

Give regular expressions for each of the following languages over the alphabet  $\{0, 1\}$ .

1. All strings containing the substring  $000$ .

**Solution:**  $(0 + 1)^*000(0 + 1)^*$  ■

2. All strings *not* containing the substring  $000$ .

**Solution:**  $(1 + 01 + 001)^*(\epsilon + 0 + 00)$  ■

**Solution:**  $(\epsilon + 0 + 00)(1(\epsilon + 0 + 00))^*$  ■

3. All strings in which every run of  $0$ s has length at least 3.

**Solution:**  $(1 + 0000^*)^*$  ■

**Solution:**  $(\epsilon + 1)((\epsilon + 0000^*)1)^*(\epsilon + 0000^*)$  ■

4. All strings in which every substring  $000$  appears after every  $1$ .

**Solution:**  $(1 + 01 + 001)^*0^*$  ■

5. All strings containing at least three  $0$ s.

**Solution:**  $(0 + 1)^*0(0 + 1)^*0(0 + 1)^*0(0 + 1)^*$  ■

**Solution (clever):**  $1^*01^*01^*0(0 + 1)^*$  or  $(0 + 1)^*01^*01^*01^*$  ■

6. Every string except  $000$ . [Hint: Don't try to be clever.]

**Solution:** Every string  $w \neq 000$  satisfies one of three conditions: Either  $|w| < 3$ , or  $|w| = 3$  and  $w \neq 000$ , or  $|w| > 3$ . The first two cases include only a finite number of strings, so we just list them explicitly. The last case includes *all* strings of length at least 4.

$$\begin{aligned} & \epsilon + 0 + 1 + 00 + 01 + 10 + 11 \\ & + 001 + 010 + 011 + 100 + 101 + 110 + 111 \\ & + (1 + 0)(1 + 0)(1 + 0)(1 + 0)(1 + 0)^* \end{aligned}$$

**Solution (clever):**  $\epsilon + 0 + 00 + (1 + 01 + 001 + 000(1 + 0))(1 + 0)^*$  ■

7. All strings  $w$  such that *in every prefix of  $w$* , the number of 0s and 1s differ by at most 1.

**Solution:** Equivalently, strings that alternate between 0s and 1s:  $(01+10)^*(\epsilon+0+1)$  ■

\*8. All strings containing at least two 0s and at least one 1.

**Solution:** There are three possibilities for how such a string can begin:

- Start with 00, then any number of 0s, then 1, then anything.
- Start with 01, then any number of 1s, then 0, then anything.
- Start with 1, then a substring with exactly two 0s, then anything.

All together:  $000^*1(0+1)^* + 011^*0(0+1)^* + 11^*01^*0(0+1)^*$

Or equivalently:  $(000^*1 + 011^*0 + 11^*01^*0)(0+1)^*$  ■

**Solution:** There are three possibilities for how the three required symbols are ordered:

- Contains a 1 before two 0s:  $(0+1)^*1(0+1)^*0(0+1)^*0(0+1)^*$
- Contains a 1 between two 0s:  $(0+1)^*0(0+1)^*1(0+1)^*0(0+1)^*$
- Contains a 1 after two 0s:  $(0+1)^*0(0+1)^*0(0+1)^*1(0+1)^*$

So putting these cases together, we get the following:

$$\begin{aligned} & (0+1)^*1(0+1)^*0(0+1)^*0(0+1)^* \\ & + (0+1)^*0(0+1)^*1(0+1)^*0(0+1)^* \\ & + (0+1)^*0(0+1)^*0(0+1)^*1(0+1)^* \end{aligned}$$
 ■

**Solution (clever):**  $(0+1)^*(101^*0 + 010 + 01^*01)(0+1)^*$  ■

\*9. All strings  $w$  such that *in every prefix of  $w$* , the number of 0s and 1s differ by at most 2.

**Solution:**  $(0(01)^*1 + 1(10)^*0)^* \cdot (\epsilon + 0(01)^*(0+\epsilon) + 1(10)^*(1+\epsilon))$  ■

- ★10. All strings in which the substring  $000$  appears an even number of times.  
(For example,  $0001000$  and  $0000$  are in this language, but  $00000$  is not.)

**Solution:** Every string in  $\{0, 1\}^*$  alternates between (possibly empty) blocks of  $0$ s and individual  $1$ s; that is,  $\{0, 1\}^* = (0^*1)^*0^*$ . Trivially, every  $000$  substring is contained in some block of  $0$ s. Our strategy is to consider which blocks of  $0$ s contain an even or odd number of  $000$  substrings.

Let  $X$  denote the set of all strings in  $0^*$  with an even number of  $000$  substrings. We easily observe that  $X = \{0^n \mid n = 1 \text{ or } n \text{ is even}\} = 0 + (00)^*$ .

Let  $Y$  denote the set of all strings in  $0^*$  with an *odd* number of  $000$  substrings. We easily observe that  $Y = \{0^n \mid n > 1 \text{ and } n \text{ is odd}\} = 000(00)^*$ .

We immediately have  $0^* = X + Y$  and therefore  $\{0, 1\}^* = ((X + Y)1)^*(X + Y)$ .

Finally, let  $L$  denote the set of all strings in  $\{0, 1\}^*$  with an even number of  $000$  substrings. A string  $w \in \{0, 1\}^*$  is in  $L$  if and only if an even number of blocks of  $0$ s in  $w$  are in  $Y$ ; the remaining blocks of  $0$ s are all in  $X$ .

$$L = ((X1)^*Y1 \cdot (X1)^*Y1)^*(X1)^*X$$

Plugging in the expressions for  $X$  and  $Y$  gives us the following regular expression for  $L$ :

$$\left( \left( (0 + (00)^*)1 \right)^* \cdot 000(00)^*1 \cdot \left( (0 + (00)^*)1 \right)^* \cdot 000(00)^*1 \right)^* \cdot \left( (0 + (00)^*)1 \right)^* \cdot (0 + (00)^*)$$

Whew! ■