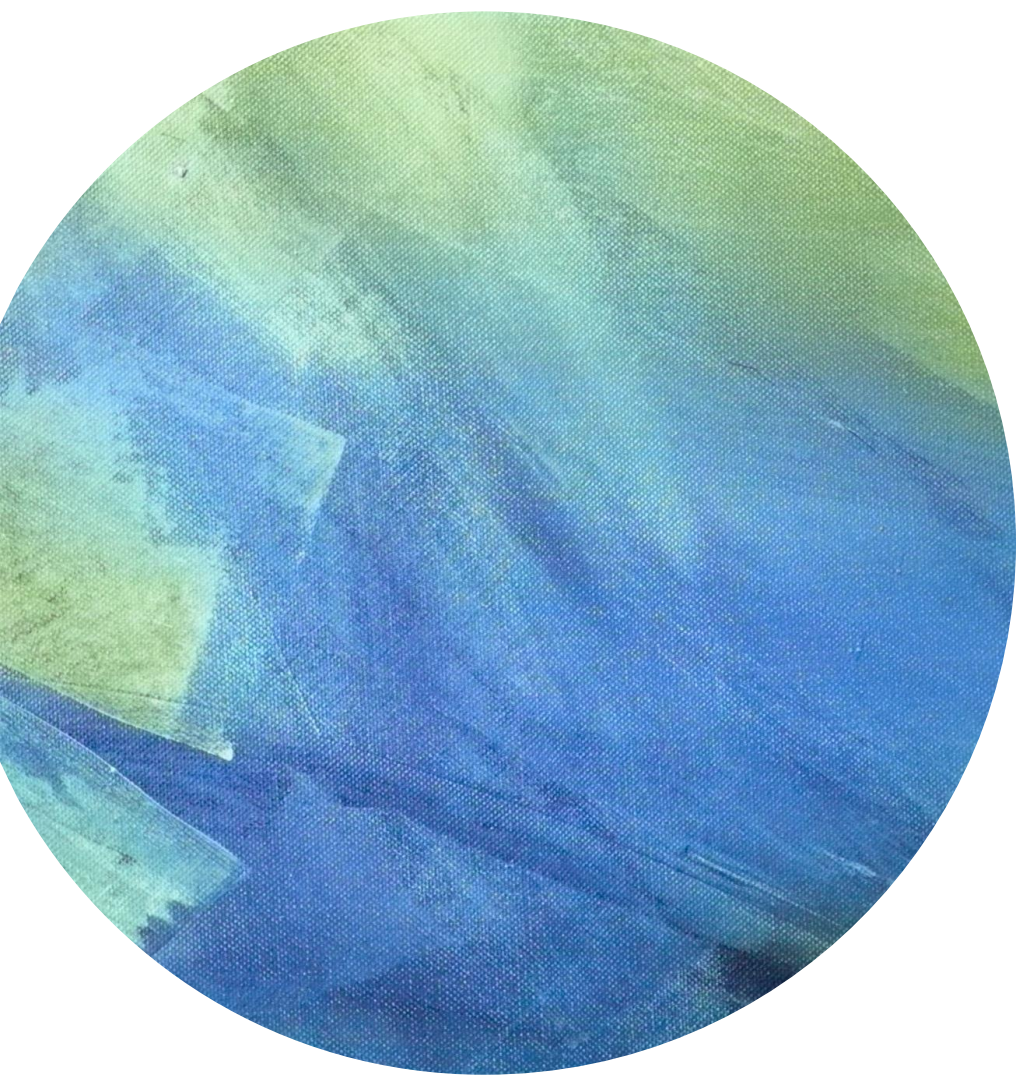


The background of the slide is an abstract composition of broad, textured brushstrokes. The top half features a mix of light and dark green strokes, while the bottom half is dominated by various shades of blue, from deep navy to a lighter, more vibrant blue. The strokes are layered and overlapping, creating a sense of depth and movement.

Lecture 19



Outline



Wrap-up commitments



NIZKs from pairings

Pedersen Commitments

Pedersen Commitments

- Unconditionally hiding
 - Given a commitment c , every value x is equally likely to be the value committed in c .
 - For example, given x, r , and any x' , there exists r' such that $g^x h^r = g^{x'} h^{r'}$, in fact $r = (x - x')a^{-1} + r \bmod q$.

Pedersen Commitments

- Computationally binding
 - Suppose committer sent $g^x h^r \bmod p$ for some (x, r)
 - Now it finds $x' \neq x$ and r' such that $c = g^{x'} h^{r'}$.
 - This means that the sender ``knows'' $\log_g(h) = (x' - x) \cdot (r - r')^{-1}$.
 - This means: assuming DL is hard, the sender cannot open the commitment to a different value.

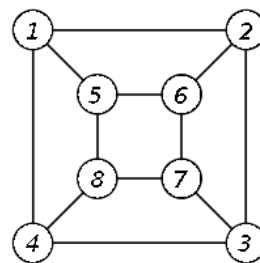
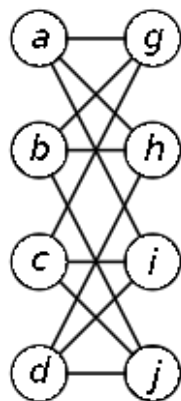
Application: Coin Tossing

- Alice and Bob want to decide on something by tossing a coin over a phone. How to do this securely?
- Solution: Alice commits to a random bit $b_A \leftarrow \{0, 1\}$, and sends $\text{Com}(b_A; r)$ to Bob
- Bob selects a random bit $b_B \leftarrow \{0, 1\}$ and sends it to Alice
- Alice decommits b_A
- Alice and Bob output $b_A \text{ xor } b_B$

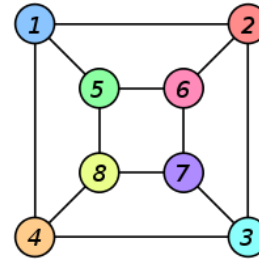
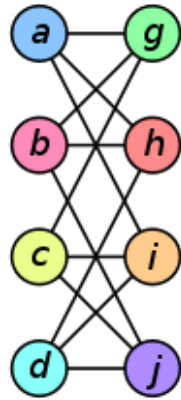
Zero-Knowledge

Problems in NP

Graph Isomorphism



Graph Isomorphism



Real World

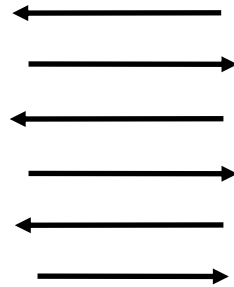
Prover



NP Statement x

Witness that x is true

Verifier



Real World

Prover



NP Statement x

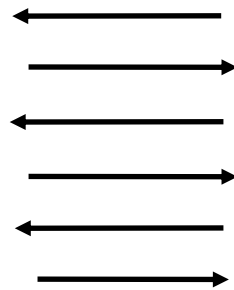
Witness that x is true

Verifier

Didn't learn
witness



Outputs
view



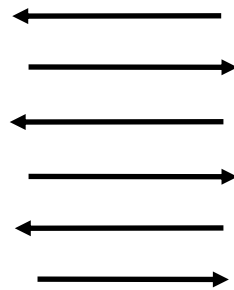
Real World

Prover



NP Statement x

Witness that x is true



Verifier

Didn't learn
witness



Outputs
view

Ideal World (Proof)

Simulator



NP Statement x

No witness

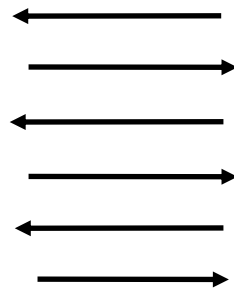
Real World

Prover



NP Statement x

Witness that x is true



Verifier

Didn't learn
witness



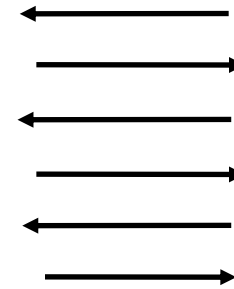
Outputs
view

Ideal World (Proof)

Simulator



NP Statement x
No witness



Verifier



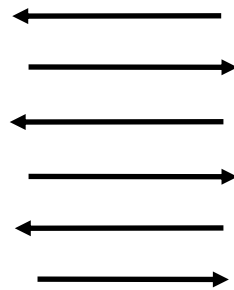
Real World

Prover



NP Statement x

Witness that x is true



Verifier

Didn't learn
witness



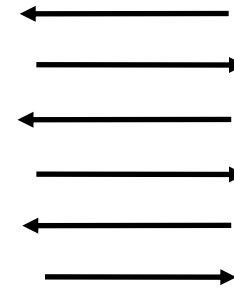
Outputs
view

Ideal World (Proof)

Simulator



NP Statement x
No witness



Verifier

Didn't learn
witness



Outputs similar
view

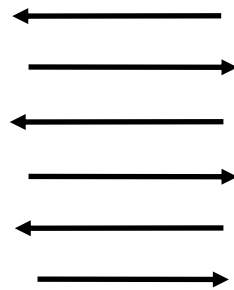
Real World

Prover



NP Statement x

Witness that x is true



Verifier

Didn't learn
witness



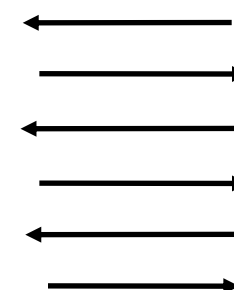
Outputs
view

Ideal World (Proof)

Simulator



NP Statement x
No witness



Verifier

Didn't learn
witness

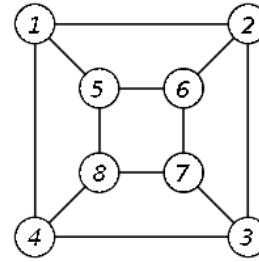
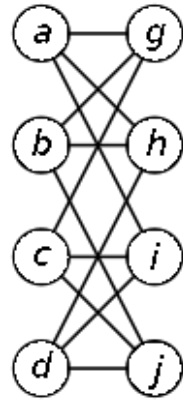


Outputs similar
view

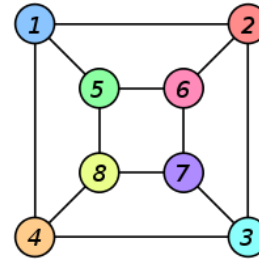
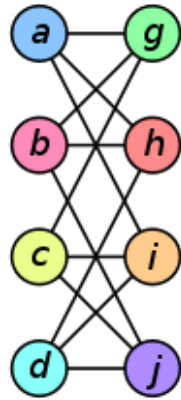
D

Cannot distinguish the two

Graph Isomorphism



Graph Isomorphism

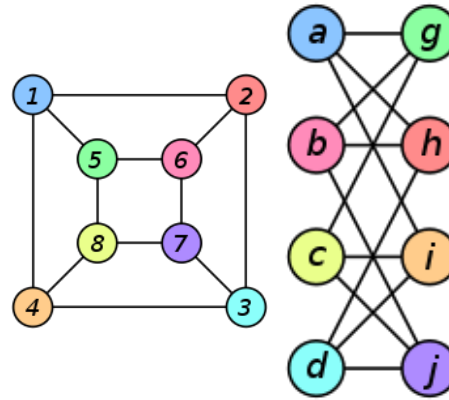


Graph Isomorphism

Prover

$X = (A, B)$

Knows η s.t.
 $A = \eta(B)$



Verifier



Graph Isomorphism

Prover

$X = (A, B)$

Knows η s.t.
 $A = \eta(B)$

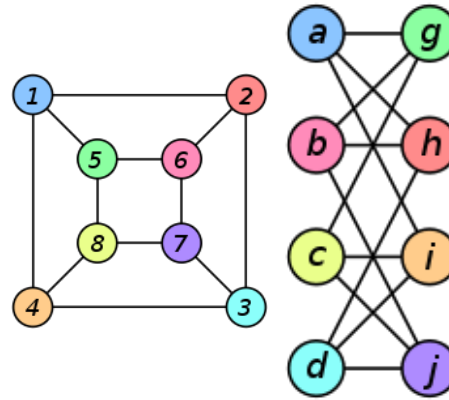


A

			1	1	1	
			1	1		1
			1		1	1
				1	1	1
1	1	1				
1	1		1			
1		1	1			
	1	1	1			

$G = \varphi(A)$

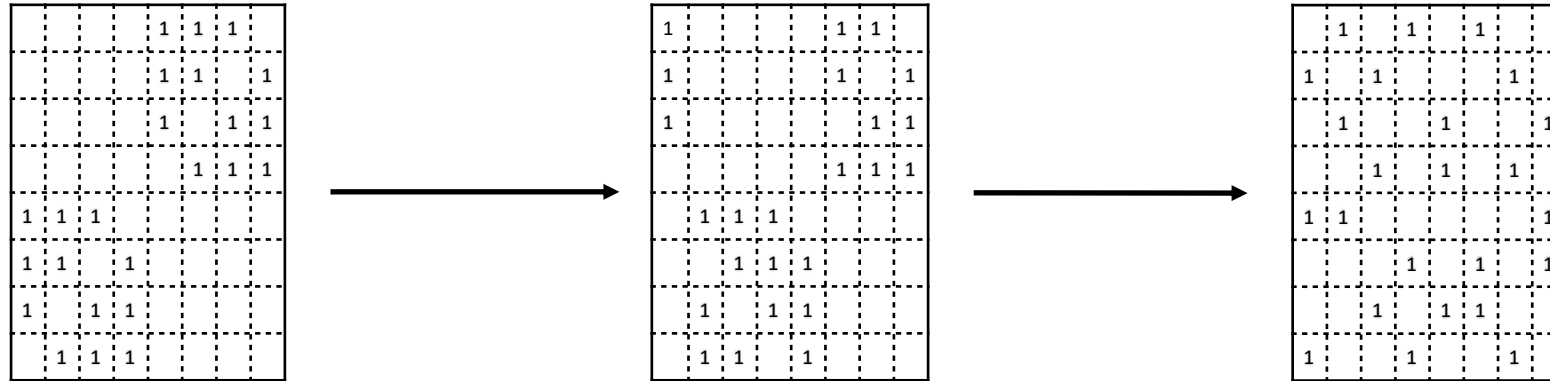
	1		1		1	
1		1				1
	1			1		1
		1		1		1
1	1					1
			1		1	1
		1		1	1	
1			1			1



Verifier



Permuting the Graph



$$G = \varphi(A)$$

Graph Isomorphism

Prover

$X = (A, B)$

Knows η s.t.
 $A = \eta(B)$

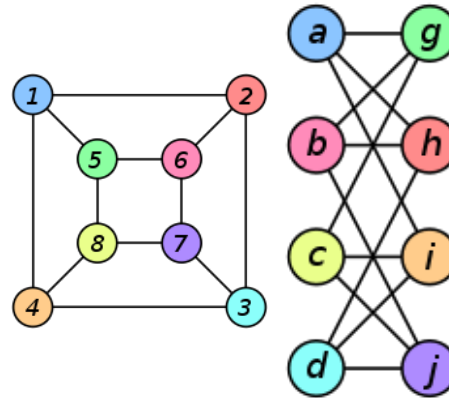


A

			1	1	1
			1	1	
			1		1
				1	1
1	1	1			
1	1		1		
1		1	1		
	1	1	1		

$G = \varphi(A)$

	1		1		1
1		1			1
	1		1		1
		1	1		
1	1				1
		1	1	1	1
			1	1	1
1			1		1



Verifier

G

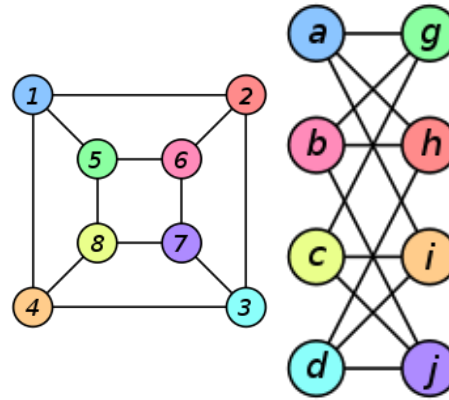
$c = A \text{ or } B$



Graph Isomorphism

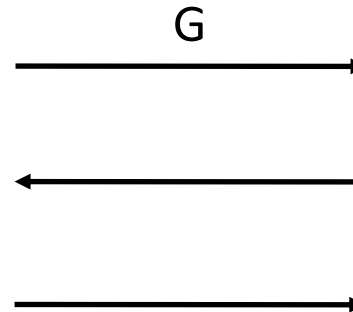
Simulator

$X = (A, B)$



Verifier

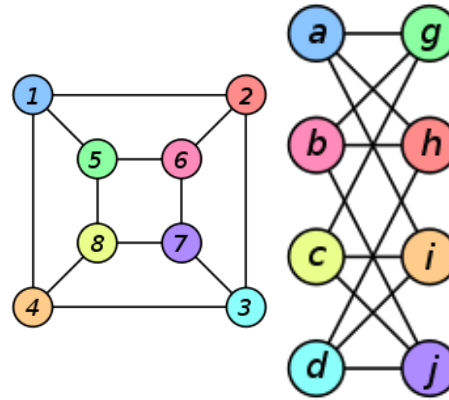
Knows c in advance



Graph Isomorphism

Simulator

$X = (A, B)$

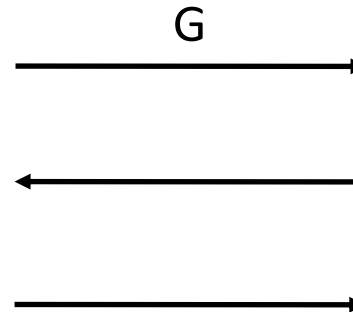


Verifier

Knows c in advance



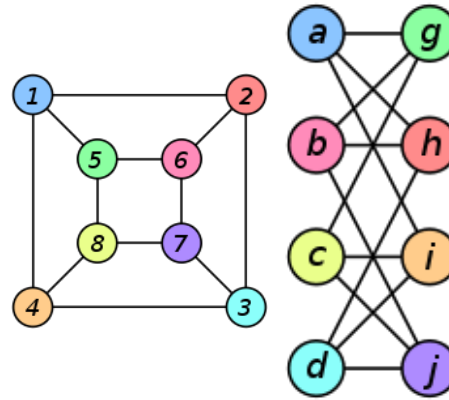
$G = \varphi(c)$



Graph Isomorphism

Simulator

$X = (A, B)$



Verifier

Knows c in advance



$G = \varphi(c)$

