## Programming Languages and Compilers (CS 421)

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https://courses.engr.illinois.edu/cs421/fa2024/CS421C

Based on slides by Elsa Gunter, which are based in part on previous slides by Mattox Beckman and updated by Vikram Adve and Gul Agha

# **CPS Transformation**

- Step 1: Add continuation argument to any function definition:
  - let f arg = e  $\Rightarrow$  let f arg k = e
  - Idea: Every function takes an extra parameter saying where the result goes
- Step 2: A simple expression in tail position should be passed to a continuation instead of returned:
  - return a  $\Rightarrow$  k a
  - Assuming a is a constant or variable.
  - "Simple" = "No available function calls."

## **CPS Transformation**

- Step 3: Pass the current continuation to every function call in tail position
  - return f arg  $\Rightarrow$  f arg k
  - The function "isn't going to return," so we need to tell it where to put the result.

## **CPS Transformation**

- Step 4: Each function call not in tail position needs to be converted to take a new continuation (containing the old continuation as appropriate)
  - return op (f arg)  $\Rightarrow$  f arg (fun r -> k(op r))

op represents a primitive operation

• return  $f(g arg) \Rightarrow g arg (fun r-> f r k)$ 

- Step 1: Add continuation argument to any function definition
- Step 2: A simple expression in tail position should be passed to a continuation instead of returned
- Step 3: Pass the current continuation to every function call in tail position
- **Step 4:** Each function call not in tail position needs to be converted to take a new continuation (containing the old continuation as appropriate)

#### **Before:**

#### After:

```
let rec add_list lst = let rec add_listk lst k =
                                            (* rule 1 *)
                          match lst with
match lst with
                          [ ] -> k Ø (* rule 2 *)
  []->0
                          0 :: xs -> add_listk xs k
0 :: xs -> add_list xs
                                            (* rule 3 *)
 x :: xs -> (+) x
                          x :: xs -> add listk xs
  (add_list xs);;
                                     (fun r -> k ((+) x r));;
                                            (* rule 4 *)
                   Same as:
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                   fun r -> addk x r k
                                                          6
```

#### **Before:**

```
let rec mem (y,lst) =
match lst with
  [ ] -> false
| x :: xs ->
  if (x = y)
   then true
   else mem(y,xs);;
```

# After: let rec memk (y,lst) k = (\* rule 1 \*)

#### **Before:**

```
let rec mem (y,lst) = _____
match lst with
 [ ] -> false
| x :: xs ->
 if (x = y)
 then true
 else mem(y,xs);;
```

```
k true (* rule 2 *)
```

#### **Before:**

```
let rec mem (y,lst) =
match lst with
  [ ] -> false
| x :: xs ->
  if (x = y)
   then true
   else mem(y,xs);;
```

# 

memk (y, xs) k (\* rule 3

#### **Before:**

let rec mem (y,lst) =
match lst with
 [ ] -> false
| x :: xs ->
 if (x = y)
 then true
 else mem(y,xs);;

#### **After:** let rec memk (y,lst) k = (\* rule 1 \*) k false (\* rule 2 \*) eqk (x, y) (fun b -> b (\* rule 4 \*) k true (\* rule 2 \*) memk (y, xs) (\* rule 3 \*

#### **Before:**

```
let rec mem (y,lst) =
match lst with
  [ ] -> false
| x :: xs ->
  if (x = y)
   then true
   else mem(y,xs);;
```

#### **After:** let rec memk (y,lst) k = (\* rule 1 \*) k false (\* rule 2 \*) eqk (x, y) (fun b $\rightarrow$ if b (\* rule 4 \*) then k true (\* rule 2 \*) else memk (y, xs) (\* rule 3 \*

#### **Before:**

let rec mem (y,lst) =
match lst with
 [ ] -> false
| x :: xs ->
 if (x = y)
 then true
 else mem(y,xs);;

#### After:

## **Other Uses for Continuations**

- CPS designed to preserve evaluation order
- Continuations used to express order of evaluation

- Can also be used to change order of evaluation
- Implements:
  - Exceptions and exception handling
  - Co-routines
  - (pseudo, aka green) threads

#### **Exceptions - Example**

```
# exception Zero;;
exception Zero
```

```
# let rec list_mult_aux list =
   match list with
   [ ] -> 1
   | x :: xs ->
        if x = 0 then raise Zero
        else x * list_mult_aux xs;;
val list_mult_aux : int list -> int = <fun>
```

#### **Exceptions - Example**

```
# let list_mult list =
    try list_mult_aux list with Zero -> 0;;
val list_mult : int list -> int = <fun>
```

```
# list_mult [3;4;2];;
```

-: int = 24

```
# list_mult [7;4;0];;
```

-: int = 0

```
# list_mult_aux [7;4;0];;
Exception: Zero.
```

#### Exceptions

# When an exception is raised

- The current computation is aborted
- Control is "thrown" back up the call stack until a matching handler is found
- All the intermediate calls waiting for a return values are thrown away

## **Implementing Exceptions**

```
# let multkp (m, n) k =
    let r = m * n in
        ( print_string "product result: ";
        print_int r; print_string "\n";
        k r);;
val multkp : int ( int -> (int -> 'a) -> 'a =
        <fun>
```

## **Implementing Exceptions**

# let rec list\_multk\_aux list k kexcp =
 match list with
 [ ] -> k 1
 | x :: xs -> if x = 0 then kexcp 0
 else
 list\_multk\_aux xs
 (fun r -> multkp (x, r) k)
 kexcp;;

```
# let rec list_multk list k =
    list_multk_aux list k
        (fun x -> print_string "nil\n");;
```

# **Implementing Exceptions**

```
# list_multk [3;4;2] report;;
product result: 2
product result: 8
product result: 24
24
- : unit = ()
```

```
# list_multk [7;4;0] report;;
nil
```

```
- : unit = ()
```

# Advanced: Using CPS as Compiler Intermediate Representation



Ocaml compiler (latest version) uses CPS:

- Blog: <u>https://discuss.ocaml.org/t/blog-the-flambda2-snippets-by-ocamlpro/14331</u>
- Tutorial: <u>https://www.youtube.com/watch?v=eI5GBpT2Brs</u>

#### Various discussions in research literature:

- With? <u>https://www.microsoft.com/en-us/research/wp-</u> <u>content/uploads/2007/10/compilingwithcontinuationscontinued.pdf</u>
- Without? <u>https://pauldownen.com/publications/pldi17.pdf</u>
- Whatever? <u>https://dl.acm.org/doi/10.1145/3341643</u>

Intermediate representations CPS for functional vs SSA for imperative

https://dl.acm.org/doi/10.1145/202530.202532

## Data type in Ocaml: lists

- Frequently used lists in recursive program
- Matched over two structural cases
  - [] the empty list
  - (x :: xs) a non-empty list
- Covers all possible lists
- type `a list = [] | (::) of `a \* `a list
  - Not quite legitimate declaration because of special syntax

Variants - Syntax (slightly simplified)

- type  $name = C_1 [of ty_1] | ... | C_n [of ty_n]$
- Introduce a type called name
- (fun x -> C<sub>i</sub>x): ty<sub>1</sub> -> name
- *C<sub>i</sub>* is called a *constructor*, if the optional type argument is omitted, it is called a *constant*
- Constructors are the basis of almost all pattern matching

## **Enumeration Types as Variants**

An enumeration type is a collection of distinct values



In C and Ocaml they have an order structure; order by order of input

## **Enumeration Types as Variants**

# type weekday = Monday | Tuesday | Wednesday | Thursday | Friday | Saturday | Sunday;; type weekday = Monday | Tuesday | Wednesday | Thursday | Friday | Saturday

| Sunday

#### **Functions over Enumerations**

# let day after day = match day with Monday -> Tuesday Tuesday -> Wednesday Wednesday -> Thursday Thursday -> Friday Friday -> Saturday Saturday -> Sunday Sunday -> Monday;; val day after : weekday -> weekday = <fun>

# Functions over Enumerations

Write a function days\_later n day that computes a day which is n days away from the day. Note that n can be greater than 7 (more than one week) and also negative (meaning a day before

# let rec days later n day = match n with 0 -> day -> if n > 0 then day after (days later (n - 1) day) else days\_later (n + 7) day;; val days later : int -> weekday -> weekday=<fun> 9/23/2024 31

## **Functions over Enumerations**

- # days\_later 2 Tuesday;;
- : weekday = Thursday
- # days\_later (-1) Wednesday;;
- : weekday = Tuesday
- # days\_later (-4) Monday;;
- : weekday = Thursday

#### Problem:

- # type weekday = Monday | Tuesday | Wednesday | Thursday | Friday | Saturday | Sunday;;
- Write function is\_weekend : weekday -> bool

#### Problem:

# type weekday = Monday | Tuesday | Wednesday | Thursday | Friday | Saturday | Sunday;;

## **Example Enumeration Types**

- # type bin\_op = IntPlusOp | IntMinusOp | EqOp | CommaOp | ConsOp
- # type mon\_op = HdOp | TlOp | FstOp | SndOp

# **Disjoint Union Types**

Disjoint union of types, with some possibly occurring more than once



We can also add in some new singleton elements

# **Disjoint Union Types**

```
type id = DriversLicense of int |
 SocialSecurity of int | Name of string
# let check id id =
   match id with
     DriversLicense num ->
      not (List.mem num [13570; 99999])
   SocialSecurity num -> num < 90000000</pre>
   Name str -> not (str = "John Doe");;
val check id : id -> bool = <fun>
```

## Problem

#### Create a type to represent the currencies for US, UK, Europe and Japan

Hint: Dollar, Pound, Euro, Yen

## Problem

 Create a type to represent the currencies for US, UK, Europe and Japan

type currency =
 Dollar of int
 Pound of int
 Euro of int
 Yen of int

## Example Disjoint Union Type

# type const = BoolConst of bool IntConst of int FloatConst of float StringConst of string NilConst UnitConst

## Example Disjoint Union Type

How to represent 7 as a const?
Answer: IntConst 7

**Polymorphism in Variants** 

The type 'a option gives us something to represent non-existence or failure

# type 'a option = Some of 'a | None;;
type 'a option = Some of 'a | None

Used to encode partial functions
Often can replace the raising of an exception
# Functions producing option # type 'a option = Some of 'a None;;

- # first (fun x -> x > 3) [1;3;4;2;5];;
- : int option = Some 4
- # first (fun x -> x > 5) [1;3;4;2;5];;
- : int option = None

## Functions over option

# type 'a option =
 Some of 'a
 None;;

# let result\_ok r =
 match r with None -> false
 Some \_ -> true;;
val result\_ok : 'a option -> bool = <fun>

# result\_ok (first (fun x -> x > 3) [1;3;4;2;5]);;

- : bool = true

# result\_ok (first (fun x -> x > 5) [1;3;4;2;5]);;

- : bool = false

## Problem

# type 'a option =
 Some of 'a
 None;;

Write a hd and tl on lists that doesn't raise an exception and works at all types of lists.

## Problem

# type 'a option = Some of 'a None;;

- Write a hd and tl on lists that doesn't raise an exception and works at all types of lists.
- let hd list = match list with [] -> None (x::xs) -> Some x let tl list = match list with [] -> None (x::xs) -> Some xs 9/23/2024

## Mapping over Variants

```
# let optionMap f opt =
    match opt with
    None -> None
    | Some x -> Some (f x);;
val optionMap : ('a -> 'b) -> 'a option -> 'b
    option = <fun>
```

# optionMap
 (fun x -> x - 2)
 (first (fun x -> x > 3) [1;3;4;2;5]);;
- : int option = Some 2

## Folding over Variants

# let optionFold someFun noneVal opt =
 match opt with
 None -> noneVal
 Some x -> someFun x;;
val optionFold : ('a -> 'b) -> 'b -> 'a option
 -> 'b = <fun>

# let optionMap f opt =
 optionFold (fun x -> Some (f x)) None opt;;
val optionMap : ('a -> 'b) -> 'a option -> 'b
 option = <fun>

## **Recursive Types**

The type being defined may be a component of itself



- # type int\_Bin\_Tree =
   Leaf of int
   Node of (int\_Bin\_Tree \* int\_Bin\_Tree);;
- type int\_Bin\_Tree = Leaf of int | Node of (int\_Bin\_Tree \* int\_Bin\_Tree)

**Recursive Data Type Values** 

# let bin\_tree =
 Node(Node(Leaf 3, Leaf 6),Leaf (-7));;

val bin\_tree : int\_Bin\_Tree = Node (Node
 (Leaf 3, Leaf 6), Leaf (-7))

**Recursive Data Type Values** 

# let bin\_tree =
 Node(Node(Leaf 3, Leaf 6),Leaf (-7));;



### **Recursive Functions**

# let rec first\_leaf\_value tree =
 match tree with
 (Leaf n) -> n
 Node (left\_tree, right\_tree) ->
 first\_leaf\_value left\_tree;;

# let left = first\_leaf\_value bin\_tree;;
val left : int = 3

```
# type exp =
    VarExp of string
   ConstExp of const
   MonOpAppExp of mon op * exp
    BinOpAppExp of bin op * exp * exp
    IfExp of exp* exp * exp
   | AppExp of exp * exp
    FunExp of string * exp
```

How to represent 6 as an exp?

- How to represent 6 as an exp?Answer: ConstExp (IntConst 6)

- How to represent (6, 3) as an exp?

- How to represent (6, 3) as an exp?
  BinOpAppExp (CommaOp, ConstExp (IntConst 6), ConstExp (IntConst 3))

- - BinOpAppExp of bin\_op \* exp \* exp | ...
- How to represent [(6, 3)] as an exp?

 BinOpAppExp (ConsOp, BinOpAppExp (CommaOp, ConstExp (IntConst 6), ConstExp (IntConst 3)), ConstExp NilConst))));;

## Problem

type int\_Bin\_Tree =Leaf of int
 Node of (int\_Bin\_Tree \* int\_Bin\_Tree);;

- Write sum\_tree : int\_Bin\_Tree -> int
- Adds all ints in tree
- let rec sum\_tree t =

## Solution

- type int\_Bin\_Tree =Leaf of int
   Node of (int\_Bin\_Tree \* int\_Bin\_Tree);;
- Write sum\_tree : int\_Bin\_Tree -> int
- Adds all ints in tree
- let rec sum\_tree t =
  - match t with
    - Leaf n -> n
    - Node(t1,t2) -> sum\_tree t1 + sum\_tree t2

### **Recursion over Recursive Data Types**

```
# type exp = VarExp of string
    | ConstExp of const
    | BinOpAppExp of bin_op * exp * exp
    | FunExp of string * exp
    | AppExp of exp * exp
```

How to count the number of variables in an exp?

#### **Recursion over Recursive Data Types**

- # type exp = VarExp of string | ConstExp of const | BinOpAppExp of bin\_op \* exp \* exp | FunExp of string \* exp | AppExp of exp \* exp
- How to count the number of variables in an exp?

```
# let rec varCnt exp =
  match exp with
    VarExp x ->
    ConstExp c ->
    BinOpAppExp (b, e1, e2) ->
    FunExp (x,e) ->
    AppExp (e1, e2) ->
```

#### **Recursion over Recursive Data Types**

- # type exp = VarExp of string | ConstExp of const | BinOpAppExp of bin\_op \* exp \* exp | FunExp of string \* exp | AppExp of exp \* exp
- How to count the number of variables in an exp?

```
# let rec varCnt exp =
   match exp with
      VarExp x -> 1
      ConstExp c -> 0
      BinOpAppExp (b, e1, e2) -> varCnt e1 +varCnt e2
      FunExp (x,e) -> 1 + varCnt e
      AppExp (e1, e2) -> varCnt e1 + varCnt e2
      66
```

## Mapping over Recursive Types

# let rec ibtreeMap f tree =
 match tree with
 (Leaf n) ->
 Node (left\_tree, right\_tree) ->

## Mapping over Recursive Types

```
# let rec ibtreeMap f tree =
    match tree with
      (Leaf n) -> Leaf (f n)
    Node (left tree, right_tree) ->
           Node (ibtreeMap f left tree,
                 IbtreeMap f right tree);;
val ibtreeMap : (int -> int) -> int_Bin_Tree ->
 int Bin Tree = <fun>
```

## Mapping over Recursive Types

# let bin\_tree =
Node(Node(Leaf 3, Leaf 6),Leaf (-7));;

- # ibtreeMap ((+) 2) bin\_tree;;
- : int\_Bin\_Tree = Node (Node (Leaf 5, Leaf 8), Leaf (-5))

## Summing up Elements of a Tree

# let rec tree\_sum\_0 tree =
 match tree with
 Leaf n ->

Node (left\_tree, right\_tree) ->

## Folding over Recursive Types

#### val ibtreeFoldRight : (int -> 'a) -> ('a -> 'a -> 'a) -> int\_Bin\_Tree -> 'a = <fun>

## Folding over Recursive Types

# let rec ibtreeFoldRight leafFun nodeFun tree =
 match tree with
 Leaf n -> leafFun n

Node (left\_tree, right\_tree) ->

nodeFun

(ibtreeFoldRight leafFun nodeFun left\_tree)
(ibtreeFoldRight leafFun nodeFun right tree);

val ibtreeFoldRight : (int -> 'a) -> ('a -> 'a -> 'a) -> int\_Bin\_Tree -> 'a = <fun>

## Folding over Recursive Types

# let tree\_sum =
 ibtreeFoldRight (fun x -> x) (+);;
val tree\_sum : int\_Bin\_Tree -> int = <fun>

- # tree\_sum bin\_tree;;
- -: int = 2

## **Mutually Recursive Types**

```
# type 'a tree =
        TreeLeaf of 'a
      | TreeNode of 'a treeList
and
   'a treeList =
        Last of 'a tree
      More of ('a tree * 'a treeList);;
```

type 'a tree = TreeLeaf of 'a | TreeNode of 'a treeList and 'a treeList = Last of 'a tree | More of ('a tree \* 'a treeList) **Mutually Recursive Types - Values** 

# let tree =
 TreeNode
 (More (TreeLeaf 5,
 (More (TreeNode
 (More (TreeLeaf 3,
 Last (TreeLeaf 2))),
 Last (TreeLeaf 7))));;

## **Mutually Recursive Types - Values**

```
val tree : int tree =
 TreeNode
  (More
    (TreeLeaf 5,
     More
      (TreeNode (More (TreeLeaf 3,
       Last (TreeLeaf 2))),
   Last (TreeLeaf 7))))
```

**Mutually Recursive Types - Values** 



## **Mutually Recursive Types - Values**

#### A more conventional picture



### **Mutually Recursive Functions**

```
# let rec fringe tree =
    match tree with
       (TreeLeaf x) \rightarrow [x]
  (TreeNode list) -> list_fringe list
and list fringe tree list =
    match tree list with
       (Last tree) -> fringe tree
  (More (tree,list)) ->
       (fringe tree) @ (list_fringe list);;
```

val fringe : 'a tree -> 'a list = <fun>
val list\_fringe : 'a treeList -> 'a list = <fun>

## **Mutually Recursive Functions**

- # fringe tree;;
- : int list = [5; 3; 2; 7]
# type 'a tree = TreeLeaf of 'a | TreeNode of 'a treeList and 'a treeList = Last of 'a tree | More of ('a tree \* 'a treeList);;

Define tree\_size

# type 'a tree = TreeLeaf of 'a | TreeNode of 'a treeList and 'a treeList = Last of 'a tree | More of ('a tree \* 'a treeList);;

Define tree\_size
 let rec tree\_size t =
 match t with TreeLeaf \_ ->
 | TreeNode ts ->

# type 'a tree = TreeLeaf of 'a | TreeNode of 'a treeList and 'a treeList = Last of 'a tree | More of ('a tree \* 'a treeList);;

- Define tree\_size
- let rec tree\_size t =

match t with TreeLeaf \_ -> 1
| TreeNode ts -> treeList\_size ts

# type 'a tree = TreeLeaf of 'a | TreeNode of 'a treeList

and 'a treeList = Last of 'a tree | More of ('a tree \* 'a treeList);;

Define tree\_size and treeList\_size

- let rec tree\_size t =
  - match t with TreeLeaf \_ -> 1
    - TreeNode ts -> treeList\_size ts

and treeList\_size ts =

# type 'a tree = TreeLeaf of 'a | TreeNode of 'a treeList and 'a treeList = Last of 'a tree | More of ('a tree \* 'a treeList);;

Define tree\_size and treeList\_size let rec tree size t =match t with TreeLeaf -> 1 TreeNode ts -> treeList size ts and treeList size ts = match ts with Last t -> | More t ts'  $\rightarrow$ 

# type 'a tree = TreeLeaf of 'a | TreeNode of 'a treeList and 'a treeList = Last of 'a tree | More of ('a tree \* 'a treeList);;

Define tree\_size and treeList\_size let rec tree\_size t = match t with TreeLeaf \_ -> 1 | TreeNode ts -> treeList\_size ts

and treeList\_size ts =

match ts with Last t -> tree\_size t

| More t ts' -> tree\_size t + treeList\_size ts'

# type 'a tree = TreeLeaf of 'a | TreeNode of 'a treeList and 'a treeList = Last of 'a tree | More of ('a tree \* 'a treeList);;

Define tree\_size and treeList\_size

- let rec tree\_size t =
  - match t with TreeLeaf \_ -> 1
  - | TreeNode ts -> treeList\_size ts

and treeList\_size ts =

match ts with Last t -> tree\_size t

More t ts' -> tree\_size t + treeList\_size ts'

## **Nested Recursive Types**

# type 'a labeled\_tree =
 TreeNode of ('a \* 'a labeled\_tree list);;

#### type 'a labeled\_tree = TreeNode of ('a \* 'a labeled\_tree list)

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# **Nested Recursive Type Values**

# let ltree =
 TreeNode(5,
 [TreeNode (3, []);
 TreeNode (2, [TreeNode (1, []);
 TreeNode (7, [])]);
 TreeNode (5, [])]);;

# **Nested Recursive Type Values**



#### **Nested Recursive Type Values**



**Mutually Recursive Functions** 

# let rec flatten\_tree labtree =
 match labtree with
 TreeNode (x,treelist) ->
 x::flatten\_tree\_list treelist

and flatten\_tree\_list treelist =
 match treelist with
 [] -> []
 labtree::labtrees ->
 flatten\_tree labtree
 @ (flatten\_tree\_list labtrees);;

# **Mutually Recursive Functions**

val flatten\_tree : 'a labeled\_tree -> 'a list = <fun>
val flatten\_tree\_list : 'a labeled\_tree list -> 'a list =
 <fun>

- # flatten\_tree ltree;;
- : int list = [5; 3; 2; 1; 7; 5]

#### Nested recursive types lead to mutually recursive functions