

# F'18 Mid-term Problem 3:

Given  $L$ ,  
 $\text{prefix}(L) = \{ u \mid uv \in L \}$   $(u, v \in \Sigma^*)$

To show:  $L$  is regular  $\Rightarrow$   $\text{prefix}(L)$  is regular.

By regular-expressions.

(or DFAs/NFAs)

$r = \emptyset \Rightarrow \emptyset$

$r = \epsilon \Rightarrow \epsilon$

$r = a \ (a \in \Sigma) \Rightarrow a + \epsilon$

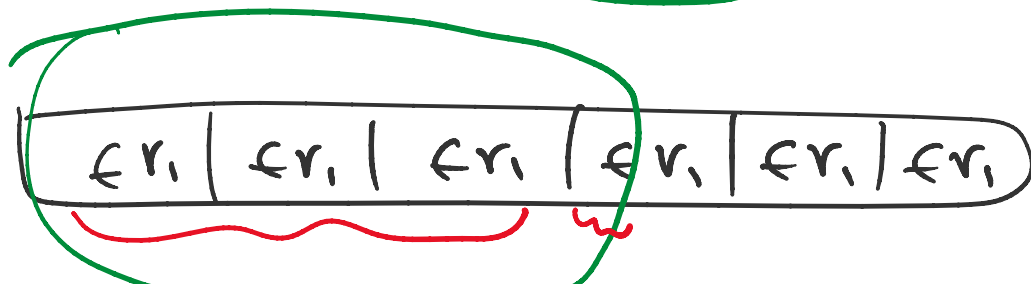
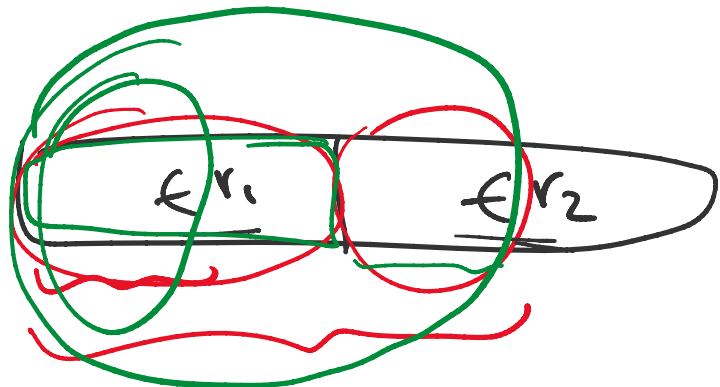
Same  $\delta$   
 $A' = \{ q \in Q : \exists \text{ path from } q \text{ to a state in } A \text{ in DFA} \}$

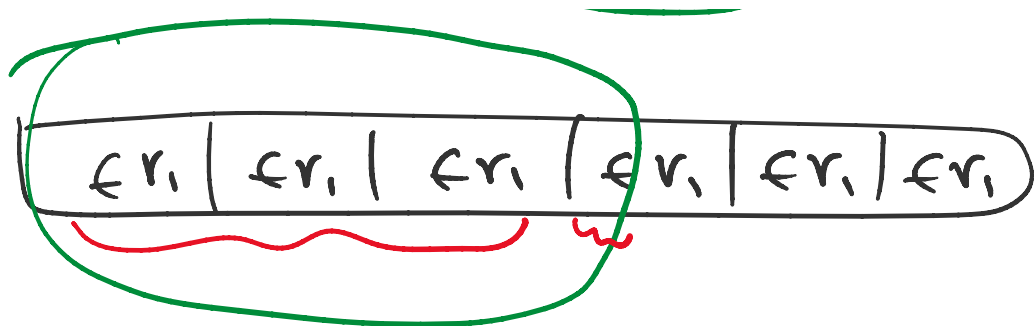
$r = r_1 + r_2$ : SUPPOSE  $S_1$  for  $\text{prefix}(L(r_1))$   
 $S_2$  for  $\text{prefix}(L(r_2))$

$\Rightarrow S_1 + S_2$

$r = r_1 r_2 \Rightarrow (S_1) + (r_1)(S_2)$

$r = r_1^*$





$$\Rightarrow \underline{r_1^*} \underline{S_1}$$

35.  $\{ 0^{F_n} \mid n \geq 0 \}$  is not regular

$F_n =$  Fibonacci numbers  
 $(F_n = F_{n-1} + F_{n-2}$   
 $F_0 = 0, F_1 = 1).$

By fooling set method.

Choose  $F = \{ 0^{F_n} \mid n \geq 3 \}$ .

Let  $x, y$  be 2 arbitrary distinct strings from  $F$ .

Then  $x = 0^{F_i}$  and  $y = 0^{F_j}$   
 for some  $i, j \geq 3, i \neq j$ .

W.l.o.g.  $i > j$ . (otherwise, can swap  $i < j$ ).

Choose  $z = 0^{F_{i+1}}$

$\Rightarrow xz = 0^{F_i + F_{i+1}} = 0^{F_{i+2}} \in L$

$$\Rightarrow xz = 0^{F_i + F_{i+1}} = 0^{i+2} \in L$$

$$yz = 0^{\underbrace{F_j + F_{i+1}}} \notin L.$$

Since  $F_j + F_{i+1}$  is not Fibonacci number.

$$F_{\underbrace{i+1}} < F_j + F_{i+1} < F_i + F_{i+1} = F_{i+2}$$

ok

Thus,  $F$  is a fooling set.

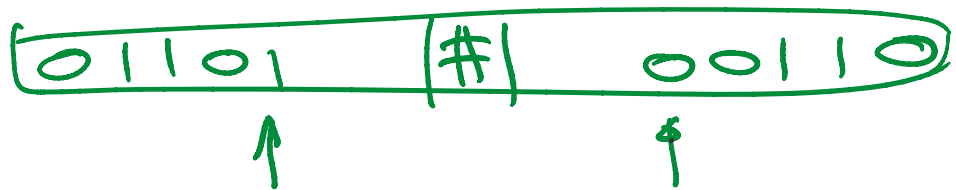
Since  $F$  is infinite,  $L$  is not regular.

□

93.  $\{x\#y^R \mid x, y \in \{0,1\}^*, x \neq y\}$

Give CFG.

$$\Sigma = \{0, 1, \#\}$$



$$S \rightarrow 0S0 \mid \underline{1S1}$$

$$1A0 \mid 0A1 \mid \underline{1B\#}$$

$$0B\# \mid \#B1 \mid \#B0$$

11#1  
~~1(1)#1~~

$$A \rightarrow 0A \mid 1A \mid \underline{\#B}$$

! ! ! !

$$A \rightarrow 0A \mid 1A \mid \underline{\#} \mid \epsilon$$

$$B \rightarrow 0B \mid \underline{1B} \mid \epsilon$$

(A generates all strings with one #).

~~Alternative:  
A → 0A | 1A | A0 | A1 | #~~

60. {all strings in  $\{0,1\}^*$  st. # 0's is even

or # 1's is divisible by 3  
or length is divisible by 5 }

$$\text{Let } M = (Q, \Sigma, s, \delta, A)$$

$$\Sigma = \{0,1\}$$

$$Q = \{ (i,j,k) : \begin{array}{l} i \in \{0,1\} \\ j \in \{0,1,2\} \\ k \in \{0,1,2,3,4\} \end{array} \}$$
$$= \{0,1\} \times \{0,1,2\} \times \{0,1,2,3,4\}$$

$$s = (0,0,0)$$

$A = \{ (i,j,k) : i=0 \text{ or } j=0 \text{ or } k=0 \}$

$$\text{~~A = \{ (0,0,0) \}}~~$$

$$\delta((i,j,k), 0) = ((i+1) \bmod 2, j, (k+1) \bmod 5)$$

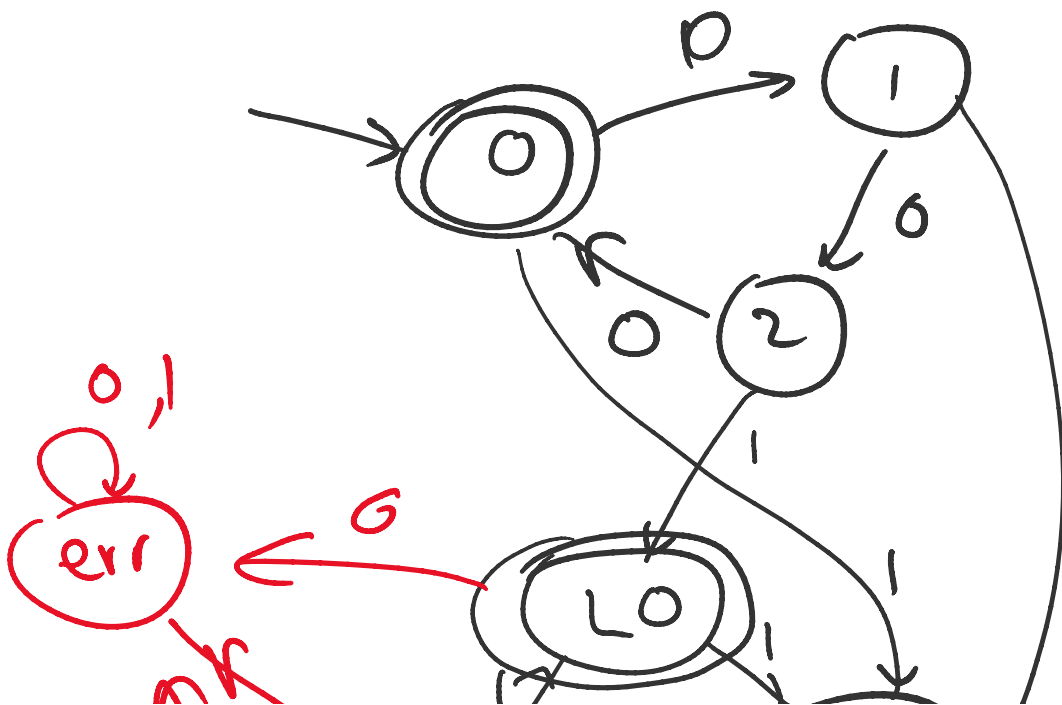
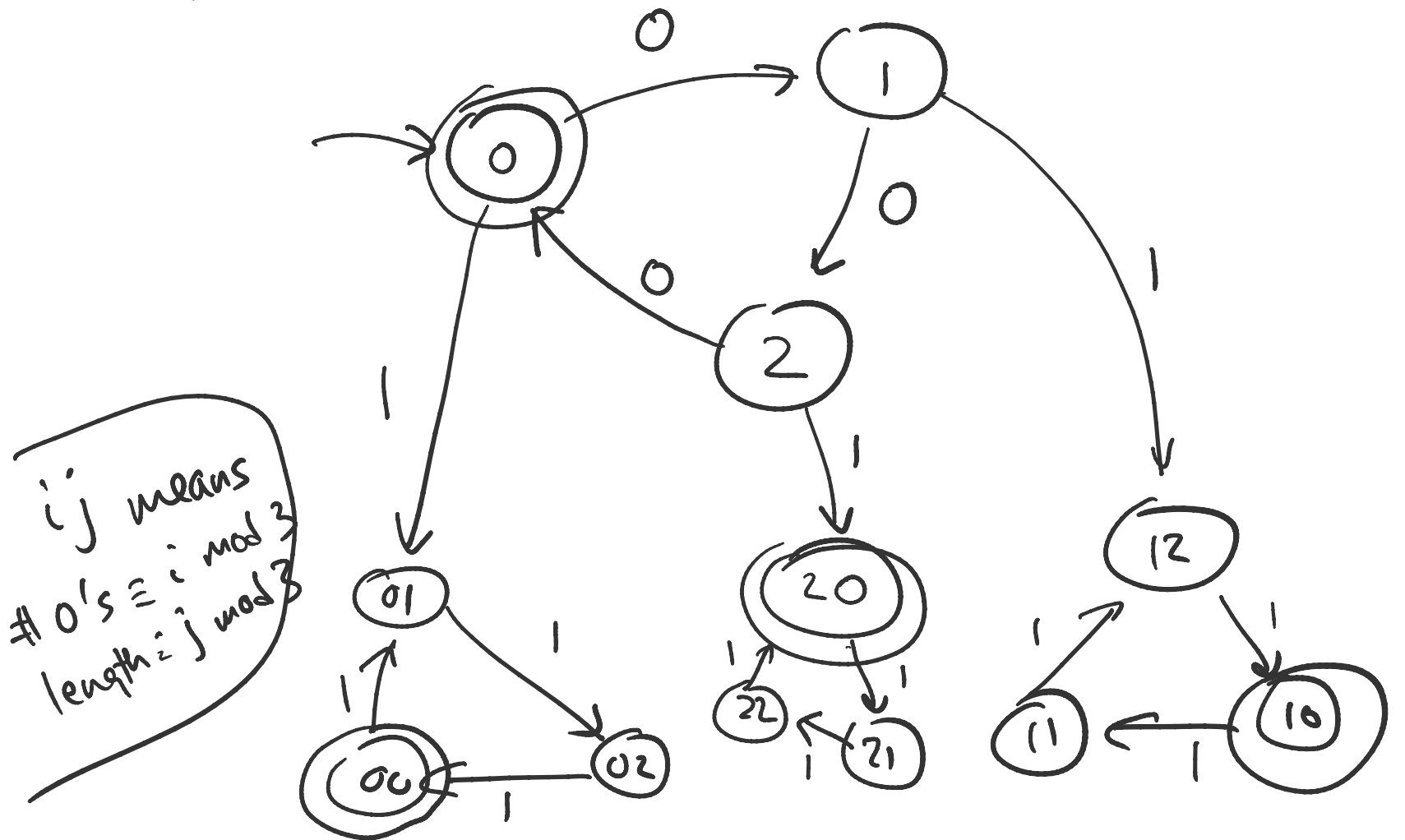
$$\delta((i,j,k), 1) = (i, (j+1) \bmod 3, (k+1) \bmod 5)$$

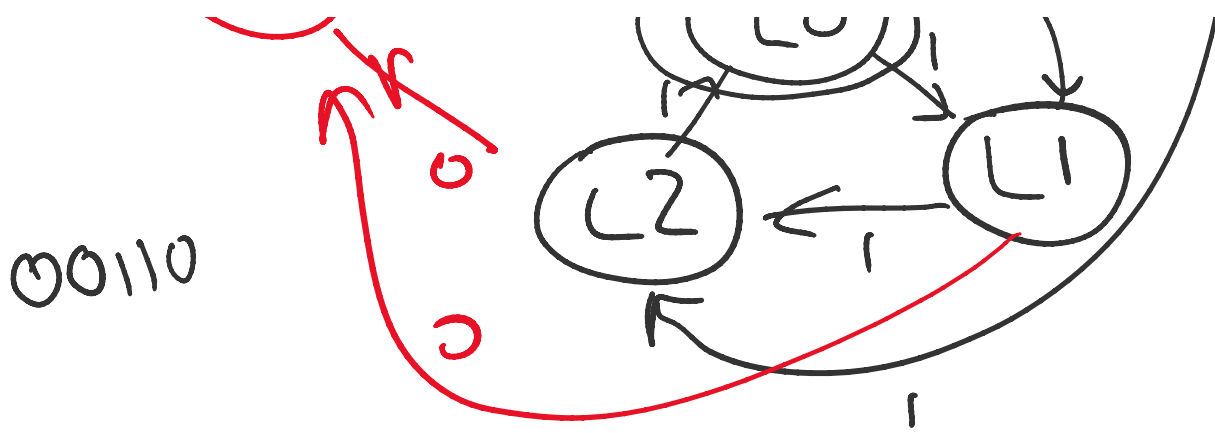
$0 \dots (k+1) \pmod{5}$

□

26. all strings in  $0^*1^*$  with length div by 3

DFA:





reg expr:  $(000)^* (111)^*$

+  $(000)^* 011 (111)^*$

+  $00(000)^* 1(111)^*$

$$(000)^* 0 \equiv 0(000)^*$$

### F'18 Midterm Problem 6:

Given CFGs  $G_i = (V_i, T, S_i, P_i)$   $i=1,2,3$ ,  
describe  
new CFG  
for  $L_1$ ,

$G' = (V', T, S', P')$  for  $L_1 \cup L_2 L_3^*$

$$P' = \left\{ \begin{array}{l} S' \rightarrow \underline{S_1} \mid S_2 S'' \\ \underline{S''} \rightarrow \underline{S_3 S''} \mid \epsilon \end{array} \right\}$$

$$\cup P_1 \cup P_2 \cup P_3.$$

$$V' = V_1 \cup V_2 \cup V_3 \cup \{S', S''\}$$

(disjoint unions)