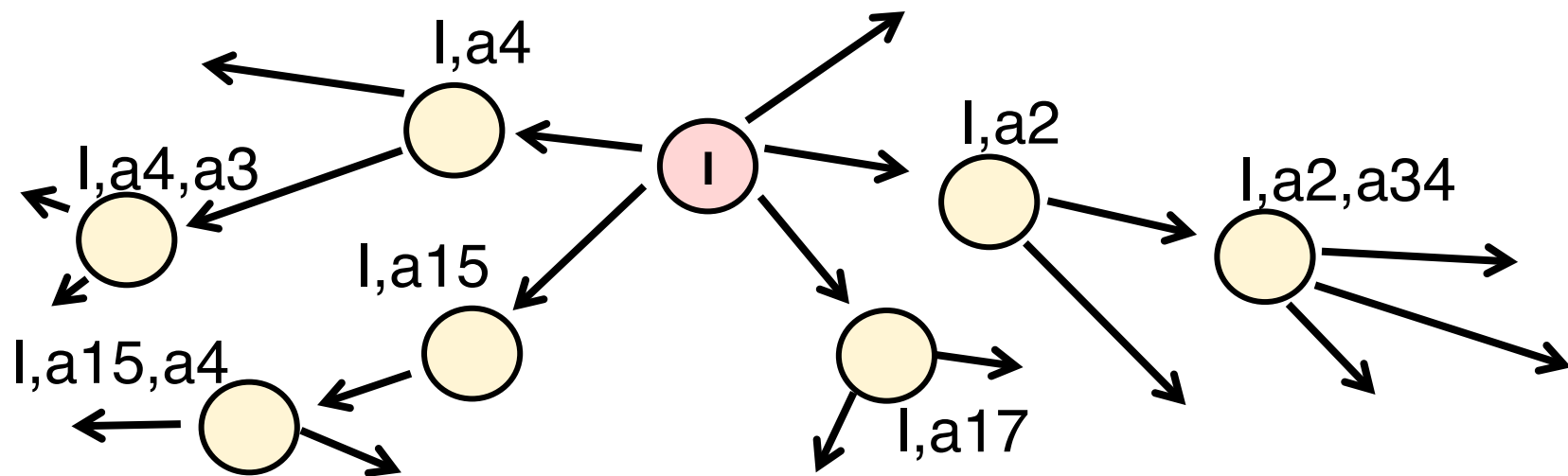
Planning graph: levels = literals and actions

State space search

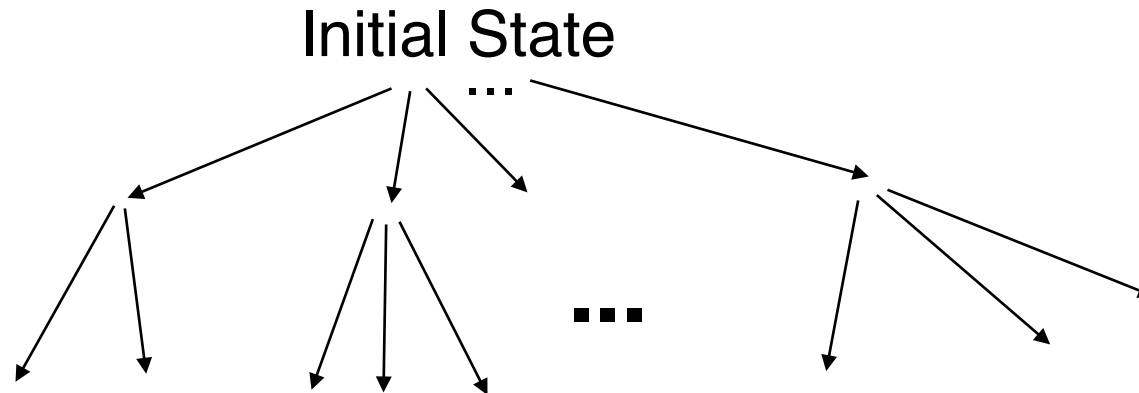
Planning as state space search

Search tree:

- Nodes: states
- Root: initial state
- Edges: actions (ground instances of operators)
- Solutions: paths from initial state to goal.



Forward search



Breadth-first forward search is sound and complete, but may require lots of memory

Depth-first forward search can be better in practice (needs graph-search to be complete)

Problem: branching factor is very large (need good heuristic: which actions may lead to goal?)

DFS and loops: iterative deepening

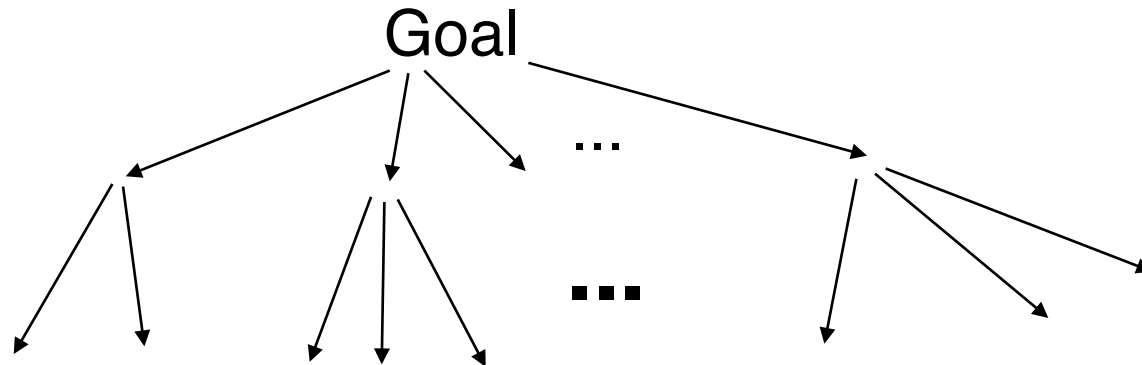
Loops ($s_i \rightarrow \dots \rightarrow s_i$) in the search graph lead to **infinite branches** in the search tree.

The tree-search variant of DFS never terminates if it goes down an infinite branch

Remedy (**iterative deepening**):

- Try to find solution of **length l** with DFS
- If this fails, **$l := l + \Delta$** ; try again.

Backward search



Start with goal; ‘undo’ actions until initial state

Edges = inverse actions a^{-1} :

$$Result(a^{-1}, s) = s \setminus effects(a) \cup precondition(a)$$

Large branching factor: many possible actions;
not every $Result(a^{-1}, s)$ leads back to initial state

Plan-space search

State space search is inefficient

Observation 1: many actions are **independent**.

e.g.: $move(A,B,C)$ and $move(D,E,F)$

We don't want to have to commit to specific order.

Observation 2: naïve backward search requires **fully instantiated actions**. We often don't know/care how to instantiate variables.

$move(A,B,?C)$ is prerequisite for $move(B,D,E)$, but we don't care which C A moves to.

Searching plan-space

States in plan-space are **partial plans**: = sets of partially instantiated actions with **constraints** on **precedence** and variable **(in)equality**

$$\textit{move}(A,B,x) < \textit{move}(B,y,z); \quad z \neq A$$

Solution is (partially ordered) complete plan with fully instantiated actions

Searching plan-space

1. Start with the **empty plan**
 $= \{start\ state, goal\ state\}$
2. Iteratively refine current plan to **resolve flaws**
(refine = add new actions and constraints)
3. **Solution** = a plan without flaws

Flaws: open goals

Open goal: a precondition that is not yet met

$move(A,B,x)$ requires $clear(x)$

Solution:

add new action and precedence constraint:

$move(A,B,x)$ needs to be preceded by
some new action A with effect $clear(x)$.

Flaws: threats

C is a threat if A is a precondition for B and C undoes the effect of A .

$move(A,B,x)$ establishes precondition $clear(B)$ for $move(C,D,B)$. $move(E,F,B)$ undoes $clear(B)$

Solution: **add new precedence constraint:** $move(E,F,B)$ has to precede $move(A,B,x)$, or follow $move(C,D,B)$.

SATplan

SATplan: basic idea

We can encode a plan of fixed length n (n time steps required for solution) as a formula in propositional logic.

There is a solution if this formula is satisfiable. Use existing tools (SAT solvers)

If there is no solution of length n , try $n+1$.

From plans to propositional logic

– **Fluents** are ground literals:

$clear(B)^t$: block B is clear at time t

– **Actions** are ground implications:

$(preconditions^t \wedge action^t) \rightarrow effect^{t+1}$

Operator $move(x,y,z)$:

PRE: $on(x,y), clr(z)$ *EFFECT*: $on(x,z), clr(y)$

Action $move(A,B,C)^{23}$

$(on(A,B)^{23} \wedge clr(C)^{23} \wedge move(A,B,C)^{23}) \rightarrow on(A,C)^{24} \wedge clr(B)^{24}$

Graphplan

Basic idea

1. Create an easier, relaxed problem P' :
 - can be solved in polynomial time
 - relax = remove some restrictions
 - Solutions to $P' \supset$ solutions to P
2. Solve this relaxed problem P'
3. Search among solutions to P' for solutions to P

Iterative deepening: try to solve P' in $1, \dots, n$ steps

Basic idea

Relaxed problem P' : can we satisfy **some necessary precondition** for the goal of P **in k steps**?
(N.B.: this only solves P if *all* necessary preconditions can be achieved in k steps)

Solve P' : build a **planning graph** of depth $(2)k$

Solve P : do **backward search on the planning graph** to extract solution to P . If this fails, set k to $k+1$.

Planning graph: nodes

Two kinds of levels alternate:

State level S_i :

Node at level S_i : a ground **fluent** (pos. or neg. literal) which may hold i actions after S_0

Action level A_i :

Node at level A_i : a ground **action** (incl. *noop*) whose preconditions might be satisfied at S_i

Have your cake and eat it too!

Init: *Have(Cake)*

Goal: *Have(Cake) \wedge Eaten(Cake)*

Action *Eat(Cake)*:

PRECOND: *Have(Cake)*

EFFECT: \neg *Have(Cake)* \wedge *Eaten(Cake)*

Action *Bake(Cake)*:

PRECOND: \neg *Have(Cake)*

EFFECT: *Have(Cake)*

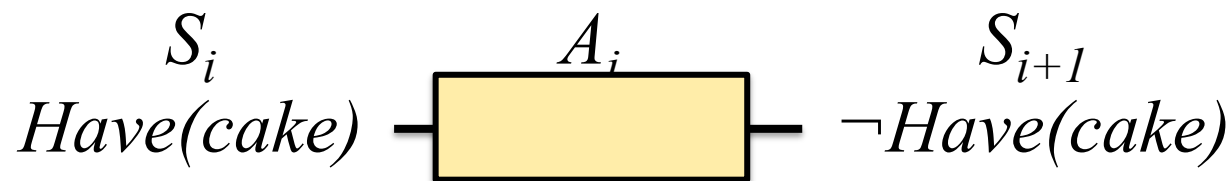
Preconditions/effects: edges *between* levels

From state level S_i to action level A_i :

the node (fluent) at level S_i is a **precondition** for the node (action) at level A_i

From action level A_i to state level S_{i+1}

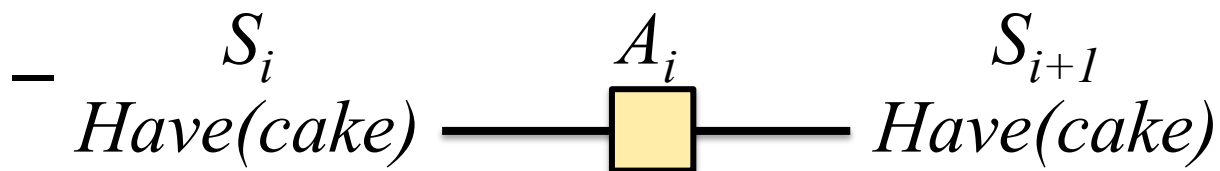
the node (fluent) at level S_{i+1} is an **effect** of the node (action) at level A_i



Persistence (no-op) actions

For each literal C_i at state level S_i :

- add a **persistence action Noop_C** at action level A_i
- Link from **C_i to Noop_C**
- Also add **$C_{(i+1)}$** to state level S_{i+1}
- Link from **Noop_C to C_{i+1}**



Mutually exclusive actions: *mutex links within action levels*

–Inconsistent effects:

effect of A_1 negates effect of A_2
eat(cake) negates *bake(cake)*

–Interference:

effect of A_1 negates precondition of A_2
eat(cake) interferes with no-op for *have(cake)*

–Competing needs:

preconditions of A_1 , A_2 are *mutex*
eat(cake) competes with *bake(cake)*

Mutually exclusive fluents: edges *within* state levels

Mutex within state level S_i : $F1$ and $F2$ cannot hold at the same time.

– Negation:

have(Cake) negates $\neg have(Cake)$

– Inconsistent support:

all actions $a1$ to achieve $F1$ are mutex with any action $a2$ that achieves $F2$.

have(Cake) and eaten(Cake)

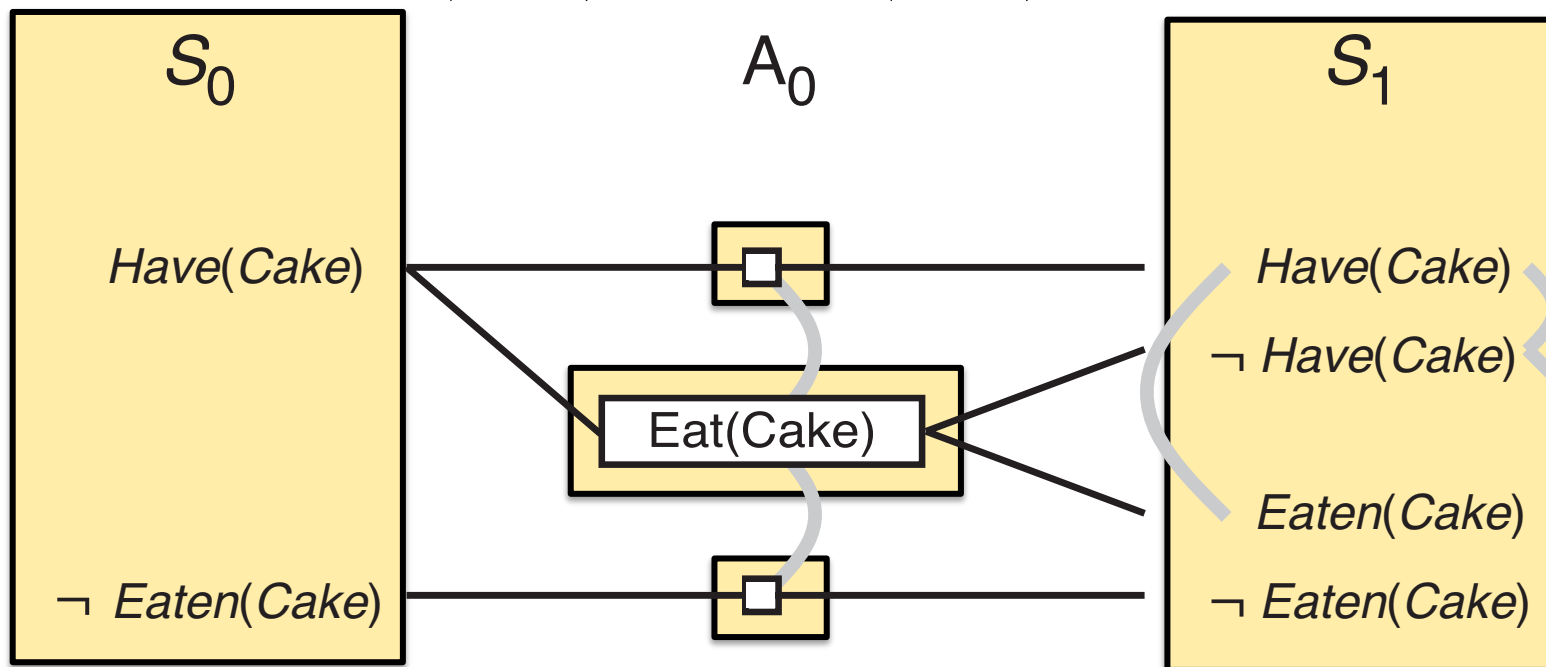
The initial planning graph

Action *Eat(Cake)*

PC: *Have(Cake)* EF: $\neg \text{Have(Cake)} \wedge \text{Eaten(Cake)}$

Action *Bake(Cake)*

PC.: $\neg \text{Have(Cake)}$ EF: *Have(Cake)*



The size of the planning graph

The original planning problem P
has l literals and a actions

Size of planning graph with n levels:
polynomial in l and a

- State levels: at most l nodes and l^2 mutex links
- Action levels: at most $a+l$ nodes, $(a+l)^2$ mutex;

Solution extraction

If all literals in $goal(P)$ hold at S_n and are not *mutex* with each other, we may be able to find a solution by backward search:

- States in backward search tree:
conjunction of (a subset of the) literals at S_i
- Actions in backward search tree:
conflict-free subset of actions at A_{i-1}
- Goals of S_n : all literals in $goal(P)$
- At each state level S_i , select

The search tree for solution extraction

- **State:** conjunction of (subset of the) literals at S_i
- **Initial state:** all literals in $goal(P)$
- **Goal state:** all literals in $init(P)$

- **Actions:** a conflict-free subset \mathbf{a} of actions at A_{i-1}
- **Result of a set of actions from A_{i-1} :**
the conjunction of all literals in S_{i-1}
that are preconditions for some action in \mathbf{a}

Heuristics for planning

Using the planning graph to estimate heuristics

What is the **cost of a goal literal g** ?

Minimum cost (level cost)

= minimum number of steps required to achieve g

= first state level at which g

What is the **cost of a conjunction of goal literals**?

Max level cost: max. of level costs of each literal

Level sum: sum of level costs (inadmissible)

Set level cost: level at which all are true and each pair of literals is non *mutex*

Today's key concepts

State space search (DFS, BFS, etc.)

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Compute $h()$ using relaxed version of the problem

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