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# Programming Languages and Compilers (CS 421)

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<https://courses.grainger.illinois.edu/cs421/fa2023/>

Based heavily on slides by Elsa Gunter, which were based in part on slides by Mattox Beckman, as updated by Vikram Adve and Gul Agha



## Objectives for Today

---

- We will look at another example of the **CPS Transformation** that we saw last week
- Then, taking a step back—how would we actually **automate** transforming programs like this?
- We need a way to **represent** the syntax of our language that allows us to (1) **construct** a representation of a new (transformed) program, and (2) **match** over the syntax of the original
- We've seen something like this for lists—if we generalize, we get **datatypes**
- We'll cover **many kinds** of datatypes



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Please post questions on Piazza!

---



# CPS Transformation Example

---



# CPS Example: List Membership

---

## Before:

```
let rec mem (y, lst) =  
  match lst with  
  | [ ] -> false  
  | x :: xs ->  
    if (x = y) then  
      true  
    else  
      mem (y, xs)
```

CPS Transformation Example



# CPS Example: List Membership

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match lst with
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```
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CPS Transformation Example





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CPS Transformation Example



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let rec **mem** (y, lst) =

match lst with

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| **x :: xs** ->

if (x = y) then

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**else**

**mem** (y, xs)

CPS Transformation Example



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let rec memk (y, lst) k = (* rule 1 *)
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CPS Transformation Example



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CPS Transformation Example





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let rec memk (y, lst) k = (* rule 1 *)  
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  eqk (x, y) (fun b -> b  
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CPS Transformation Example

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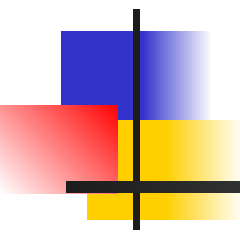
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How to implement automatically in  
compiler, rather than by hand?

---



How do we even represent the syntax of our language, and map over it to transform programs?

---



# Datatypes

---



# OCaml Datatype You've Seen: lists

---

- Frequently used lists in recursive program
- Matched over two structural cases
  - `[]` - the empty list
  - `(x :: xs)` a non-empty list
- Covers all possible lists
- `type 'a list = [] | (::) of 'a * 'a list`
  - Not quite legitimate declaration because of special syntax



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# OCaml Datatypes in General

- $\text{type } name = C_1 [\text{of } ty_1] \mid \dots \mid C_n [\text{of } ty_n]$
- Introduce a type called *name*
- $(\text{fun } x \rightarrow C_i x) : ty_1 \rightarrow name$
- $C_i$  is called a *constructor*; if the optional type argument is omitted, it is called a *constant*
- Constructors are the basis of almost all *pattern matching* (alt. *destruction* or, with caveats, *elimination*)

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# Enumeration Types

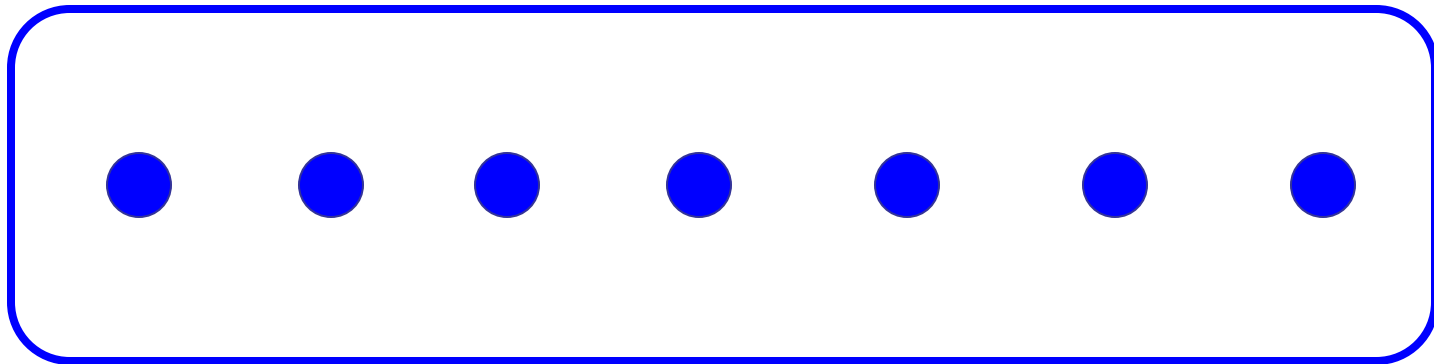
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# OCaml Variants

- $\text{type name} = C_1 [\text{of } ty_1] \mid \dots \mid C_n [\text{of } ty_n]$
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- Constructors are the basis of almost all *pattern matching* (alt. *destruction* or, with some extra machinery, *elimination*)

# Enumeration Types as Variants

An enumeration type is a collection of distinct values



In C and Ocaml they have an order structure;  
order by order of input



# Enumeration Types as Variants

---

```
# type weekday = Monday | Tuesday | Wednesday  
| Thursday | Friday | Saturday | Sunday;;
```

```
type weekday =  
    Monday  
| Tuesday  
| Wednesday  
| Thursday  
| Friday  
| Saturday  
| Sunday
```

Enumeration Types



# Functions over Enumerations

---

```
# let day_after day =  
  match day with  
  | Monday -> Tuesday  
  | Tuesday -> Wednesday  
  | Wednesday -> Thursday  
  | Thursday -> Friday  
  | Friday -> Saturday  
  | Saturday -> Sunday  
  | Sunday -> Monday;;
```

```
val day_after : weekday -> weekday = <fun>
```

Enumeration Types



# Functions over Enumerations

---

```
# let rec days_later n day =  
  match n with  
  | 0 -> day  
  | _ ->  
    if n > 0 then  
      day_after (days_later (n - 1) day)  
    else  
      days_later (n + 7) day;;  
val days_later : int -> weekday -> weekday = <fun>
```

Enumeration Types





# Functions over Enumerations

---

```
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```

```
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```

```
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Enumeration Types



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Enumeration Types



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val days_later : int -> weekday -> weekday = <fun>
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Enumeration Types



# Problem:

---

```
# type weekday = Monday | Tuesday | Wednesday  
| Thursday | Friday | Saturday | Sunday;;
```

- Write function `is_weekend : weekday -> bool`

```
let is_weekend day =
```

**Your turn!**



# Problem:

---

```
# type weekday = Monday | Tuesday | Wednesday  
| Thursday | Friday | Saturday | Sunday;;
```

- Write function `is_weekend : weekday -> bool`

```
let is_weekend day =  
  match day with
```

**Weekend days?**



# Problem:

---

```
# type weekday = Monday | Tuesday | Wednesday  
| Thursday | Friday | Saturday | Sunday;;
```

- Write function `is_weekend : weekday -> bool`

```
let is_weekend day =  
  match day with  
  | _ -> true
```

**In a better world ...**



# Problem:

---

```
# type weekday = Monday | Tuesday | Wednesday  
| Thursday | Friday | Saturday | Sunday;;
```

- Write function `is_weekend : weekday -> bool`

```
let is_weekend day =  
  match day with  
  | Saturday -> true  
  | Sunday -> true
```

**Other days?**



# Problem:

---

```
# type weekday = Monday | Tuesday | Wednesday  
| Thursday | Friday | Saturday | Sunday;;
```

- Write function `is_weekend : weekday -> bool`

```
let is_weekend day =  
  match day with  
  | Saturday -> true  
  | Sunday -> true  
  | Monday -> false  
  | Tuesday -> false ...
```

**More concisely?**

Enumeration Types





# Problem:

---

```
# type weekday = Monday | Tuesday | Wednesday  
| Thursday | Friday | Saturday | Sunday;;
```

- Write function `is_weekend : weekday -> bool`

```
let is_weekend day =  
  match day with  
  | Saturday -> true  
  | Sunday -> true  
  | _ -> false
```

**Yay**



# Enumeration Types in Languages!

---

# (\* Binary operators \*)

```
type bin_op = IntPlusOp | IntMinusOp  
            | EqOp | CommaOp | ConsOp
```

# (\* Unary operators \*)

```
type mon_op = HdOp | TlOp | FstOp | SndOp
```

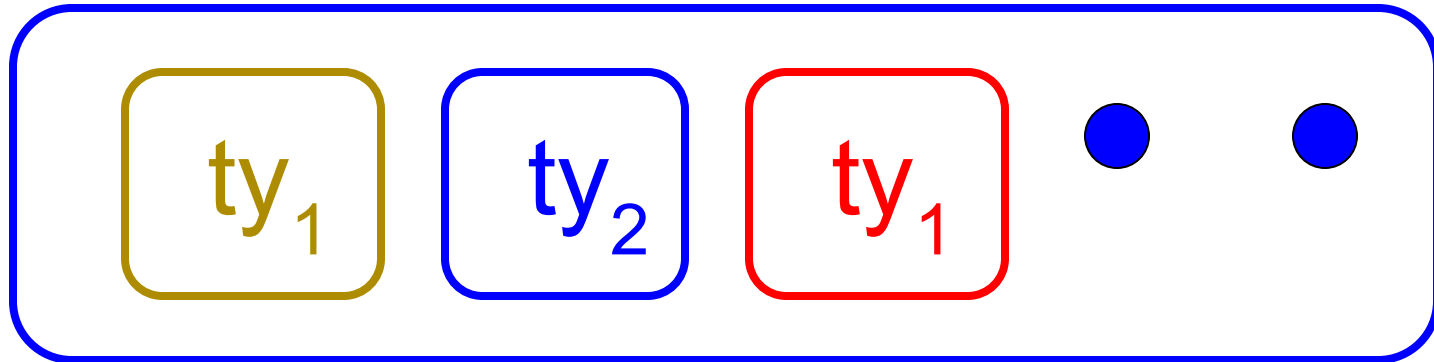


# Disjoint Union Types

---

# Disjoint Union Types as Variants

- Disjoint union of types, with some possibly occurring more than once



- We can also add in some new singleton elements



# Disjoint Union Types

---

(\* Different forms of identification \*)

```
type id = DriversLicense of int  
       | SocialSecurity of int | Name of string
```

```
let check_id id =  
  match id with  
  | DriversLicense num ->  
    not (List.mem num [13570; 99999])  
  | SocialSecurity num -> num < 900000000  
  | Name str -> not (str = "John Doe")
```

Disjoint Union Types



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Disjoint Union Types



# Problem

---

- Create a type to represent the currencies for US, UK, Europe and Japan

**Your turn!**



# Problem

---

- Create a type to represent the currencies for US, UK, Europe and Japan

type currency =

**How many constructors?**





# Problem

---

- Create a type to represent the currencies for US, UK, Europe and Japan

type currency =

```
| (* US *)  
| (* UK *)  
| (* Europe *)  
| (* Japan *)
```

**What currencies?**

Disjoint Union Types



# Problem

---

- Create a type to represent the currencies for US, UK, Europe and Japan

type currency =

| Dollar        (\* US \*)  
| Pound        (\* UK \*)  
| Euro         (\* Europe \*)  
| Yen          (\* Japan \*)

**How to store values?**

Disjoint Union Types



# Problem

---

- Create a type to represent the currencies for US, UK, Europe and Japan

type currency =

| Dollar of int (\* US \*)  
| Pound of int (\* UK \*)  
| Euro of int (\* Europe \*)  
| Yen of int (\* Japan \*)

Disjoint Union Types



# Disjoint Unions in Languages!

---

```
type const =  
  BoolConst of bool  
| IntConst of int  
| FloatConst of float  
| StringConst of string  
| NilConst  
| UnitConst
```

Disjoint Union Types



# Disjoint Unions in Languages!

---

```
type const =  
  BoolConst of bool  
| IntConst of int  
| FloatConst of float  
| StringConst of string  
| NilConst  
| UnitConst
```

- How to represent 7 as a const?



# Disjoint Unions in Languages!

---

```
type const =  
  BoolConst of bool  
| IntConst of int  
| FloatConst of float  
| StringConst of string  
| NilConst  
| UnitConst
```

- How to represent 7 as a const?
- Answer: `IntConst 7`

Disjoint Union Types



Please post questions on Piazza!

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# Polymorphic Datatypes

---





# Polymorphism in Variants

---

- Variants can also be **polymorphic**
- For example, the type `'a option` gives us something to represent non-existence or failure

```
# type 'a option = Some of 'a | None;;
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```
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- Used to encode partial functions
- Often can replace the raising of an exception

Polymorphic Datatypes



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Polymorphic Datatypes



# Functions producing option

---

```
# let rec first p list =  
  match list with  
  | [] -> None  
  | (x :: xs) -> if p x then Some x else first p xs;;  
val first : ('a -> bool) -> 'a list -> 'a option = <fun>  
# first (fun x -> x > 3) [1; 3; 4; 2; 5];;  
- : int option = Some 4  
# first (fun x -> x > 5) [1; 3; 4; 2; 5];;  
- : int option = None
```



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```



# Functions over option

---

```
# let result_ok r =  
  match r with  
  | None -> false  
  | Some _ -> true;;
```

```
val result_ok : 'a option -> bool = <fun>
```

```
# result_ok (first (fun x -> x > 3) [1; 3; 4; 2; 5]);;
```

```
- : bool = true
```

```
# result_ok (first (fun x -> x > 5) [1; 3; 4; 2; 5]);;
```

```
- : bool = false
```

Polymorphic Datatypes





# Functions over option

---

```
# let result_ok r =  
  match r with  
  | None -> false  
  | Some _ -> true;;
```

```
val result_ok : 'a option -> bool = <fun>
```

```
# result_ok (first (fun x -> x > 3) [1; 3; 4; 2; 5]);;
```

```
- : bool = true
```

```
# result_ok (first (fun x -> x > 5) [1; 3; 4; 2; 5]);;
```

```
- : bool = false
```

Polymorphic Datatypes



# Functions over option

---

```
# let result_ok r =  
  match r with  
  | None -> false  
  | Some _ -> true;;
```

```
val result_ok : 'a option -> bool = <fun>
```

```
# result_ok (first (fun x -> x > 3) [1; 3; 4; 2; 5]);;
```

```
- : bool = true
```

```
# result_ok (first (fun x -> x > 5) [1; 3; 4; 2; 5]);;
```

```
- : bool = false
```

Polymorphic Datatypes



# Problem

---

- Write a `hd` on lists that doesn't raise an exception and works at all types of lists.

**Your turn!**



# Problem

---

- Write a `hd` on lists that doesn't raise an exception and works at all types of lists.

let `hd list =`

`match list with`

**Nil case?**



# Problem

---

- Write a `hd` on lists that doesn't raise an exception and works at all types of lists.

let `hd list =`

`match list with`

`| [] -> None`

**Cons case?**



# Problem

---

- Write a `hd` on lists that doesn't raise an exception and works at all types of lists.

```
let hd list =  
  match list with  
  | [] -> None  
  | (x :: xs) -> Some x
```



# Mapping over Variants

---

```
# let optionMap f opt =  
  match opt with  
  | None -> None  
  | Some x -> Some (f x);;
```

```
val optionMap :
```

```
('a -> 'b) -> 'a option -> 'b option = <fun>
```

```
# optionMap
```

```
(fun x -> x - 2)
```

```
(first (fun x -> x > 3) [1; 3; 4; 2; 5]);;
```

```
- : int option = Some 2
```

Polymorphic Datatypes

# Mapping over Variants

```
# let optionMap f opt =  
  match opt with  
  | None -> None  
  | Some x -> Some (f x);;
```

```
val optionMap :
```

```
('a -> 'b) -> 'a option -> 'b option = <fun>
```

```
# optionMap
```

```
(fun x -> x - 2)
```

```
(first (fun x -> x > 3) [1; 3; 4; 2; 5]);;
```

```
- : int option = Some 2
```

Polymorphic Datatypes





# Folding over Variants

---

```
# let optionFold someFun noneVal opt =  
  match opt with  
  | None -> noneVal  
  | Some x -> someFun x;;
```

```
val optionFold :
```

```
('a -> 'b) -> 'b -> 'a option -> 'b = <fun>
```

```
# let optionMap f opt =
```

```
  optionFold (fun x -> Some (f x)) None opt;;
```

```
val optionMap :
```

```
('a -> 'b) -> 'a option -> 'b option = <fun>
```

Polymorphic Datatypes



# Folding over Variants

---

```
# let optionFold someFun noneVal opt =  
  match opt with  
  | None -> noneVal  
  | Some x -> someFun x;;
```

```
val optionFold :
```

```
('a -> 'b) -> 'b -> 'a option -> 'b = <fun>
```

```
# let optionMap f opt =  
  optionFold (fun x -> Some (f x)) None opt;;
```

```
val optionMap :
```

```
('a -> 'b) -> 'a option -> 'b option = <fun>
```

Polymorphic Datatypes



Please post questions on Piazza!

---

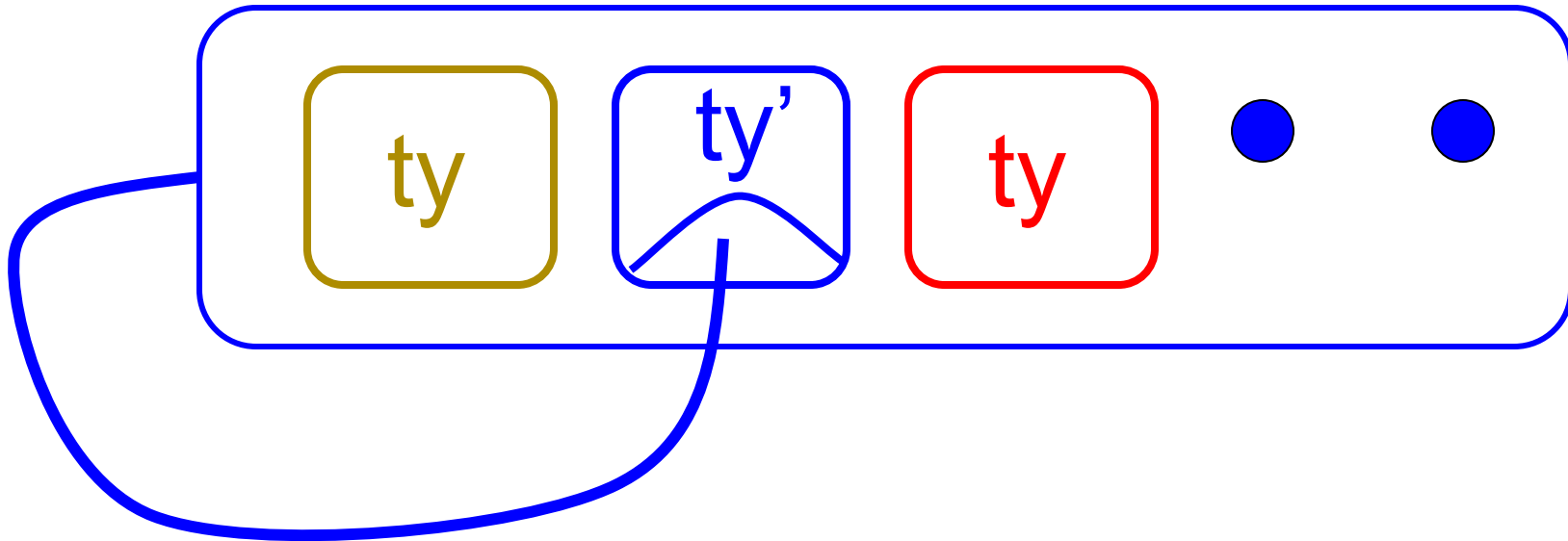


# Preview: Recursive Datatypes

---

# Recursive Types as Variants

- The type being defined may be a component of itself





# Recursive Data Types

---

```
type int_Bin_Tree =  
  Leaf of int | Node of (int_Bin_Tree * int_Bin_Tree)
```

```
let my_tree =  
  Node (Node (Leaf 3, Leaf 6), Leaf (-7))
```



# Recursive Data Types

---

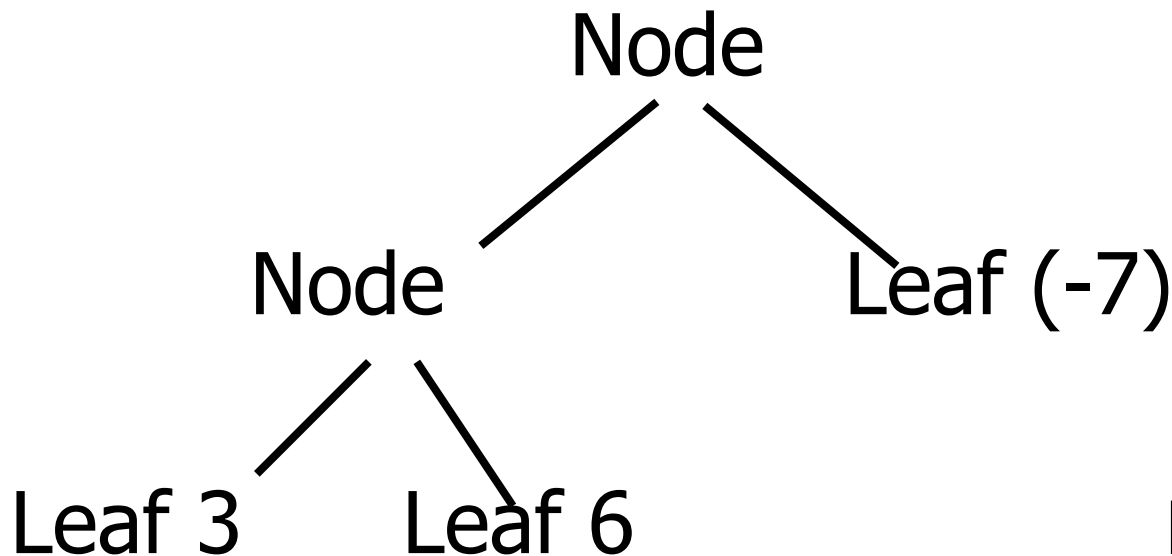
```
type int_Bin_Tree =  
  Leaf of int | Node of (int_Bin_Tree * int_Bin_Tree)
```

```
let my_tree =  
  Node (Node (Leaf 3, Leaf 6), Leaf (-7))
```

# Recursive Data Types

```
type int_Bin_Tree =  
  Leaf of int | Node of (int_Bin_Tree * int_Bin_Tree)
```

```
let my_tree =  
  Node (Node (Leaf 3, Leaf 6), Leaf (-7))
```







# Recursive Data Types in Languages!

---

```
# type exp =  
  | VarExp of string  
  | ConstExp of const  
  | MonOpAppExp of mon_op * exp  
  | BinOpAppExp of bin_op * exp * exp  
  | IfExp of exp * exp * exp  
  | AppExp of exp * exp  
  | FunExp of string * exp
```



How do we even represent the syntax of our language, and map over it to transform programs?

---



How to implement automatically in compiler, rather than by hand?

---



Please post questions on Piazza!

---



# Takeaways

---

- **Variants** let us represent custom **datatypes**
  - Can be **polymorphic**
  - Can be **recursive**
  - Can represent **lists** and **trees**
  - Can represent **language syntax!**
- Can do **two things** with them:
  - **construct**
  - **deconstruct** (match, eliminate)
- Can write **program transformations**, **interpreters**, and **compilers** this way :)



## Next Class

---

- **I will be back!** Lecture will happen in person
- **EC1** is due, if interested (**extra credit**)
- **WA3XC** also due, if interested (**extra credit**)
- **MP4** will be due next Tuesday
- All deadlines can be found on **course website**
- Use **office hours** and **class forums** for help