



NFA → DFA

Thm If L is accepted by some NFA M ,
then L is accepted by some DFA M' .

Pf: **idea** - remember a subset of states

Given NFA $M = (Q, \Sigma, s, \delta, A)$, $\delta: Q \times (\Sigma \cup \{\epsilon\}) \rightarrow \underline{\underline{Q}}$

Construct DFA $M' = (Q', \Sigma, s', \delta', A')$

$$\delta': \underline{\underline{Q'}} \times \Sigma \rightarrow \underline{\underline{Q'}}$$

where $Q' = \underline{\underline{2^Q}}$ (power set of Q)
(= {all subsets of Q })

$$s' = \epsilon\text{-reach}(s)$$

$$A' = \{S \in Q' : S \cap A \neq \emptyset\}$$

Called
subset construction
(or power-set
construction)

$$\delta'(\underline{\underline{S}}, a) = \bigcup_{q \in S} \delta^*(q, a) \quad \forall S \in Q', \forall a \in \Sigma$$

Lemma $\delta'^*(S, x) = \bigcup_{q \in S} \delta^*(q, x) \quad \forall S \in Q', \forall x \in \Sigma^*$

Pf: by induction (skipped). \square

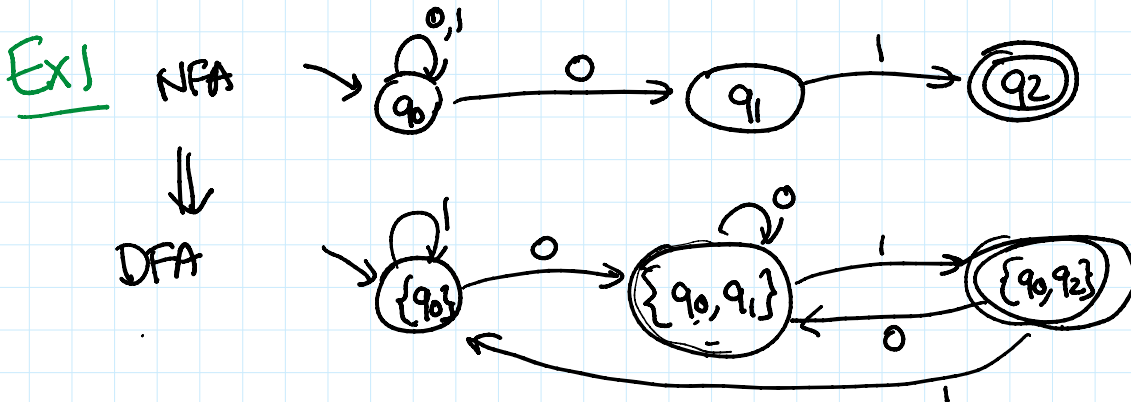
$$\left(\text{Then } x \in L(M') \iff \delta'^*(s', x) \in A' \right.$$

def of
accept of
DFA

$$\iff \bigcup_{q \in \epsilon\text{-reach}(s)} \delta^*(q, x) \in A'$$

$$\begin{aligned} &\Leftrightarrow \delta^*(s, x) \in A' \\ &\stackrel{\text{def of } A'}{\Leftrightarrow} \delta^*(s, x) \cap A \neq \emptyset \\ &\stackrel{\text{def of accept in NFA}}{\Leftrightarrow} x \in L(M). \end{aligned}$$

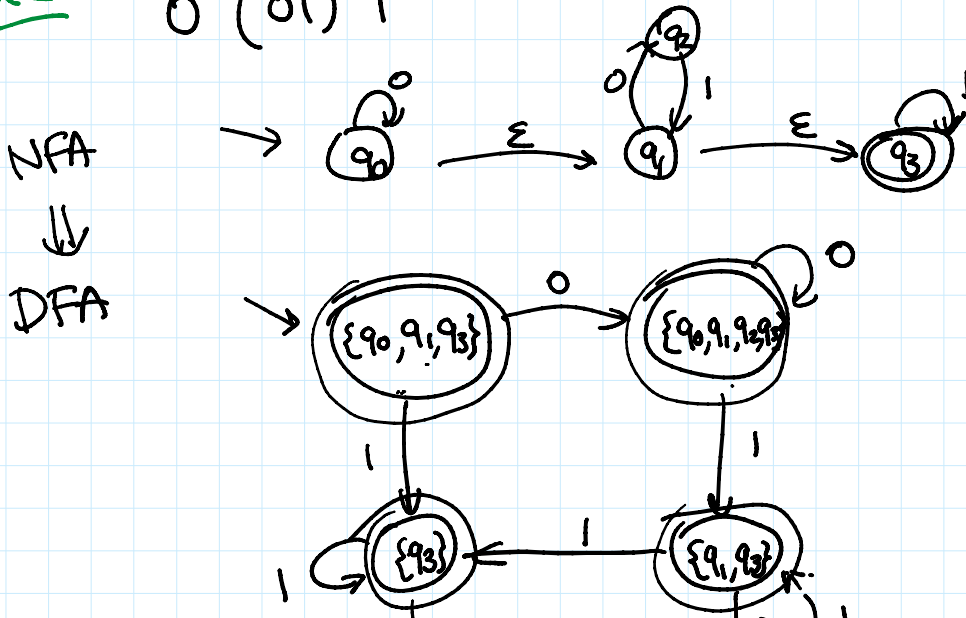
$$L(M') = L(M). \quad \square$$

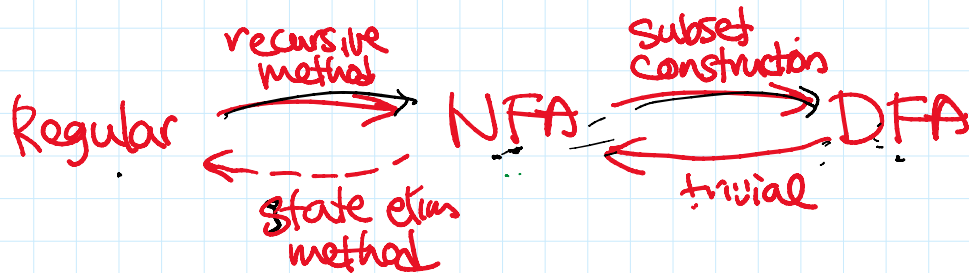
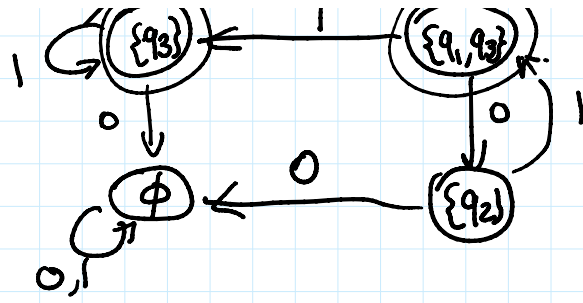


$$\begin{aligned} \delta'(\{q_0, q_1\}, 0) &= \delta(q_0, 0) \cup \delta(q_1, 0) \\ &= \{q_0, q_1\} \cup \emptyset = \{q_0, q_1\}. \\ \delta'(\{q_0, q_1\}, 1) &= \delta(q_0, 1) \cup \delta(q_1, 1) \\ &= \{q_0\} \cup \{q_2\} = \{q_0, q_2\} \end{aligned}$$

Ex2

$$0^*(01)^*1^*$$





NFA \rightarrow Regular

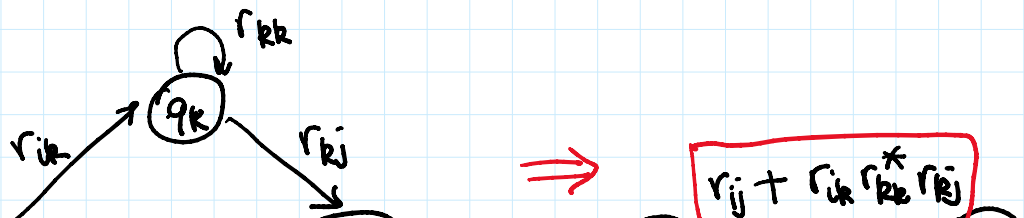
Thm If L is accepted by NFA M then L is regular.

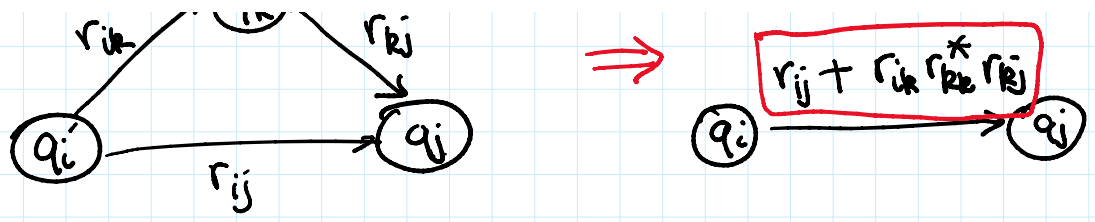
Pf. Sketch: Given $M = (Q, \Sigma, s, \delta, A)$
 $Q = \{q_0, q_1, \dots, q_{n-1}\}$

state elimination method

0. add new state s', f' with ϵ -transition from s' to s
 & from A to f'

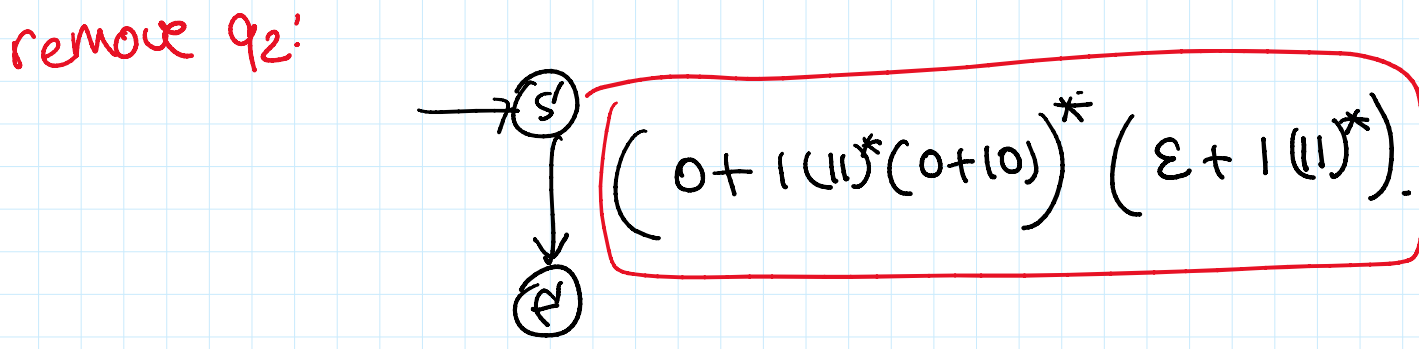
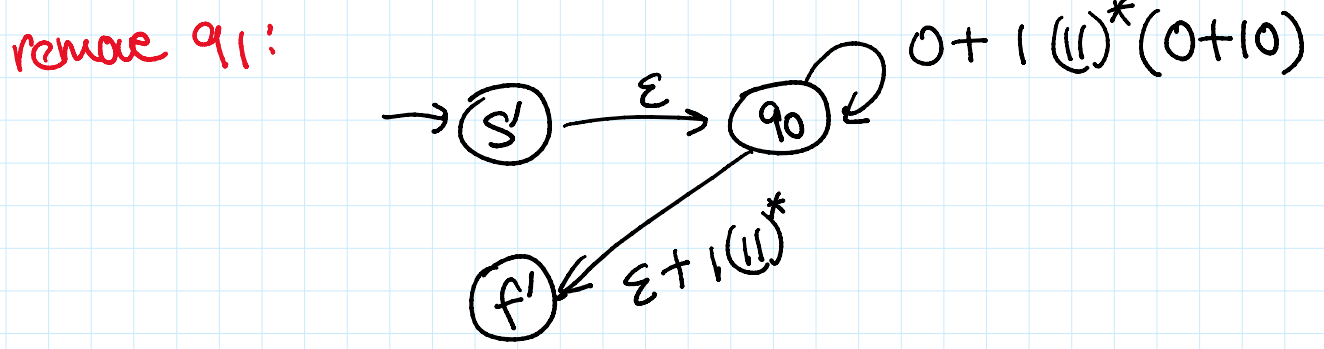
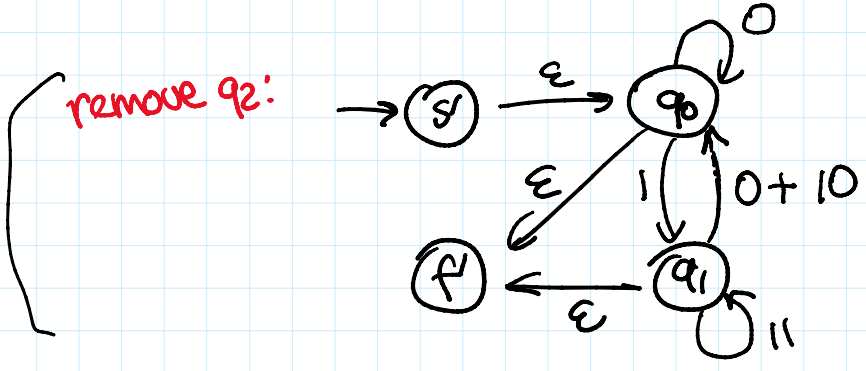
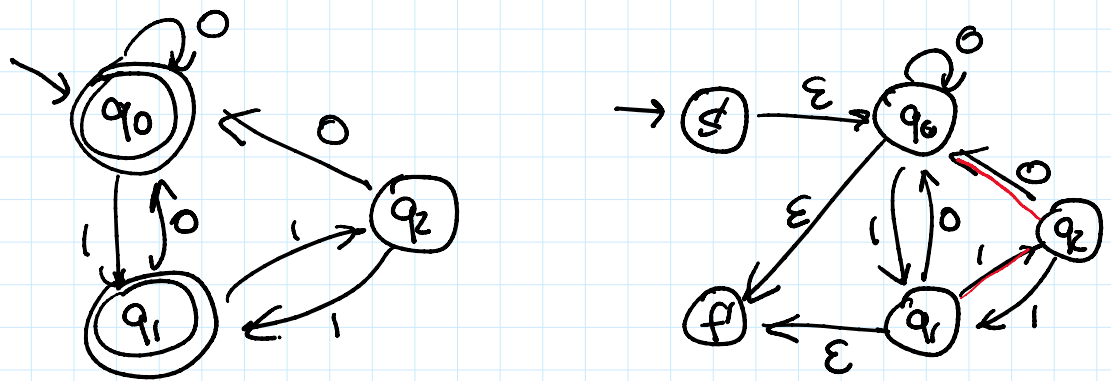
1. for $k = n-1$ to 0:
 remove q_k & apply the rule below:





2. return label from s' to f' . □

Ex



Kleene's Thm (1956)

L is regular iff L is accepted by some DFA.

Corollaries

If L is regular, then \bar{L} is also regular.

If L_1, L_2 are regular, then
so is $L_1 \cap L_2$.

If L is accepted by some DFA,
then so is L^* .
is L^R .

many closure props . . .