

12.4.1

Running time analysis

Running time of LIS($[1..n]$)

```
LIS_smaller( $A[1..n]$ ,  $x$ ):  
  if ( $n = 0$ ) then return 0  
   $m = \mathbf{LIS\_smaller}(A[1..(n - 1)], x)$   
  if ( $A[n] < x$ ) then  
     $m = \mathbf{max}(m, 1 + \mathbf{LIS\_smaller}(A[1..(n - 1)], A[n]))$   
  Output  $m$ 
```

```
LIS( $A[1..n]$ ):  
  return  $\mathbf{LIS\_smaller}(A[1..n], \infty)$ 
```

Running time of LIS([1..n])

Lemma

LIS_smaller runs in $O(2^n)$ time.

Improvement: From $O(n2^n)$ to $O(2^n)$.

....one can do much better using memoization!

Running time of LIS([1..n])

Lemma

LIS_smaller runs in $O(2^n)$ time.

Improvement: From $O(n2^n)$ to $O(2^n)$.

....one can do much better using memoization!

Running time of LIS([1..n])

Lemma

LIS_smaller runs in $O(2^n)$ time.

Improvement: From $O(n2^n)$ to $O(2^n)$.

....one can do much better using memoization!

THE END

...

(for now)