

Kleene's Thm ('51)

L is regular $\Leftrightarrow L$ is accepted by some DFA.


Corollary

If L is regular, then \bar{L} is regular.

If L_1, L_2 are regular,
then $L_1 \cap L_2$ is regular.

If L is accepted by some DFA,
then so is L^R .

Proof Outline:

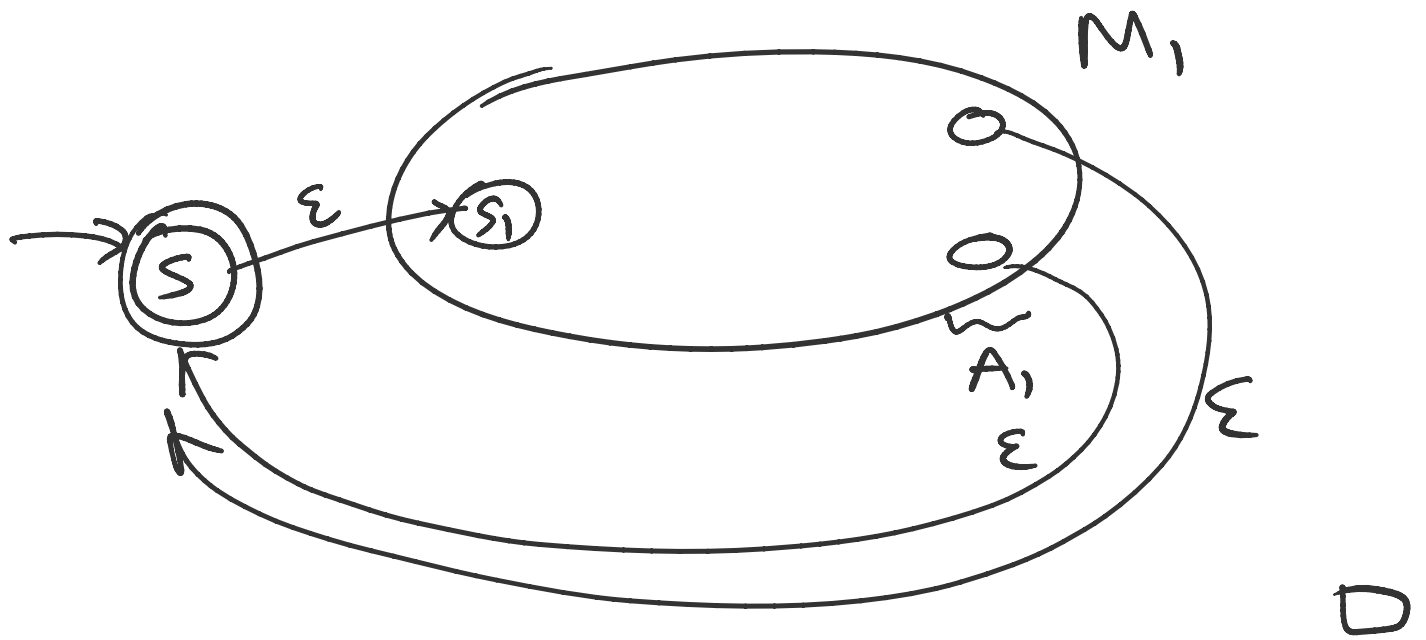
regular \rightarrow NFA \rightarrow DFA


Regular \rightarrow NFA

Lemma If L is regular,
then L is accepted by some NFA.

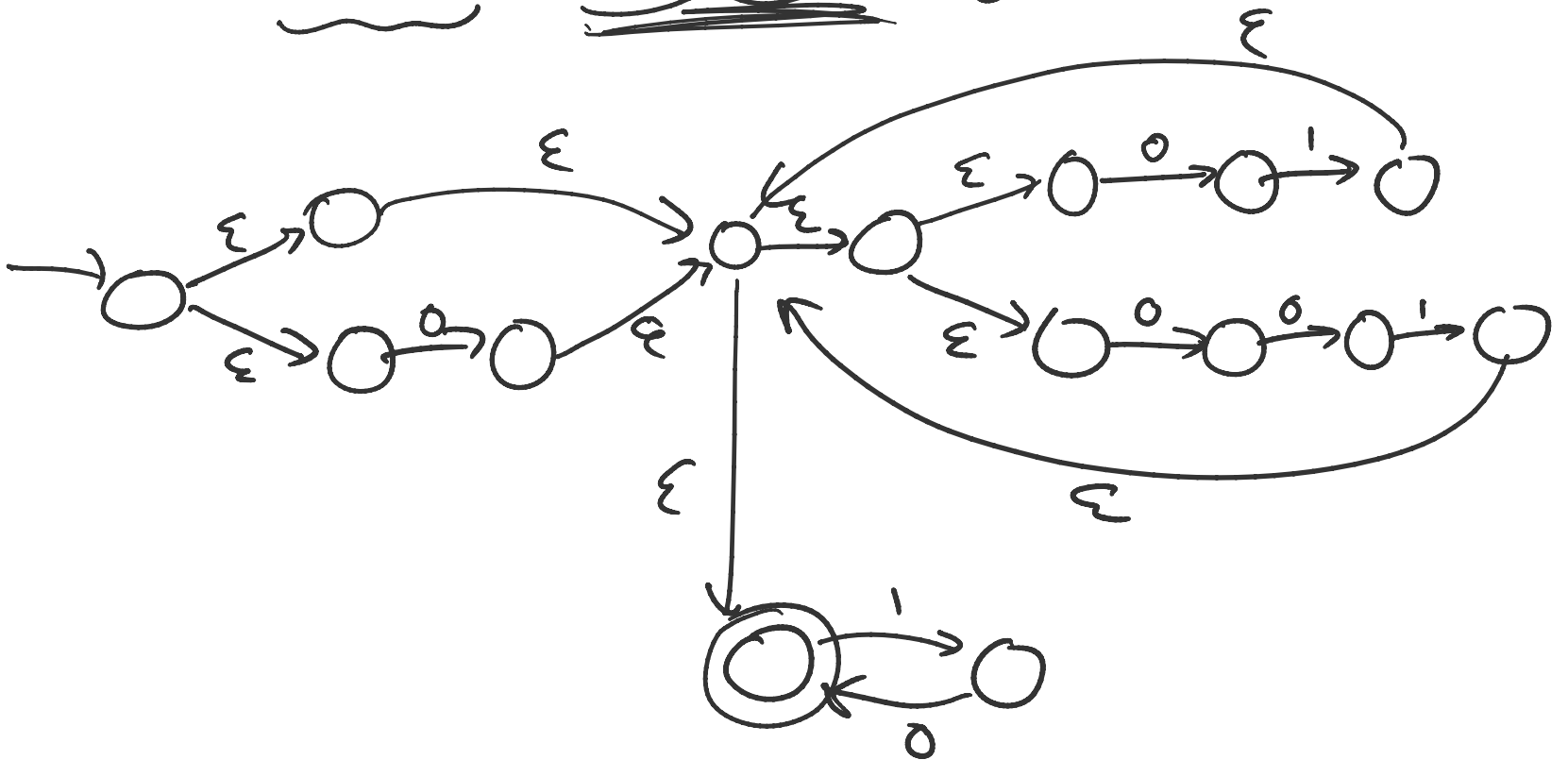
Pf Sketch:

method - recursion (i.e. induction)



Ex

$(\epsilon + 0)$ $(01 + 001)^*$ $(10)^*$



NFA \rightarrow DFA

Lemma If L is accepted by NFA M , then L is accepted by some DFA M' .

Pf: method - Subset construction (i.e. power set)

idea - at any time,

idea - at any time, remember subset of all states you can be in

Given NFA $M = (\Sigma, Q, s, A, \delta)$

where $\delta: Q \times (\Sigma \cup \{\epsilon\}) \rightarrow \mathcal{P}(Q)$,

Construct DFA $M' = (\Sigma, Q', s', A', \delta')$

where $\delta': Q' \times \Sigma \rightarrow Q'$:

$$Q' = \mathcal{P}(Q)$$

$$(|Q'| = 2^{|Q|})$$

$$s' = \epsilon\text{-reach}(s)$$

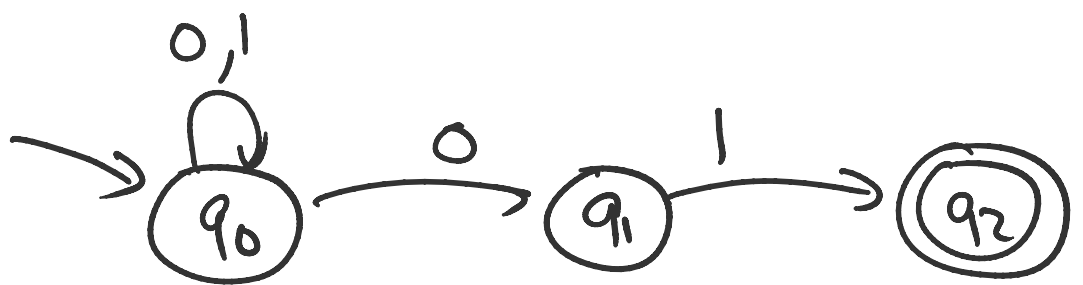
$$A' = \{S \in Q' \mid S \cap A \neq \emptyset\}$$

$$\forall S \in Q', a \in \Sigma, \delta'(S, a) = \bigcup_{q \in S} \delta^*(q, a)$$

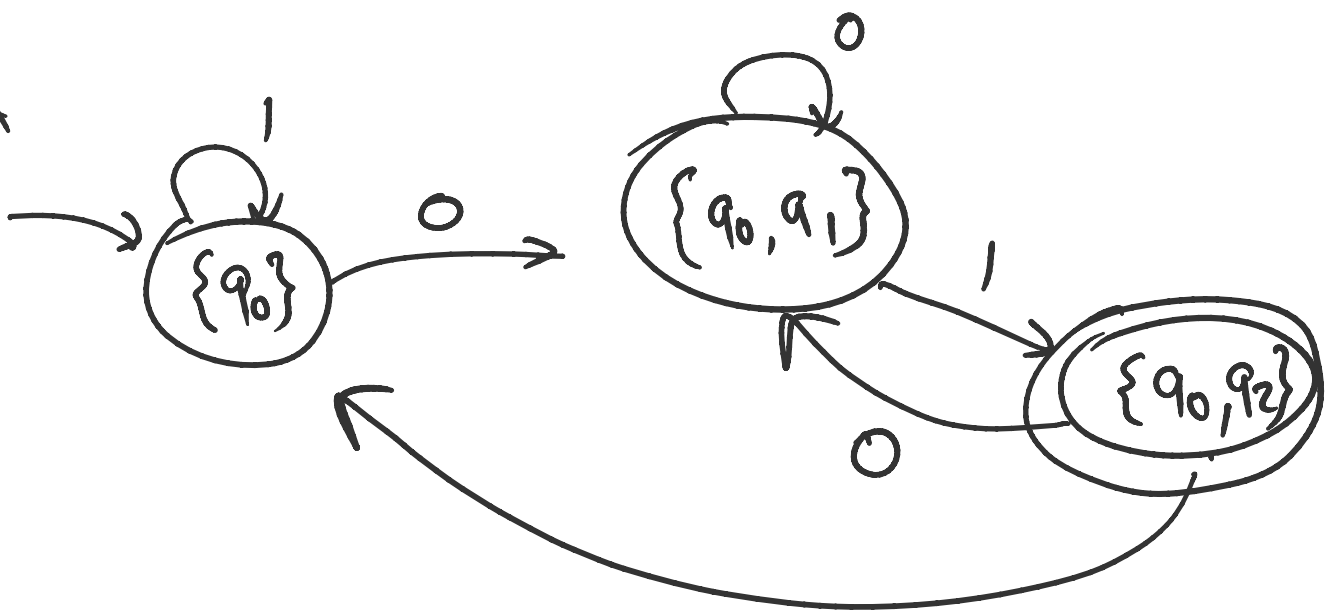
□

Ex

NFA



DFA



DFA \rightarrow Regular

Lemma If L is accepted by DFA M , then L is regular.

Pf sketch: method - state elimination

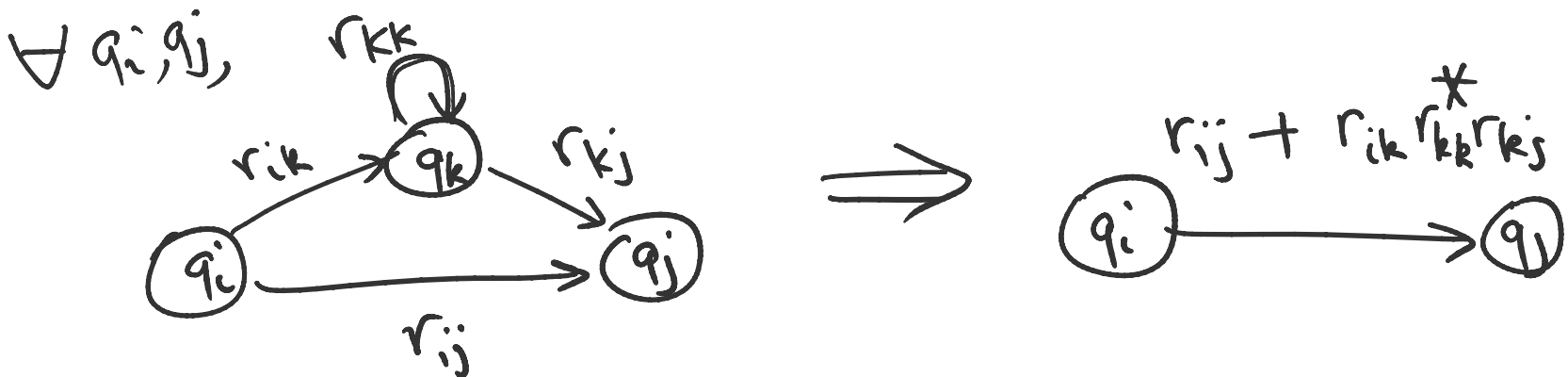
Given DFA $M = (\Sigma, Q, s, A, \delta)$,

add new start state s' , new final state f'

for each $q_k \in Q - \{s', f'\}$,

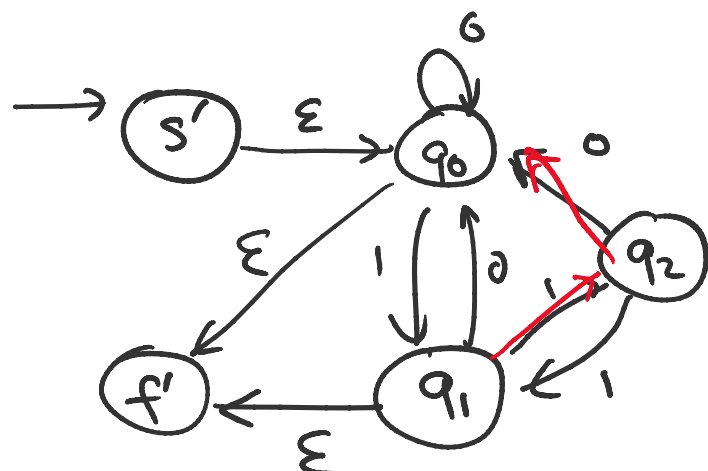
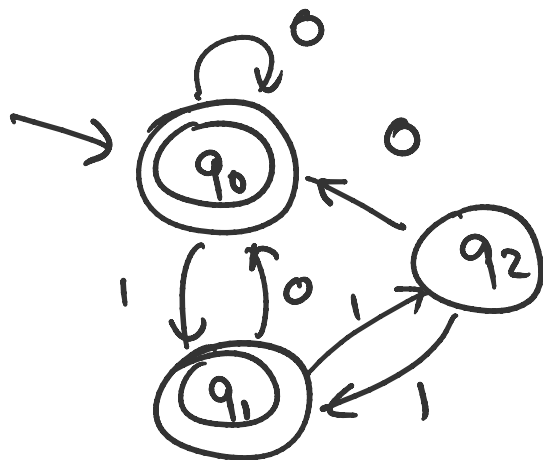
remove q_k &

relabel transitions by this rule:

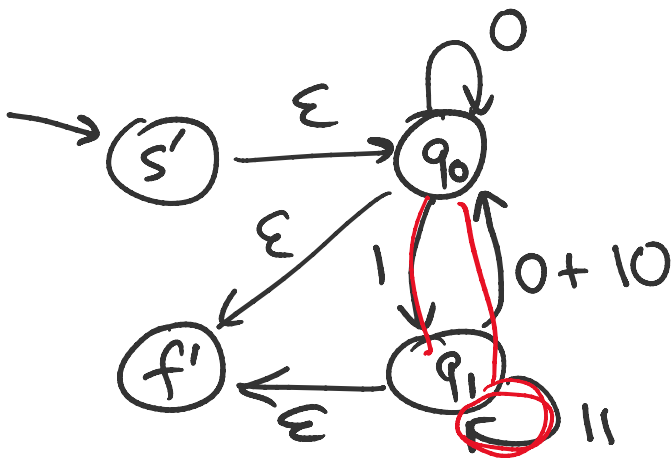


return label from s' to f' . \square

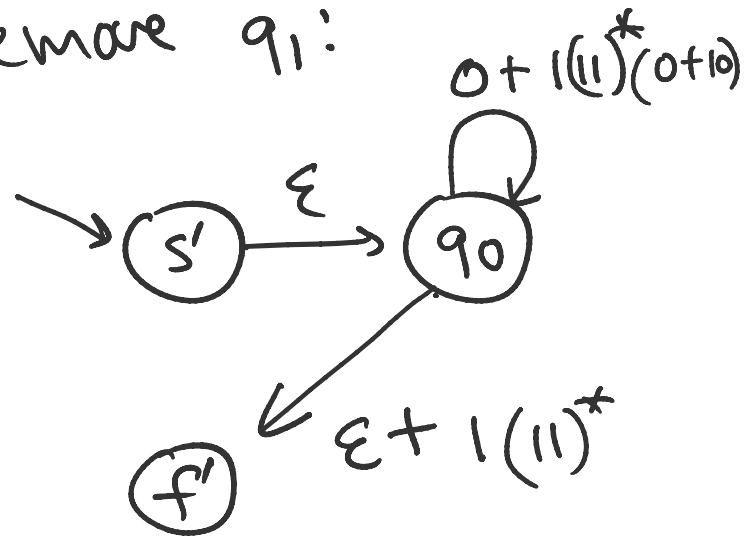
Ex



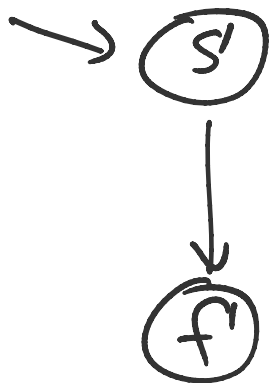
remove q_2 :



remove q_1 :



remove q_0 :



$$(0 + 1(11)^*(0+10))^* (\epsilon + 1(11)^*)$$