Algorithms \& Models of Computation

## CS/ECE 374, Fall 2020

### 24.4.2

The consistency of execution

## The variables of $\varphi$

## Variables:

$\left\langle\boldsymbol{q}_{\boldsymbol{j}}, \boldsymbol{b}_{j}, \boldsymbol{q}_{\boldsymbol{j}}^{\prime}, \boldsymbol{b}_{\boldsymbol{j}}^{\prime}, \boldsymbol{d}_{\boldsymbol{j}}\right\rangle: \boldsymbol{j}$ th instruction of $\boldsymbol{M}$ $\boldsymbol{I}(\boldsymbol{j}, \boldsymbol{i})$ : Instruction $\boldsymbol{j}$ was issued at time $\boldsymbol{i}$. $\boldsymbol{H}(\boldsymbol{h}, \boldsymbol{i})$ : The head is at location $\boldsymbol{h}$ at time $\boldsymbol{i}$. $\boldsymbol{T}(\boldsymbol{c}, \boldsymbol{h}, \boldsymbol{i})$ : The tape at location $\boldsymbol{h}$ at time $\boldsymbol{i}$ stored the character $\boldsymbol{c}$.

## $\varphi_{1}$ : The input is encoded correctly

$\varphi_{1}$ asserts (is true eff) the variables are set $\mathrm{T} / \mathrm{F}$ indicating that $\boldsymbol{M}$ starts in state $\boldsymbol{q}_{0}$ at time $\mathbf{0}$ with tape contents containing $\boldsymbol{x}$ followed by blanks. Let $\boldsymbol{x}=x_{1} x_{2} \ldots x_{n}$

$$
\begin{aligned}
& \varphi_{1}=\boldsymbol{S}\left(\boldsymbol{q}_{\mathbf{0}}, \mathbf{0}\right) \quad / / \text { state at time } \mathbf{0} \text { is } \boldsymbol{q}_{\mathbf{0}} \\
& \bigwedge_{h=1} \boldsymbol{T}\left(x_{h}, \boldsymbol{h}, \mathbf{0}\right) \quad / / \text { at time } \mathbf{0} \text { cells } \mathbf{1} \text { to } \boldsymbol{n} \text { have value } x_{1} \text { to } x_{n} \\
& \wedge \bigwedge_{h=n+1}^{p(n)} T(\sqcup, h, 0) \\
& \text { // all remaining cells are blank } \\
& \wedge \boldsymbol{H}(\mathbf{1}, \mathbf{0}) \quad / / \text { The head is at time } \mathbf{0} \text { at start of tape }
\end{aligned}
$$

## $\varphi_{2}: M$ is in exactly one state at any point in time

$\varphi_{2}$ asserts $M$ in exactly one state at any time $\boldsymbol{i}$ :

$$
\varphi_{2}=\bigwedge_{i=0}^{p(|x|)}\left(\oplus\left(S\left(q_{0}, i\right), S\left(q_{1}, i\right), \ldots, S\left(q_{|Q|}, i\right)\right)\right)
$$

## Variables:

$\left\langle\boldsymbol{a}_{j}, \boldsymbol{b}_{\boldsymbol{j}}, \boldsymbol{q}_{\boldsymbol{j}}^{\prime}, \boldsymbol{b}_{j}^{\prime}, \boldsymbol{d}_{\boldsymbol{j}}\right\rangle: j$ th instruction of $\boldsymbol{M}$
$\boldsymbol{I}(\boldsymbol{j}, \boldsymbol{i})$ : Instruction $\boldsymbol{j}$ was issued at time $\boldsymbol{i}$.
$\boldsymbol{H}(\boldsymbol{h}, \boldsymbol{i})$ : The head is at location $\boldsymbol{h}$ at time $\boldsymbol{i}$.
$\boldsymbol{T}(\boldsymbol{c}, \boldsymbol{h}, \boldsymbol{i})$ : The tape at location $\boldsymbol{h}$ at time $\boldsymbol{i}$ stored the character $\boldsymbol{c}$.

## $\varphi_{3}:$ Each tape cell holds a unique symbol at any time

$\varphi_{3}$ asserts that each tape cell holds a unique symbol at any given time.

$$
\varphi_{3}=\bigwedge_{i=0}^{p(|x| \mid)} \bigwedge_{h=1}^{p(|x|)} \oplus\left(T\left(b_{1}, \boldsymbol{h}, \boldsymbol{i}\right), T\left(\boldsymbol{b}_{2}, \boldsymbol{h}, \boldsymbol{i}\right), \ldots, T\left(\boldsymbol{b}_{|\Gamma|}, \boldsymbol{h}, \boldsymbol{i}\right)\right)
$$

For each time $\boldsymbol{i}$ and for each cell position $\boldsymbol{h}$ exactly one symbol $\boldsymbol{b} \in \Gamma$ at cell position $\boldsymbol{h}$ at time $\boldsymbol{i}$

## Variables:

$\left\langle\boldsymbol{q}_{\boldsymbol{j}}, \boldsymbol{b}_{\boldsymbol{j}}, \boldsymbol{q}_{\boldsymbol{j}}^{\prime}, \boldsymbol{b}_{\boldsymbol{j}}^{\prime}, \boldsymbol{d}_{\boldsymbol{j}}\right\rangle: \boldsymbol{j}$ th instruction of $\boldsymbol{M}$
$\boldsymbol{I}(\boldsymbol{j}, \boldsymbol{i})$ : Instruction $\boldsymbol{j}$ was issued at time $\boldsymbol{i}$.
$\boldsymbol{H}(\boldsymbol{h}, \boldsymbol{i})$ : The head is at location $\boldsymbol{h}$ at time $\boldsymbol{i}$.
$\boldsymbol{T}(\boldsymbol{c}, \boldsymbol{h}, \boldsymbol{i})$ : The tape at location $\boldsymbol{h}$ at time $\boldsymbol{i}$ stored the character $\boldsymbol{c}$.

## $\varphi_{4}$ : tape head of $\boldsymbol{M}$ is in exactly one position at any time $\boldsymbol{i}$

 $\varphi_{4}$ asserts that the read/write head of $\boldsymbol{M}$ is in exactly one position at any time $\boldsymbol{i}$$$
\varphi_{4}=\bigwedge_{i=0}^{p(|x|)}(\oplus(\boldsymbol{H}(\mathbf{1}, \boldsymbol{i}), \boldsymbol{H}(2, \boldsymbol{i}), \ldots, \boldsymbol{H}(\boldsymbol{p}(|x|), \boldsymbol{i})))
$$

## Variables:

$\left\langle\boldsymbol{q}_{j}, \boldsymbol{b}_{\boldsymbol{j}}, \boldsymbol{q}_{\boldsymbol{j}}^{\prime}, \boldsymbol{b}_{\boldsymbol{j}}^{\prime}, \boldsymbol{d}_{\boldsymbol{j}}\right\rangle: \boldsymbol{j}$ th instruction of $\boldsymbol{M}$
$\boldsymbol{I}(\boldsymbol{j}, \boldsymbol{i})$ : Instruction $\boldsymbol{j}$ was issued at time $\boldsymbol{i}$.
$\boldsymbol{H}(\boldsymbol{h}, \boldsymbol{i})$ : The head is at location $\boldsymbol{h}$ at time $\boldsymbol{i}$.
$\boldsymbol{T}(\boldsymbol{c}, \boldsymbol{h}, \boldsymbol{i})$ : The tape at location $\boldsymbol{h}$ at time $\boldsymbol{i}$ stored the character $\boldsymbol{c}$.

## $\varphi_{5}: M$ accepts the input

$\varphi_{5}$ asserts that $M$ accepts

- Let $\boldsymbol{q}_{a}$ be unique accept state of $\boldsymbol{M}$
- without loss of generality assume $\boldsymbol{M}$ runs all $\boldsymbol{p}(|\boldsymbol{x}|)$ steps

$$
\varphi_{5}=\boldsymbol{S}\left(\boldsymbol{q}_{a}, \boldsymbol{p}(|\boldsymbol{x}|)\right)
$$

State at time $\boldsymbol{p}(|\boldsymbol{x}|)$ is $\boldsymbol{q}_{\boldsymbol{a}}$ the accept state.
If we don't want to make assumption of running for all steps

$$
\varphi_{5}=\bigvee_{i=1}^{p(|x|)} S\left(q_{a}, i\right)
$$

which means $M$ enters accepts state at some time.

## $\varphi_{6}: M$ executes a unique instruction at each time

$\varphi_{6}$ asserts that $M$ executes a unique instruction at each time

$$
\varphi_{6}=\bigwedge_{i=0}^{p(|x|)} \oplus(I(1, i), I(2, i), \ldots, I(m, i))
$$

where $\boldsymbol{m}$ is max instruction number.

## Variables:

$\left\langle\boldsymbol{q}_{\boldsymbol{j}}, \boldsymbol{b}_{j}, \boldsymbol{q}_{\boldsymbol{j}}^{\prime}, \boldsymbol{b}_{\boldsymbol{j}}^{\prime}, \boldsymbol{d}_{\boldsymbol{j}}\right\rangle: \boldsymbol{j}$ th instruction of $\boldsymbol{M}$
$\boldsymbol{I}(\boldsymbol{j}, \boldsymbol{i})$ : Instruction $\boldsymbol{j}$ was issued at time $\boldsymbol{i}$.
$\boldsymbol{H}(\boldsymbol{h}, \boldsymbol{i})$ : The head is at location $\boldsymbol{h}$ at time $\boldsymbol{i}$.
$\boldsymbol{T}(\boldsymbol{c}, \boldsymbol{h}, \boldsymbol{i})$ : The tape at location $\boldsymbol{h}$ at time $\boldsymbol{i}$ stored the character $\boldsymbol{c}$.

## $\varphi_{7}$ : Tape changes only because of the head writing something

$\varphi_{7}$ ensures that variables don't allow tape to change from one moment to next if the read/write head was not there.
"If head is not at position $\boldsymbol{h}$ at time $\boldsymbol{i}$ then at time $\boldsymbol{i}+\mathbf{1}$ the symbol at cell $\boldsymbol{h}$ must be unchanged"

$$
\varphi_{7}=\bigwedge_{i} \bigwedge_{h} \bigwedge_{b \neq c}(\overline{H(h, i)} \Rightarrow \overline{T(b, h, i) \bigwedge T(c, h, i+1)})
$$

since $\boldsymbol{A} \Rightarrow \boldsymbol{B}$ is same as $\neg \boldsymbol{A} \vee \boldsymbol{B}$, rewrite above in CNF form

$$
\varphi_{7}=\bigwedge_{i} \bigwedge_{h} \bigwedge_{b \neq c}(H(h, i) \vee \neg T(b, h, i) \vee \neg T(c, h, i+1))
$$

## $\varphi_{8}:$ Transitions are done from correct states

$\boldsymbol{j}$ th instruction of $\boldsymbol{M}:<\boldsymbol{q}_{j}, \boldsymbol{b}_{\boldsymbol{j}}, \boldsymbol{q}_{\boldsymbol{j}}^{\prime}, \boldsymbol{b}_{\boldsymbol{j}}^{\prime}, \boldsymbol{d}_{\boldsymbol{j}}>$

$$
\varphi_{8}=\bigwedge_{i} \bigwedge_{j}\left(I(j, i) \Rightarrow S\left(q_{j}, i\right)\right)
$$

If instruction $\boldsymbol{j}$ is executed at time $\boldsymbol{i}$ then state at time $\boldsymbol{i}$ must be $\boldsymbol{q}_{\boldsymbol{j}}$.

## Variables:

$\left\langle\boldsymbol{q}_{j}, \boldsymbol{b}_{j}, \boldsymbol{q}_{\boldsymbol{j}}^{\prime}, \boldsymbol{b}_{\boldsymbol{j}}^{\prime}, \boldsymbol{d}_{j}\right\rangle: \boldsymbol{j}$ th instruction of $\boldsymbol{M}$
$\boldsymbol{I}(\boldsymbol{j}, \boldsymbol{i})$ : Instruction $\boldsymbol{j}$ was issued at time $\boldsymbol{i}$.
$\boldsymbol{H}(\boldsymbol{h}, \boldsymbol{i})$ : The head is at location $\boldsymbol{h}$ at time $\boldsymbol{i}$.
$\boldsymbol{T}(\boldsymbol{c}, \boldsymbol{h}, \boldsymbol{i})$ : The tape at location $\boldsymbol{h}$ at time $\boldsymbol{i}$ stored the character $\boldsymbol{c}$.

## $\varphi_{9}:$ Transitions are done into correct state

$j$ th instruction of $\boldsymbol{M}:<\boldsymbol{q}_{j}, \boldsymbol{b}_{j}, \boldsymbol{q}_{\boldsymbol{j}}^{\prime}, \boldsymbol{b}_{j}^{\prime}, \boldsymbol{d}_{j}>$

$$
\varphi_{9}=\bigwedge_{i} \bigwedge_{j}\left(I(j, i) \Rightarrow S\left(q_{j}^{\prime}, i+1\right)\right)
$$

If instruction $\boldsymbol{j}$ was performed at time $\boldsymbol{i}$, then state at time $\boldsymbol{i}+\mathbf{1}$ must be $\boldsymbol{q}_{\boldsymbol{j}}^{\prime}$.

## Variables:

$\left\langle\boldsymbol{q}_{j}, \boldsymbol{b}_{\boldsymbol{j}}, \boldsymbol{a}_{\boldsymbol{j}}^{\prime}, \boldsymbol{b}_{j}^{\prime}, \boldsymbol{d}_{\boldsymbol{j}}\right\rangle: \boldsymbol{j}$ th instruction of $\boldsymbol{M}$
$\boldsymbol{I}(\boldsymbol{j}, \boldsymbol{i})$ : Instruction $\boldsymbol{j}$ was issued at time $\boldsymbol{i}$.
$\boldsymbol{H}(\boldsymbol{h}, \boldsymbol{i})$ : The head is at location $\boldsymbol{h}$ at time $\boldsymbol{i}$.
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## $\varphi_{10}:$ The character written on tape that triggered an

 instruction, is the correct one$$
\varphi_{10}=\bigwedge_{i} \bigwedge_{h} \bigwedge_{j}\left[(I(j, i) \bigwedge H(h, i)) \Rightarrow T\left(\boldsymbol{b}_{j}, \boldsymbol{h}, \boldsymbol{i}\right)\right]
$$

If instruction $\boldsymbol{j}$ was executed at time $\boldsymbol{i}$ and head was at position $\boldsymbol{h}$, then cell $\boldsymbol{h}$ has the symbol needed to issue instruction $\boldsymbol{j}$ is written under the head location on the tape.

## Variables:

$\left\langle\boldsymbol{q}_{j}, \boldsymbol{b}_{j}, \boldsymbol{q}_{\boldsymbol{j}}^{\prime}, \boldsymbol{b}_{j}^{\prime}, \boldsymbol{d}_{\boldsymbol{j}}\right\rangle: \boldsymbol{j}$ th instruction of $\boldsymbol{M}$
$\boldsymbol{I}(\boldsymbol{j}, \boldsymbol{i})$ : Instruction $\boldsymbol{j}$ was issued at time $\boldsymbol{i}$.
$\boldsymbol{H}(\boldsymbol{h}, \boldsymbol{i})$ : The head is at location $\boldsymbol{h}$ at time $\boldsymbol{i}$.
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## $\varphi_{11}$ : The correct symbol was written to the tape at time $\boldsymbol{i}$

$$
\varphi_{11}=\bigwedge_{i} \bigwedge_{j} \bigwedge_{h}\left[(I(j, i) \wedge H(h, i)) \Rightarrow T\left(b_{j}^{\prime}, h, i+1\right)\right]
$$

If instruction $\boldsymbol{j}$ was executed time $\boldsymbol{i}$ with head at $\boldsymbol{h}$, then at next time step symbol $\boldsymbol{b}_{\boldsymbol{j}}^{\prime}$ was written in position $\boldsymbol{h}$

## Variables:

$\left\langle\boldsymbol{q}_{j}, \boldsymbol{b}_{\boldsymbol{j}}, \boldsymbol{q}_{\boldsymbol{j}}^{\prime}, \boldsymbol{b}_{\boldsymbol{j}}^{\prime}, \boldsymbol{d}_{\boldsymbol{j}}\right\rangle: \boldsymbol{j}$ th instruction of $\boldsymbol{M}$
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$\boldsymbol{H}(\boldsymbol{h}, \boldsymbol{i})$ : The head is at location $\boldsymbol{h}$ at time $\boldsymbol{i}$.
$\boldsymbol{T}(\boldsymbol{c}, \boldsymbol{h}, \boldsymbol{i})$ : The tape at location $\boldsymbol{h}$ at time $\boldsymbol{i}$ stored the character $\boldsymbol{c}$.

## $\varphi_{12}:$ Head was moved in the right direction at time $\boldsymbol{i}$

$$
\varphi_{12}=\bigwedge_{i} \bigwedge_{j} \bigwedge_{h}\left[(I(\boldsymbol{j}, \boldsymbol{i}) \wedge \boldsymbol{H}(\boldsymbol{h}, \boldsymbol{i})) \Rightarrow \boldsymbol{H}\left(\boldsymbol{h}+\boldsymbol{d}_{j}, \boldsymbol{i}+\mathbf{1}\right)\right]
$$

The head is moved properly according to instr $\boldsymbol{j}$.

## Variables:

$\left\langle\boldsymbol{q}_{\boldsymbol{j}}, \boldsymbol{b}_{\boldsymbol{j}}, \boldsymbol{q}_{\boldsymbol{j}}^{\prime}, \boldsymbol{b}_{j}^{\prime}, \boldsymbol{d}_{\boldsymbol{j}}\right\rangle: \boldsymbol{j}$ th instruction of $\boldsymbol{M}$
$\boldsymbol{I}(\boldsymbol{j}, \boldsymbol{i})$ : Instruction $\boldsymbol{j}$ was issued at time $\boldsymbol{i}$.
$\boldsymbol{H}(\boldsymbol{h}, \boldsymbol{i})$ : The head is at location $\boldsymbol{h}$ at time $\boldsymbol{i}$.
$\boldsymbol{T}(\boldsymbol{c}, \boldsymbol{h}, \boldsymbol{i})$ : The tape at location $\boldsymbol{h}$ at time $\boldsymbol{i}$ stored the character $\boldsymbol{c}$.

## THE END

(for now)

