

# CS 439: Wireless Networking

IoT in Public Spaces

# Bridging the Gap with IoT



**Incentives and Privacy  
Exposure vs. Benefits**

**IoT**



# Designing an IoT Ecosystem

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## Context Discovery

**“What environment  
am I in?”**



**The CoffeeShop**

# Designing an IoT Ecosystem

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Context Discovery

User Identity



**“What identity  
should I expose?”**

**FrenchRoast99**



# Designing an IoT Ecosystem

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Context Discovery

User Identity

User Identity Use and Reuse



**FrenchRoast99**

**Location-based  
pseudonyms**



**Espresso42**



# Designing an IoT Ecosystem

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Context Discovery

User Identity

User Identity Use and Reuse

Sharing/Querying



**The CoffeeShop**

**Enable social networking  
based recommendations**



**A friend**

# Designing an IoT Ecosystem

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Context Discovery

User Identity

User Identity Use and Reuse

Sharing/Querying

**Secure Validation and Privacy**



**The CoffeeShop**

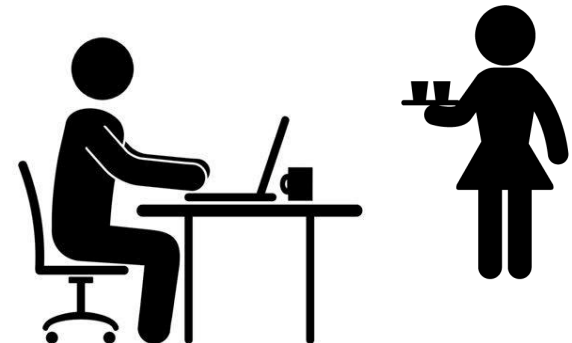
**Prevent impersonation  
and  
unauthorized access**



# Incognito

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- ▶ A privacy-preserving IoT ecosystem architecture
  - ▶ Users share any desired part of their identity within an environment
- ▶ User-managed identities
  - ▶ cid: context-based identities





# Incognito

- ▶ User and Infrastructure Devices
  - ▶ Wi-Fi to infrastructure



# Incognito

## ▶ User and Infrastructure Devices

- ▶ Wi-Fi to infrastructure
- ▶ Bluetooth LE for local environment

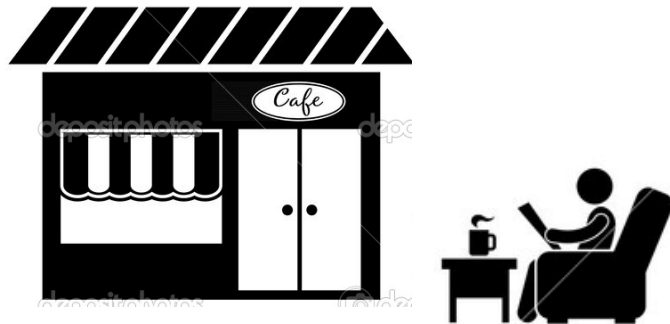


# Incognito

## ▶ Environments

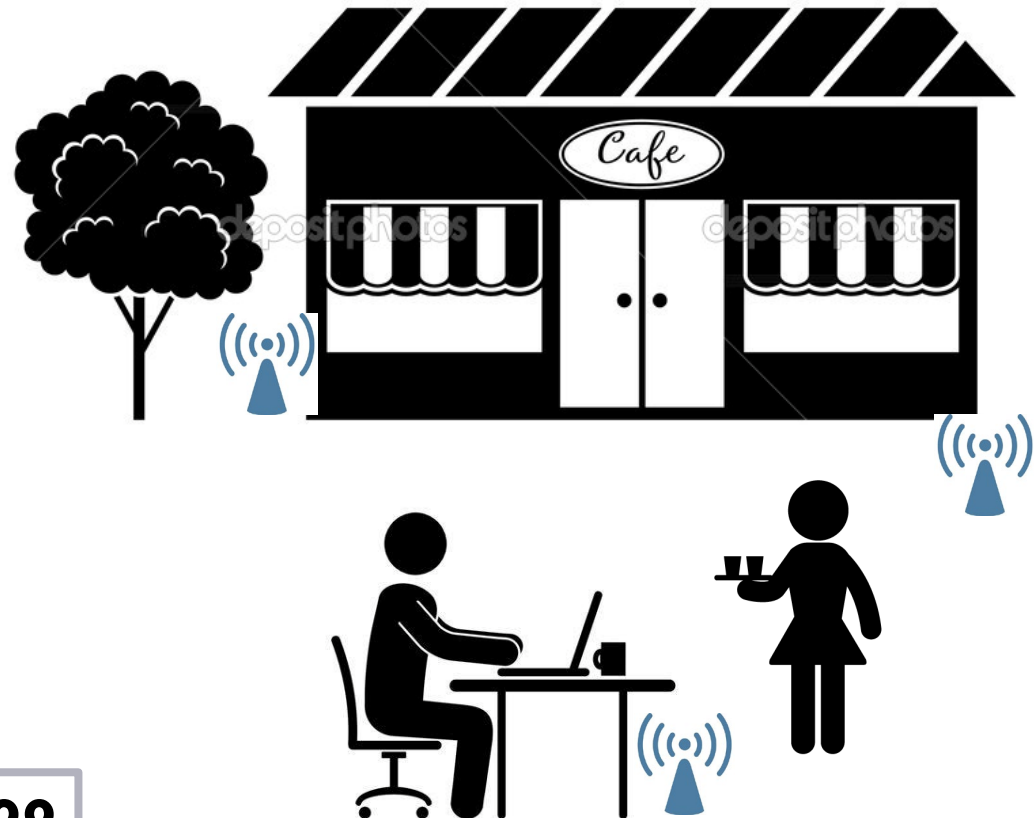
- ▶ Organizations can share data across locations

### The CoffeeShop II



### FrenchRoast99

### The CoffeeShop

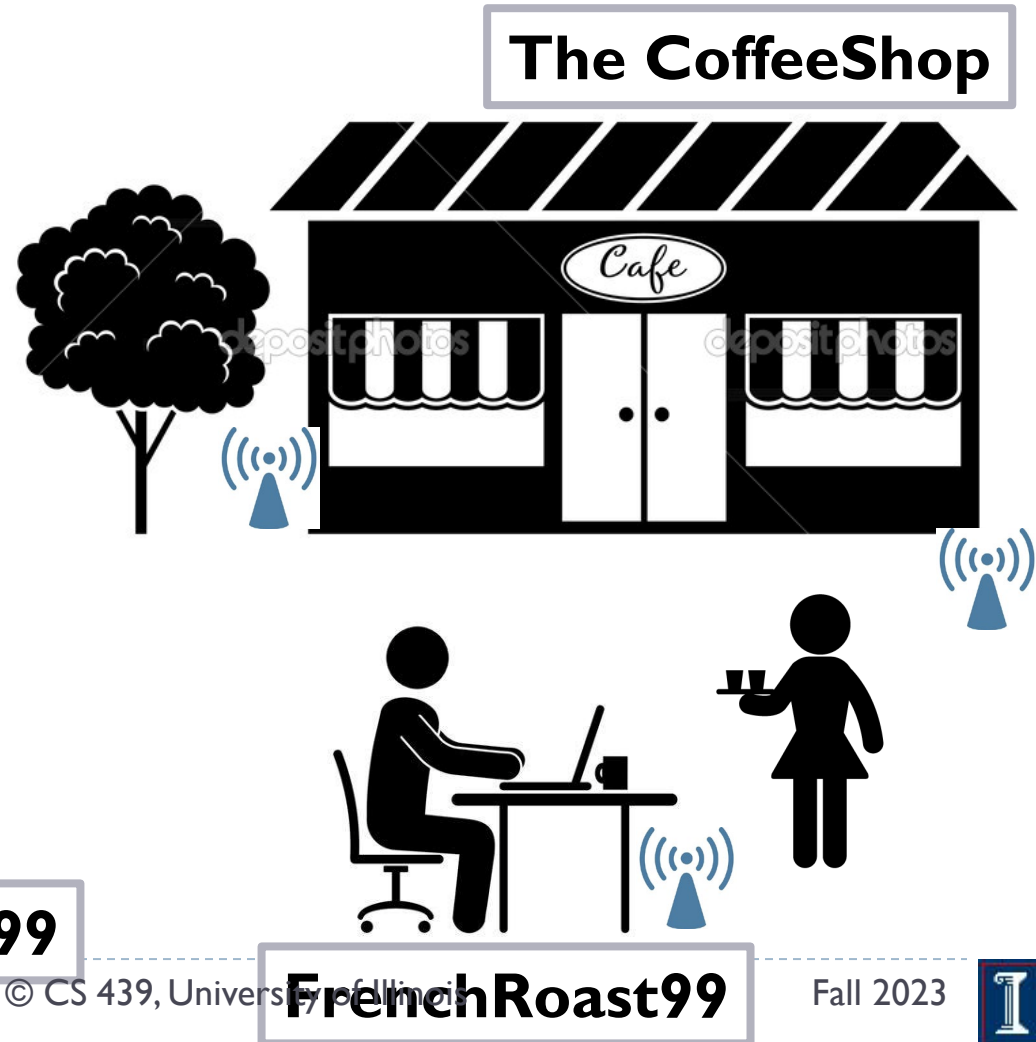
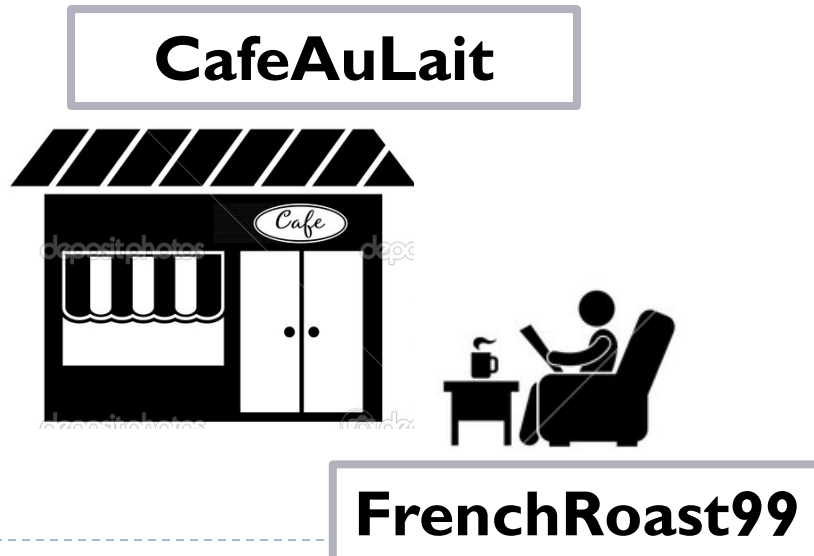


### FrenchRoast99

# Incognito

## ▶ Environments

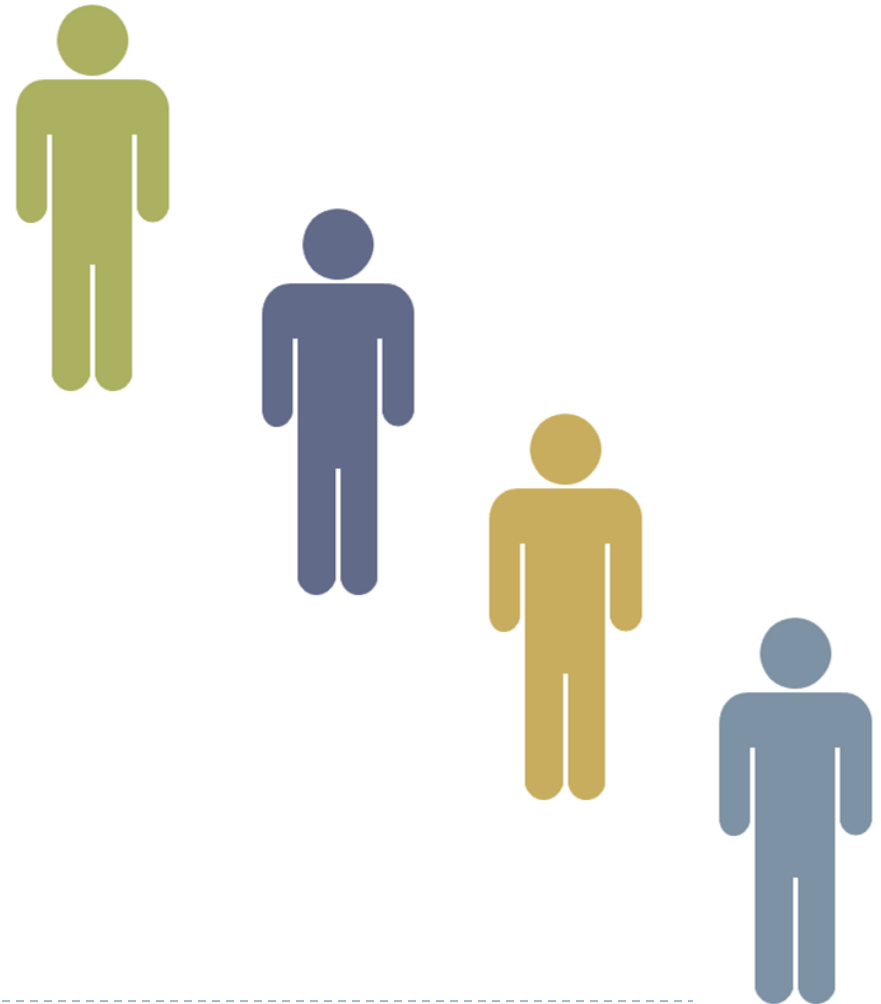
- ▶ Users can enable sharing of data within categories



# User Identities

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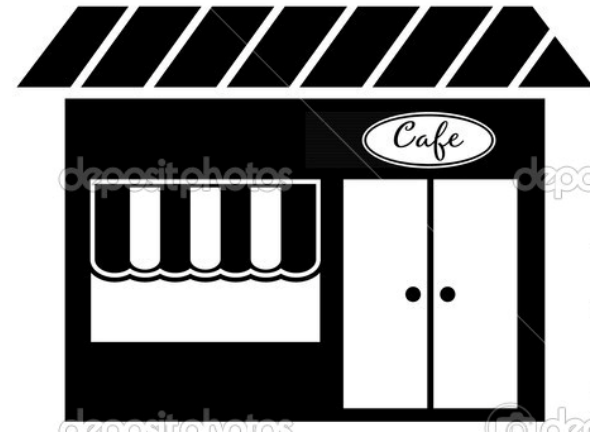
- ▶ **Anonymous**
  - ▶ Random cid every packet
  - ▶ For every message, the user appears as someone new
  - ▶ No tracking from environment



# User Identities

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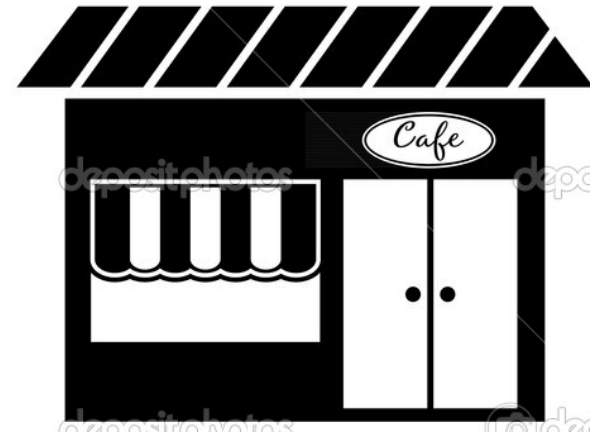
- ▶ **Local-One-Time**
  - ▶ Random cid per session



# User Identities

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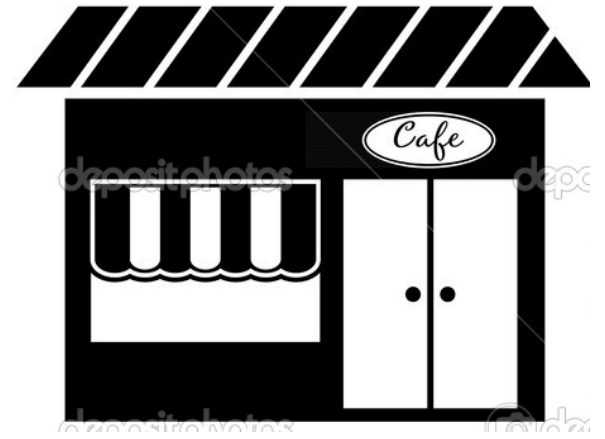
- ▶ **Local-One-Time**
  - ▶ Random cid per session
  - ▶ No connection between multiple sessions from the same user in the same environment



# User Identities

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- ▶ **Local**
  - ▶ Random cid per environment



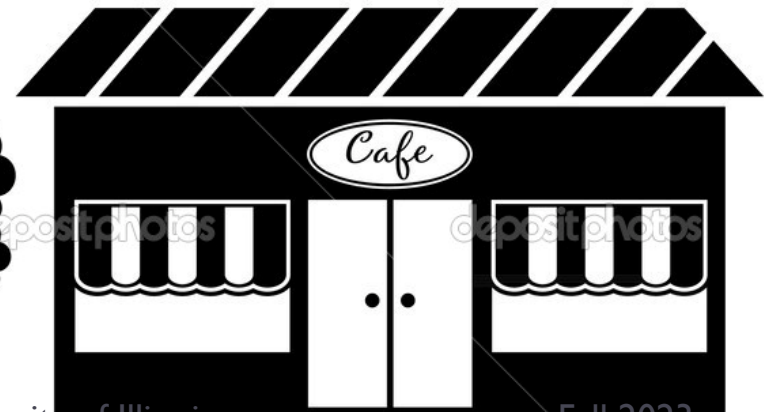
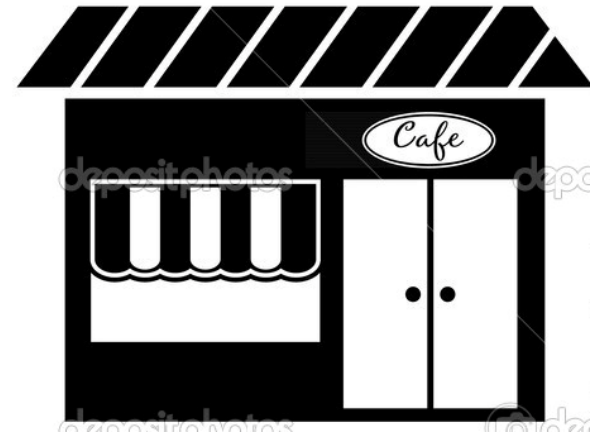


# User Identities

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## ▶ Local

- ▶ Random cid per environment
- ▶ No connection to the same user in a different environment

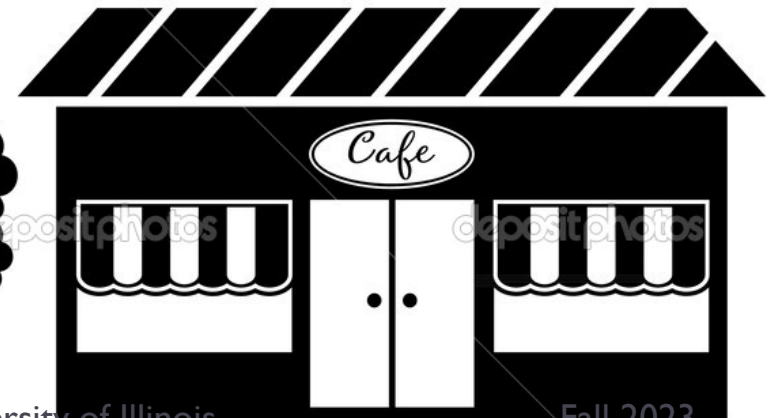
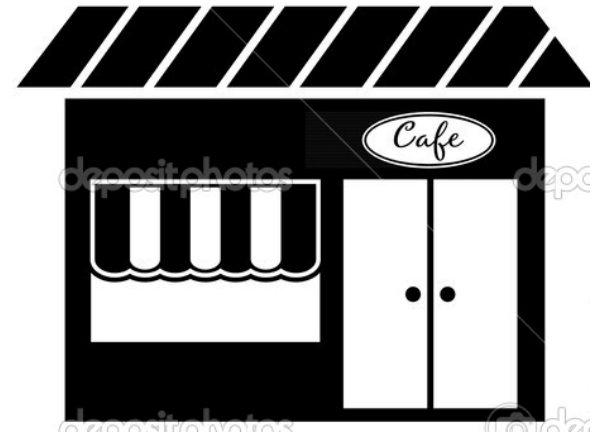


# User Identities

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- ▶ **Cross-Domain**

- ▶ Random cid per environment class
- ▶ Track and share user information within an environment class



# User Identities

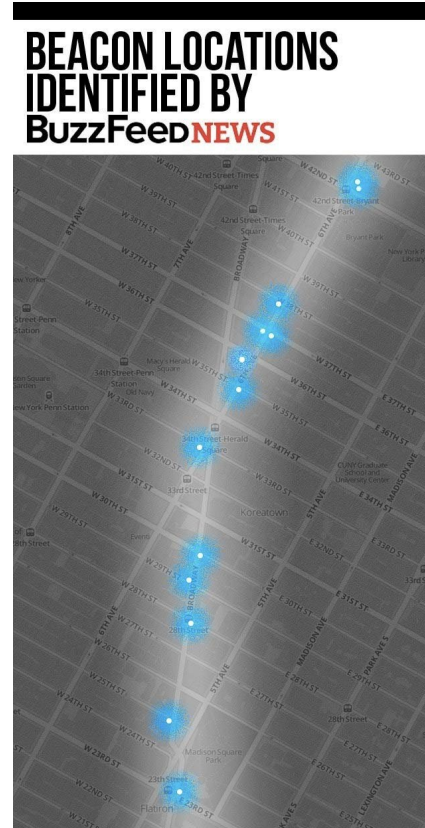
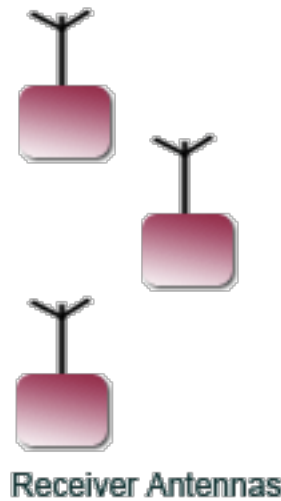
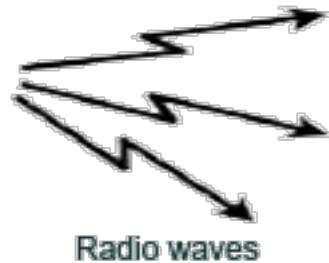
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- ▶ **Global**
  - ▶ Global cid
  - ▶ User exposes an identity all of the time



# Managing Exposure

**All transmissions contain the identity of the sender (MAC address)**



**Anyone can listen and track the user**



# Managing Exposure

---

**Cycle through random  
MAC address**



**No one knows who  
you are!**



# Managing Exposure

---

**Cycle through random  
MAC address**



**No one knows who  
you are!**

**Add cid into beacon  
message**



**Uses limited 31B of  
data every time**



# Managing Exposure

---

**Cycle through random  
MAC address**



**No one knows who  
you are!**

**Add cid into beacon  
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**Uses limited 31B of  
data every time**

**Incognito: User-managed  
MAC addresses**



# Managing Exposure

---

**Cycle through random  
MAC address**



**No one knows who  
you are!**

**Add cid into beacon  
message**



**Uses limited 31B of  
data every time**

**Incognito: User-managed  
MAC addresses  
BLE and WiFi!**





# Managing Exposure

---

**Cycle through random  
MAC address**



**No one knows who  
you are!**

**Add cid into beacon  
message**



**Uses limited 31B of  
data every time**

**Both MAC addresses  
are set to current cid**

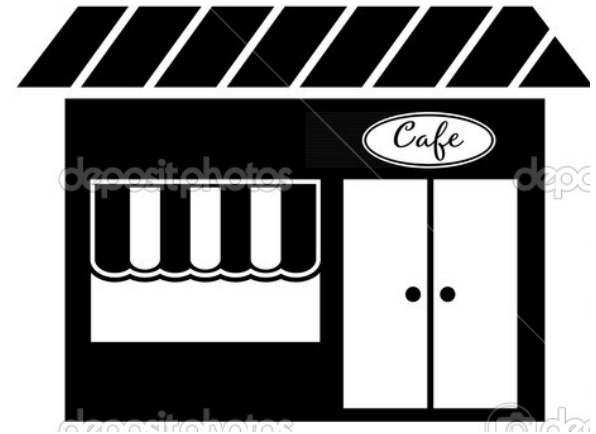


# User-to-Environment Sharing

## User registers with environment



MAC = cid  
Data = (PublicKey,  
PrivateKeyEncrypti(cid))



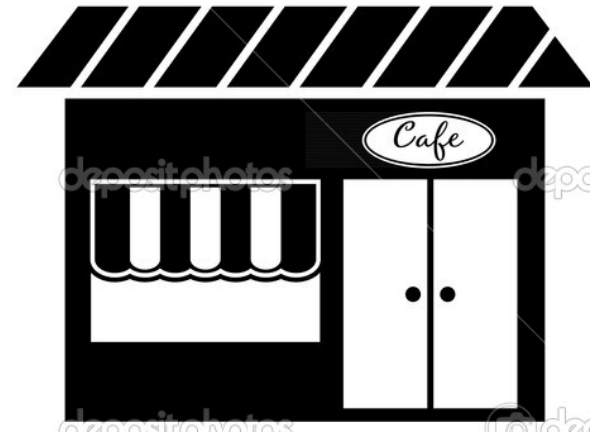
**Environment specific  
public-private key pair**

Conte xt ID	Public Key	Encrypted CID
cid	PublicKey	PrivateKey(cid)
cid-a	PublicKey-a	PrivateKey-a(cid-a)



# User-to-Environment Sharing

## User advertises presence



MAC = cid  
Data = (Hash(TimeStamp),  
PrivateKeyEncrypt(TimeStamp))

**Timestamp added to prevent impersonation**

Context ID	Public Key	Encrypted CID
cid	PublicKey	PrivateKey(cid)
cid-a	PublicKey-a	PrivateKey-a(cid-a)

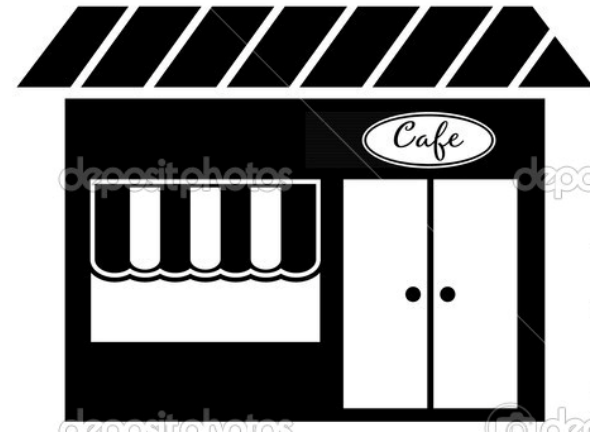


# User-to-User Sharing

## User enables sharing with friends



$(cid, PrivateKeyEncrypt(cid))$



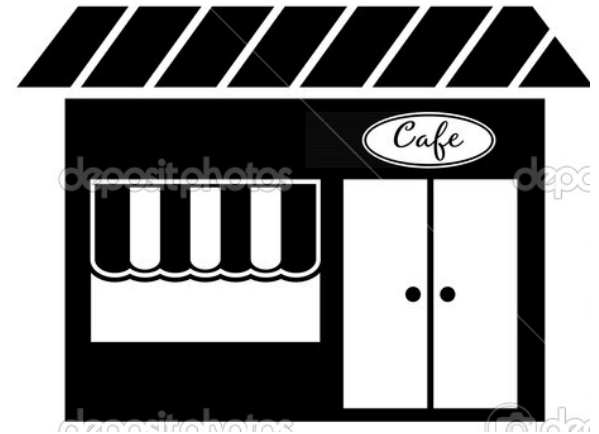
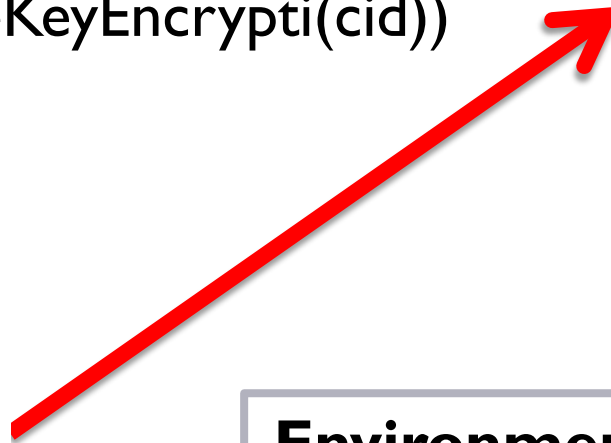
**Share cid and encrypted cid with second user**

# User-to-User Sharing

## Friend queries environment data



(cid,  
PrivateKeyEncrypt(cid))



**Environment can validate  
cid without any  
knowledge of second user**

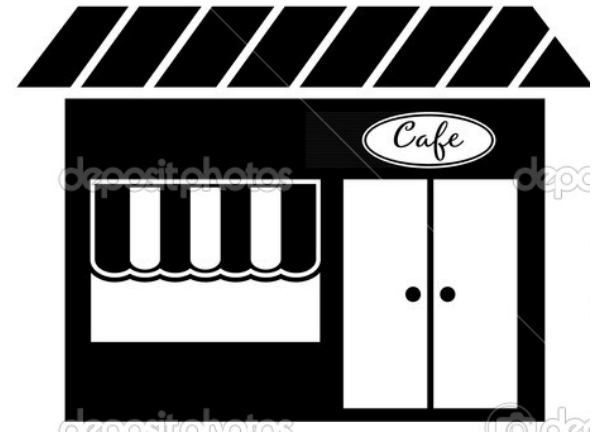
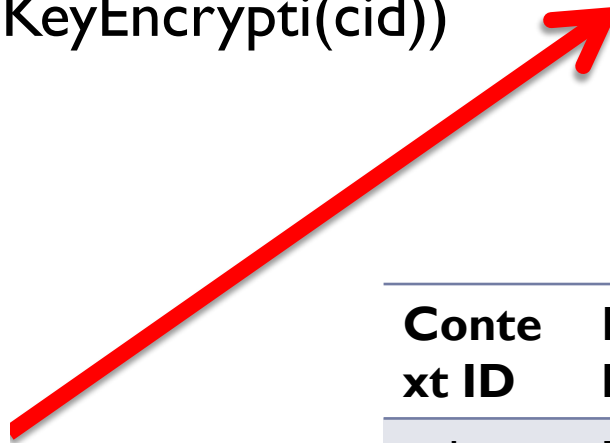


# User-to-User Sharing

## Friend queries environment data



(cid,  
PrivateKeyEncrypt(cid))



Context ID	Public Key	Encrypted CID
cid	PublicKey	PrivateKey(cid)
cid-a	PublicKey-a	PrivateKey-a(cid-a)

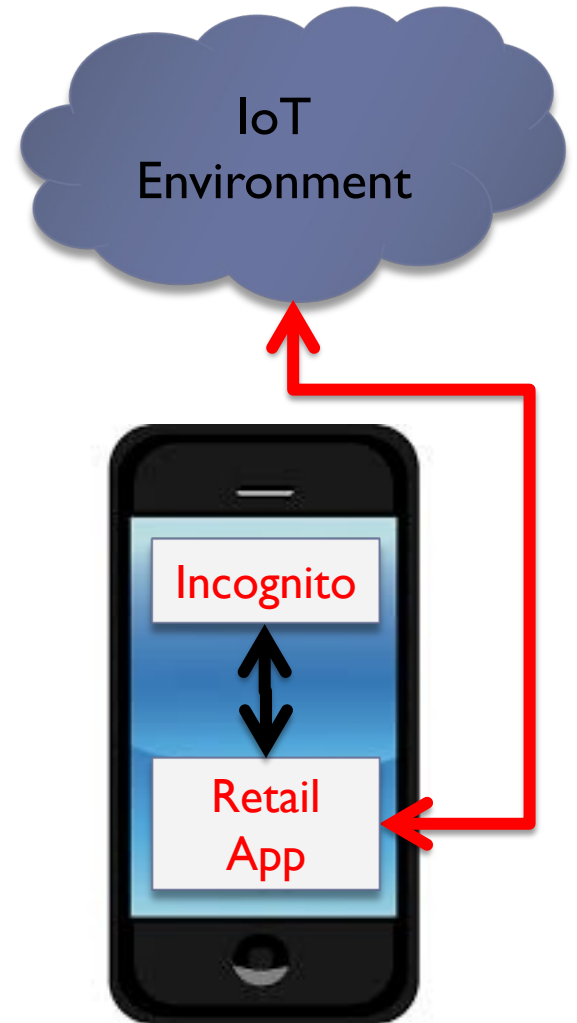


# Realizing Incognito

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## Ensuring privacy from malicious apps

**Challenge:  
External apps could  
leak cid**

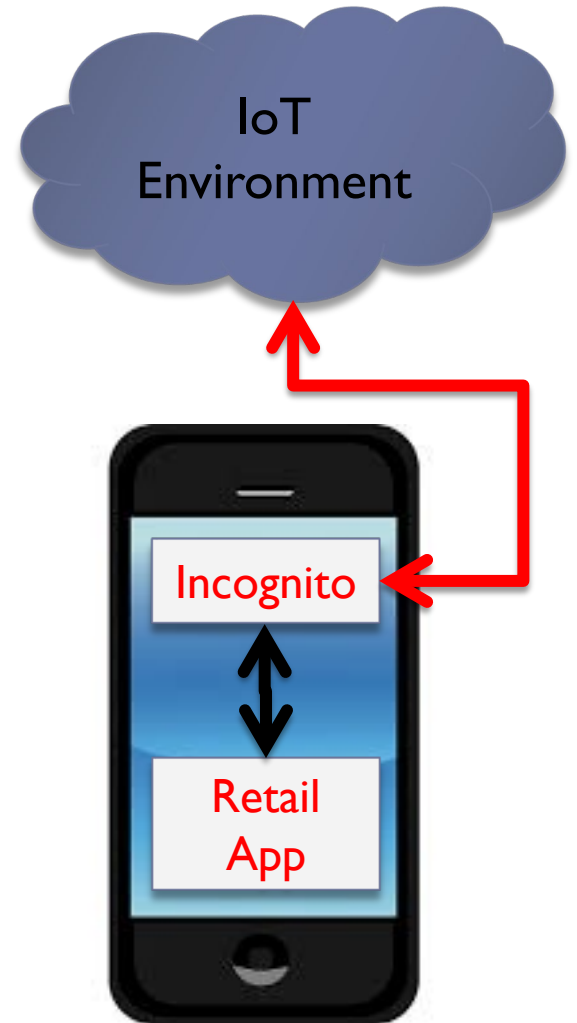


# Realizing Incognito

## Ensuring privacy from malicious apps

**Challenge:**  
External apps could leak cid

**Solution:**  
Sandbox IoT apps  
All communication with IoT infrastructure through incognito





# Realizing Incognito

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## CID Lifetime and Exposure

**Challenge:**  
**One environment could**  
**snoop into the other**

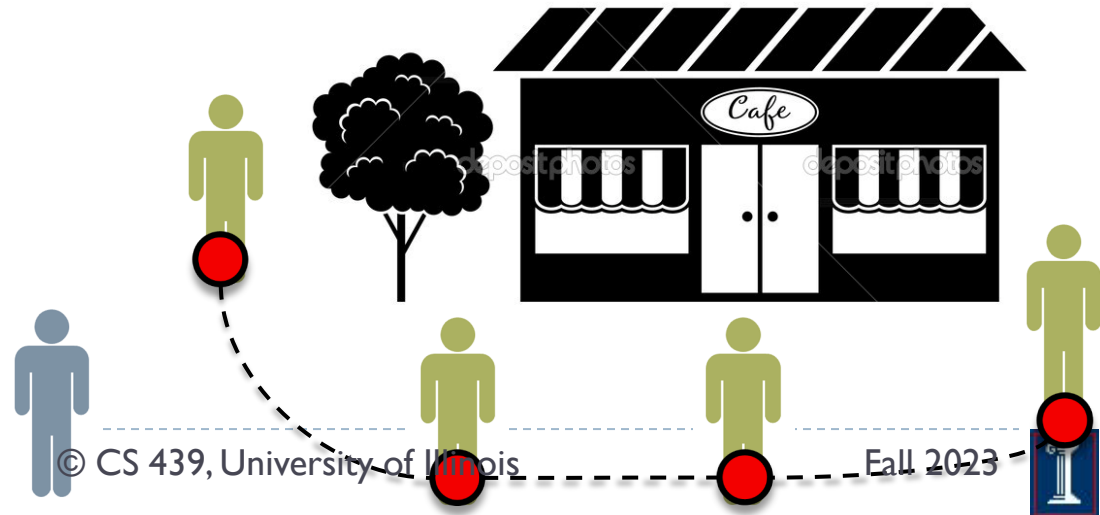


# Realizing Incognito

## CID Lifetime and Exposure

**Challenge:**  
One environment could  
snoop into the other

**Challenge:**  
A snooper can track the  
path of a cid



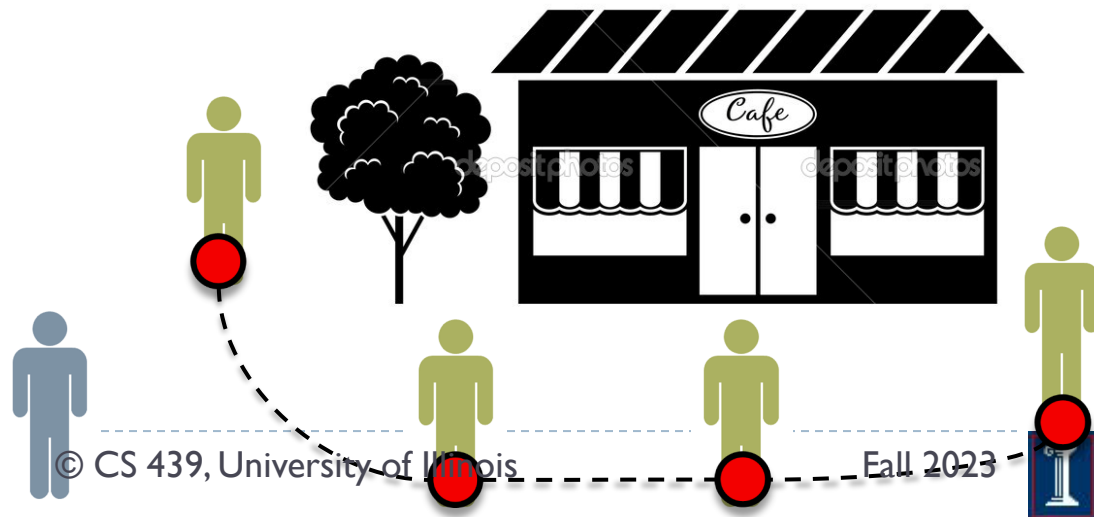
# Realizing Incognito

## CID Lifetime and Exposure

**Challenge:**  
One environment could snoop into the other

**Challenge:**  
A snooper can track the path of a cid

Can we bring back random MAC addresses and still maintain cid as a location based identifier?



# Anonymizing MAC Addresses

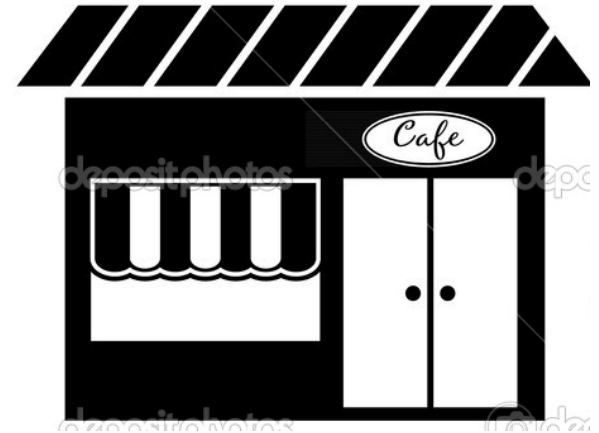


Registration



MAC = cid

Data = (PublicKey,  
PrivateKeyEncrypt(cid))

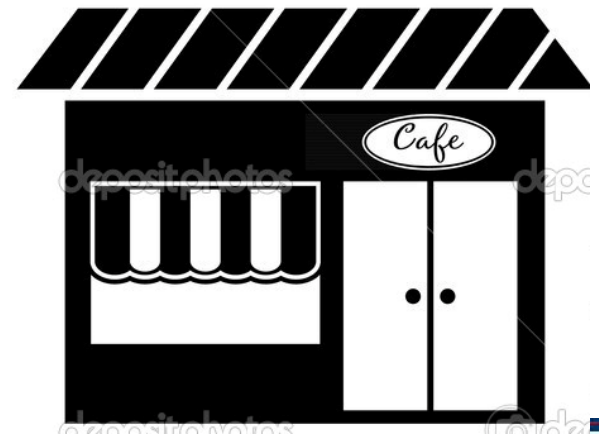


Advertising



MAC = cid

Data = (Hash(TimeStamp),  
PrivateKeyEncrypt(TimeStamp))



# Anonymizing MAC Addresses

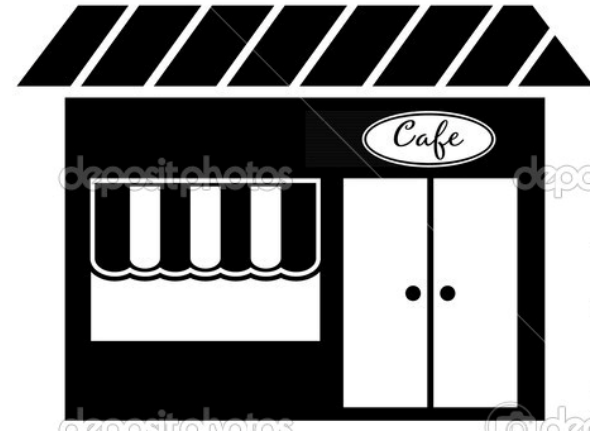


Registration



MAC = cid

Data = (PublicKey,  
PrivateKeyEncrypt(cid))



**Only prevents  
impersonation, not tracking**

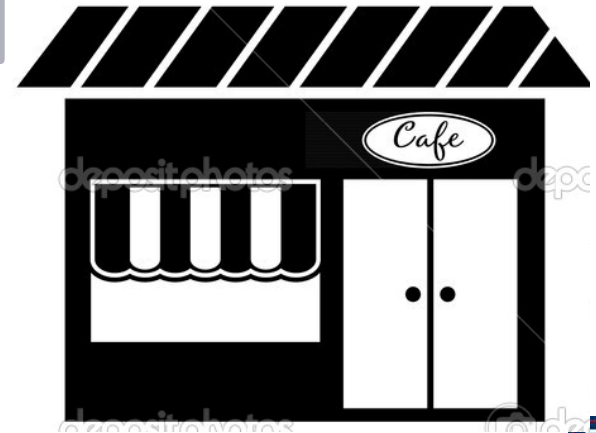


Advertising



MAC = cid

Data = (Hash(TimeStamp),  
PrivateKeyEncrypt(TimeStamp))



# Anonymizing MAC Addresses



Registration



**PrivateKeyEncrypt** during advertising is expensive

**High computation and energy**

**Large packets**

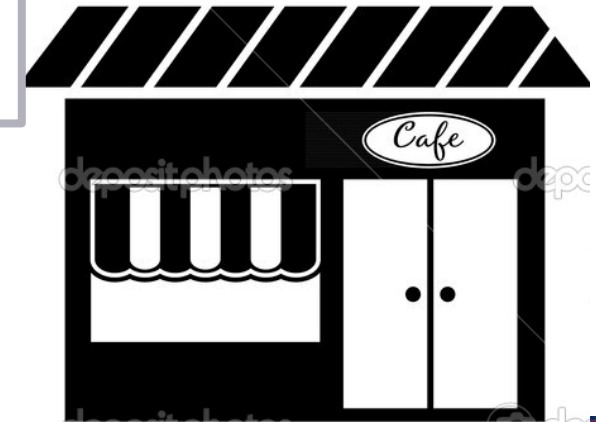
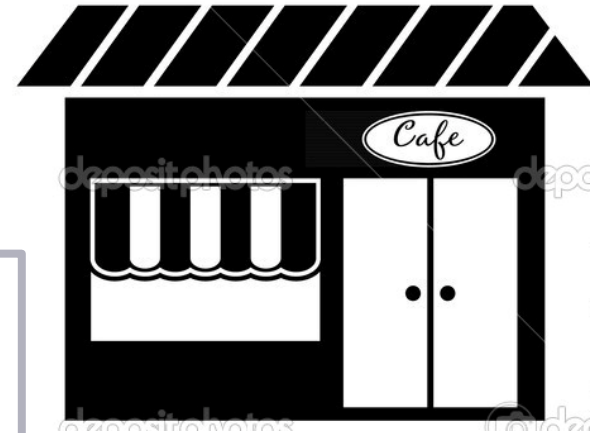
Advertising



MAC = cid

Data = (Hash(TimeStamp),

**PrivateKeyEncrypt(TimeStamp))**



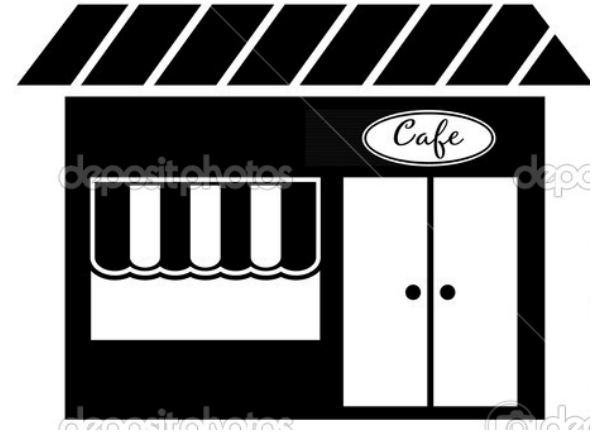
# Lamina

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## User registers with environment

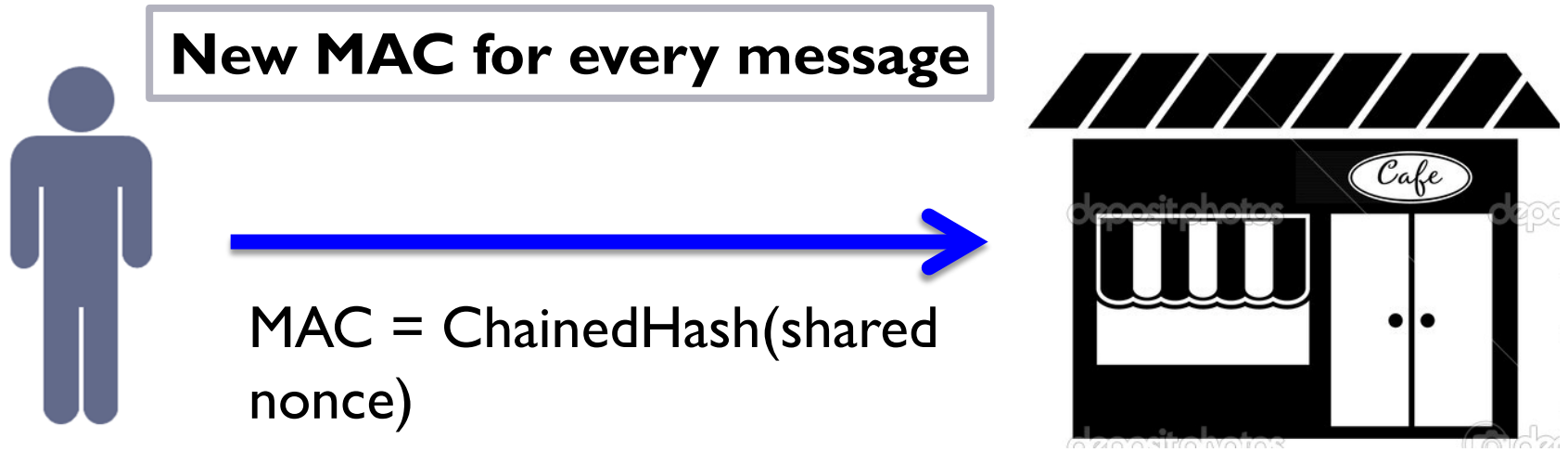


(cid , PublicKey,  
PrivateKeyEncrypti(cid),  
**shared nonce**)



# Lamina

## User advertises presence





# Lamina

## User advertises presence



### Problem

**Many encryption protocols  
assume a reliable channel**

# Lamina

## User advertises presence



### Problem

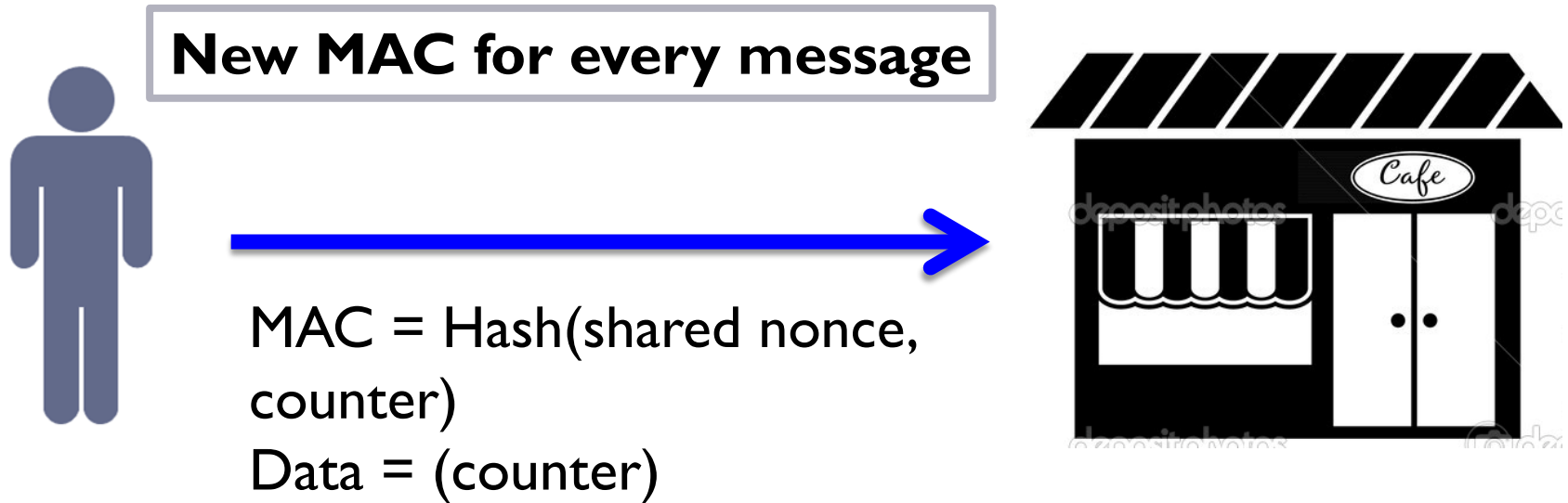
Many encryption protocols  
assume a reliable channel

Loss → protocol desync,  
further decryption of the  
stream impossible

# Lamina

---

## User advertises presence



**Use AES Counter mode**

# Lamina

---

## User advertises presence

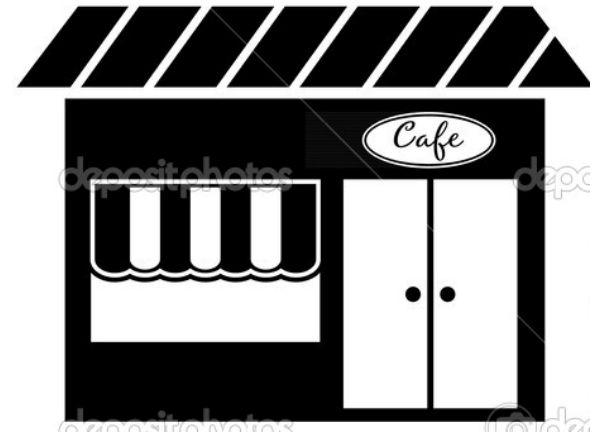


**New MAC for every message**



MAC = Hash(shared nonce,  
**counter**)

Data = (**counter**)



**Use AES Counter mode**

**Self-synchronizing cipher  
maintains encrypted  
channel with packet loss**

# Lamina

## User advertises presence

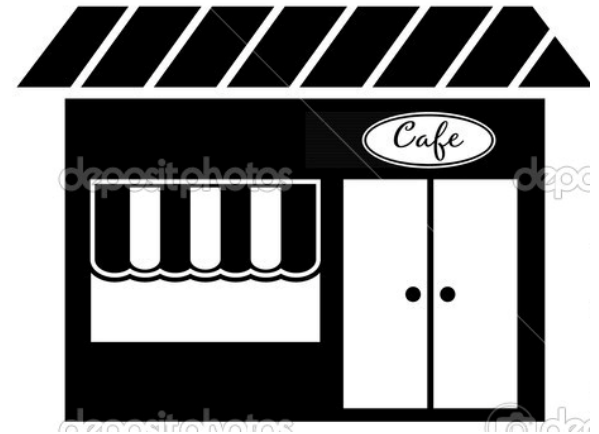


**New MAC for every message**



MAC = Hash(shared nonce,  
counter)

Data = (counter)



**Problem**  
Incremental counters can  
leak information



# Lamina

## User advertises presence

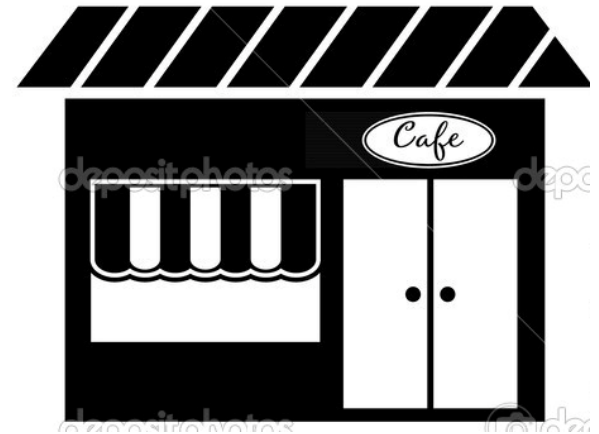


**New MAC for every  
message**



MAC = Hash(shared nonce,  
counter)

Data = (counter)



**Counter does not have to  
be an incremental  
sequence, it just has to  
change**

# Lamina

## User advertises presence

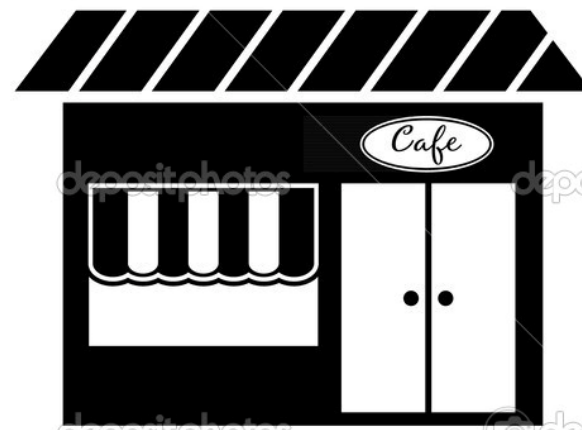


**New MAC for every message**



$\text{MAC} = \text{Hash}(\text{shared nonce}, \text{Hash}(\text{shared nonce}, \text{counter}))$

$\text{Data} = \text{Hash}(\text{shared nonce}, \text{counter})$



**Counter does not have to be an incremental sequence, it just has to change**

**Use shared sequence of randomized counters**



# Lamina

## User advertises presence

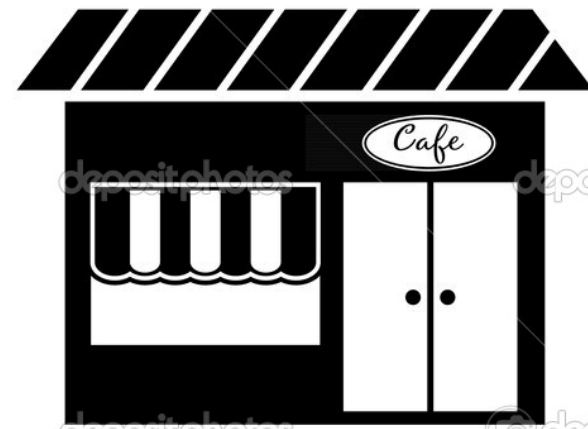


**New MAC for every message**



$MAC = \text{Hash}(\text{shared nonce}, \text{Hash}(\text{shared nonce}, \text{counter}))$

$\text{Data} = \text{Hash}(\text{shared nonce}, \text{counter})$



**Counter does not have to be an incremental sequence, it just has to change**

**Hashed Counter =  $\text{Hash}(\text{shared nonce}, \text{counter})$**





# Lamina

## User advertises presence

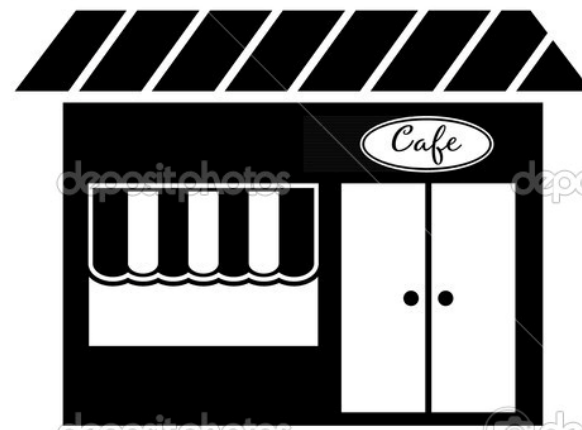


**New MAC for every message**



$\text{MAC} = \text{Hash}(\text{shared nonce}, \text{Hash}(\text{shared nonce}, \text{counter}))$

$\text{Data} = \text{Hash}(\text{shared nonce}, \text{counter})$



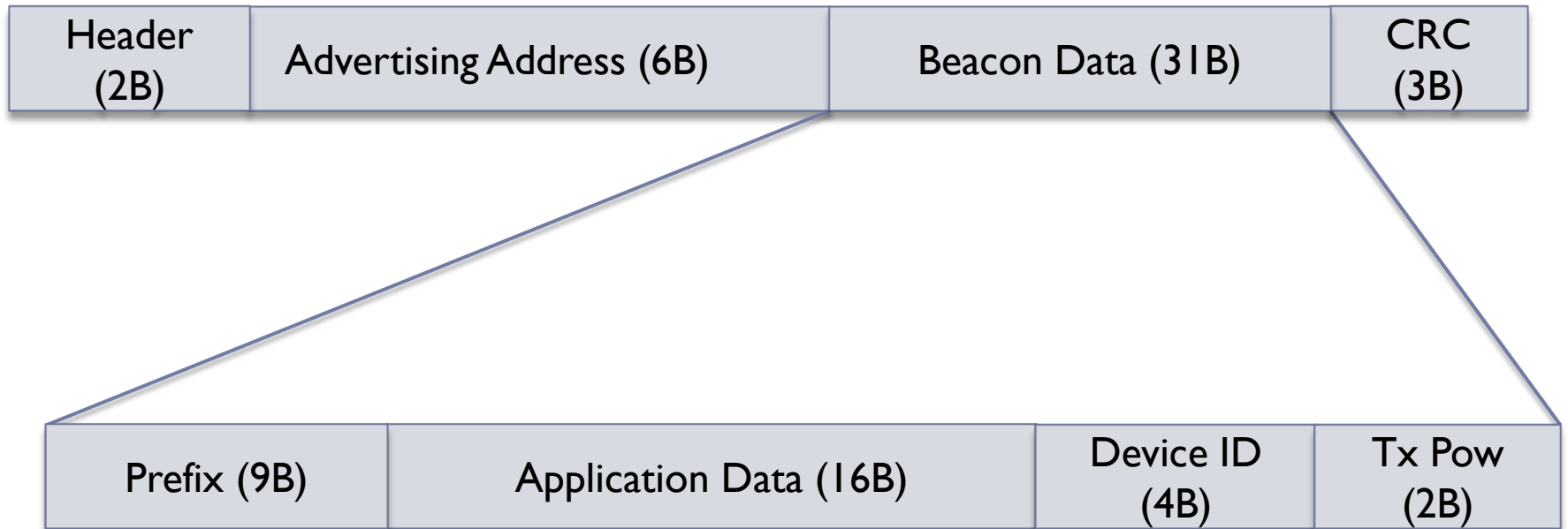
**Counter does not have to be an incremental sequence, it just has to change**

**Counter sequence can be pre-loaded and pre-hashed for low-power and low-computation devices**

# Lamina and BLE

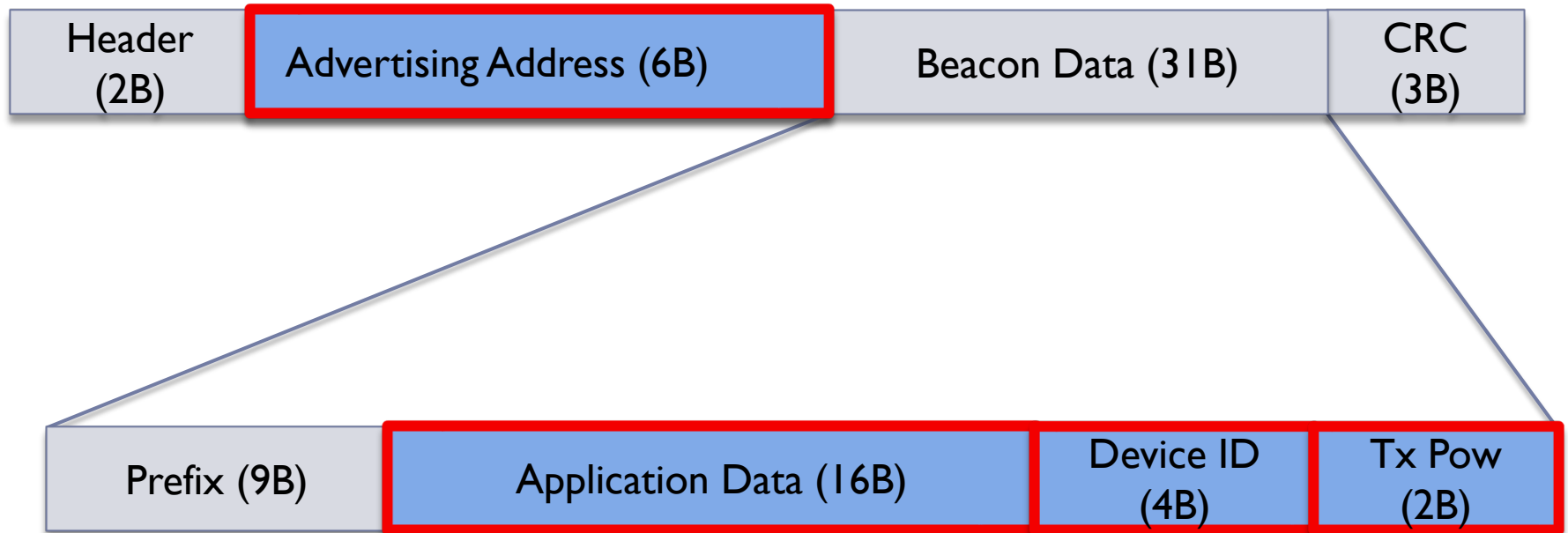
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## BLE Packet Format



# Lamina and BLE

## BLE Packet Format

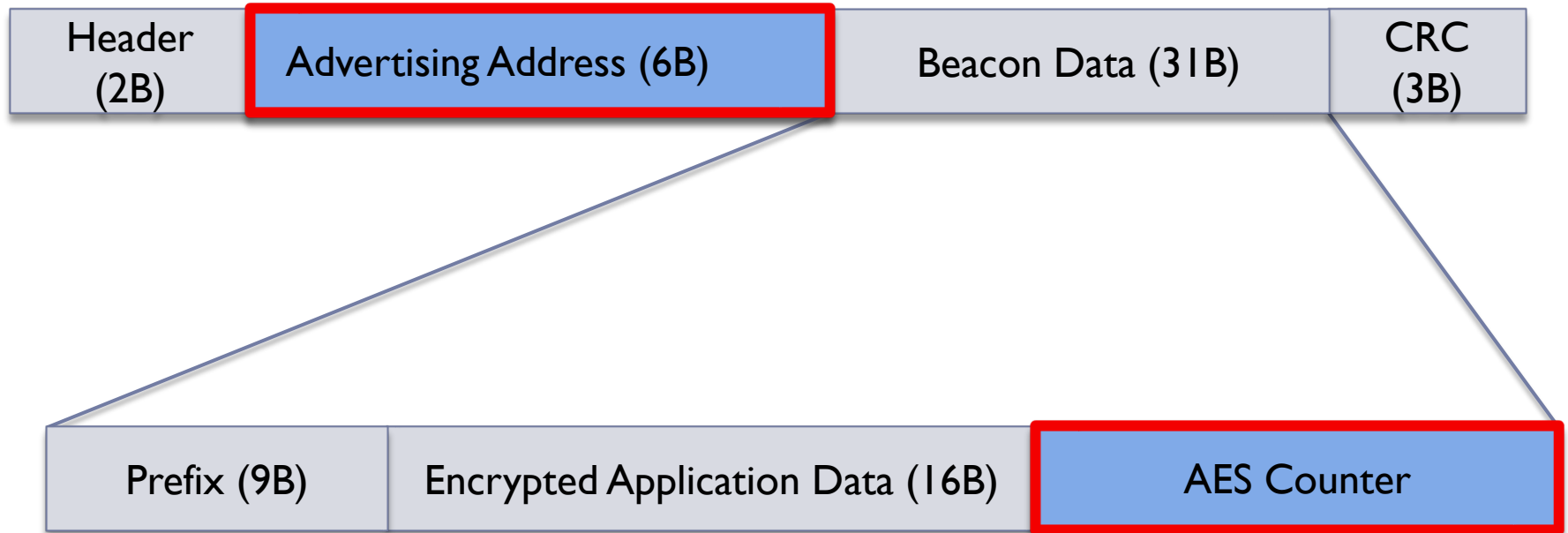


(Potentially privacy-leaking fields)

# Lamina and BLE

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## Lamina BLE Packet Format



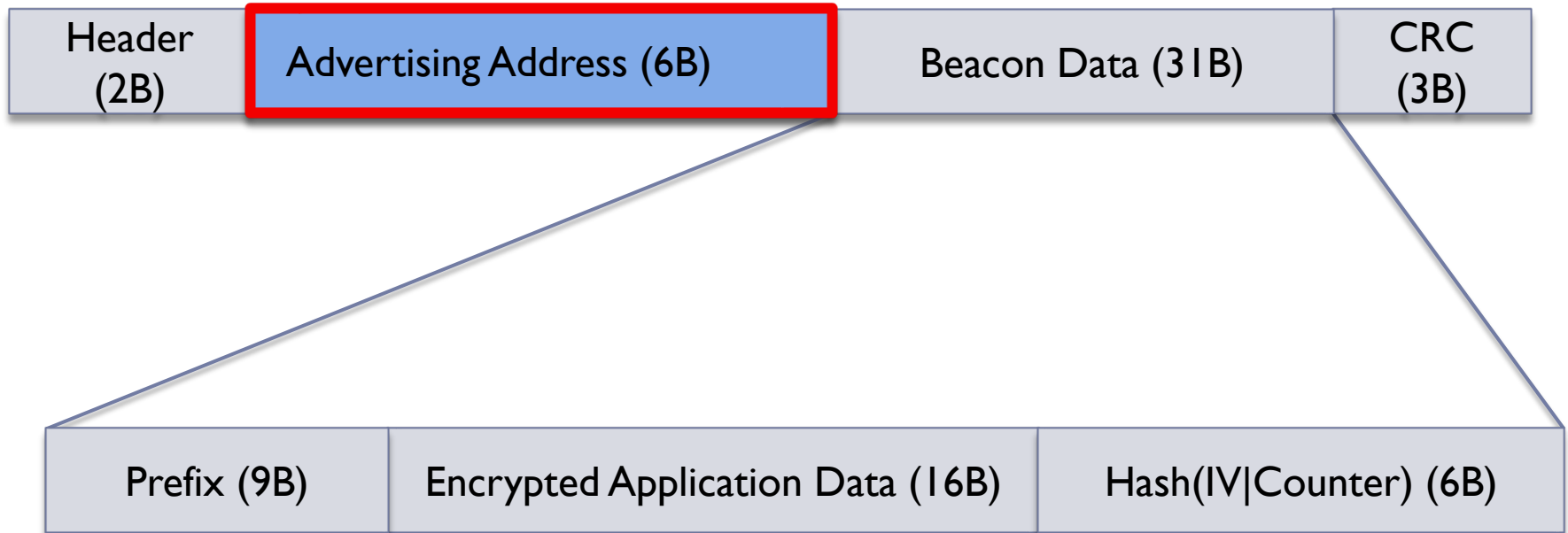
(Potentially privacy-leaking fields)



# Lamina and BLE

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## Lamina BLE Packet Format



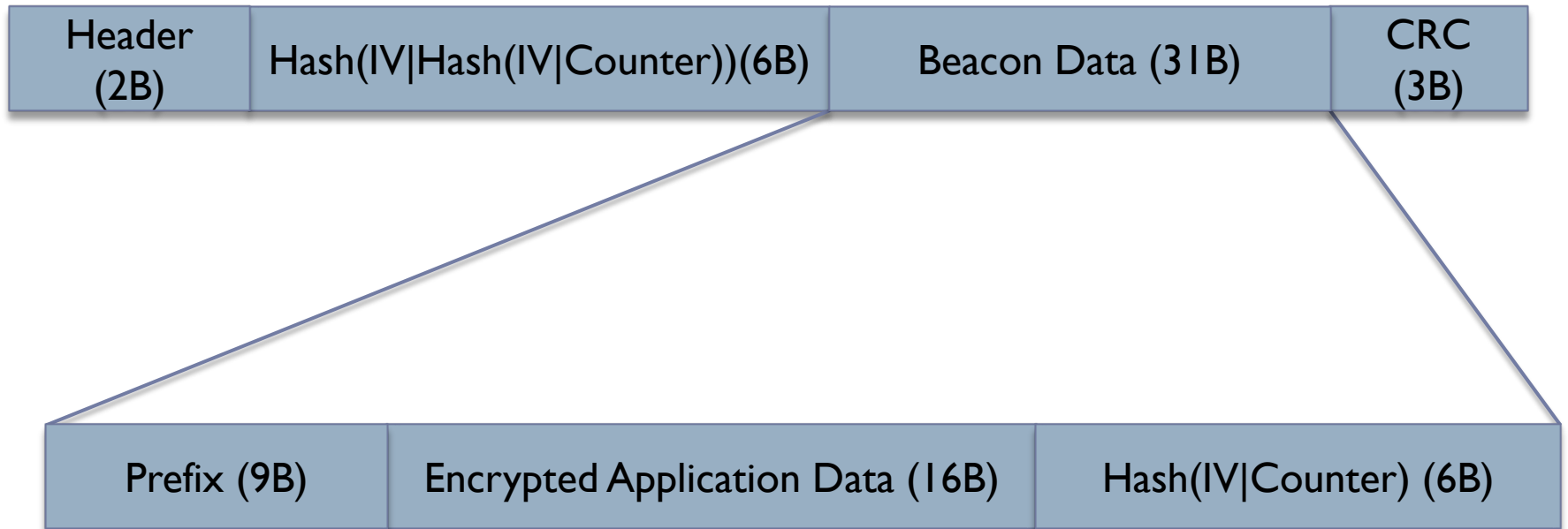
(Potentially privacy-leaking fields)



# Lamina and BLE

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## Lamina BLE Packet Format

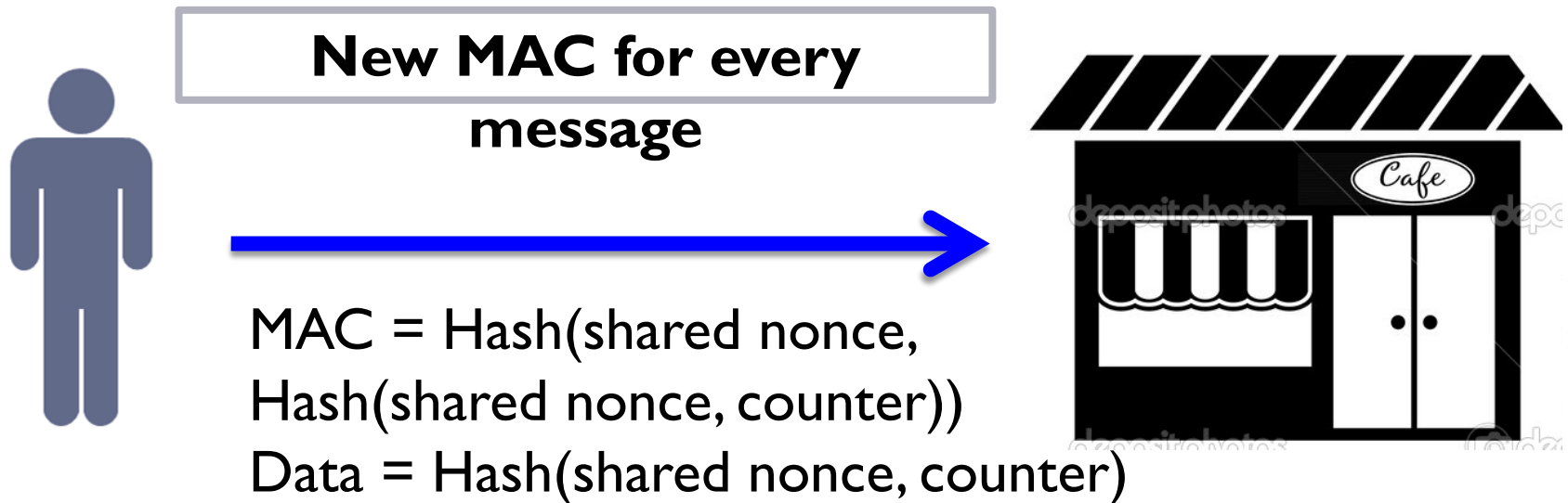


(Potentially privacy-leaking fields)



# Lamina

## User advertises presence



**Unique MAC address  
identifiable by IoT**

**No public key cryptography**

**Loss tolerant**

# Incognito + Lamina

