

# Disambiguating Grammars

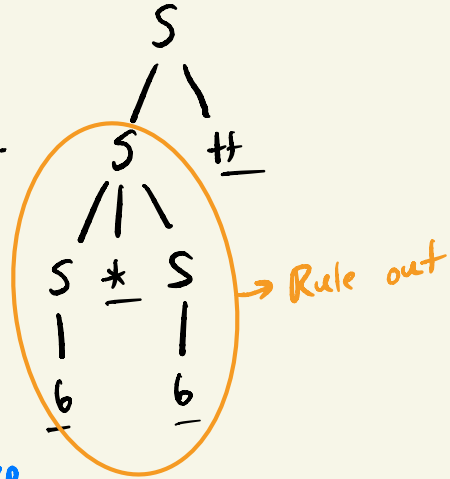
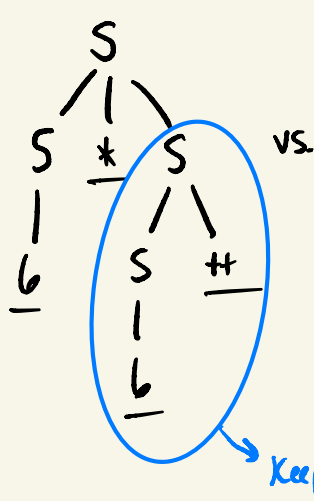
[No algorithm in general]  
[Requires creativity.]

## Example 1

Starting grammar :

$$S ::= S * S \mid \underline{S \#} \mid 6$$

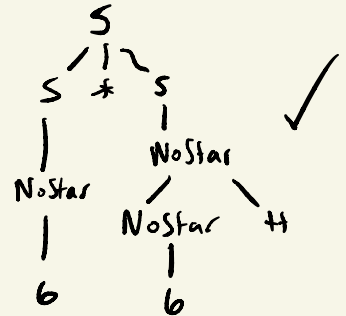
6 \* 6 #



$$S ::= S * S \mid \text{NoStar}$$

$$\text{NoStar} ::= \text{NoStar} \# \mid 6$$

Check 6 \* 6 #

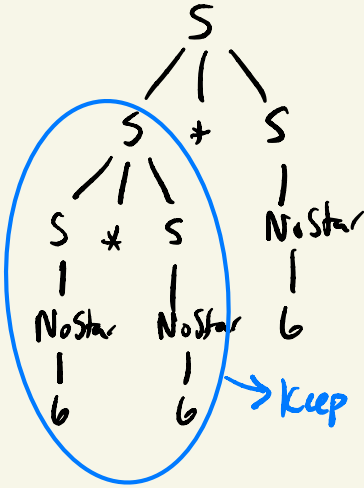


Is this unambiguous

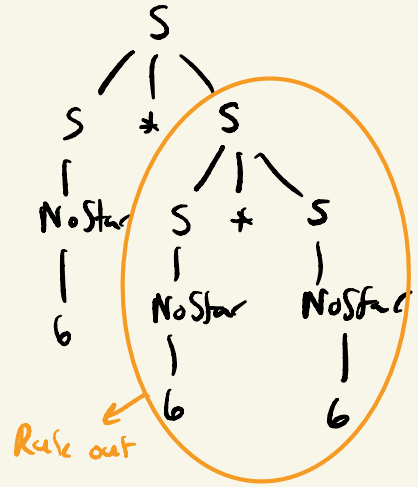
$S ::= S * \underline{S} \mid \text{NoStar}$

$\text{NoStar} ::= \text{NoStar} \# \mid 6$

$6 * 6 * 6$



vs.



left associativity

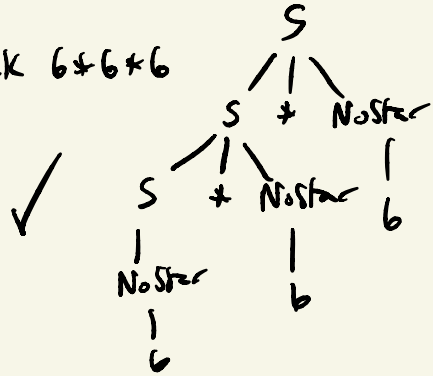
$S ::= S * \text{NoStar} \mid \text{NoStar}$

$\text{NoStar} ::= \text{NoStar} \# \mid 6$

Unambiguous? Yes.

Prove it by induction on the depth of parse trees.

check  $6 * 6 * 6$



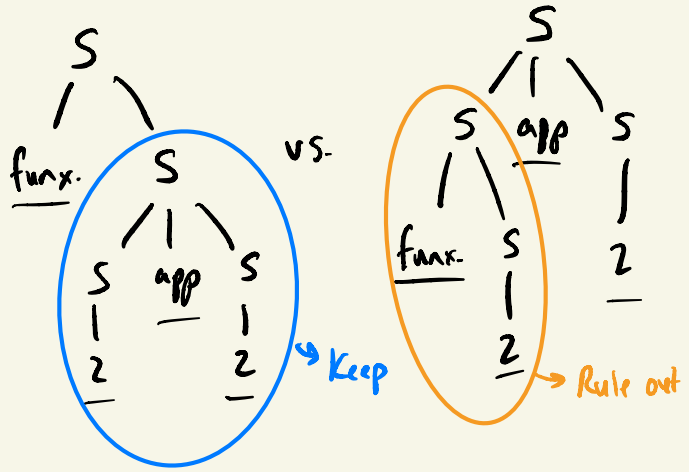
## Example 2

Starting grammar:

$S ::= \text{fun } x. S \mid S \text{ app } S \mid \_ \mid (S)$

Ambiguous?

$\text{fun } x. \_ \text{ app } \_$



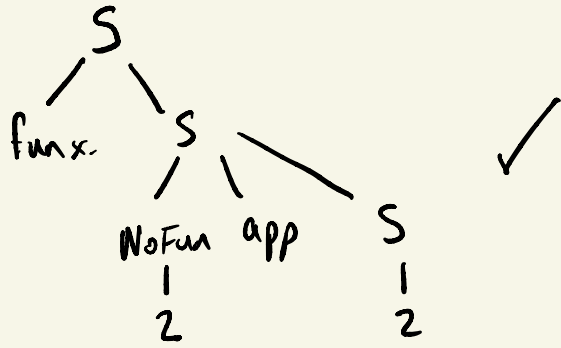
$S ::= \text{fun } x. S \mid \text{NoFun app } S \mid \_ \mid (S)$

$\text{NoFun} ::= \text{NoFun app NoFun} \mid \_ \mid (S)$

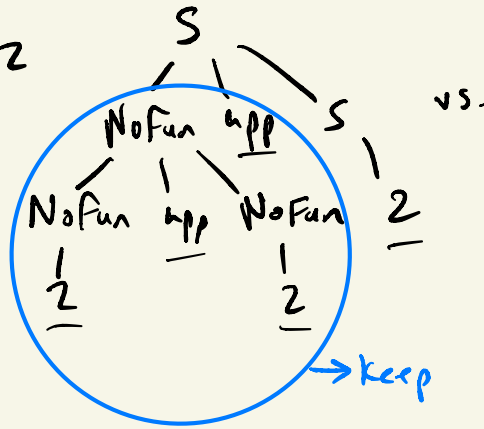
$S ::= \text{fun } x. S \mid \text{NoFun app } \underline{S} \mid \underline{2} \mid (S)$

$\text{NoFun} ::= \text{NoFun app NoFun} \mid \underline{2} \mid (S)$

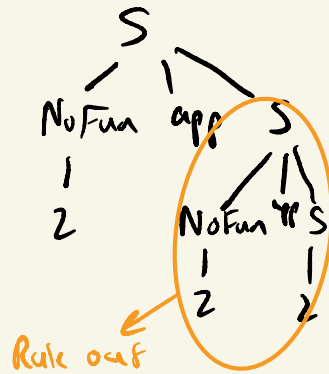
fun x. 2 app 2



2 app 2 app 2



vs.



$S ::= \text{fun } x. S \mid \text{NoFun app NoApp} \mid \text{NoFun NoApp}$

$\text{NoFun} ::= \text{NoFun app NoFun NoApp} \mid \text{NoFun NoApp}$

$\text{NoApp} ::= \text{fun } x. S \mid \text{NoFun NoApp}$

$\text{NoFun NoApp} ::= \underline{2} \mid (S)$

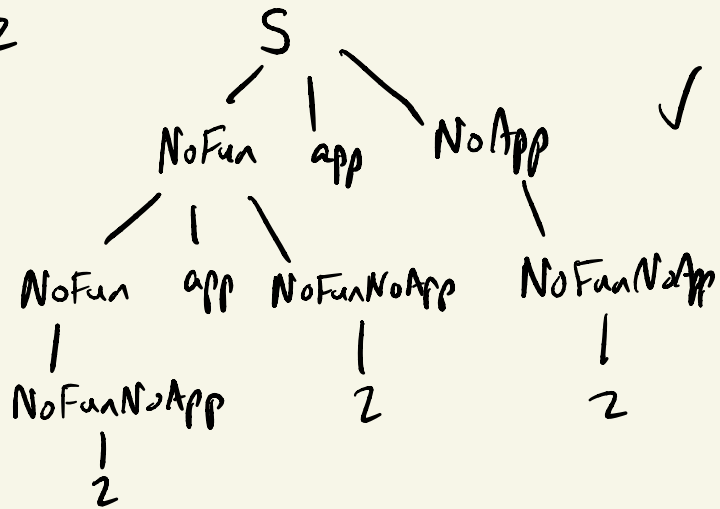
$S ::= \text{fun } x. S \mid \text{NoFun app NoApp} \mid \text{NoFunNoApp}$

$\text{NoFun} ::= \text{NoFun app NoFunNoApp} \mid \text{NoFunNoApp}$

$\text{NoApp} ::= \text{fun } x. S \mid \text{NoFunNoApp}$

$\text{NoFunNoApp} ::= \lambda \mid (S)$

$\lambda \text{ app } \lambda \text{ app } \lambda$



### Example 3

Starting grammar:

$$S ::= \text{fun } x. S \mid S * S \mid S + S \mid Z$$

### Precedence & Associativity Specification

right \*                      highest precedence  
left +  
fun x.                      lowest precedence

Idea: Start from lowest precedence, and work your way up to higher precedence.

$$\begin{aligned} S &::= \text{fun } x. S \mid \text{Plus} \\ \text{Plus} &::= \text{Plus} + \text{Star} \mid \text{Star} \\ \text{Star} &::= \text{Atom} * \text{Star} \mid \text{Atom} \\ \text{Atom} &::= Z \end{aligned}$$

Missing things like  $Z * \text{fun } x. Z$

$S ::= \text{fun } x. S \mid \text{Plus}$  Missing things

$\text{Plus} ::= \text{Plus} + \text{Star} \mid \text{Star}$  like

$\text{Star} ::= \text{Atom} * \text{Star} \mid \text{Atom}$   $2 * \text{fun } x. 2$

$\text{Atom} ::= 2$

$S ::= \text{NoFun} \mid \text{Fun}$

$\text{NoFun} ::= \text{NoFun} + \text{NoPlusNoFun} \mid \text{NoPlusNoFun}$

$\text{NoPlusNoFun} ::= \text{NoPlusNoFunNoStar} * \text{NoPlusNoFun}$

$\text{NoPlusNoFunNoStar} ::= 2$

$\text{Fun} ::= \text{NoFun} + \text{NoPlus} \mid \text{NoPlus}$

$\text{NoPlus} ::= \text{NoPlusNoFunNoStar} * \text{NoPlus} \mid \text{FinallyFun}$

$\text{FinallyFun} ::= \text{fun } x. S$

$S ::= \underline{\text{NoFun}} \mid \text{Fun}$

$\text{NoFun} ::= \text{NoFun} + \underline{\text{NoPlusNoFun}} \mid \underline{\text{NoPlusNoFun}}$

$\text{NoPlusNoFun} ::= \text{NoPlusNoFunNoStar} * \text{NoPlusNoFun} \mid \text{NoPlusNoFunNoStar}$

$\text{NoPlusNoFunNoStar} ::= 2$

$\text{Fun} ::= \text{NoFun} + \underline{\text{NoPlus}} \mid \text{NoPlus}$

$\text{NoPlus} ::= \text{NoPlusNoFunNoStar} * \text{NoPlus} \mid \text{FinallyFun}$

$\text{FinallyFun} ::= \text{fun } x. S$

$2 * \text{fun } x. 2$

