Computer Science 425 Distributed Systems

CS 425 / CSE 424 / ECE 428

> Klara Nahrstedt August 25, 2009

Acknowledgement

- The slides during this semester are based on ideas and material from the following sources:
 - Slides prepared by Professors M. Harandi, J. Hou, I. Gupta, N. Vaidya, Y-Ch. Hu, S. Mitra.
 - Slides from Professor S. Gosh's course at University o lowa.

Course Staff – how to meet us

In person (Office Hours):

- Klara Nahrstedt: 3104 Siebel Center
 - Every Wednesday 9-10am, and Every Thursday 9-10am
- TA, Ying Huang: 207 Siebel Center
 - Every Monday: 3-4pm, and Every Thursday 3:30-4:30pm

Virtually:

- Newsgroup: class.cs425 (most preferable, monitored daily)
- Email (turnaround time may be longer than newsgroup)
 - Klara Nahrstedt: <u>klara@illinois.edu</u>
 - Ying Huang: huang23@illinois.edu

How will you Learn?

Take a look at handout "Course Information and Schedule"

- http://www.cs.uiuc.edu/class/fa09/cs425
- Text: Colouris, Dollimore and Kindberg (4th edition)
- Lectures
- Homework Sets (4 Sets)
 - Approx. one every two weeks
 - Solutions need to be typed, figures can be hand-drawn
- Programming assignments (3 MPs)
 - Incremental, in 3 stages
 - We will build a peer to peer system for mobile phones
 - Form a group up to 3 members for each programming assignment
 - Programming Language needed: Java
- Exams
 - Midterm Tuesday, October 13, 2-3:15pm in class
 - Final Exam Wednesday, December 16, 7-10pm, location TBD

On the Textbook

- Text: Colouris, Dollimore and Kindberg (4th edition). "White book".
- The 3rd edition will suffice for most material too. However, we will refer to section, chapter, and problem numbers only in the 4th edition.
 - The 3rd edition may have a different numbering for some HW problems (that we give from the textbook). Make sure you solve the right problem; the responsibility is yours (no points for solving the wrong problem!)

What assistance is available to you?

Lectures

- lecture slides will be placed online at course website
 - » "Tentative" version before lecture
 - » "Final" version after lecture
- Homework office hours to help you (without giving you the solution)
 - 2 Homework Assignments before Midterm
 - 2 Homework Assignments before Final
- Course Prerequisite: Operating Systems/Systems
 Programming (CS 241 or CS 423 or instructor permission)
 and Java programming language
- Compass tool as grade-book where all grades will be posted.

Programming Assignments

Server

MP1:

Peer Registration to Server P2P Socket Programming Peer-Server Socket Programming GUI Design (register, view other Members send/receive msg)

We will use Eclipse Emulator

MP2·

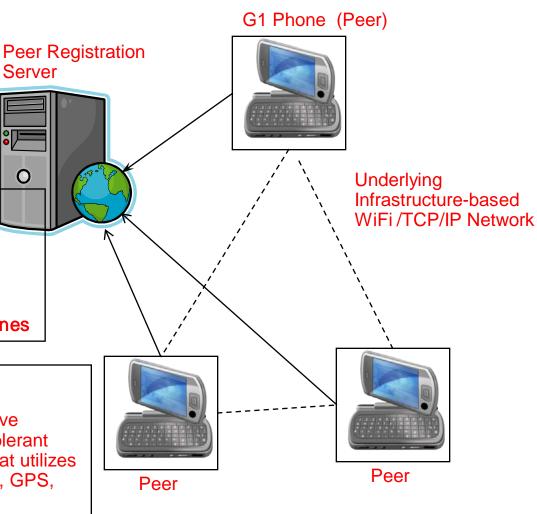
Build P2P File Sharing with Membership Service, Functions: File Search, Insert, Delete, Browse

We will use Eclipse Emulator and G1 Phones

MP3:

Design and Implement a P2P application (e.g., find a group member, SVN: cooperative Editing, chat: group conversation log, fault-tolerant P2P file sharing using coding, application that utilizes G1 phone hardware such as WiFi, Bluetooth, GPS, accelerometers

We will use Eclipse and G1 Phones



Lecture 1-7

Course Logistics

- Exams, Homework and Machine Problems (Programming Assignments) deadlines on the website
 - http://www.cs.uiuc.edu/class/fa09/cs425/lectures.html
- Grade Distribution:
 - Homework: 20% (each HW 5% of the grade)
 - Programming Assignments: 40% (MP1:10%, MP2: 16%, MP3: 14%)
 - Midterm: 10%
 - Final Exam: 30% (comprehensive)

To Do List

- Create groups of 2-3 students for machine problems programming assignments (use newsgroup, TA to find a group member)
- Deadline for group creation September 1
 (Tuesday) email TA the names of your group members
 - Based on the group names TSG will create group directories where the group can build the MPs
- Check carefully the class website will all the deadlines for homework, machine problems, exams

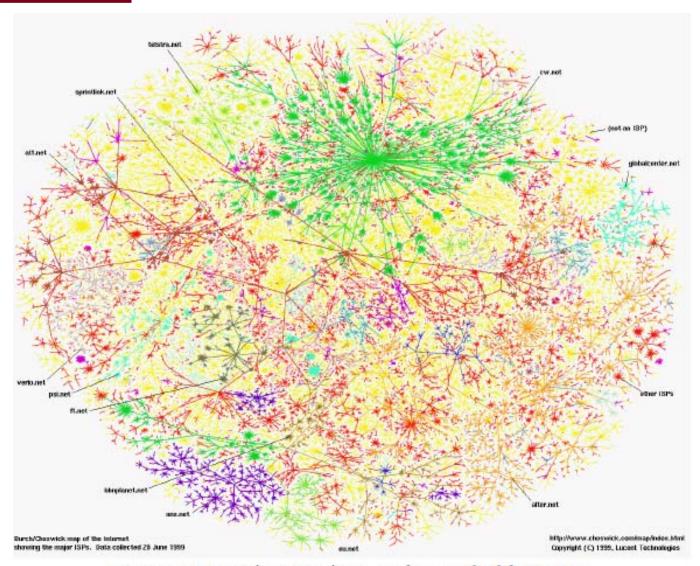
Our Only Goal Today

To Define the Term Distributed System	

Can you name some examples of Distributed Systems?

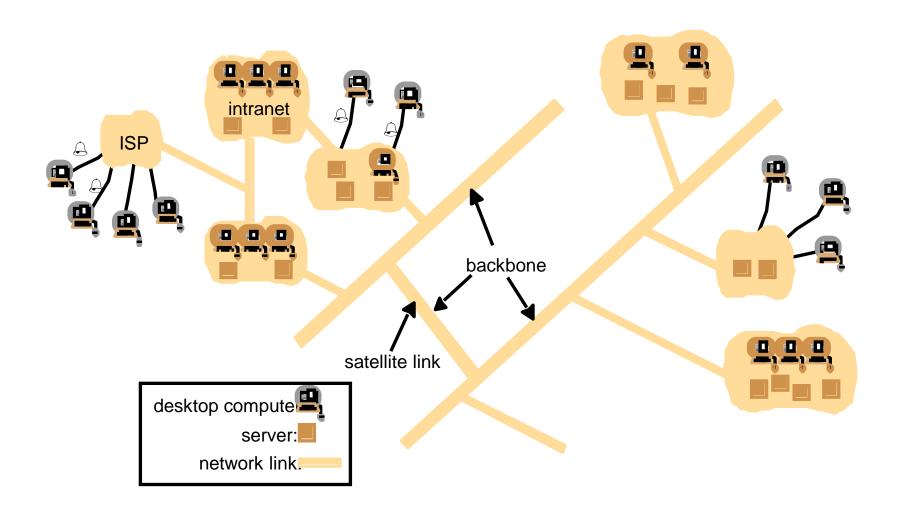


1. Internet

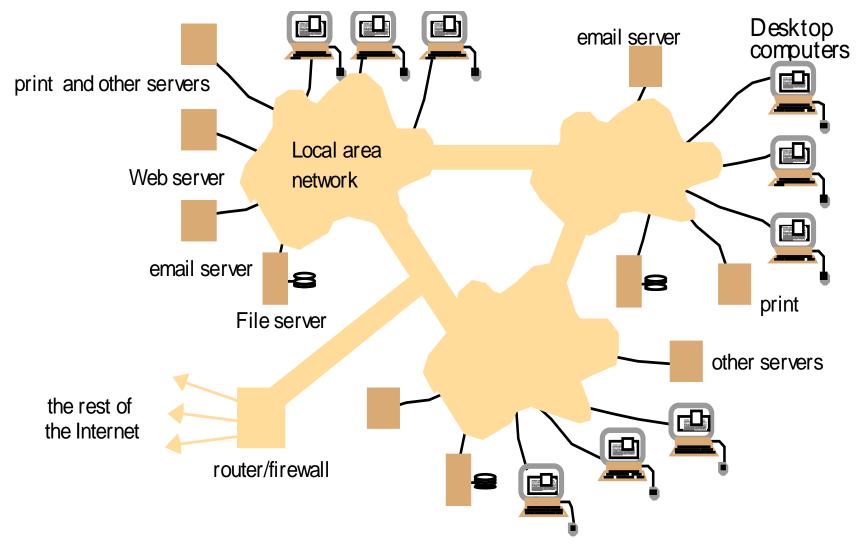


Internet Mapping Project, color coded by ISPs

A typical portion of the Internet



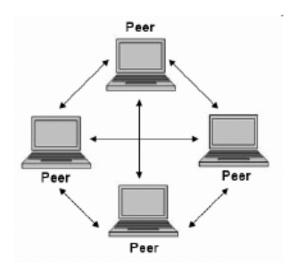
A typical intranet



2. Peer-to-Peer Overlays



client-server e.g., NFS



peer-to-peer e.g., KaZaa, napster



3. Sensor Network



4. Distributed Mobile Robots



Kiva robot

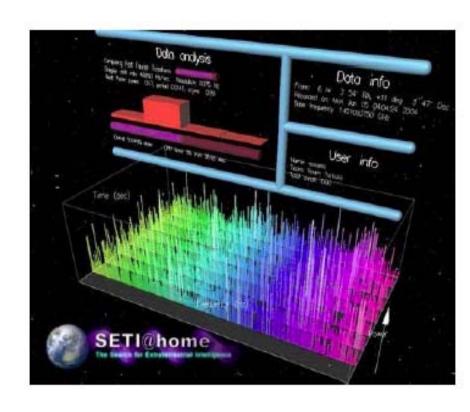
500 robots coordinate to manage inventory in a warehouse





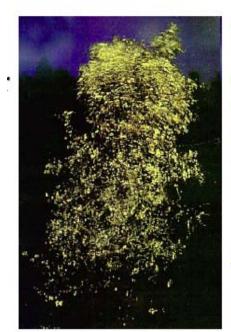
5. Distributed Computation Grids

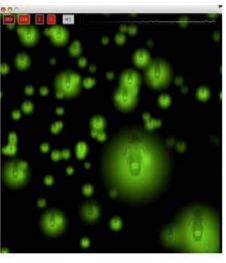
- Seti@Home
- Uses Internet-connected computers in the Search for Extraterrestrial Intelligence
- Downloads and analyzes telescope data chunks on your computer and sends back results to Seti servers

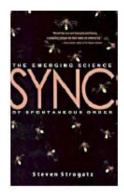


6. Distributed Structures in Nature

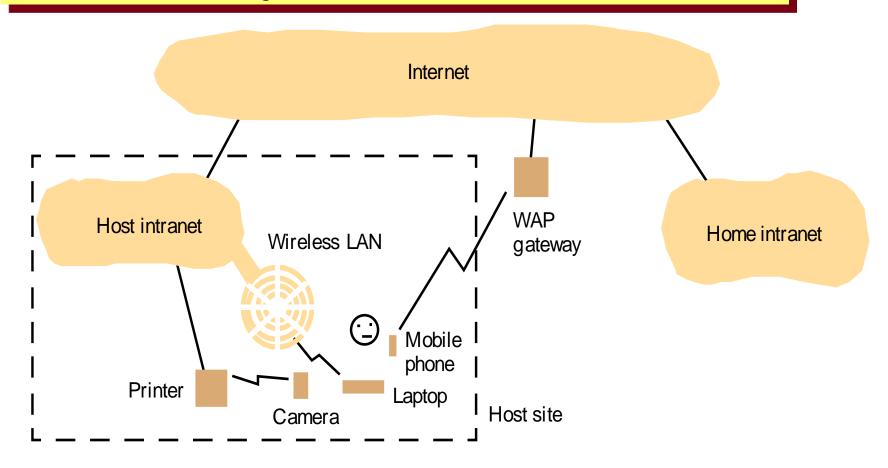
- Tropical Fireflies synchronize their flashes precisely among large groups
- Example of biological synchronicity







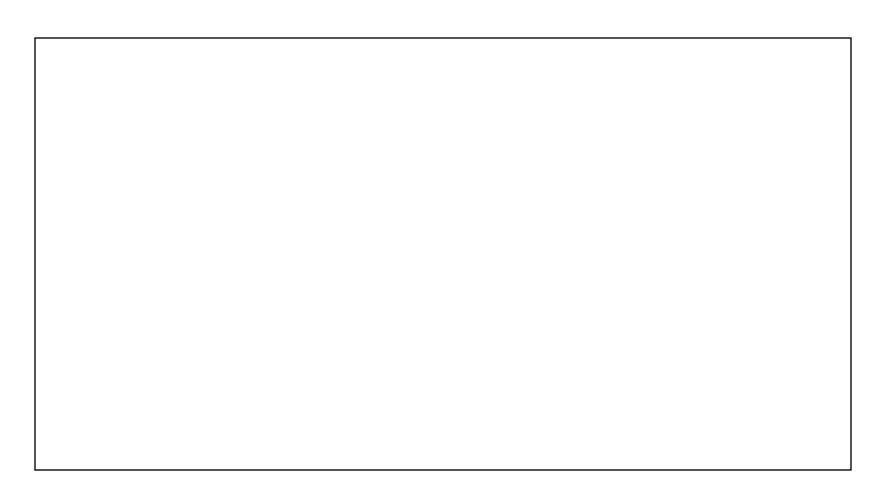
7. Portable and handheld devices in a distributed system



Can you name some examples of Distributed Systems?

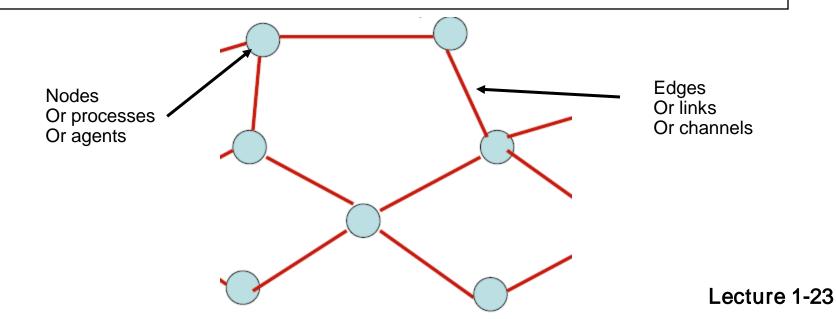
- Client-Server
- Web
- Internet
- Mobile Cell-Phone Network
- Sensor Network
- Network of Robots
- DNS (Domain Name Service)
- Napster, KaZaA (peer to peer file sharing overlayss)
- Synchronized Fireflies
- Computing Grid

What is a Distributed System?



What is a Distributed System?

- It is s collection of entities (computers, robots, fireflies,..) where
 - Each of them is autonomous, asynchronous and [possibly] failure-prone
 - Communicating through [possibly] unreliable channels (Ethernet, wireless network, vision,...)
 - To perform some common function (detection, computation, communication)



Why Distributed Systems?

- Geographic distributed of processes
- Resource sharing (as used in P2P networks)
- Computing speed up (as applied in grid and clusters)
 - MapReduce distributed computing framework to parallelize operation on large datasets (webpages)
 - Seti@Home
- Fault tolerance (as applied in Google Distributed File System (GFS)
 - On top of GFS is BigTable (Google Distributed Database)

CS 425

- CS 425 is an introduction to the key concepts for
 - Designing
 - Analyzing and
 - Implementing

Distributed Systems

Key Design Issues

- Knowledge is local
- Clocks are not synchronized
- No shared address space
- Topology and routing
- Scalability and adaptability
- Faulty nodes and channels
 - Processes may crash, processes may misbehave
 - Variable latency and bandwidth

Fundamental Problems in Distributed Systems and their Analysis

- Time synchronization
- Leader election
- Mutual exclusion
- Distributed snapshot
- Routing
- Consensus
- Replica management
- Transactions

We will study and analyze algorithms for these types of problems

Approaching Implementation of Distributed Systems

- In implementing distributed systems we have to know the nature of underlying computing and networking infrastructure
- What can we assume?
 - How large are message delays and latencies?
 - Are processor clocks synchronized?
 - Do processors and links fail?
 - Could processes become malicious?
 - Are channels secure from eavesdropping and corruption?
 - What are the allowed primitive HW/OS operations? What are the atomic operations?
- Different answers give rise to different models for a distributed system
- We will study algorithms for fundamental problems in the context of certain models

Lecture 1-28

Problems and Models

- Time synchronization
- Leader election
- Mutual exclusion
- Distributed snapshot
- Routing
- Consensus
- Replica management
- Transactions

- Message Passing Model
 - Synchronous (with bounded delays)
 - Asynchronous (with unbounded delays)
 - Partially synchronous (with bounded but unknown delays)
- Shared memory model

Example: Simple Communication Model Message Passing

- Assume distributed system topology as a graph G=(V,E)
 - V = set of nodes (sequential processes)
 - E = set of edges (links or channels which could unidirectional or bidirectional)
- Actions by a process
 - Internal action: sequential process computes
 - Send action: sends a message (put a message in channel) and perform computation
 - Receive action: receives a message (message taken out from channel) and performs computation

Weak vs Strong Models

- One object (or operation) of a strong model - more than one objects(or operations) of a weaker model
- Often, weaker models are synonymous with fewer restrictions
- One can add layers (additional restrictions) to create a stronger model from weaker one

- High Level Language model is stronger than assembly language model
- Asynchronous is weaker than synchronous model
- Bounded delay is stronger than unbounded delay (channel)

Model Transformation and Implementation

- Stronger models
 - Simplify reasoning, but
 - Need extra work to implement
- Weaker models
 - Are easier to implement
 - Have a closer relationship with real world

- Can model X be implemented using model Y?
- Examples
 - Transformation of message passing to share memory
 - Transformation non-FIFO to FIFO channel

Distributed Systems Design Goals

Common Goals:

- Heterogeneity can the system handle different types of PCs and devices?
- Robustness is the system resilient to host crashes and failures, and to the network dropping messages?
- Availability are data, services always there for clients?
- Transparency can the system hide its internal workings from the users?
- Concurrency can the server handle multiple clients simultaneously?
- Efficiency is it fast enough?
- Scalability can it handle 100 million nodes?
 (nodes=clients and/or servers)
- Security can the system withstand hacker attacks?
- Openness is the system extensible?

Readings

For today's lecture: Chapter 1

- For next Thursday's lecture:
 - Read sections 11.1-11.4
 - Fill out and return Student InfoSheet

Additional Slides



Heterogeneity

Variety and differences among

- Networks, hardware, Operating Systems
- Programming languages
- Implementations
- Middleware software layer that provides programming abstract as well as
 - Masking heterogeneity of the underlying networks, hardware, operating systems and programming languages
 - Providing uniform computational models, e.g., remote object invocation, remote event notification

Virtual machine

 Compiler for a particular language generates code for VMs, instead for particular hardware code to assist execution of mobile code

Openness

- Degree to which new resource-sharing services can be added and used
- Requires publication of interfaces for access to shared resources
- Requires uniform communication mechanism
- Conformance of each component to the published standard must be tested and verified

Security

- Security has three components
 - Confidentiality (protection against disclosure to unauthorized individuals)
 - Integrity (protection against alternation or corruption)
 - Availability (protection against interference with means to access the resources)
- Two security challenges (examples):
 - Denial of service attacks
 - Mobile code security

Scalability

- System is said to be scalable if it will remain effective when there is a significant increase in the number of resources and users
 - Controlling cost of resources
 - Controlling performance loss
 - Preventing software resources running out (e.g., IP addresses)
 - Avoiding performance bottlenecks

Failure Handling

- Failure in DS is partial some components fail while the rest is functional
 - Detecting failures (remote site crash or delay in message transmission)
 - Masking failures (message retransmission, file replication)
 - Tolerance for failure (clients give up after a predetermined number of attempts and take other actions)
 - Failure recovery (check-point and rollback recovery)
 - Redundancy (multi-path routing, replicated database, replicated DNS)

Concurrency

- Is a problem when two or more users access to the same resource by two at the same time
 - Each resource is encapsulated as an object and invocations are executed in concurrent threads
 - Concurrency can be maintained by use of semaphores and other mutual exclusion mechanisms

Transparency

- Concealment of the separation of components from users:
 - Access transparency: local and remote resources can be accessed using identical operations
 - Location transparency: resources can be accessed without knowing their whereabouts
 - Concurrency transparency: processes can operate concurrently using shared resources without interferences
 - Failure transparency: faults can be concealed from users/applications
 - Mobility transparency: resources/users can move within a system without affecting their operations
 - Performance transparency: system can be reconfigured to improve performance
 - Scaling transparency: system can be expanded in scale without change to the applications

Transparency Examples

- Distributed File System allows access transparency and location transparency
- URLs are location transparent, but are not mobility transparent
- Message retransmission governed by TCP is a mechanism for providing failure transparency
- Mobile phone is an example of mobility transparency