
CS 421 — LL Grammar Problems

Question 1: What advantage does it give the programmer if a grammar turns out to be LL?

Question 2: What are the two conditions that will cause a grammar to **not** be LL?

Question 3: Consider the following grammar.

$$\begin{array}{l} S \rightarrow x E a \\ \quad | y E z \\ E \rightarrow x F \\ \quad | x E q \\ F \rightarrow q \\ \quad | F z \end{array}$$

This grammar is not LL. There are two reasons for that. What are they?

Question 4: Convert the following grammar into an equivalent LL grammar.

$$\begin{array}{l} S \rightarrow x E a \\ \quad | y E z \\ E \rightarrow x F \\ \quad | x E q \\ F \rightarrow q \\ \quad | F z \end{array}$$

Question 5: Convert the following non-LL grammar into an equivalent LL grammar.

$$\begin{array}{l} S \rightarrow S x \\ \quad | a E \\ E \rightarrow y a y \\ \quad | y a z \end{array}$$

Question 6: Convert the following non-LL grammar into an equivalent LL grammar.

$$\begin{array}{l} S \rightarrow y y \\ \quad | y E \\ E \rightarrow E x \\ \quad | x a \end{array}$$

Question 7: Convert the following non-LL grammar into an equivalent LL grammar.

$$\begin{array}{l} S \rightarrow S x \\ \quad | a E \\ E \rightarrow x y \\ \quad | x z \end{array}$$

Question 8: Convert the following non-LL grammar into an equivalent LL grammar.

$$\begin{array}{l} S \rightarrow S x \\ \quad | x E \\ E \rightarrow y y \\ \quad | y z \end{array}$$

Question 9: Convert the following non-LL grammar into an equivalent LL grammar.

$$\begin{array}{l} S \rightarrow y E \\ \quad | y z \\ E \rightarrow E z z \\ \quad | x a a \end{array}$$

Question 10: Convert the following non-LL grammar into an equivalent LL grammar.

$$\begin{array}{l} S \rightarrow S x \\ \quad | a E \\ E \rightarrow z b b \\ \quad | z b z \end{array}$$

1 Solutions to exercises

Solution 1 If the grammar is LL, then we can write a parser for it very simply using recursive descent.

Solution 2

1. two rules for the same symbol that have overlapping first sets (The “common prefix” problem.)
2. a left recursive rule

Solution 3 The E productions share a common prefix, x ; and one of the F rules is left recursive.

Solution 4

$$\begin{array}{lcl} S & \rightarrow & x E a \\ & | & y E z \\ E & \rightarrow & x E' \\ E' & \rightarrow & F \\ & | & E q \\ F & \rightarrow & q F' \\ F' & \rightarrow & z F' \\ & | & \epsilon \end{array}$$

Solution 5

$$\begin{array}{lcl} S & \rightarrow & a E S' \\ S' & \rightarrow & x S' \\ & | & \epsilon \\ E & \rightarrow & y a E' \\ E' & \rightarrow & y \\ & | & z \end{array}$$

Solution 6

$$\begin{array}{lcl} S & \rightarrow & y S' \\ S' & \rightarrow & y \\ & | & E \\ E & \rightarrow & x a E' \\ E' & \rightarrow & x E' \\ & | & \epsilon \end{array}$$

Solution 7

$$\begin{array}{lcl} S & \rightarrow & a E S' \\ S' & \rightarrow & x S' \\ & | & \epsilon \\ E & \rightarrow & x E' \\ E' & \rightarrow & y \\ & | & z \end{array}$$

Solution 8

$$\begin{array}{lcl} S & \rightarrow & x E S' \\ S' & \rightarrow & x S' \\ & | & \epsilon \\ E & \rightarrow & y E' \\ E' & \rightarrow & y \\ & | & z \end{array}$$

Solution 9

$$\begin{array}{lcl} S & \rightarrow & y S' \\ S' & \rightarrow & z \\ & | & E \\ E & \rightarrow & x a a E' \\ E' & \rightarrow & z z E' \\ & | & \epsilon \end{array}$$

Solution 10

$$\begin{array}{lcl} S & \rightarrow & a E S' \\ S' & \rightarrow & x S' \\ & | & \epsilon \\ E & \rightarrow & z b E' \\ E' & \rightarrow & b \\ & | & z \end{array}$$