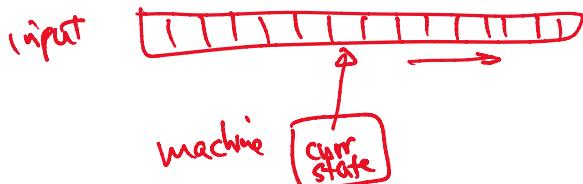
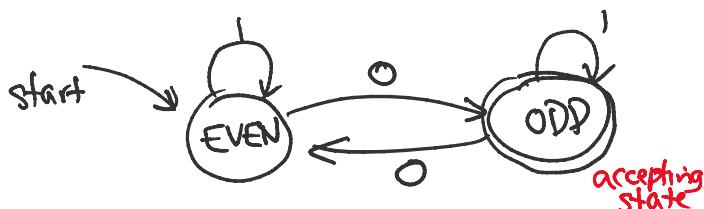


# Deterministic Finite Automata (DFA)

Intuitively,  
machine that uses const amount of memory  
program & reads input in pass  
(from left to right)  $\xrightarrow{Q}$   
const # of "states"

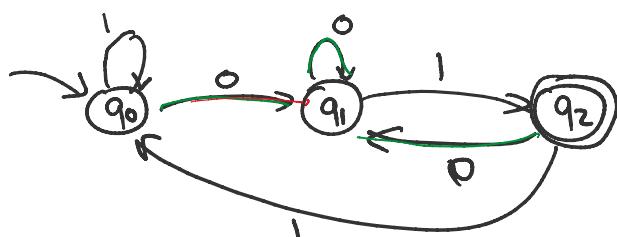


Ex0 all strings over  $\{0,1\}$  with odd # of 0's



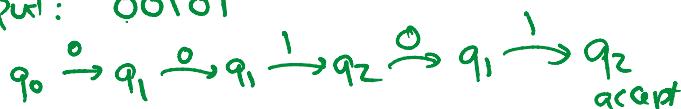
$$\begin{aligned} Q &= \{q_0, q_1, q_2\} \\ \Sigma &= \{0, 1\} \\ S &= q_0 \\ A &= \{q_2\} \end{aligned}$$

Ex1 all strings ending with 01



States:  
 $q_1$ : just seen 0  
 $q_2$ : just seen 01  
 $q_0$ : none of above

e.g. Input: 00101



$$\begin{aligned} \delta^*(q_0, 001) &= \delta^*(\delta(q_0, 0), 01) \\ &= \delta^*(q_1, 01) \\ &= \delta^*(\delta(q_1, 0), 1) \\ &= \delta^*(q_2, 1) \\ &= \delta^*(\delta(q_2, \epsilon), \epsilon) = \delta^*(q_2, \epsilon) = q_2 \end{aligned}$$

state =  $q_0$   
while (not end of input) {  
  c = next input symbol;  
  if (state ==  $q_0$  && c == 0)  
    state =  $q_1$ ;  
  else if (state ==  $q_0$  && c == 1)  
    state =  $q_2$ ;

}:  $O(n)$  time  
&  $O(1)$  space  
  if (state ==  $q_2$ )  
    output yes  
  else ... no

Formal Def'n A DFA is specified by 5 things:

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$$M = (Q, \Sigma, s, \delta, A)$$

$Q$  is a finite set of states

$\Sigma$  is finite alphabet

$s \in Q$  is the start state

$A \subseteq Q$  is the set of accepting states

$\delta: Q \times \Sigma \rightarrow Q$  is the transition fn

Curv state Curv symbol next state

Def Given transition fn  $\delta$ ,  
define its extended transition fn

$\delta^*: Q \times \Sigma^* \rightarrow Q$  inductively:

$$(i) \quad \delta^*(q, \epsilon) = q$$

$$(ii) \quad \delta^*(q, x) = \delta^*(\delta(q, a), y) \text{ if } x = ay \text{ with } a \in \Sigma, y \in \Sigma^*$$

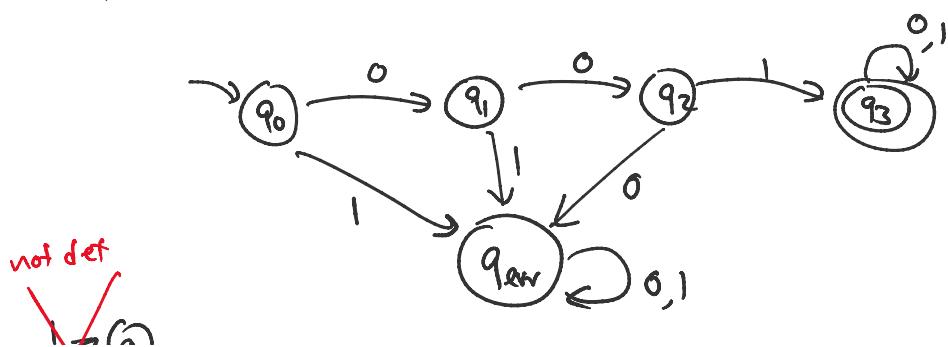
Def  $M$  accepts  $x$  iff  $\delta^*(s, x) \in A$

Define  $L(M) = \{x \in \Sigma^*: M \text{ accepts } x\}$

Lang. accepted by  $M$

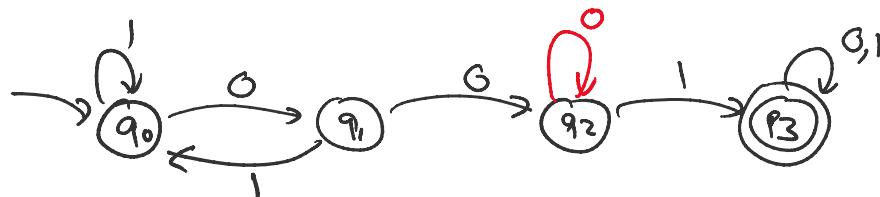
Exs ( $\Sigma = \{0, 1\}$ )

a) all strings beginning with 001





b) all strings containing 001 as a substring



$q_3$ : found 001

$q_2$ : just seen 00 but not found 001 yet

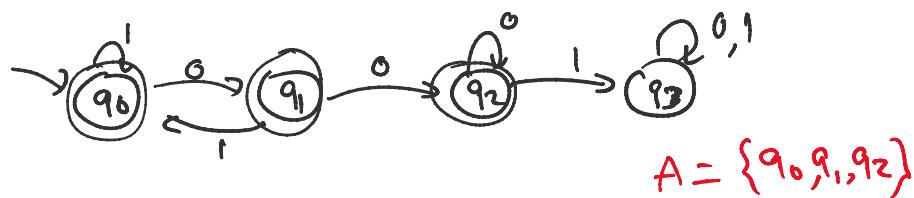
$q_1$ : just seen 0 but not in  $q_2, q_3$

$q_0$ : none of above

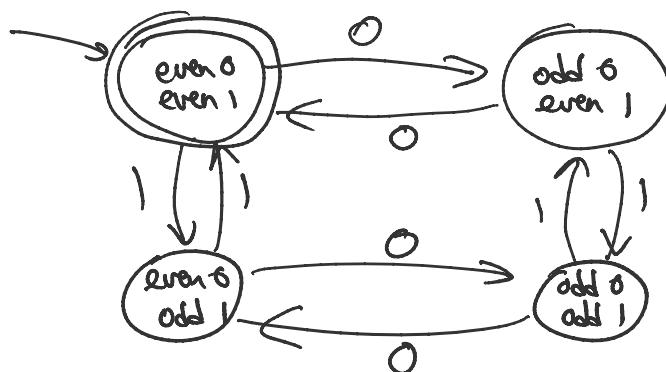
⇒ generalizes to other pattern strings

⇒ pattern matching alg'm with  $O(n)$  time

c) all strings not containing 001



d) all strings with even # of 0's  
and even # of 1's



e) strings with length divisible by 5

