Computer Science 425 Distributed Systems

CS 425 / CSE 424 / ECE 428 Fall 2012

Indranil Gupta (Indy)
September 25, 2012
Lecture 9
Peer-to-peer Systems I

Reading: Gnutella paper on website

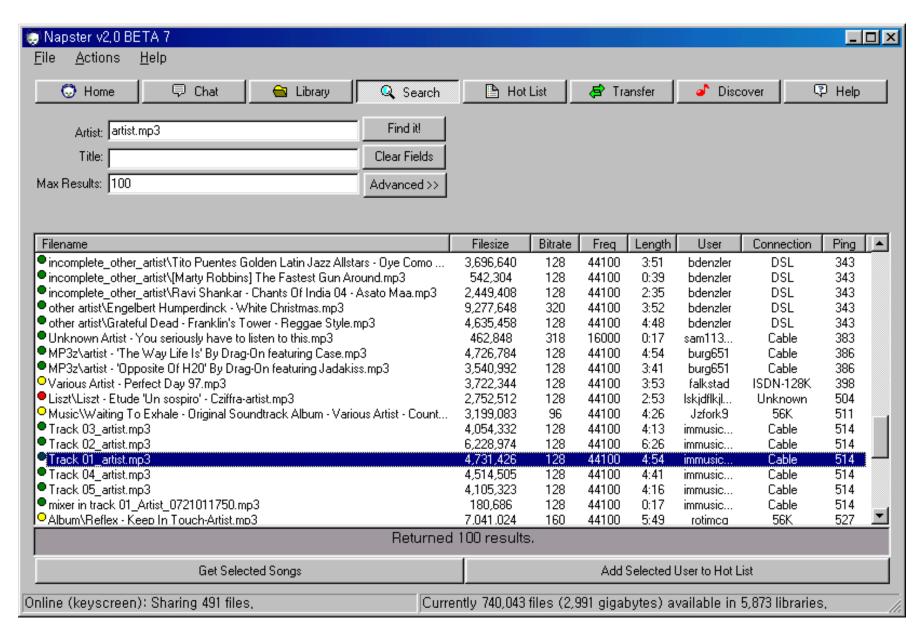
Why Study Peer to peer (P2P) systems?

- To understand how they work
- To understand the **techniques and principles** within them
- To modify, adapt, reuse these techniques and principles in other related areas
 - Cloud computing: key-value stores borrow heavily from p2p systems
 - To build your own p2p system
- To grow the body of knowledge about distributed systems

Some Questions

- Why do people get together?
 - to share information
 - to share and exchange resources they have
 - books, class notes, experiences, videos, music cd's
- How can computers help people
 - find information
 - find resources
 - exchange and share resources

- Existing technologies: The Web!
 - Search engines
 - Forums: chat rooms, blogs, ebay
 - Online business
- But, the web is heavy weight if you want specific resources: say a Beatles' song "PennyLane"
- A search engine will give you their bio, lyrics, chords, articles on them, and then perhaps the mp3
- But you want only the song, nothing else!
- If you can find a peer who has a copy of the Beatles song (mp3), perhaps in exchange for your UIUC Homeocoming videos, that would be great!
 - Napster: a solution light weight that was lighter than the Web



A Brief History

- [6/99] Shawn Fanning (freshman Northeastern U.) releases Napster online music service
- [12/99] RIAA sues Napster, asking \$100K per download
- [3/00] 25% UWisc traffic Napster, many universities ban it
- [00] 60M users
- [2/01] US Federal Appeals Court: users violating copyright laws, Napster is abetting this
- [9/01] Napster decides to run paid service, pay % to songwriters and music companies
- [Today] Napster protocol is open, people free to develop opennap clients and servers

http://opennap.sourceforge.net

Napster Structure Filename Info about Store a directory, i.e., filenames with peer pointers PennyLane.mp3 Beatles, @ 128.84.92.23:1006 napster.com Servers Client machines ("Peers") Store their own files

Napster Operations

Client

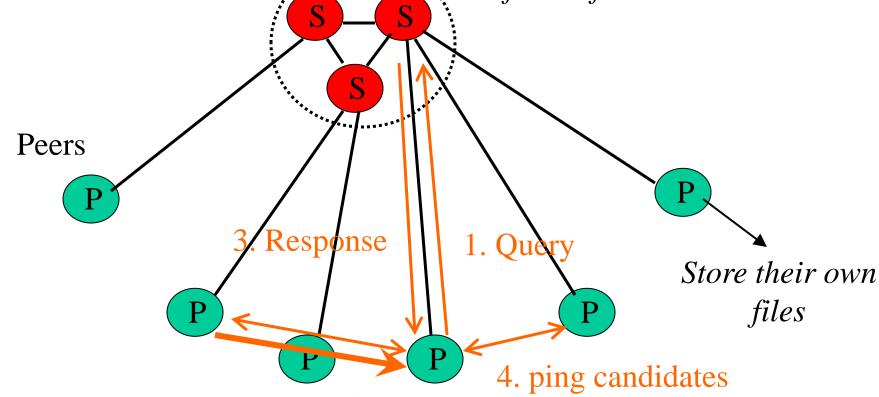
- Connect to a Napster server
- Upload list of music files that you want to share
 - Server maintains list of <filename, ip_address,
 portnum> tuples. Server stores no files.

Search

- Send server keywords to search with
- (Server searches its list with the keywords)
- Server returns a list of hosts <ip_address, portnum> tuples - to client
- Client pings each host in the list to find transfer rates
- Client fetches file from best host
- All communication uses TCP

Napster Search

2. All servers search their lists (ternary tree algo.) Store peer pointers napster.com Servers **▼** for all files



5. download from best host

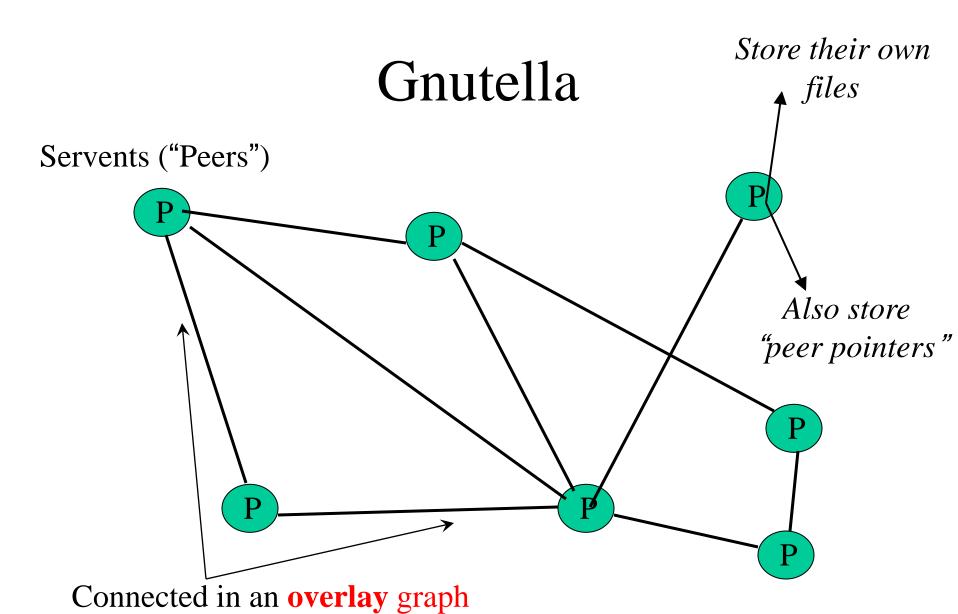
Problems

- Centralized server a source of congestion
- Centralized server single point of failure
- No security: plaintext messages and passwds
- Courts declared napster.com responsible for users' copyright violation
 - "Indirect infringement"

Gnutella

- Eliminate the servers
- Client machines search and retrieve amongst themselves
- Clients act as servers too, called **servents**
- [3/00] release by AOL, 88K users by 3/03
- Original design underwent several modifications
- Available as an open protocol today

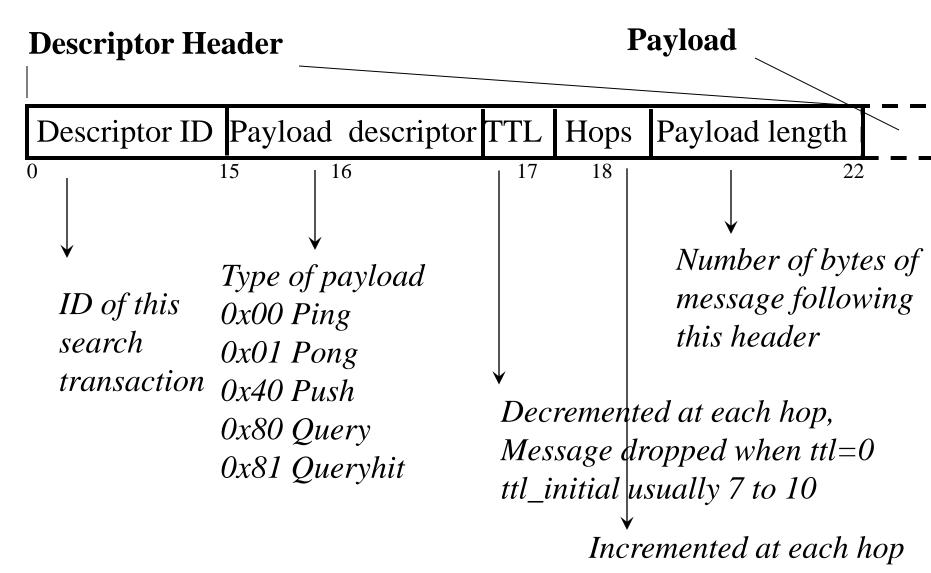
http://www.limewire.com



(== each link is an implicit Internet path)

How do I search for my Beatles file?

- Gnutella *routes* different messages within the overlay graph
- Gnutella protocol has 5 main message types
 - Query (search)
 - QueryHit (response to query)
 - Ping (to probe network for other peers)
 - Pong (reply to ping, contains address of another peer)
 - Push (used to initiate file transfer)
- We'll go into the message structure and protocol now (note: all fields except IP address are in little-endian format)



Gnutella Message Header Format

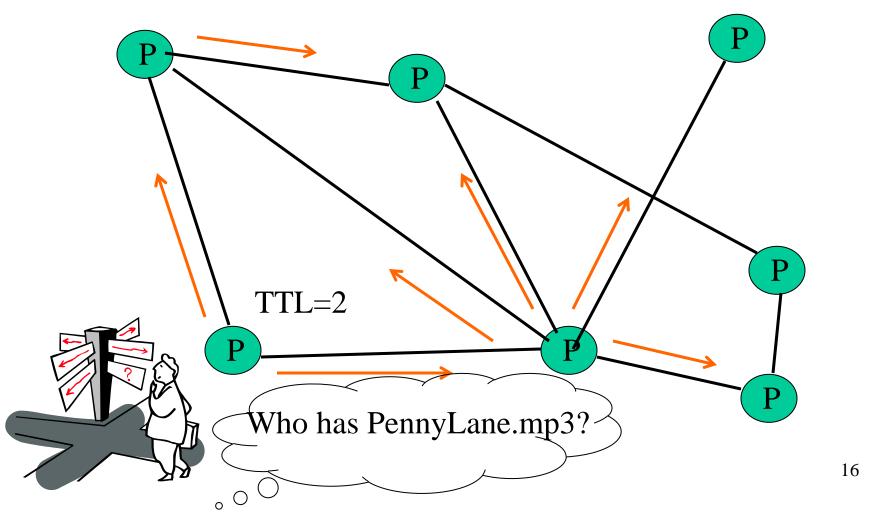
Query (0x80)

Minimum Speed Search criteria (keywords)

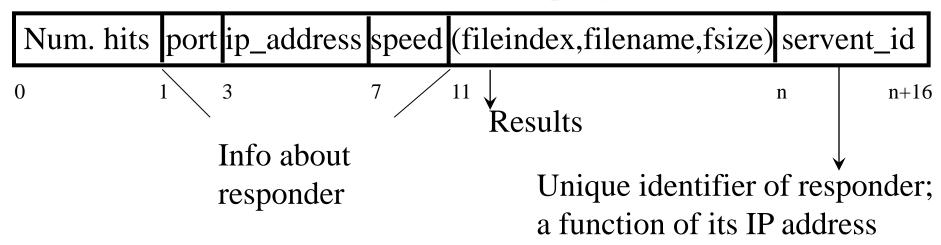
Payload Format in Gnutella Query Message

Gnutella Search

Query's flooded out, ttl-restricted, forwarded only once



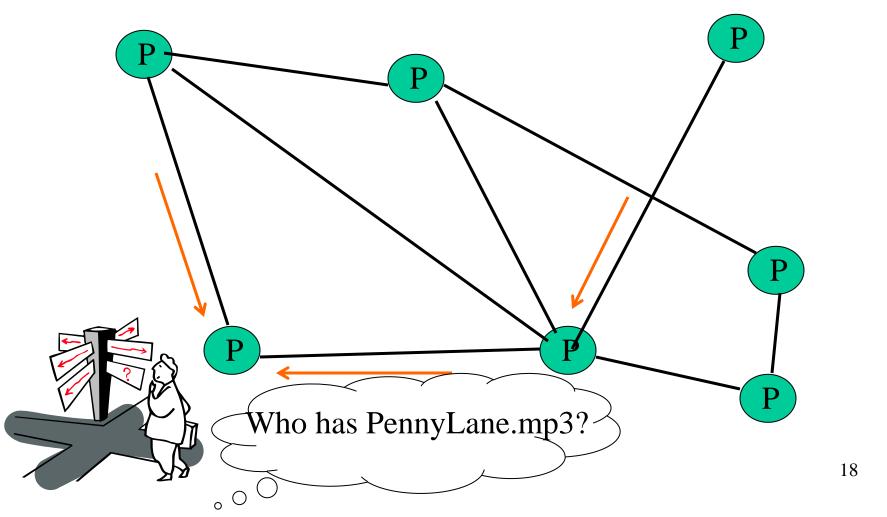
QueryHit (0x81): successful result to a query



Payload Format in Gnutella Query Reply Message

Gnutella Search

Successful results QueryHit's routed on reverse path



Avoiding excessive traffic

To avoid duplicate transmissions, each peer maintains a list of recently received messages

- Query forwarded to all neighbors except peer from which received
- Each Query (identified by DescriptorID) forwarded only once
- QueryHit routed back only to peer from which Query received with same DescriptorID
 - If neighbor does not exist anymore, drop QueryHit
- Duplicates with same DescriptorID and Payload descriptor (msg type) are dropped
- QueryHit with DescriptorID for which Query not seen is dropped

After receiving QueryHit messages

- Requestor chooses "best" QueryHit responder
 - Initiates HTTP request directly to responder's ip+port

```
GET /get/<File Index>/<File Name>/HTTP/1.0\r\n Connection: Keep-Alive\r\n Range: bytes=0-\r\n User-Agent: Gnutella\r\n \r\n
```

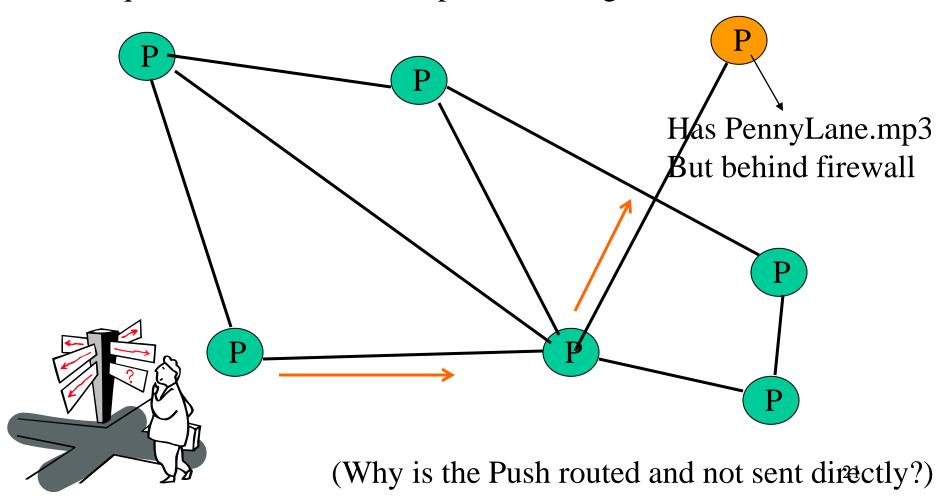
 Responder then replies following start message, followed by packets containing file:

```
HTTP 200 OK\r\n
Server: Gnutella\r\n
Content-type:application/binary\r\n
Content-length: 1024 \r\n
\r\n
```

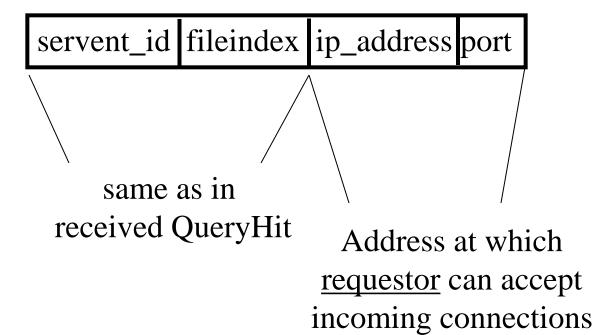
- HTTP is the file transfer protocol. Why?
- Why the "range" field in the GET request?
- What if responder is behind firewall that disallows incoming connections?

Dealing with Firewalls

Requestor sends Push to responder asking for file transfer



Push (0x40)



• Responder establishes a TCP connection at ip_address, port specified. Sends

GIV <File Index>:<Servent Identifier>/<File Name>\n\n

- Requestor then sends GET to responder (as before) and file is transferred
- What if requestor is behind firewall too?
 - Gnutella gives up
 - Can you think of an alternative solution?

Ping-Pong

```
Ping (0x00)
no payload
Pong (0x01)
```

•P2P systems have churn – peers continuously joining, leaving, and failing

Port | ip_address | Num. files shared | Num. KB shared

- Peers initiate Ping's periodically
- •Ping's flooded out like Query's, Pong's routed along reverse path (like QueryHit's)
- Pong replies used to update set of neighboring peers
 to keep neighbor lists fresh in spite of churn

Gnutella Summary

- No servers
- Peers/servents maintain "neighbors", this forms an overlay graph
- Peers store their own files
- Queries flooded out, ttl restricted
- QueryHit (replies) reverse path routed
- Supports file transfer through firewalls
- Periodic Ping-pong to continuously refresh neighbor lists
 - List size specified by user at peer: heterogeneity means some peers may have more neighbors
 - Gnutella found to follow **power law** distribution:

$$P(\#links = L) \sim L^{-k}$$
 (k is a constant)

Problems

- Ping/Pong constituted 50% traffic
 - Solution: Multiplex, cache and reduce frequency of pings/pongs
- Repeated searches with same keywords
 - Solution: Cache Query, QueryHit messages
- Modem-connected hosts do not have enough bandwidth for passing Gnutella traffic
 - Solution: use a central server to act as proxy for such peers
 - Another solution:
 - → FastTrack System (in a few slides)

Problems (contd.)

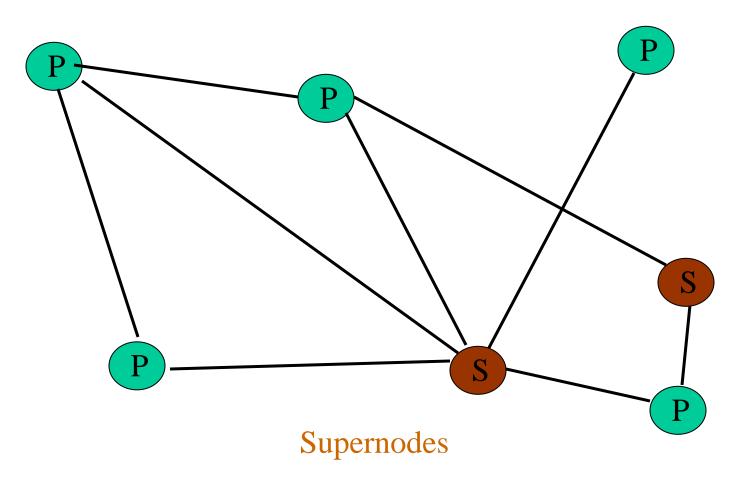
- Large number of *freeloaders*
 - 70% of users in 2000 were freeloaders
 - Only download files, never upload own files
 - Endemic to all p2p systems in deployment
- Flooding causes excessive traffic
 - Is there some way of maintaining meta-information about peers that leads to more intelligent routing?
 - → Structured Peer-to-peer systems e.g., Chord System (next lecture)

FastTrack

- Hybrid between Gnutella and Napster
- Takes advantage of "healthier" participants in the system
- Underlying technology in Kazaa, KazaaLite, Grokster
- Proprietary protocol, but some details available
- Like Gnutella, but with some peers designated as supernodes

A FastTrack-like System

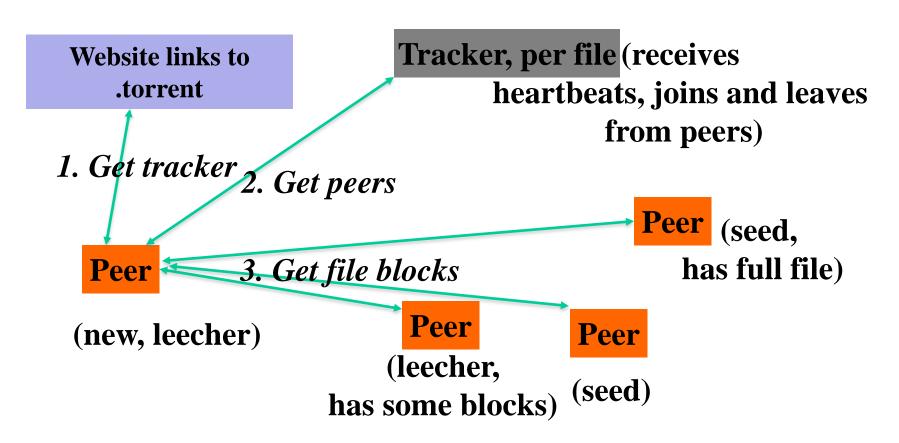
Peers



FastTrack (contd.)

- A supernode stores a directory listing (<filename,peer pointer>), similar to Napster servers
- A peer searches by contacting a nearby supernode
- Supernode membership changes over time
- Any peer can become (and stay) a supernode, provided it has earned enough *reputation*
 - Kazaalite: participation level (=reputation) of a user between 0 and 1000. Initially 10, then affected by length of periods of connectivity and total number of uploads.
 - More sophisticated Reputation schemes invented, especially based on economics

BitTorrent – A Quick Overview



BitTorrent – A Quick Overview (2)

- File split into blocks (32 KB 256 KB)
- Download Local Rarest First block policy: prefer early download of blocks that are least replicated among neigh bors
 - Exception: New node allowed to pick one random neighbor: helps in bootstrapping
- Tit for tat bandwidth usage: Provide blocks to neighbors that provided it the best download rates
 - Incentive for nodes to provide good download rates
 - Seeds do the same too
- Choking: Limit number of neighbors to which concurrent uploads <= a number (5), i.e., the "best" neighbors
 - Everyone else choked
 - Periodically re-evaluate this set (e.g., 10 s)
 - Optimistic unchoke: periodically (e.g., ~30 s), unchoke a random neigbhor – helps keep unchoked set fresh

Wrap-up Notes

Applies to all p2p systems

- How does a peer join the system
 - Send an http request to well-known url for that P2P
 service http://www.myp2pservice.com
 - Message routed (after DNS lookup) to a well known server which then initializes new peers' neighbor table
 - Server only maintains a partial list of online clients
- Lookups can be speeded up by having each peer cache:
 - Queries and their results that it sees
 - All directory entries (filename,host) mappings that it sees
 - The files that pass through it

Summary

- Napster: protocol overview, more details available on webpage
- Gnutella protocol
- FastTrack protocol
- Protocols continually evolving, software for new clients and servers conforming to respective protocols: developer forums at
 - Napster: http://opennap.sourceforge.net
 - Gnutella: http://www.limewire.com
- Others
 - Peer to peer working groups: http://p2p.internet2.edu

For Next Lecture

- Read "Chord" paper from website
 - Sections 1-4, 6-7

- MP2 and HW2 out
 - By now, you should have an initial design for MP2.