

17.3.5

The basic algorithm: Find the i th closest vertex

A Basic Strategy

Explore vertices in increasing order of distance from s :

(For simplicity assume that nodes are at different distances from s and that no edge has zero length)

```
Initialize for each node  $v$ ,  $\text{dist}(s, v) = \infty$   
Initialize  $X = \{s\}$ ,  
for  $i = 2$  to  $|V|$  do  
    (* Invariant:  $X$  contains the  $i - 1$  closest nodes to  $s$  *)  
    Among nodes in  $V - X$ , find the node  $v$  that is the  
         $i$ th closest to  $s$   
    Update  $\text{dist}(s, v)$   
     $X = X \cup \{v\}$ 
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How can we implement the step in the for loop?

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How can we implement the step in the for loop?

Finding the i th closest node

- 1 X contains the $i - 1$ closest nodes to s
- 2 Want to find the i th closest node from $V - X$.

What do we know about the i th closest node?

Claim

Let P be a shortest path from s to v where v is the i th closest node. Then, all intermediate nodes in P belong to X .

Proof.

If P had an intermediate node u not in X then u will be closer to s than v . Implies v is not the i th closest node to s - recall that X already has the $i - 1$ closest nodes. \square

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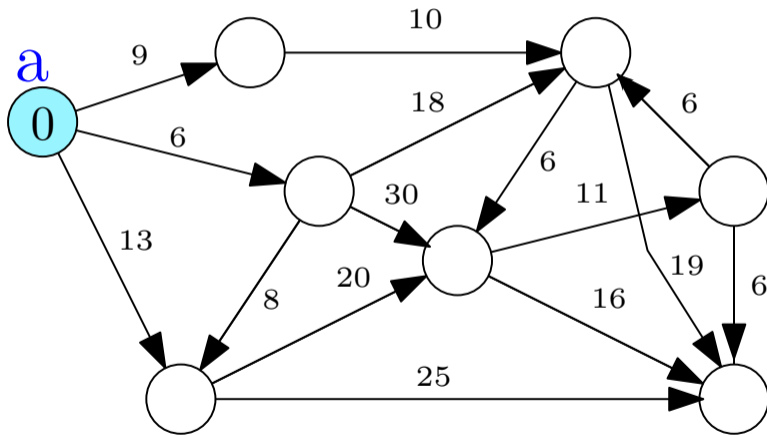
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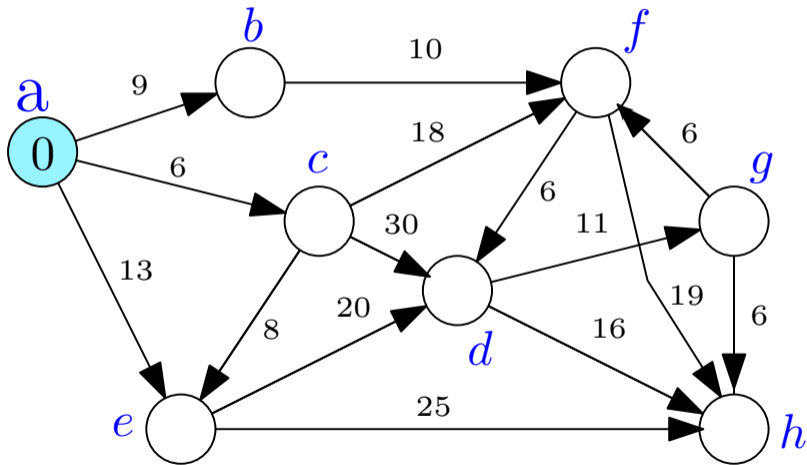
Finding the i th closest node repeatedly

An example



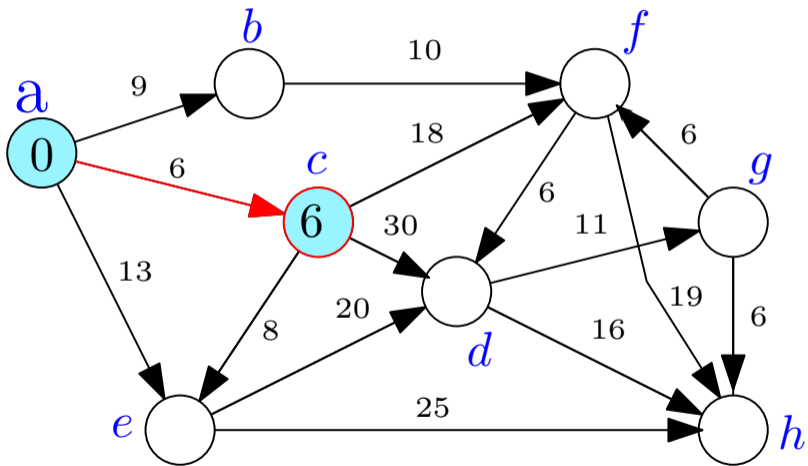
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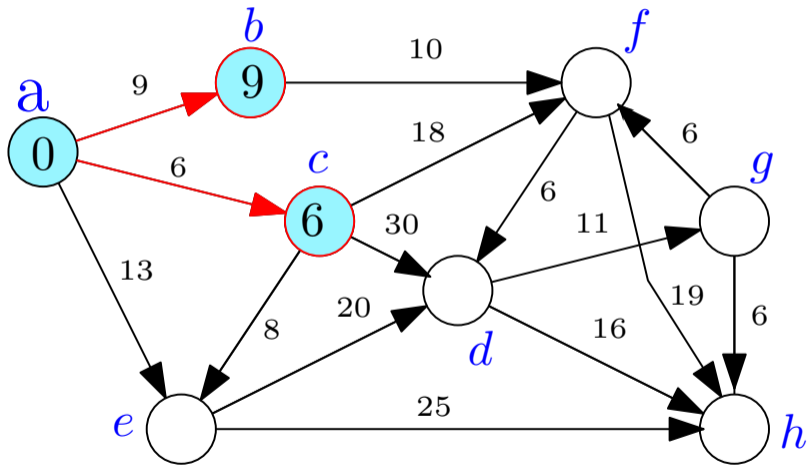
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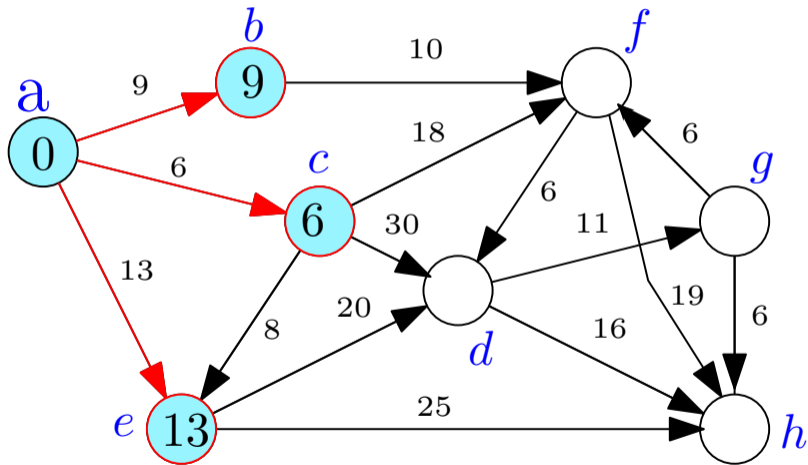
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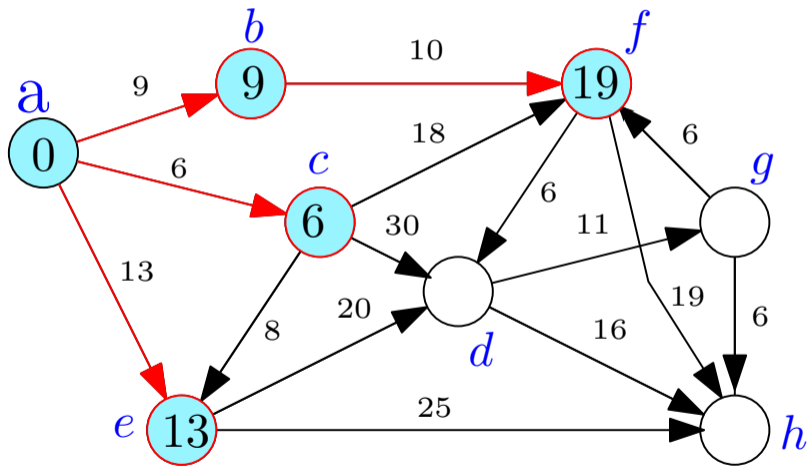
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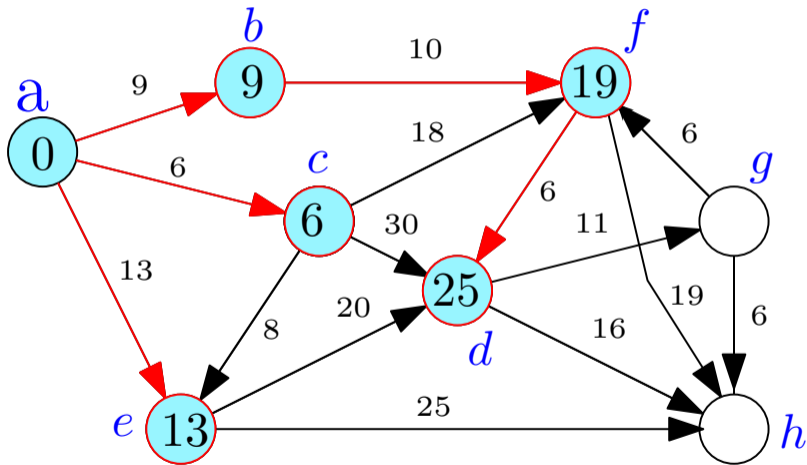
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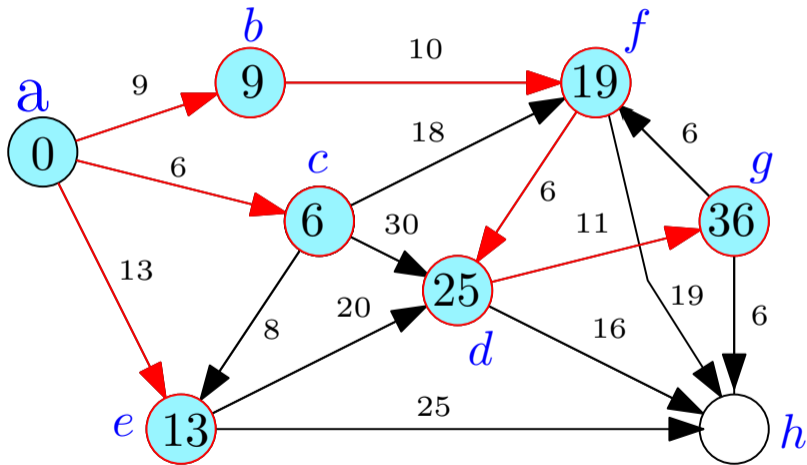
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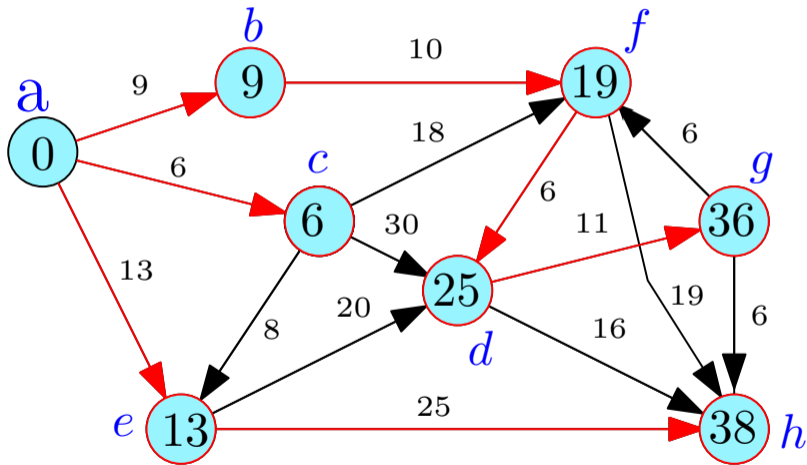
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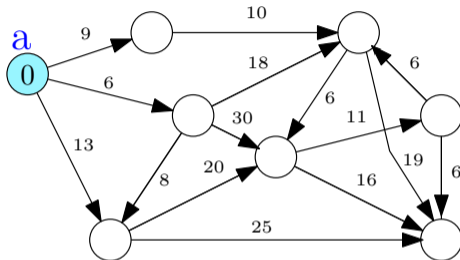


Finding the i th closest node repeatedly

An example



Finding the i th closest node



Corollary

The i th closest node is adjacent to X .

Summary

Proved that the basic algorithm is (intuitively) correct...

...but is missing details

...and how to implement efficiently?

THE END

...

(for now)