# CS 421 — Small Step Semantics Activity

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## The Rules

$$<\operatorname{skip},\sigma>\to < E,\sigma> \\ < u:=t,\sigma>\to < E,\sigma[u:=\sigma(t)]> \\ \frac{< S_1,\sigma>\to < S_2,\tau>}{< S_1;S,\sigma>\to < S_2;S,\tau>} \\ E;S\equiv S \\ <\operatorname{if}\ B\ \operatorname{then}\ S_1\ \operatorname{else}\ S_2\ \operatorname{fi}\ ,\sigma>\to < S_1,\sigma> \ \operatorname{where}\ \sigma\models B \\ <\operatorname{if}\ B\ \operatorname{then}\ S_1\ \operatorname{else}\ S_2\ \operatorname{fi}\ ,\sigma>\to < S_2,\sigma> \ \operatorname{where}\ \sigma\models \neg B \\ <\operatorname{while}\ B\ \operatorname{do}\ S_1\ \operatorname{od}\ ,\sigma>\to < S_1; \ \operatorname{while}\ B\ \operatorname{do}\ S_1\ \operatorname{od}\ ,\sigma> \ \operatorname{where}\ \sigma\models B \\ <\operatorname{while}\ B\ \operatorname{do}\ S_1\ \operatorname{od}\ ,\sigma>\to < E,\sigma> \ \operatorname{where}\ \sigma\models \neg B \\ <\operatorname{while}\ B\ \operatorname{do}\ S_1\ \operatorname{od}\ ,\sigma>\to < E,\sigma> \ \operatorname{where}\ \sigma\models \neg B \\ <\operatorname{while}\ B\ \operatorname{do}\ S_1\ \operatorname{od}\ ,\sigma>\to < E,\sigma> \ \operatorname{where}\ \sigma\models \neg B \\ <\operatorname{hile}\ B\ \operatorname{do}\ S_1\ \operatorname{od}\ ,\sigma>\to < E,\sigma> \ \operatorname{where}\ \sigma\models \neg B \\ <\operatorname{hile}\ S_1\ \operatorname{od}\ ,\sigma>\to < E,\sigma> \ \operatorname{where}\ \sigma\models \neg B \\ <\operatorname{hile}\ S_1\ \operatorname{od}\ ,\sigma>\to < E,\sigma> \ \operatorname{where}\ \sigma\models \neg B \\ <\operatorname{hile}\ S_1\ \operatorname{od}\ ,\sigma>\to < E,\sigma> \ \operatorname{where}\ \sigma\models \neg B \\ <\operatorname{hile}\ S_1\ \operatorname{od}\ ,\sigma>\to < E,\sigma> \ \operatorname{where}\ \sigma\models \neg B \\ <\operatorname{hile}\ S_1\ \operatorname{od}\ ,\sigma>\to < E,\sigma> \ \operatorname{where}\ \sigma\models \neg B \\ <\operatorname{hile}\ S_1\ \operatorname{od}\ ,\sigma>\to < E,\sigma> \ \operatorname{where}\ \sigma\models \neg B \\ <\operatorname{hile}\ S_1\ \operatorname{od}\ ,\sigma>\to < E,\sigma> \ \operatorname{where}\ \sigma\models \neg B \\ <\operatorname{hile}\ S_1\ \operatorname{od}\ ,\sigma>\to < E,\sigma> \ \operatorname{hile}\ S_1\ \operatorname{od}\ ,\sigma>\to < E,\sigma> \ \operatorname{hile}\ S_1\ \operatorname{od}\ ,\sigma>\to < E,\sigma> \ \operatorname{hile}\ S_1\ \operatorname{od}\ ,\sigma>$$

## Reductions

Reduce the following programs according to the semantic rules given.

### Problem 1)

< if x>y then m:=x else skip fi; if x<y then m:=y else skip fi;,  $\{x:=10,y:=30\}$ >

### Problem 2)

```
n := 0; while x > 1 do x := x/2; n := n+1 od, \{x := 8\}
```

#### Problem 3)

(Don't spend too much time on this one.)

```
p:=1; n:=3; while n>1 do p:=p*x od, {x:=3}>
```

# Make your own rules!

## Problem 4)

Write a rule to explain the when  $\, \, {\tt B} \, \, {\tt S} \, \, {\tt Statement}. \, {\tt It} \, \, {\tt executes} \, S \, {\tt only} \, {\tt if} \, B \, {\tt is} \, \, {\tt true}.$ 

## Problem 5)

Write a rule for do S while B od. It is like while, but executes S at least one time.

## Church Rosser

**Problem 6)** Consider this semantic rule:

$$x_1 \circ x_2 \circ \cdots \circ x_i \circ x_{i+1} \circ \cdots x_n \to x_1 \circ x_2 \circ \cdots \circ (x_i * x_{i+1}) \circ \cdots x_n$$

Does it have the Church-Rosser property? Try to prove it.

**Problem 7)** Consider this semantic rule:

$$x_1 \circ x_2 \circ \cdots \circ x_i \circ x_{i+1} \circ \cdots x_n \to x_1 \circ x_2 \circ \cdots \circ (x_i - x_{i+1}) \circ \cdots x_n$$

Does it have the Church-Rosser property? Try to prove it.